

# PAR: A New Representation Scheme for Rotational Parts \*

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## **Abstract**

A new representation scheme PAR (Principal Axis Representation) for rotational parts is proposed as an internal representation scheme for Constructive Solid Geometry (CSG). The key idea of PAR is to represent an object by its principal axis and a set of boundary curves. Based on a mathematical framework, an algorithm is designed to convert a CSG tree into a PAR, which represents the same object as the CSG tree does but is in an evaluated form. Geometrical properties of parts can then be computed more directly and efficiently from this evaluated representation than from the original CSG tree. In addition to its computational efficiency, the PAR is a *unique* representation scheme.

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## 1 Introduction

Solid modeling techniques have been considered as one of the keys to the integration of computer-aided design and manufacturing (CAD/CAM)[Bee82]. Two of the prevailing geometric schemes in solid modeling are *Boundary Representation* (B-rep) scheme and *Constructive Solid Geometry* (CSG) scheme[ReV82]. B-rep represents an object by segmenting its boundary into a finite number of bounded subsets usually called faces or patches, and representing each face by (for example) its bounding edges and vertices. CSG scheme represents objects as constructions or combinations, via the regularized set operators, of solid components. While B-rep in general is in an evaluated form, which explicitly shows what and where each geometric entity is, CSG is not. The much more concise CSG is easy to construct and to check the validity of the object being constructed. The tradeoff between the evaluatedness and conciseness seems to be clear but the choice between B-rep and CSG has never been easy.

As a means to integrate design and manufacturing, solid modeling techniques not only display objects but also compute important properties of objects. A more comprehensive understanding of the geometric representation by the computer becomes so important that shape features can then be extracted, parts classified, processes planned, and etc.. Several interesting researches have been done along the direction of enhancing machine's understanding of part geometric representations. Woo[Woo77] focused on cavity recognition in studying problems of transforming volumetric designs of parts into numerical control (NC) descriptions. Staley et al.[StHA83], without specific application emphasized, also dealt with cavity recognition but using syntactic pattern recognition techniques. Grayer[Gra76] compares the part in B-rep with its initial work part in order to compute NC tool paths. In addition to NC path generation, Armstrong[ArCP84] also considered the determination of fixture orientation. For global understanding of part geometry, Woo[Woo82] suggested a convex hull technique which transforms a boundary representation into an expression of alternating sum of volumes. To generate a part code for group technology, Kyprianou[Kyp80] used syntactic pattern recognition techniques to extract features characterized by protru-

sions and depressions. Jakubowski[Jak82] also used syntactic methods to develop a part description language primarily for the determination of part shapes. Henderson[Hen84] was able to use PROLOG and expert system techniques to extract features and organize them in a feature graph so as to facilitate automatic process planning.

While most of the work mentioned above is based on B-rep, very little has been done on CSG. In CSG, a machine part is represented as a tree with some primitive solids as terminal nodes and the regularized set operations (*union* and *difference*) and *movement* operations as non-terminal nodes. Properties of the machine part can be computed by evaluating the CSG tree. The evaluation procedure is basically a tree traversal routine that combines CSG subtrees into higher level CSG subtrees. The algorithm is simple but the computation involved in calculating the intersection (the major computation of *union* and *difference* operations) of two solids are complex and time-consuming.

Without evaluation, Lee and Fu[LeK86] proposed an approach to the problem of feature extraction. The approach is based on the spatial relationships among axes of primitive solids. They also suggested a unification process using tree reconstruction to unify feature representations so that properties can be computed more easily. As more general cases for this approach are yet to be studied, there are algorithms that convert CSG into B-rep[BoG82]. Not surprisingly, such algorithms are often called for when certain properties must be computed from a CSG tree[KaO84]. However, the conversion often appears to be computationally expensive or, being worse, a waste to simply compute the specific property.

There are part families for which CSG should be an excellent representation scheme and, furthermore, a transformation into B-rep can actually be avoided. A such example, rotational parts, is investigated in this study. Instead of B-rep, a new representation scheme PAR (Principal Axis Representation) is proposed as an internal representation scheme for rotational parts that are initially described by CSG. The key idea of PAR is to represent an object by a principal axis and a set of boundary curves. Based on a mathematical framework, an algorithm is designed and implemented to convert a CSG tree into a PAR, which represents the same object as the CSG tree does but is in an evaluated form. From PAR, the profile of the part can be efficiently computed. Other properties such as length

and maximum diameter can also be obtained easily.

The conversion algorithm is much simpler than the one converting CSG into B-rep. The main reason is that one dimensional curves (e.g. lines and arcs) are used to characterize the rotational parts and the evaluation of PAR is performed on two dimensional space, e.g. the intersection of two lines on the same plane rather than that of two solids on a three dimensional space. In addition to its computational ease, the PAR is proved to be a unique representation scheme. Being unique, it facilitates feature definition and thus simplifies the tasks of feature extraction. The possible extension of PAR is discussed in Section 6.

## **2 The Problem and Proposed Approach**

### **2.1 Problem**

Rotational parts include all parts that are symmetrical with respect to their principal axes. In this study, however, the primitive solids are limited to cylinders, cones, and tori only. The problem is formulated as follows:

Given a CSG tree of a machine part, which is

1. axis-symmetrical, and
2. constructed from cylinders, cones and tori in an arbitrary order of combination such as union and difference.

can we:

1. by utilizing the property of axis-symmetry, efficiently derive its profile and some geometric properties such as length and diameter so that part classification and process planning can possibly be supported?
2. develop an internal scheme which not only supports the above computation but also represents the part uniquely?

## 2.2 Basic Ideas

If a machine part is axis-symmetrical, it can be viewed as a  $2\frac{1}{2}$ D object and represented by the union of components, each being generated by rotating its boundary curve segments with respect to the principal axis. For example, a cylinder can be represented by rotating a line segment around its axis where the axis and the line segment are parallel to each other and their distance is the radius of the cylinder. Other primitive solids like cones and tori can also be represented in a similar way. Since the machine part is axis-symmetrical, each component solid depicted by a CSG subtree can also be represented by a collection of principal axis segments, each being associated with several pairs of boundary curves. For each boundary curve pair, rotating it around the associated axis segment would generate a volume layer for the corresponding component solid.

This idea of representing axis-symmetrical machine parts by a principal axis and a set of bounded curves can be exploited to develop efficient algorithms to evaluate CSG trees and deduce their geometrical properties because the computational complexity of the algorithms can be reduced significantly. For example, in the course of evaluating CSG trees, instead of testing the intersection of two three dimensional solids, we need only to test the intersection of a few one dimensional curves on the same plane, where the number of the 1D curves needed to be tested depends on the complexity of the two original 3D solids involved.

To evaluate a CSG tree and convert it into this kind of representation, new operations such as *union*, *difference*, and *movement* operated on this new representation scheme must be defined in order to maintain the actual semantics as their counterpart operations do in the CSG tree representation scheme. To deal with the *movement* operation is easy; if all the coordinates of the segments and boundary curves are relative to a principal axis, then the result of the *movement* operation is merely applying the *movement* transformation to the principal axis, and others remain unchanged. But the *union* and *difference* operations are not so simple to deal with, as more elaborate processing would be required.

After a CSG tree is converted to the new representation, the profile and other geometric properties of the corresponding part can be easily computed. These properties may be



computed all at the same time after the CSG tree is converted into the new representation. To compute the desired geometric properties directly from the CSG tree, on the other hand, is complex and time-consuming.

In summary, the proposed approach includes the following:

1. To define a new representation scheme (Principal Axis Representation) for axis-symmetrical machine parts
2. To convert CSG trees into their new representations.
3. To derive desired geometrical properties from this new representation, and it is believed to be easier and more efficient.
4. To prove that the Principal Axis Representation is a unique representation scheme for axis-symmetrical machine parts.

### 3 Formulation of Principal Axis Representation

#### 3.1 Terminology

**Definition: Principal Axis**

A *Principal Axis* A of an axis-symmetrical object O is a line segment in the three dimensional space such that the object O is represented in a way that starts from and ends at the two end points of A, and O is symmetrical with respect to A.

**Definition: Axis Segment**

An *Axis Segment* S is a line segment ( i.e. subset ) of a Principal Axis which has two end points.

**Definition: Principal Axis Coordinate System**

A curve C and a principal axis A form a coordinate system if they are co-plane and, on this plane, the horizontal axis is the line containing A and its vertical axis is the line

perpendicular to A and passing through the start point of A. This 2-D coordinate system is called *principal axis coordinate system*.

**Definition: Bounded Curve**

A curve C is bounded at [a b] with respect to a principal axis A iff in the principal axis coordinate system formed from C and A, there exist two real numbers a and b ( $a \leq b$ ) such that the curve C is differentiable within the interval (a b). a and b are called the *bound points* of the curve C, and [a b] is called a *pair of bounds* of C.

**Definition: Range of a Bounded Curve**

If C is a bounded curve within [a b] with respect to a principal axis A, the *range* of C at  $x, x \in [a b]$ , on the principal axis coordinate system is  $C(x)$ , i.e. the mapping of C from x.

**Definition: Arc**

An *Arc* is a bounded curve which is also a subset of a circle.

Note that : a line segment is a bounded curve.

**Definition: Bounded Curve Set**

A *Bounded Curve Set* is a set of *Arcs* and/or *Line Segments*.

**Definition: Principal Axis Representation ( PAR )**

If an object O is an axis-symmetrical machine part which can be characterized by a Principal Axis A, then its *Principal Axis Representation*  $PAR(O, A)$  can be defined as a set of tuples  $(S_i, C_i)$ ,  $i = 1, \dots, n$ ,  $n \in \mathbb{N}$ , such that

1.  $S_i$  is an Axis Segment, and  $S_i \subseteq A$ ,  $\bigcup_{i=1}^n S_i = A$ ,  $\exists i, j \leq n$ ,  $S_i \cup S_j$  is either a point (i.e. their common bound point) if  $j = i + 1$  or  $\emptyset$ , otherwise; that is,  $\{ S_i \}$  is a partition of A.
2.  $C_i = \{ C_{ij} \mid j = 1, \dots, 2m_i \}$  is a Bounded Curve Set.

3.  $\exists a, b \in R$  such that all  $C_{ij}$  's and  $S_i$  have the same bound  $[a, b]$  with respect to  $A$ .
4. all  $C_{ij}$  's do not intersect within their bound interval  $(a, b)$ .
5. all  $C_{ij}$  's are differentiable (i.e. their first derivatives exist ) within  $(a, b)$ , and  $\exists$  a curve  $C_{ik} \in C_i$  such that  $C_{ik}$  is not differentiable at the bound point  $a$ , and  $\exists$  a curve  $C_{il} \in C_i$  such that  $C_{il}$  is not differentiable at the bound point  $b$ .

**Theorem 1:**

For a Principal Axis Representation of an axis-symmetrical machine part  $O$  and its Principal Axis  $A$ ,  $PAR(O, A) = \{ (S_i, C_i) \mid i=1, \dots, n \}$ .  $\forall i \leq n$ , all  $C_{ij}$  's ( $C_{ij} \in C_i$ ,  $j = 1, \dots, 2r_i$ ), within  $(a, b)$  which is the bound of  $S_i$  with respect to  $A$ , form a *total ordering*.

Proof: let's denote  $C_{ij}(x)$  to be the range of the bounded curve  $C_{ij}$  at  $x$ ,  $x \in (a, b)$ .

For  $C_{ik}$  and  $C_{im} \in C_i$ ,  $k \neq m$ ,  $\exists x \in (a, b)$ ,  $C_{ik}(x)$  and  $C_{im}(x) \in R$ , and  $C_{ik}(x) \neq C_{im}(x)$ ; otherwise  $C_{ik}$  and  $C_{im}$  intersect at  $x$  within  $(a, b)$  and this contradicts to the definition of  $PAR(O, A)$ . Therefore either  $C_{ik}(x) < C_{im}(x)$  or  $C_{ik}(x) > C_{im}(x)$  is true. And we prove that all  $C_{ij}(x)$ ,  $C_{ij} \in C_i$ , form a total ordering at some  $x$ ,  $x \in (a, b)$ .

Next, we will prove that the total orderings of all  $C_{ij}$  ( $j = 1, 2, \dots, 2r_i$ ) at all  $x$ ,  $x \in (a, b)$  are consistent.

For  $x_1$  and  $x_2 \in (a, b)$ ,  $x_1 \neq x_2$ , if the total ordering of all  $C_{ij}$  ( $j = 1, 2, \dots, 2r_i$ ) at  $x_1$  and that of all  $C_{ij}$  ( $j = 1, 2, \dots, 2r_i$ ) at  $x_2$  are not consistent, then there must exist two different curves,  $C_{ik}$  and  $C_{il}$ , such that  $C_{ik}(x_1) < C_{il}(x_1)$  and  $C_{ik}(x_2) > C_{il}(x_2)$ . It means that there must exist some value  $x_3$  between  $x_1$  and  $x_2$  such that  $C_{ik}$  and  $C_{il}$  intersect at  $x_3$ . This contradicts to the definition of  $PAR$ . Therefore the total orderings of all  $C_{ij}$  ( $j = 1, 2, \dots, 2r_i$ ) at all  $x$ ,  $x \in (a, b)$ , must be consistent, and we may conclude that all  $C_{ij}$  ( $j = 1, 2, \dots, 2r_i$ ) within  $(a, b)$  form a unique total ordering. **Q.E.D.**

Note: this total ordering will be used to determine the relation among those curves within some bound.

Using the concept of Theorem 1, we may define the concept of a layer.

**Definition: Layer and Layer Set**

For  $\text{PAR}(O, A)$ ,  $S_i$  is an Axis Segment and  $C_i$  is its associated bounded curve set.  $C_{ij}$ 's ( $j = 1, \dots, 2m_i$ ) are those bounded curves in  $C_i$  and they are sorted in such a sequence :  $C_{i1} > C_{i2} > C_{i3} > \dots > C_{i2m_i}$ . Then a *layer* within this Axis Segment  $S$  is defined as a pair of curves  $C_{ik}$  and  $C_{i(k+1)}$ , denoted as  $[C_{ik}C_{i(k+1)}]$ ,  $k = 1, 3, \dots, 2m_i - 1$ . The set of layers  $\{ [C_{ik}C_{i(k+1)}] \mid k = 1, 3, \dots, 2m_i - 1 \}$  is called a *layer set*.

Note: a layer with its bound actually generates a volume of the object by rotating the curves in the layer with respect to its corresponding Axis Segment.

### 3.2 Represent Primitive Solids by PAR

#### 1. Use PAR to represent cylinder

$\text{PAR}(\text{cylinder}, A) = \{ (S, C) \mid C = \{ C_1, C_2 \}, C_1 \text{ and } C_2 \text{ are line segments}; S = A = C_2; C_1 \parallel C_2; |C_1 - C_2| = \text{radius of the cylinder}; [a \ b] \text{ is the bound for } S \text{ and } C \text{ where } a \text{ and } b \text{ are the horizontal coordinates of the starting and ending points, respectively, of } A \text{ in the Principal Axis Coordinate System formed by } A \text{ and } C_1. \}$

Notice that the statement " $C_1 \parallel C_2$ " means that the line segment  $C_1$  is parallel to the line segment  $C_2$ , and the statement " $|C_1 - C_2|$ " denotes the distance between  $C_1$  and  $C_2$ .

#### 2. Use PAR to represent cone

$\text{PAR}(\text{cone}, A) = \{ (S, C) \mid C = \{ C_1, C_2 \}, C_1 \text{ and } C_2 \text{ are line segment}; S = A = C_2; [a \ b] \text{ is the bound for } S \text{ and } C \text{ where } a \text{ and } b \text{ are the horizontal coordinates of the starting and ending points, respectively, of } A \text{ in the Principal Axis Coordinate System formed by } A \text{ and } C_1; C_1 \text{ and } C_2 \text{ intersect at } b, \text{ and the distance of } C_1 \text{ and } C_2 \text{ at } a \text{ is the radius of the cone. } \}$

#### 3. Use PAR to represent torus

$\text{PAR}(\text{torus}, A) = \{ (S, C) \mid S = A; C = \{ C_1, C_2 \}, [a \ b] \text{ is the bound for } S \text{ and } C \text{ where } a \text{ and } b \text{ are the horizontal coordinates of the starting and ending points, respectively, of } A \text{ in the Principal Axis Coordinate System formed by } A \text{ and } C; C_1, C_2 \text{ are Arcs of the circle with the radius equal to the } \textit{local radius} \text{ of the torus and with its center located at } (x$

y) in the Principal Axis Coordinate System where  $x = (a + b)/2$ , and y is the *global radius* of the torus,  $C_1$  is the upper half and  $C_2$  is the lower half of the circle. }

### 3.3 Operations on PAR

#### Definition: Overlap of Two Layers

The *overlap* of two layers  $[c_{11}c_{12}]$  and  $[c_{21}c_{22}]$  with respect to some bound  $[a b]$  is defined as either a layer  $[c_{31}c_{32}]$  with respect to  $[a b]$  if  $c_{31} = \min(c_{11}, c_{21})$ ,  $c_{32} = \max(c_{12}, c_{22})$  and  $c_{31} > c_{32}$ ; or  $\emptyset$ , otherwise.

#### Definition: Layer-Union of Two Layers

The *layer-union* of two layers  $[c_{11}c_{12}]$  and  $[c_{21}c_{22}]$  with respect to some bound  $[a b]$  is defined as either the layer  $[c_{31}c_{32}]$  where  $c_{31} = \max(c_{11}, c_{21})$ ,  $c_{32} = \min(c_{12}, c_{22})$  if the overlap of  $[c_{11}c_{12}]$  and  $[c_{21}c_{22}]$  is not  $\emptyset$ ; or the two layers  $[c_{11}c_{12}]$  and  $[c_{21}c_{22}]$ , otherwise. The result layer or layers are with respect to the same bound  $[a b]$ .

#### Definition: Layer-Union of Two Layer Sets

The *layer-union* of two layer sets  $L_1$  and  $L_2$  is defined as a layer set which is the *set union* of all layer-union's of two layers, one from  $L_1$  and the other from  $L_2$ , respectively.  $L_1$ ,  $L_2$  and the resulting layer set are all with respect to the same bound.

#### Definition: Maximum Union-able Layer Set

The *maximum union-able layer set* of two layer sets  $L_1$  and  $L_2$  is defined as a layer set which is the transitive closure of applying the layer-union operation to the resulting layers of the layer-union of  $L_1$  and  $L_2$ ; that is, if  $L_3$  is the layer-union of  $L_1$  and  $L_2$ , then the *maximum union-able layer set* of  $L_1$  and  $L_2$  is the results of repeatedly applying the layer-union operation to the layers in  $L_3$  until the resulting set is no more changed.  $L_1$ ,  $L_2$  and the maximum union-able layer set are all with respect to the same bound.

#### Definition: Union of Two PAR's

For two PAR's,  $\text{PAR}(O_1, A_1)$  and  $\text{PAR}(O_2, A_2)$ , if  $A_1$  and  $A_2$  are subsets of a line, then the *union* of  $\text{PAR}(O_1, A_1)$  and  $\text{PAR}(O_2, A_2)$  is defined as a mapping from  $\text{PAR} \times \text{PAR}$  to  $\text{PAR}$  such that if  $\text{PAR}(O_1, A_1) = \{ (S_i^1, C_i^1) \mid i = 1, \dots, n_1, [a_i^1 b_i^1] \text{ is the bound of } S_i^1 \}$ ,  $\text{PAR}(O_2, A_2) = \{ (S_i^2, C_i^2) \mid i = 1, \dots, n_2, [a_i^2 b_i^2] \text{ is the bound of } S_i^2 \}$ , and there exists a  $\text{PAR}(O_3, A_3) = \{ (S_i^3, C_i^3) \mid i = 1, \dots, n, [a_i^3 b_i^3] \text{ is the bound of } S_i^3 \}$ , then

1.  $O_3 = O_1 \cup O_2$ ;  $A_3 = A_1 \cup A_2$ ;
2.  $S_i^3$  is a line segment of  $A_3$  with bounds  $[a_i^3 b_i^3]$ ,  $i=1, \dots, n$ ; and there exists a set of bounds  $[d_m^3 e_m^3]$ ,  $m = 1, 2, \dots, n_3$  ( $n_3 \geq n$ ) which is a partition of the set of bounds  $[a_i^3 b_i^3]$  and there exists a bounded curve set  $K_m^3$  for each bound  $[d_m^3 e_m^3]$  where
  - $K_m^3 = C_j^1$  with adjusted bound  $[d_m^3 e_m^3]$  if  $\exists [a_l^2 b_l^2] \supseteq [d_m^3 e_m^3]$ ,  $l = 1, \dots, n_2$
  - $K_m^3 = C_k^2$  with adjusted bound  $[d_m^3 e_m^3]$  if  $\exists [a_l^1 b_l^1] \supseteq [d_m^3 e_m^3]$ ,  $l = 1, \dots, n_1$
  - $K_m^3 = \text{maximum union-able layer set of } C_j^1 \text{ and } C_k^2 \text{ with respect to a bound } [d_m^3 e_m^3] \text{ if } [a_j^1 b_j^1] \supseteq [d_m^3 e_m^3], \text{ and } [a_k^2 b_k^2] \supseteq [d_m^3 e_m^3] \text{ for some } j, k.$
3.  $C_i^3 = \bigcup_{m=x_i}^{y_i} K_m^3$ , ( $x_i, y_i$  are integers) if  $[d_m^3 e_m^3]$  for  $m = x_i, \dots, y_i$  is a partition of  $[a_i^3 b_i^3]$

NOTE: This mapping is from  $\text{PAR} \times \text{PAR}$  to  $\text{PAR}$ , the non-differentiable properties of curves across the bound  $[a_i^3 b_i^3]$  for the resulting PAR should be preserved by the definition of PAR.

#### Definition: Layer-Difference of Two Layers

The *layer-difference* of two layers  $[c_{11}c_{12}]$  and  $[c_{21}c_{22}]$  with respect to some bound  $[a b]$  is defined as

1. the layer  $[c_{11}c_{12}]$  if  $c_{11} \leq c_{22}$  or  $c_{12} \geq c_{21}$ ;
2. the layer  $[c_{22}c_{12}]$  if  $c_{21} \geq c_{11} \geq c_{22} > c_{12}$ ;
3. the layer  $[c_{11}c_{21}]$  if  $c_{11} > c_{21} \geq c_{12} \geq c_{22}$ ;
4. the layers  $[c_{11}c_{21}]$  and  $[c_{22}c_{12}]$  if  $c_{11} > c_{21} > c_{22} > c_{12}$ ;

5.  $\emptyset$  if  $c_{21} \geq c_{11} > c_{12} \geq c_{22}$ ;

The result layer or layers is with respect to the same bound [a b].

**Definition: Layer-Difference of Two Layer Sets**

The *layer-difference* of two layer sets  $L_1$  and  $L_2$  (assume  $L_1 = \{ [c_i c_{i+1}] \mid i = 1, 3, \dots, 2n-1 \}$  is the first operand and  $L_2 = \{ [d_i d_{i+1}] \mid i = 1, 3, \dots, 2m-1 \}$  is the second operand) is defined as a layer set which is the *set union* of all the set  $OV_k$  ( $k = 1, 2, \dots, n$ ) where each  $OV_k$  is the overlap of all the layers  $l_{kj}$  ( $j = 1, 2, \dots, m$ ),  $l_{kj}$  = the layer-difference of  $[c_{2k-1} c_{2k}]$  to  $[d_{2j-1} d_{2j}]$ .  $L_1$ ,  $L_2$  and the resulting layer set are all with respect to the same bound.

**Definition: Minimum Difference-able Layer Set**

The *minimum difference-able layer set* of two layer sets  $L_1$  and  $L_2$  with respect to some bound [a b] is the *layer-difference* of  $L_1$  to  $L_2$  with respect to the same bound [a b].

**Definition: Difference of Two PAR's**

For two PAR's,  $PAR(O_1, A_1)$  and  $PAR(O_2, A_2)$ , if  $A_1$  and  $A_2$  are subsets of a line, then the *difference* of  $PAR(O_1, A_1)$  and  $PAR(O_2, A_2)$  is defined as a mapping from  $PAR \times PAR$  to  $PAR$  such that if  $PAR(O_1, A_1) = \{ (S_i^1, C_i^1) \mid i = 1, \dots, n_1, [a_i^1 b_i^1] \text{ is the bound of } S_i^1 \}$ ,  $PAR(O_2, A_2) = \{ (S_i^2, C_i^2) \mid i = 1, \dots, n_2, [a_i^2 b_i^2] \text{ is the bound of } S_i^2 \}$ , and there exists a  $PAR(O_3, A_3) = \{ (S_i^3, C_i^3) \mid i = 1, \dots, n, [a_i^3 b_i^3] \text{ is the bound of } S_i^3 \}$ , then

1.  $O_3 = O_1 - O_2$ ;  $A_3 \subseteq A_1$ ;
2.  $S_i^3$  is a line segment of  $A_3$  with bounds  $[a_i^3 b_i^3]$ ,  $i=1, \dots, n$ ; and there exists a set of bounds  $[d_m^3 e_m^3]$ ,  $m = 1, 2, \dots, n_3$  ( $n_3 \geq n$ ) which is a partition of the set of bounds  $[a_i^3 b_i^3]$  and there exists a bounded curve set  $K_m^3$  for each bound  $[d_m^3 e_m^3]$  where
  - $K_m^3 = C_j^1$  with adjusted bound  $[d_m^3 e_m^3]$  if  $B [a_i^2 b_i^2] \supseteq [d_m^3 e_m^3]$ ,  $l = 1, \dots, n_2$
  - $K_m^3 = \emptyset$  if  $B [a_l^1 b_l^1] \supseteq [d_m^3 e_m^3]$ ,  $l = 1, \dots, n_1$
  - $K_m^3 =$  minimum difference-able layer set of  $C_j^1$  and  $C_k^2$  with respect to a bound  $[d_m^3 e_m^3]$  if  $[a_j^1 b_j^1] \supseteq [d_m^3 e_m^3]$ , and  $[a_k^2 b_k^2] \supseteq [d_m^3 e_m^3]$  for some  $j, k$ .

3.  $C_i^3 = \bigcup_{m=x_i}^{y_i} K_m^3$ , ( $x_i, y_i$  are integers) if  $[d_m^3 e_m^3]$  for  $m = x_i, \dots, y_i$  is a partition of  $[a_i^3 b_i^3]$ .

**Definition: Movement of a PAR**

The *movement* of a PAR(O, A) is defined as a mapping from PAR to PAR such that the new PAR(O,B) after movement is the same as PAR(O,A) except that A is transformed to B in 3-D space by applying the movement operation specified.

Note: The bounds in PAR are with respect to A, and therefore movement-invariant.

### 3.4 Relation between PAR and CSG

In this subsection, the relation between CSG and PAR representation schemes is investigated. It is the basis of our algorithm to convert a CSG tree into a PAR.

**Theorem 2.**

For a CSG tree with *union* as its root, the object represented by this CSG tree is identical to that represented by a PAR which is the *union* of two PAR's, each corresponding to either the left- or right- subtree of the original CSG tree.

Proof : Assume that  $O_3 = O_1 \cup O_2$  in CSG domain where  $O_3$  is the CSG tree, and  $O_1$  and  $O_2$  are its left- and right- subtrees, respectively. In the PAR domain, assume  $O_1$  and  $O_2$  are represented by  $\text{PAR}(O_1, A_1)$  and  $\text{PAR}(O_2, A_2)$ , respectively. Also assume  $A_3 = A_1 \cup A_2$ .

Now we want to prove  $\text{PAR}(O_1, A_1) \cup \text{PAR}(O_2, A_2)$  represents  $O_3$ . In this proof, we may view an object as a set of points in 3-D space and use the notation  $V(\text{layers})$  to denote the volume generated by rotating the layers with respect to a principal axis in 3-D space.

Step 1: To prove  $O_3 \subseteq V(\text{layers of } (\text{PAR}(O_1, A_1) \cup \text{PAR}(O_2, A_2)))$ .

For every point  $p$  of  $O_3$ ,  $p \in O_1$  or  $p \in O_2$  is true in CSG domain because  $O_3 = O_1 \cup O_2$ . But  $O_1$  and  $O_2$  can be represented by  $\text{PAR}(O_1, A_1)$  and  $\text{PAR}(O_2, A_2)$ , respectively. Therefore,  $p \in V(\text{layers of } (\text{PAR}(O_1, A_1)))$  or  $p \in V(\text{layers of } (\text{PAR}(O_2, A_2)))$ . More specifically, on the principal axis coordinate system, the horizontal coordinate of  $p$  must lie



5.  $\emptyset$  if  $c_{21} \geq c_{11} > c_{12} \geq c_{22}$ ;

The result layer or layers is with respect to the same bound[a b].

**Definition: Layer-Difference of Two Layer Sets**

The *layer-difference* of two layer sets  $L_1$  and  $L_2$  (assume  $L_1 = \{ [c_i c_{i+1}] \mid i = 1, 3, \dots, 2n-1 \}$  is the first operand and  $L_2 = \{ [d_i d_{i+1}] \mid i = 1, 3, \dots, 2m-1 \}$  is the second operand) is defined as a layer set which is the *set union* of all the set  $OV_k$  ( $k = 1, 2, \dots, n$ ) where each  $OV_k$  is the overlap of all the layers  $l_{kj}$  ( $j = 1, 2, \dots, m$ ),  $l_{kj}$  = the layer-difference of  $[c_{2k-1} c_{2k}]$  to  $[d_{2j-1} d_{2j}]$ .  $L_1$ ,  $L_2$  and the resulting layer set are all with respect to the same bound.

**Definition: Minimum Difference-able Layer Set**

The *minimum difference-able layer set* of two layer sets  $L_1$  and  $L_2$  with respect to some bound [a b] is the *layer-difference* of  $L_1$  to  $L_2$  with respect to the same bound [a b].

**Definition: Difference of Two PAR's**

For two PAR's,  $PAR(O_1, A_1)$  and  $PAR(O_2, A_2)$ , if  $A_1$  and  $A_2$  are subsets of a line, then the *difference* of  $PAR(O_1, A_1)$  and  $PAR(O_2, A_2)$  is defined as a mapping from  $PAR \times PAR$  to  $PAR$  such that if  $PAR(O_1, A_1) = \{ (S_i^1, C_i^1) \mid i = 1, \dots, n_1, [a_i^1 b_i^1] \text{ is the bound of } S_i^1 \}$ ,  $PAR(O_2, A_2) = \{ (S_i^2, C_i^2) \mid i = 1, \dots, n_2, [a_i^2 b_i^2] \text{ is the bound of } S_i^2 \}$ , and there exists a  $PAR(O_3, A_3) = \{ (S_i^3, C_i^3) \mid i = 1, \dots, n, [a_i^3 b_i^3] \text{ is the bound of } S_i^3 \}$ , then

1.  $O_3 = O_1 - O_2$ ;  $A_3 \subseteq A_1$ ;
2.  $S_i^3$  is a line segment of  $A_3$  with bounds  $[a_i^3 b_i^3]$ ,  $i=1, \dots, n$ ; and there exists a set of bounds  $[d_m^3 e_m^3]$ ,  $m = 1, 2, \dots, n_3$  ( $n_3 \geq n$ ) which is a partition of the set of bounds  $[a_i^3 b_i^3]$  and there exists a bounded curve set  $K_m^3$  for each bound  $[d_m^3 e_m^3]$  where
 
$$K_m^3 = C_j^1 \text{ with adjusted bound } [d_m^3 e_m^3] \text{ if } \exists [a_l^2 b_l^2] \supseteq [d_m^3 e_m^3], l = 1, \dots, n_2$$

$$K_m^3 = \emptyset \text{ if } \exists [a_l^1 b_l^1] \supseteq [d_m^3 e_m^3], l = 1, \dots, n_1$$

$$K_m^3 = \text{minimum difference-able layer set of } C_j^1 \text{ and } C_k^2 \text{ with respect to a bound } [d_m^3 e_m^3]$$
 if  $[a_j^1 b_j^1] \supseteq [d_m^3 e_m^3]$ , and  $[a_k^2 b_k^2] \supseteq [d_m^3 e_m^3]$  for some  $j, k$ .

3.  $C_i^3 = \bigcup_{m=x_i}^{y_i} K_m^3$ , ( $x_i, y_i$  are integers) if  $[d_m^3 e_m^3]$  for  $m = x_i, \dots, y_i$  is a partition of  $[a_i^3 b_i^3]$ .

**Definition: Movement of a PAR**

The *movement* of a PAR(O, A) is defined as a mapping from PAR to PAR such that the new PAR(O,B) after movement is the same as PAR(O,A) except that A is transformed to B in 3-D space by applying the movement operation specified.

Note: The bounds in PAR are with respect to A, and therefore movement-invariant.

### 3.4 Relation between PAR and CSG

In this subsection, the relation between CSG and PAR representation schemes is investigated. It is the basis of our algorithm to convert a CSG tree into a PAR.

**Theorem 2.**

For a CSG tree with *union* as its root, the object represented by this CSG tree is identical to that represented by a PAR which is the *union* of two PAR's, each corresponding to either the left- or right- subtree of the original CSG tree.

Proof : Assume that  $O_3 = O_1 \cup O_2$  in CSG domain where  $O_3$  is the CSG tree, and  $O_1$  and  $O_2$  are its left- and right- subtrees, respectively. In the PAR domain, assume  $O_1$  and  $O_2$  are represented by  $\text{PAR}(O_1, A_1)$  and  $\text{PAR}(O_2, A_2)$ , respectively. Also assume  $A_3 = A_1 \cup A_2$ .

Now we want to prove  $\text{PAR}(O_1, A_1) \cup \text{PAR}(O_2, A_2)$  represents  $O_3$ . In this proof, we may view an object as a set of points in 3-D space and use the notation  $V(\text{layers})$  to denote the volume generated by rotating the layers with respect to a principal axis in 3-D space.

Step 1: To prove  $O_3 \subseteq V(\text{layers of } (\text{PAR}(O_1, A_1) \cup \text{PAR}(O_2, A_2)))$ .

For every point  $p$  of  $O_3$ ,  $p \in O_1$  or  $p \in O_2$  is true in CSG domain because  $O_3 = O_1 \cup O_2$ . But  $O_1$  and  $O_2$  can be represented by  $\text{PAR}(O_1, A_1)$  and  $\text{PAR}(O_2, A_2)$ , respectively. Therefore,  $p \in V(\text{layers of } (\text{PAR}(O_1, A_1)))$  or  $p \in V(\text{layers of } (\text{PAR}(O_2, A_2)))$ . More specifically, on the principal axis coordinate system, the horizontal coordinate of  $p$  must lie

in some bounds  $[a_i b_i]$ , which is a subset of  $[a_j^1 b_j^1]$  of  $\text{PAR}(O_1, A_1)$  or a subset of  $[a_k^2 b_k^2]$  of  $\text{PAR}(O_2, A_2)$  for some  $j$  or  $k$ .

1. if  $[a_i b_i]$  is a subset of  $[a_j^1 b_j^1]$  but not a subset of  $[a_k^2 b_k^2]$ , then  $p \in V$ ( layers specified by  $C_j^1$  of  $\text{PAR}(O_1, A_1)$  ).
2. if  $[a_i b_i]$  is a subset of  $[a_k^2 b_k^2]$  but not a subset of  $[a_j^1 b_j^1]$ , then  $p \in V$ ( layers specified by  $C_k^2$  of  $\text{PAR}(O_2, A_2)$  ).
3. if  $[a_i b_i]$  is both a subset of  $[a_k^2 b_k^2]$  and a subset of  $[a_j^1 b_j^1]$ , then  $p \in V$ (  $[c_m c_{m+1}]$  ) or  $p \in V$ (  $[d_n d_{n+1}]$  ) where  $[c_m c_{m+1}]$  and  $[d_n d_{n+1}]$  are some layers of  $\text{PAR}(O_1, A_1)$  and  $\text{PAR}(O_2, A_2)$ , respectively. By the definition of *layer-union* of two layers and *Maximum Union-able Layer Set*,  $p \in V$ ( layer-union of  $[c_m c_{m+1}]$  and  $[d_n d_{n+1}]$  ), and therefore  $p \in V$ ( maximum union-able layer set of  $L_1$  and  $L_2$  ) where  $L_1$  and  $L_2$  are the layer sets with the bound  $[a_i b_i]$  that contain  $[c_m c_{m+1}]$  and  $[d_n d_{n+1}]$ , respectively.

Thus, from the definition of the *union* of two PAR's,  $p \in V$ ( layers of  $(\text{PAR}(O_1, A_1) \cup \text{PAR}(O_2, A_2))$  ). And we prove that  $O_3 \subseteq V$ ( layers of  $(\text{PAR}(O_1, A_1) \cup \text{PAR}(O_2, A_2))$  ).

Step 2: To prove  $O_3 \supseteq V$ ( layers of  $(\text{PAR}(O_1, A_1) \cup \text{PAR}(O_2, A_2))$  ).

For every point  $p$ ,  $p \in V$ ( layers of  $(\text{PAR}(O_1, A_1) \cup \text{PAR}(O_2, A_2))$  ), by the definition of the *union* of two PAR's, there are three possible cases for  $p$ :

1.  $p \in V$ (  $[c_m c_{m+1}]$  ) where  $[c_m c_{m+1}]$ , which is bounded by some bound  $[a b] \subseteq [a_i^1 b_i^1]$  but  $\not\subseteq [a_j^2 b_j^2]$ , is a layer of  $\text{PAR}(O_1, A_1)$ . In this case, since  $\text{PAR}(O_1, A_1)$  represents  $O_1$ ,  $p \in O_1$  and therefore  $p \in O_1 \cup O_2 = O_3$ .
2.  $p \in V$ (  $[d_n d_{n+1}]$  ) where  $[d_n d_{n+1}]$ , which is bounded by some bound  $[a b] \subseteq [a_j^2 b_j^2]$  but  $\not\subseteq [a_i^1 b_i^1]$ , is a layer of  $\text{PAR}(O_2, A_2)$ . In this case, since  $\text{PAR}(O_2, A_2)$  represents  $O_2$ ,  $p \in O_2$  and therefore  $p \in O_1 \cup O_2 = O_3$ .
3.  $p \in V$ ( maximum union-able layer set of  $L_1$  and  $L_2$  ) where  $L_1$  and  $L_2$  are the sets of layers of  $\text{PAR}(O_1, A_1)$  and  $\text{PAR}(O_2, A_2)$  at some bounds  $[a b]$ , respectively. By the definition of the maximum union-able layer set, there are three possible cases for  $p$ :

- (a)  $p \in V( [c_m c_{m+1}] \subseteq L_1 \text{ but } \not\subseteq L_2)$ . Thus,  $p \in O_1 \subseteq (O_1 \cup O_2) = O_3$ .
- (b)  $p \in V( [d_n d_{n+1}] \subseteq L_2 \text{ but } \not\subseteq L_1)$ . Thus,  $p \in O_2 \subseteq (O_1 \cup O_2) = O_3$ .
- (c)  $p \in V( [e_l e_{l+1}] )$  where  $[e_l e_{l+1}] \subseteq (L_1 \cap L_2)$ . In this case,  $p \in (V(L_1) \cap V(L_2))$ .  
That is  $p \in (O_1 \cap O_2) \subseteq (O_1 \cup O_2) = O_3$ .

Therefore we prove that  $O_3 \supseteq V(\text{layers of } (\text{PAR}(O_1, A_1) \cup \text{PAR}(O_2, A_2)))$ .

From Step 1 and Step 2, we prove the theorem. **Q.E.D.**

Notice that the union operator on PAR performs the actual semantics of the union operator on CSG.

### Theorem 3.

For a CSG tree with *difference* as its root, the object represented by this CSG tree is identical to that represented by a PAR which is the *difference* of two PAR's, which correspond to the left- and right- subtrees of the original CSG tree.

Proof: Instead of using the concept of *maximum union-able layer set* and the definition of the *union* of two PAR's, the concept of *minimum difference-able layer set* and the definition of the *difference* of two PAR's may be used here to prove this theorem in a similar way as in Theorem 2. **Q.E.D.**

Notice that the difference operator on PAR performs the actual semantics of the difference operator on CSG.

### Theorem 4.

For a CSG tree with *movement* as its root, the object represented by this CSG tree is identical to that represented by a PAR which is the *movement* of a PAR, which corresponds to the subtree of the original CSG tree.

Proof: By the definition of *movement* of a PAR, the object represented by the PAR after applying a *movement* operation is the same as the object represented by the original PAR (i.e. the structure of the object is not changed) except that the location and orientation of the object are changed. In the CSG representation scheme, an object represented by a CSG

subtree has the same structure as the object represented by the CSG tree after applying a *movement* operation to the original CSG subtree except these two objects have different location and orientation in 3-D space.

Therefore, after the same transformation matrix (i.e. *movement* operation) is applied to both the PAR and the original CSG subtree which represent the same object, the resulting objects should be identical; that is, they not only have the same structure but also have the same location and orientation in the 3-D space. **Q.E.D.**

Notice that the movement operator on PAR performs the actual semantics of the movement operator on CSG.

### **Theorem 5.**

An axis-symmetrical machine part represented by a CSG tree can be evaluated on PAR domain and the final resulting PAR after evaluation represents the same object as the CSG tree represents.

**Proof :** A CSG tree is composed of operators (*union, difference, movement*) as non-terminal nodes and primitive solids (*cylinder, cone, torus*) as terminal nodes in this study. It is shown in Section 3.2 that these primitive solids can be represented by their corresponding PAR's. If the CSG tree is converted in bottom-up fashion from leaves to root, by Theorem 2, 3 and 4, the final resulting PAR should represent the identical object as the CSG tree represents. **Q.E.D.**

Notice that by this theorem, a given CSG tree can be evaluated in PAR domain to obtain an identical object as the CSG tree represents. This theorem is the theoretical basis of our algorithms in Section 4.

### **Theorem 6. Uniqueness**

Any axis-symmetrical machine part has a *unique* PAR representation.

**Proof :** Our theorem may be rephrased as follows:

For two representations,  $PAR(O_1, A_1)$  and  $PAR(O_2, A_2)$ , if  $O_1 = O_2$  then  $PAR(O_1, A_1) = PAR(O_2, A_2)$ .

Assume that the object  $O_1$  ( or  $O_2$  ) is in some 3-D coordinate system. Then there must be a unique Principal Axis for the object in this coordinate system, thus  $A_1 = A_2$ .

Our proof procedure consists of two steps:

Step 1: To prove  $\text{PAR}(O_1, A_1)$  and  $\text{PAR}(O_2, A_2)$  have the same Axis Segment bounds with respect to  $A_1$  and  $A_2$ , respectively. That is, if  $\text{PAR}(O_1, A_1) = \{ (S_i^1, C_i^1) \mid i = 1, \dots, n; S_i^1 \text{ is bounded by } [a_i^1 \ b_i^1] \}$ ,  $\text{PAR}(O_2, A_2) = \{ (S_i^2, C_i^2) \mid i = 1, \dots, m; S_i^2 \text{ is bounded by } [a_i^2 \ b_i^2] \}$  then we want to prove  $n = m$ , and  $a_i^1 = a_i^2$  and  $b_i^1 = b_i^2$ , for all  $i = 1, \dots, n$ .

By the definition of PAR, all the curves within a bound must be differentiable and there exists at least one curve non-differentiable at one of its lower and upper bounds. This means that the bounds on the Principal Axis are uniquely defined by the discontinuity of the first derivatives of the curves specifying the object. Because the object has fixed shape and the bounded curve to specify some part of the shape of an object can be either an arc or a line segment but not both, the bounded curve to describe the object shape and thus the discontinuity points of the bounded curve can be uniquely specified in PAR. Therefore, there is a unique way to partition the Principal Axis in PAR and this unique partition decides a unique set of bounds.

Step 2: Within each bound of  $(S_i^1, C_i^1)$  and  $(S_i^2, C_i^2)$  of  $\text{PAR}(O_1, A_1)$  and  $\text{PAR}(O_2, A_2)$ , respectively, we want to prove  $C_i^1 = C_i^2$  if  $S_i^1 = S_i^2$ .

This can be proved by contradiction. Assume  $C_i^1 \neq C_i^2$ . Since  $C_i^1$  and  $C_i^2$ , both are composed of a set of layers, if  $C_i^1 \neq C_i^2$ , there must exist, for some  $j$ , a layer  $[c_{ij}^1, c_{i,j+1}^1]$  in  $C_i^1$  and its corresponding layer  $[c_{ij}^2, c_{i,j+1}^2]$  in  $C_i^2$  such that  $[c_{ij}^1, c_{i,j+1}^1] \neq [c_{ij}^2, c_{i,j+1}^2]$ , and generate different volumes for  $O_1$  and  $O_2$ . In this case  $O_1 \neq O_2$ , which contradicts to our initial assumption  $O_1 = O_2$ . Therefore,  $C_i^1 = C_i^2$  if  $S_i^1 = S_i^2$ .

From the proofs of Step1 and Step2, in PAR, there is a unique set of bounds on the Principal Axis for an object, and on each pair of bounds there is a unique set of layers ( i.e. pairs of curves) to represent the object. Therefore, we conclude that there is only one unique PAR for an axis-symmetrical object and further PAR is a unique representation scheme. **Q.E.D.**

## 4 Algorithm to Convert CSG Representation to PAR

In this section, we will describe how to convert a CSG representation to its corresponding Principal Axis Representation (PAR). Section 4.1 first describes the data structure for PAR. The algorithm is described in Section 4.2, while Section 4.3 describes how to combine (either *union* or *difference*) two PAR's into one. The combination procedure is definitely a key component of the conversion algorithm.

### 4.1 Data Structure for PAR

Since a PAR is composed of a set of tuples  $(S_i, C_i)$ , we may use a record to represent each tuple and a link list to link these records together. Being a pair of bounds on the principal axis, each  $S_i$  can thus be represented by a pair of real numbers and stored in the tuple record. Each  $C_i$  consists of a set of curves  $C_{ij}$ , and can therefore be represented by a pointer to a link list where each element stores the parameters of one curve pair  $C_{ij}$  and  $C_{ij+1}$ , that is, a record representing a layer of the machine part.

We will assume the link list for tuples is sorted in the order of their bounds. Within each record for a tuple, the link list for the pairs of curves  $C_{ij}$  is also kept sorted in their ordering (from Theorem 1). This makes the concept of layers easy to deal with because each pair of the curves enumerating from the beginning of the sorted sequence is just a layer.

### 4.2 Conversion from CSG Representation to PAR

The conversion algorithm basically traverses the CSG tree from bottom to top. It first converts the CSG leaf nodes to their corresponding PAR's, and then combines (either *union* or *difference*) the PAR subtrees into composite PAR subtrees in higher levels. The combination procedure is repeated until the root is visited. The following recursive procedure describes this algorithm.

```
function EvaluateCSG ( T : CSG_tree ) : principal_axis_rep;
var
    P, P1, P2 : principal_axis_rep;
```

```

begin

    case T.type of

        movement : P := EvaluateCSG ( T.left_child^ );
                  Transform ( T.movement, P );
                  return ( P );

        union,

        difference: P1 := EvaluateCSG ( T.left_child^ );
                  P2 := EvaluateCSG ( T.right_child^ );
                  P := Combine ( T.type, P1, P2 );
                  return ( P );

        primitive : P := BuildAxisRep ( T );
                  return ( P );

    end; { case }

end; { EvaluateCSG }

```

To convert a CSG tree, its root can be passed to this function, which will recursively walk the whole CSG tree and transform it into a PAR. We assume that each node of the CSG tree is either a primitive solid (cylinder, cube, or torus) or an operator (*movement*, *union* or *difference*).

Primitives (i.e. the CSG-tree leaf nodes) are converted to corresponding PAR's by *BuildAxisRep* procedure, which is based on the framework of Section 3.2. Using the data structure representation in Section 4.1, the *BuildAxisRep* procedure simply creates one record to represent the primitive. In this record, the pair of bounds (in terms of the line parameter values of its principal axis) of the primitive solid are stored. A pair of curves representing a layer is also stored within the record. For a cylinder or cone, the outer curve of the pair is the line specifying the outer shape of the primitive solid and the inner curve is just the principal axis. For a torus, this pair is just the outer and inner half circles of the torus.



For the *movement* node in the CSG tree, its subtree (i.e. left-subtree) is evaluated first, then the transformation matrix stored within the *movement* node is applied to the evaluated subtree. The transformation matrix is a 4 by 4 matrix denoting the translation and rotation components of the movement. Since all the curves in PAR are specified relative to the principal axis, procedure *Transform* needs only transform the principal axis, not the whole set of curves. In fact, the transformation matrix is applied only to the two ending points of the principal axis.

To deal with the *union* or *difference* node, its two subtrees are evaluated first, then the *Combine* procedure is called to union or difference them together. To implement the definitions of union and difference of two PAR's developed in Section 3.3, the *Combine* procedure includes four passes. We will describe these four passes in the next section.

### 4.3 Union and Difference of Two PAR's

This algorithm (the procedure *Combine*) basically employs the split-and-merge paradigm, splitting the two composite axes into segments, computing the union or difference of two segments and then merging all the resulting segments into a new PAR. It has four steps (passes).

#### Step 1. Find New Pairs of Bounds

In this stage, the two operand subtrees being combined are represented in the PAR form, each being a link list of segments (records) in the order of the values of the bounded pairs. The function of this step is to compute the new bounded pairs for the resulting PAR.

Since the values of the bounds in each operand PAR are in terms of line parameters of its own principal axis, we must first convert these values so that they are in terms of the line parameters of the resulting principal axis. Depending on the operation to be performed, the resulting principal axis can be chosen either as that of the first operand for difference operation or as the union of those of the two operands for the union operation. This selection process can be easily accomplished by simply checking the two ending points of the two operand principal axes. We assume that the two operands have the principal

axes in the same or opposite direction if they are to be combined. (Note that the opposite direction may occur because *cones* are not symmetrical in its top and bottom and thus have two directions.)

After the resulting principal axis is roughly chosen and the values of the bounds are adjusted in term of this resulting axis, we may determine the pairs of all bounds for the resulting PAR. This process is simply the merge-and-sort procedure by repeatedly inputting two values of the bounds from each operand PAR and choosing the smaller value for the new bound. The state transition diagrams of this procedure are shown in Figure 1 and Figure 2.

In Figure 1 and Figure 2, we may imagine that there are two stacks for two operand PAR's. Each stack stores the bounds of one operand PAR which are in ascending order with the smallest one on the top of the stack. A pair of bounds is just an even-odd pair of values on the stacks if the top element of the stack is numbered from zero. This implies that a common bound of two neighboring segments has duplicate values on the stack. Let's also assume that *a* is the top element of the first operand stack A and *b* is the top element of the second stack B, respectively.

Now the following four states in the state transition diagrams can be defined.

- State “-0” means a upper bound of a pair in stack A and a upper bound of a pair in stack B have just been popped out.
- State “+0” means a lower bound of a pair in stack A and a upper bound of a pair in stack B have just been popped out.
- State “+1” means a lower bound of a pair in stack A and a lower bound of a pair in stack B have just been popped out.
- State “-1” means a upper bound of a pair in stack A and a lower bound of a pair in stack B have just been popped out.

The starting and final state is State “-0”. The notation { *condition/ action1; action2; ...* } is used in the state transition diagrams to indicate that, if the *condition* is satisfied, the

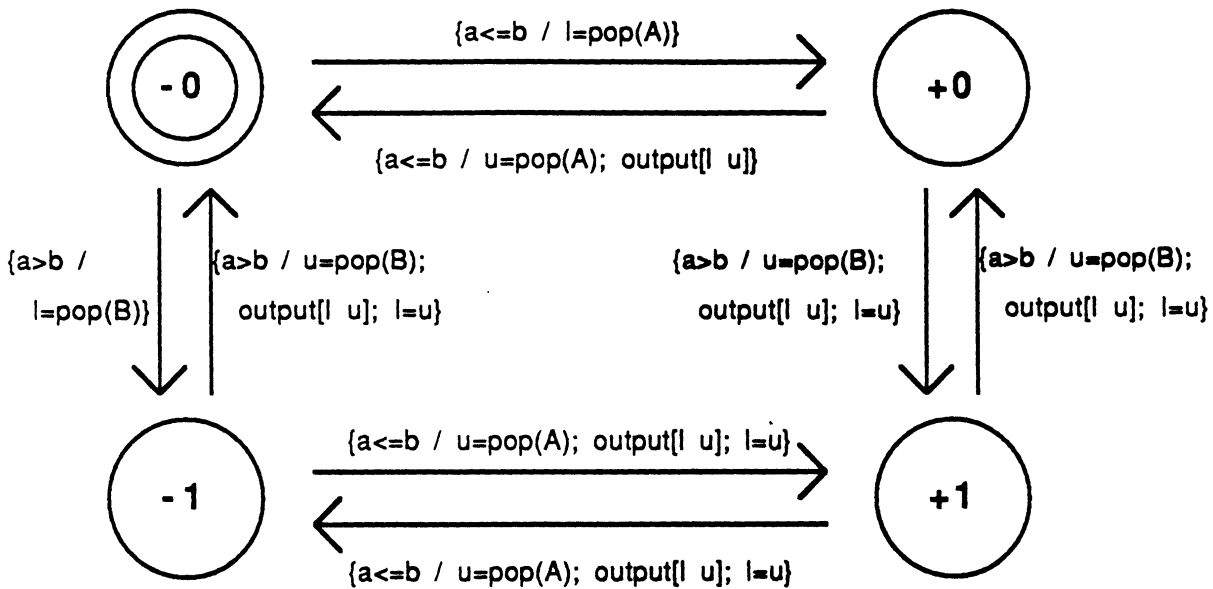


Figure 1: State Transition to Compute the New Bounds for Union.

actions specified in the bracket are executed and the new state is pointed to by the arrow sign. Let's also use  $l$  and  $u$  to denote lower bound and upper bound of a bound pair. Note that these state transition diagrams may generate null bound pairs like  $[c c]$ , where  $c$  is a real number, and they should be discarded.

**Step 2. Refine the Pairs of Bounds**

After Step 1, a set of pairs of bounds is obtained for the resulting PAR. For those pairs of bounds which are not subintervals common to both operands, the resulting shape within them can be immediately determined by the curves from the original (i.e. operand PAR) bound pairs. For those new bound pairs which are subintervals of both original operand PARs, the story is, however, not so simple. They should be refined. That is, subdivision of them must be considered because the resulting shape within a pair of bounds might be determined by curves from both original PARs.

To refine the pairs of bounds which are the subintervals of both original PARs, we compute the intersection points of the curves within them. After sorting these intersecting points, a set of refined pairs of bounds may be created. The resulting curves within these refined pairs of bounds are then determined by the subcurves of the original curves.

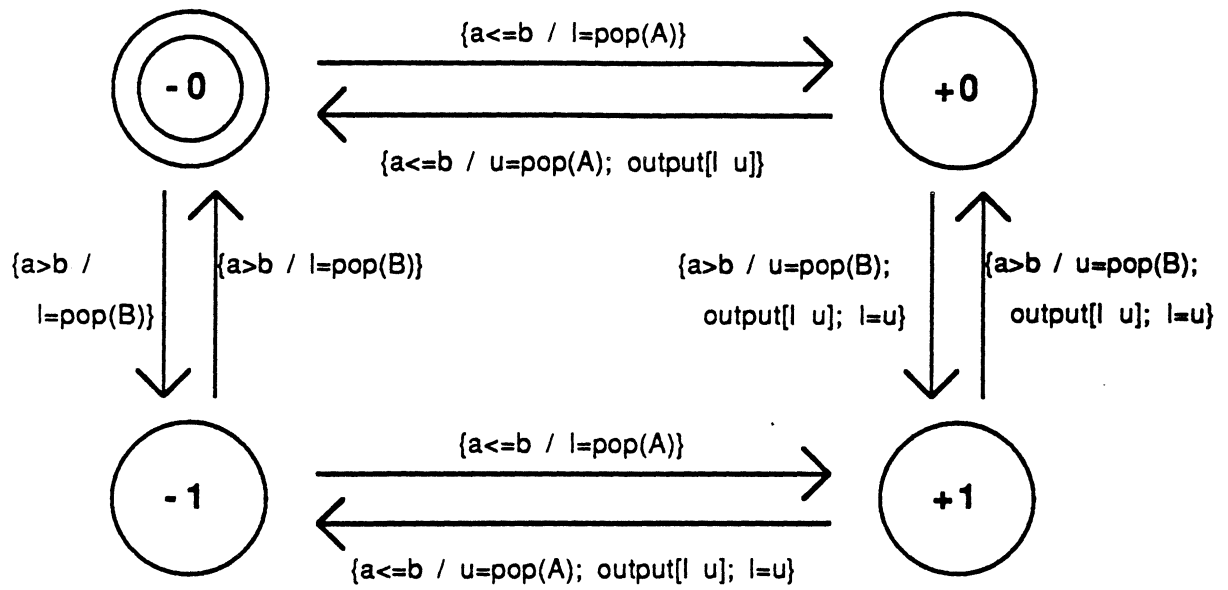


Figure 2: State Transition to Compute the New Bounds for Difference.

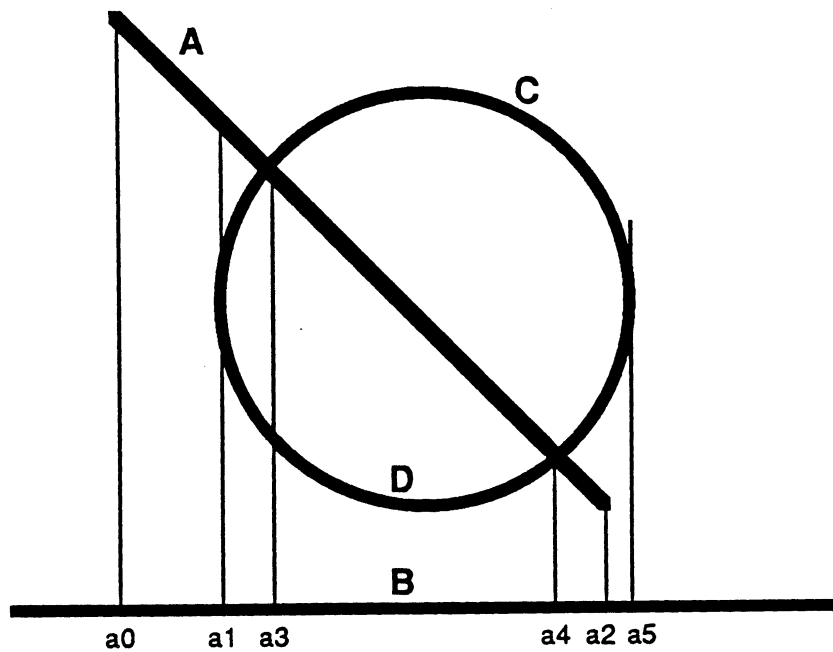


Figure 3: Refine a Pair of Bounds by Computing Intersection Points.

Consider the example of Figure 3. A cone and a torus are combined ( either *union* or *difference*). The cone is specified by lines A and B, and the torus is specified by the upper half circle C and lower half circle D within the bounds  $[a_1 a_2]$ . Since  $[a_1 a_2]$  obtained after Step 1 is both a subinterval of the  $[a_0 a_2]$  of the cone and of the  $[a_1 a_5]$  of the torus, we compute the intersection points for the curves A, B, C, D. Two intersection points  $a_2$  and  $a_3$  are obtained, then we create three refined bound pairs  $[a_1 a_3]$ ,  $[a_3 a_4]$ , and  $[a_4 a_2]$  to replace the  $[a_1 a_2]$ . The curve shape within these three refined intervals can then be uniquely determined, which depends on the operation to be performed. We will discuss it in the next step of the algorithm.

### Step 3. Apply the Union or Difference Operation

In this stage, the resulting curve shape within each pair of bounds can be determined. For those pairs of bounds created from Step 1 but not refined by Step2, that is, they are subintervals that belong to only one original pair of bounds, the curves within the original pair of bounds are directly copied into the newly created bound pair with the two ending points of the curves adjusted so as to be consistent with the new bounds. This works for the union operation. For the difference operation, the bound pairs contributed solely by the second operand are simply thrown away, but those bound pairs from the first operand should be kept and the curves defined within these bound pairs should also be copied.

For those refined bound pairs from Step 2, it is assured that there is no curve intersection within them, thus from Theorem 1, the curves within them form a total ordering. This property of total ordering together with the concepts of *maximum union-able* and *minimum difference-able layer sets* is useful for determining which curves from the original ones contribute to the resulting curves.

Here we use the concept of layer to determine the resulting curves. A layer is an area bounded by two curves within some interval. Because the curves of layers of the two operands within the bound pair form a total ordering, they are topologically equivalent to lines which are parallel to the principal axis within that bound pair. At any point within the bound pair, a line, passing this point and perpendicular to the principal axis, intersects all

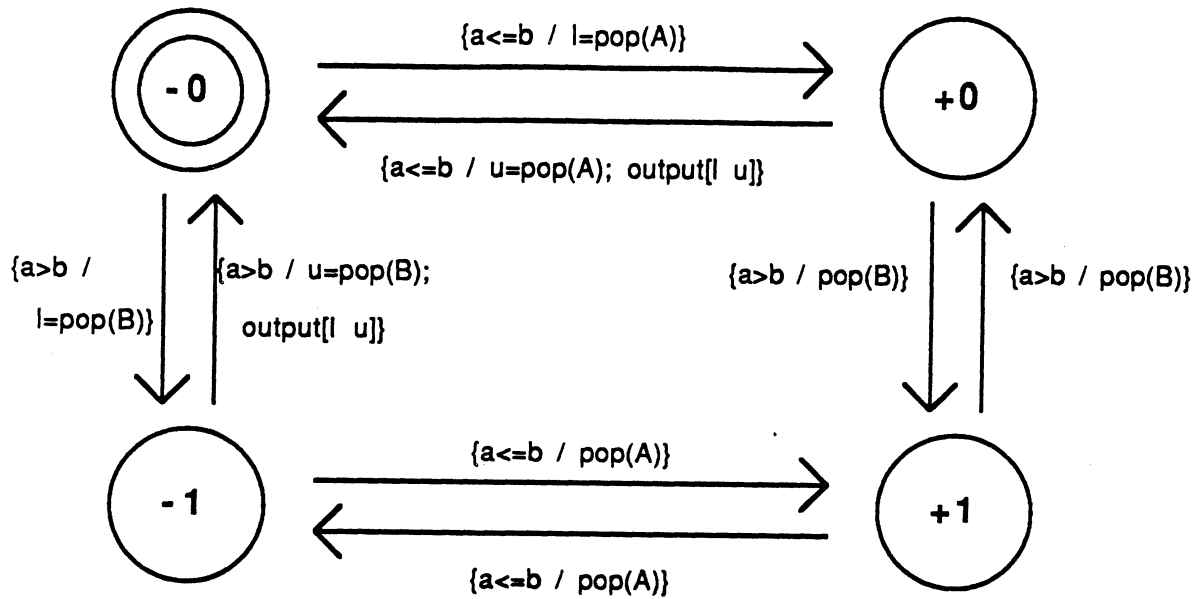


Figure 4: State Transition Diagram for Union.

the curves within that bound pair. Layers can then be represented by segments (determined by these intersection points) on this line. If the intersection point of this line and the principal axis is assumed to be the zero point, then the coordinates of the original curves and those of segments for layers can thus be uniquely determined in this line coordinate.

By using this line coordinate system, computing the union or difference of the line segments on this line is a simple task. It is similar to the procedure of merge-and-sort in Step 1. The coordinates of the line segments (for layers), which are kept sorted, from two operands are compared one by one, and depending on the operation performed (either *union* or *difference*), the resulting curves and layers can be determined. The following state transition diagrams (Figure 4 and Figure 5) demonstrate this algorithm. The meaning of the states and symbols is very similar to that in Figure 1 and Figure 2 except that the concept of bound pairs is replaced by the concept of layers and the values of the bounds are replaced by the coordinates of the curves in the line coordinate system.

Consider again the example of Figure 3. We assume the cone is operated with the torus, the cone is the first operand (i.e. left subtree) and the torus is the second operand (i.e. right subtree). Within the bounds  $[a_1, a_3]$ , lines A and B form a layer for the cone and curves C and D form another layer for the torus. Since the layer bounded by A and B

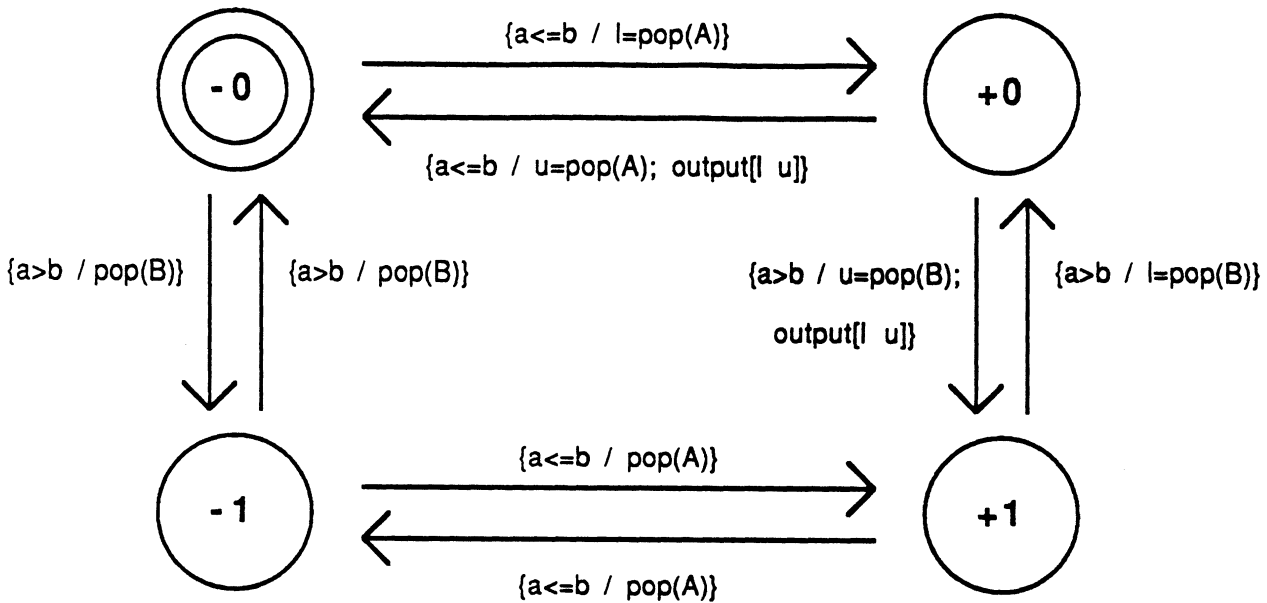


Figure 5: State Transition Diagram for Difference.

covers that bounded by C and D (this can be checked by their line coordinates), if the union operation is performed, the resulting layer would be that bounded by A and B. However, if the difference operation is executed, the results would be the layer bounded by A and C, and the layer bounded by D and B.

Similarly, within the bounds  $[a_3 a_4]$ , the result of the union operation is the layer bounded by the curves C and B, and the result of the difference operation is the layer bounded by D and B. Within the bounds  $[a_4 a_2]$ , the result of the union operation is the layer bounded by C and D and the layer bounded by A and B. The result of the difference operation would be the layer bounded by A and B.

**Step 4. Merge Neighboring Segments If Possible**

After the processing of the previous steps, the resulting shape of the combined object is obtained. However, it might not be in the PAR form because the curves at one of the two bound points might be differentiable. This may happen if some curves are separated due to the bound pair refinement and later only those separated subcurves are saved to be the resulting curves by the selection process of Step 3. In this case, the bound point at which all curves are differentiable is not a really breakpoint and it does not satisfy the

non-differentiable requirement of the PAR definition at the bound points. Therefore further processing is necessary to make the result a PAR.

The actual work in this step is to go through the link list and check if the current segment is possible to be merged with its neighboring segments. The conditions for merging two neighboring segments are that

1. two neighboring segments have one common bound point
2. all the curves in one segment are one-to-one connected to all curves in another segment at the common bound point.
3. the two curves from both segments which are connected at the common point must have the common curve type and must belong to (i.e. subcurves) a curve whose curve type is identical to the common curve type.
4. two curves which satisfy condition (3) must be differentiable at the common bound point.

After testing the merge condition, if two neighboring segments are mergeable, a segment is deleted from the list and the bound point of the extant segment is extended to cover the range of the deleted segment. Since all the curves in the original two segments are one-to-one correspondance and the corresponding curves have the same curve parameters, nothing the bounds of the curves need to be updated.

This step makes sure that the result of operating two PAR's is still a PAR. It is very important because it assures that, from Theorem 6, PAR is a unique representational scheme. The uniqueness property of a representational scheme is very useful because it minimizes the redundancy of storing data and eliminates the concern of data inconsistency in a database. Furthermore, it saves programmers' effort of writing code to analyze each of the possible representations for the same object and makes the access of objects in the database more efficient.



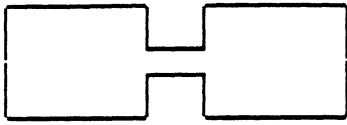


Figure 6.1

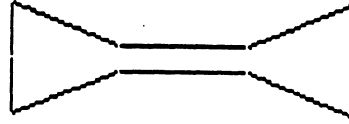


Figure 6.2

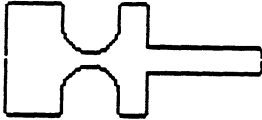


Figure 6.3

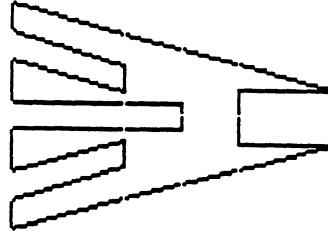


Figure 6.4

Figure 6: Examples.

## 5 Illustrative Examples

In this section, the testing results of several examples by our program are shown in Figure 6 and Figure 7. This program implements the algorithms described in Section 4. Also implemented are the algorithms to compute the length and maximum diameter of a symmetrical part and to compute and plot its profile from PAR. The whole program code, which has more than 2,500 lines Pascal code on Apollo computer, can be found in Appendix 1, while Appendix 2 shows the input data for the example "bottle" of Figure 7. Appendix 3 shows all the pictures of the machine parts in the example of Figure 6 and 7. These pictures are generated from a *ray tracing* program with CSG trees as input data.

Figure 6.1 shows the profile of a machine part which is joined together from three cylinders of different sizes. Some part of the middle cylinder originally overlapping with the other two is removed after processing, therefore it is not seen from the profile. Figure 6.2 shows the profile of a machine part which is unioned from two cones and a cylinder. Notice that the two cones are in opposite directions, that is, one of them is rotated  $180^\circ$  with respect to  $y$ -axis from its original orientation. Figure 6.3 is a profile of a neck. The neck is made from a cylinder, cut off a pipe and further a torus from its outer surface. In Figure 6.4, a pipe and two holes are drilled from a cone. The computation for this object

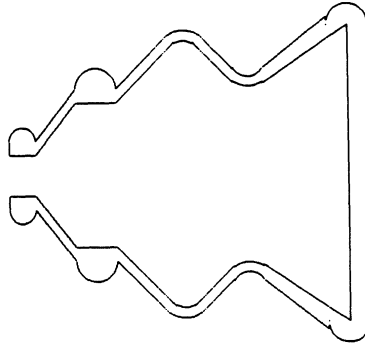


Figure 7: Example of Bottle.

involves many *layer difference* operations. Its profile clearly displays the layer information.

Figure 7 displays the profile of a “bottle”. Its length and diameter are also computed by the program (Length = 5.30 units, Diameter = 4.905 units). This “bottle” is composed of 27 primitive solids involving 8 cones, 12 cylinders and 7 tori (see Appendix 2, *primitive* file). Its CSG tree contains 53 nodes (*object* file). Since there is no movement node used in this example, its *movement* file is empty. Many *principal axis union* and *difference* operations as well as *layer union* and *layer difference* operations are involved in converting the CSG tree into its PAR representation. Its length, diameter information and profile are then computed from this PAR.

To compute a profile of a machine part, its algorithm may take advantages of the PAR representation. Since the layer information which consists of its outer, inner curves, and starting and ending points exists in the PAR representation, the profile of a machine part can be generated by drawing the four boundary curves (two vertical lines for the starting and ending line boundaries, and two for the inner and outer curves) for each layer in an *exclusive* or plotting mode. In such mode, the common boundary of two layers will disappear, rather than being plotted twice.

This algorithm is efficient and easy to implement, it has some problems, however. It

relies on the *exclusive or* operation to eliminate the common boundary of two layers, but the *exclusive or* operation also eliminates the conner points of actual boundary because the intersection point of two boundary curves is plotted twice. In general, a *right-turn* algorithm [WeA77] can be used to solve this problem. The vertices of layers and their relations in the *principal axis coordinate system* can be easily obtained from the PAR. Then the *right-turn* algorithm can be applied to these vertices to get the profile. At this time, we only implement the profile computation using the *exclusive or* algorithm. Close examination of Figure 6 and Figure 7 reveals the fact of the missing common corner points.

## 6 Conclusion and Discussion

In this section, some concluding remarks are made first in Section 6.1, then several issues about PAR and its future extension are discussed in Section 6.2.

### 6.1 Conclusion

As a new representational scheme toward *representational uniqueness*, the Principal Axis Representation (PAR) is developed for *axis symmetrical* machine parts. Based on the CSG, the machine parts can be constructed by applying the *regularized set operations union, difference* and *movement* on the primitive solids *cylinder, cone* and *torus*. The PAR can represent the machine parts composed of primitives that have the same *Principal Axis*.

The PAR is first defined in this report. The operations for operating ( combining) PAR's : *union, difference* and *movement* are then defined. Also shown is how to use PAR to represent primitive solids: cylinder, cone and torus. Based on this formulation, several theorems are proved, which show that a CSG tree can be converted into a PAR representing the same axis-symmetrical object.

An algorithm based on the mathematical formulation of PAR to convert a CSG tree has been designed and implemented in this study. It first converts the CSG leaf nodes into their PAR representation and then operates these PAR's (using *union, difference, movement* operations of PAR ) on the CSG tree from bottom to top until the root of the CSG tree is

visited. Several examples and testing results of this algorithm are shown in Section 5.

To support CAD applications, for example, generating the profile of machine parts for Numerical Control and computing the geometrical properties of a machine part such as its length and maximum diameter for part classification and CAPP (Computer Aided Process Planning), the PAR seems to be more efficient than its counterpart: the CSG scheme. This is due to the fact that, to represent an object, the CSG tree is left unevaluated while the PAR is already the result after evaluation. Algorithms that compute the profile and the length as well as diameter of a machine part from PAR have also been implemented. Its results are shown in Section 5.

In addition to its computational efficiency, the PAR is proved to possess yet another important property as a representational scheme, that is, representational uniqueness. A unique representation scheme allows much simpler *feature definition* and therefore *feature extraction* or *object recognition* because only one representation for a feature or an object is required to deal with [LeK86].

## 6.2 Discussion and Future Work

The key ideas of PAR are to represent an object by its principal axis and boundary curves, and to resolve the overlap (intersection) of two composite solids, i.e. to compute the *union* and *difference* of layers. At present, PAR works for axis-symmetrical machine parts that are constructed from primitives such as cylinder, cone and torus. To make it more general as a unique representational scheme, several extensions are currently under investigation:

- relax the assumption of common principal axis requirement. Based on the PAR, a hierarchical PAR could be defined as a set of PAR's, each having its own common principal axis. But the problem of how to characterize the overlap parts of two PAR's during *union* and *difference* operations to ensure that the generalized PAR is also a uniqueness representation needs to be studied.
- include more primitives like *sphere* and *cube*. In fact, *sphere* is a special case of *torus*. But naively integrating *sphere* into PAR would invalidate the uniqueness property

of the PAR because a *sphere* has two different representations in PAR, i.e. using *sphere* primitive or *torus* primitive. Including *cube* into PAR relaxes the constraint of axis-symmetry to a large extent. How to characterize the boundary curves, however, requires further study.

To make the PAR more useful, several areas of application are also under study:

- support feature extraction and object recognition. Since the PAR is representational uniqueness, object features such as *round*, *fillet*, *keyway*, *hole* can be defined on the PAR more easily than on the CSG tree. Developing algorithms to extract features or recognize object based on the PAR should be simpler because less cases need to be analyzed.
- support CAD (Computer Aided Design) applications. In this report, we have shown the profile generation of machine parts for NC (Numerical Control) and the length and diameter computation for machine parts to support part classification and CAPP (Computer Aided Process Planning). Other geometrical properties and geometrical codes [KaO84] could also be computed easily from the PAR. We are working on exploit it now.

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***Appendix 1. Program Listing***



```

1 (*****
2 (*
3 (* This program implements how to convert a CSG
4 (* input tree into a PAR (Principal Axis Representa-
5 (* tion) that represents the identical object as the
6 (* original CSG tree does.
7 (* All the algorithms and procedures implemented
8 (* here are based on the definitions of PAR, its
9 (* union, difference and movement operations.
10 (* Geometrical properties such as the length and
11 (* diameter of objects are computed from the result-
12 (* ing PAR to support part classification and process
13 (* planning. Profile is also generated to support
14 (* Numerical Control application.
15 (*
16 (* Author : K. F. Jack Jea
17 (* Date : 6/12/1986
18 (*
19 (*****
20
21 PROGRAM AXIS( input, output );
22
23 $nolist;
24 $include '/sys/ins/base.ins.pas';
25 $include '/sys/ins/error.ins.pas';
26 $include '/sys/ins/gpr.ins.pas';
27 $include '/sys/ins/time.ins.pas';
28 $list;
29
30 CONST
31   { part number of the csg tree root }
32   CSG_ROOT = -1,
33   PI = 3.1415926;
34
35 TYPE
36
37   { normalized parametric value on the axis }
38   t_value = real;
39
40   { value of vertical axis on principal axis coordinate }
41   Y_value = real;
42
43   { disparity of directions of two axes }
44   direction = (same_dir, opposite_dir, non_co_linear);
45
46   { basic primitive solids currently implemented }
47   solid_kind = ( cylinder, cone, torus );
48
49   { node types of the CSG tree }
50   node_kind = ( primitive, movement, union_op, diff_op );
51
52   { axis operation type of two principal axes }
53   axis_op_kind = ( union, difference );
54
55   { lower or upper bounds }
56   lower_or_upper = ( lower, upper );

```

```

57
58   { upper or lower part of a circle }
59   up_or_down = ( up, down );
60
61   { curve type }
62   curve_kind = ( line, arc );
63
64   { possible states when two axes are operated }
65   state = ( pluszero, minuszero, plusone, minusone );
66
67   { point coordinate on the principal axis coordinate }
68   D2_point = record
69     x_coord : t_value;
70     Y_coord : Y_value;
71   end;
72
73   { actual coordinate in the 3-D space }
74   D3_point = record
75     x_coord,
76     Y_coord,
77     z_coord : real;
78   end;
79
80   { chain to link intersection points of two axes }
81   intersect_rec_ptr = ^intersect_rec;
82
83   { record for storing the intersection points }
84   intersect_rec = record
85     intersection : t_value;
86     next_intersect : intersect_rec_ptr;
87   end;
88
89   { record to store a bounded curve }
90   curve = record
91     curve_start,
92     curve_end : Y_value ;
93     case curve_type : curve_kind of
94       line : ( );
95       arc : (
96         center : D2_point;
97         radius : real;
98         which_half : up_or_down );
99     end;
100
101   { pointer to a bounded curve }
102   curve_ptr = ^curve;
103
104   { pointer to a layer }
105   layer_ptr = ^layer;
106
107   { a layer which is bounded by two curves }
108   layer = record
109     next_layer : layer_ptr;
110     inner_curve,
111     outer_curve : curve_ptr;
112   end;

```

```

113 { pointer to a segment of the axis }
114 segment_ptr = ^segment;
115
116 { segment record, part of axis }
117 segment
118 = record
119   lower_bound,
120   upper_bound : t_value;
121   next_segment : segment_ptr;
122   layer_head : layer_ptr;
123   end;
124
125 { information of the principal axis }
126 principal_axis = record
127   start_point,
128   end_point : D3_point;
129   segment_head : segment_ptr;
130   end;
131
132 { transformation matrix for movement }
133 xform_matrix = record
134   translate_x,
135   translate_y,
136   translate_z : real;
137   rotate_x,
138   rotate_y,
139   rotate_z : real;
140   end;
141
142 { information of a primitive solid }
143 solid = record
144   base, top : D3_point;
145   case solid_type : solid_kind of
146     cylinder : ( radius : real );
147     cone : ( radius_b,
148             radius_t : real );
149     torus : ( inner_radius,
150             outer_radius : real );
151   end; { solid }
152
153 { pointer to a CSG tree }
154 CSG_tree_ptr = ^CSG_tree;
155
156 { CSG tree node information }
157 CSG_tree = record
158   node_type : node_kind;
159   partno : integer;
160   case node_kind of
161     primitive : ( prim_solid : solid );
162     movement : ( move : xform_matrix;
163               child : CSG_tree_ptr );
164     union_op,
165     diff_op : ( left_child,
166               right_child : CSG_tree_ptr );
167   end;
168

```

```

169 { global variables for main program use }
170 VAR
171   { original input CSG tree }
172   CSG : CSG_tree_ptr;
173   { resulting PAR corresponding to CSG above }
174   P : principal_axis;
175   { length and maximum diameter of the object }
176   L, D : real;
177   { input data for creating the CSG }
178   obj_file, mov_file, prim_file : text;
179   { starting and ending points of a combined axis }
180   min_point, max_point : D3_point;
181
182   { switches for debugging use }
183   DEBUG_INPUT : boolean;
184   DEBUG_EVALUATE_CSG : boolean;
185   DEBUG_BUILD_AXIS : boolean;
186   DEBUG_ID : boolean;
187   DEBUG_DRAW : boolean;
188   DEBUG_COMPUTE_T : boolean;
189   DEBUG_CURVE_CURVE, DEBUG_INTERSECT : boolean;
190   DEBUG_AXIS_OP, DEBUG_NORMALIZATION,
191   DEBUG_PARTITION, DEBUG_MERGE_INTERVAL : boolean;
192   DEBUG_LAYER_DIFFERENCE, DEBUG_COMPUTE_LAYER,
193   DEBUG_ADD_LAYER : boolean;
194   TRACE_AXIS : boolean;
195
196 procedure dump_axis( PX : principal_axis );
197 { dump the principal axis; for debugging purpose }
198 var
199   s : segment_ptr;
200   l : layer_ptr;
201 begin
202   writeln('@@@@@@@@ Dump Principal Axis @@@@@@');
203   writeln('start point =', PX.start_point.x_coord,
204         PX.start_point.y_coord, PX.start_point.z_coord);
205   writeln('end point =', PX.end_point.x_coord,
206         PX.end_point.y_coord, PX.end_point.z_coord);
207
208 { go thru the axis and print its content }
209 s := PX.segment_head;
210 while s <> nil do
211 begin
212   writeln('lower and upper bounds=',
213         s.lower_bound, s.upper_bound);
214
215   l := s.layer_head;
216   while l <> nil do
217 begin
218   writeln('outer curve, start and end y-value=',
219         l.outer_curve.curve_start,
220         l.outer_curve.curve_end);
221   writeln('inner curve, start and end y-value=',
222         l.inner_curve.curve_start,
223         l.inner_curve.curve_end);
224

```

```

225 l^.inner_curve^.curve_end);
226
227 l := l^.next_layer;
228 end;
229
230 s := s^.next_segment;
231 end;
232
233 writeln('@@@@@@ End of Dumping Principal Axis @@@@@@@');
234
235 end; { dump_axis }
236
237
238 function compute_t ( pt1, pt2, point : D3_point ) : t_value;
239 { calculate the t parameter for point in line segment
240 bounded by pt1 and pt2. Note: we do not check the linearity
241 of point, pt1 and pt2 in this version, it will be refined in
242 later version. }
243 const
244 { numerical error range }
245 EPS = 0.0001;
246 var
247 t : t_value;
248 begin
249
250 if DEBUG_COMPUTE_T then writeln(
251 '***** In compute_t, pt1, pt2, point=',
252 pt1.x_coord, pt1.y_coord, pt1.z_coord,
253 pt2.x_coord, pt2.y_coord, pt2.z_coord,
254 point.x_coord, point.y_coord, point.z_coord );
255
256 if abs(pt1.x_coord - pt2.x_coord) >= EPS then
257 t := (point.x_coord - pt1.x_coord) /
258 (pt2.x_coord - pt1.x_coord)
259 else if abs(pt1.y_coord - pt2.y_coord) >= EPS then
260 t := (point.y_coord - pt1.y_coord) /
261 (pt2.y_coord - pt1.y_coord)
262 else if abs(pt1.z_coord - pt2.z_coord) >= EPS then
263 t := (point.z_coord - pt1.z_coord) /
264 (pt2.z_coord - pt1.z_coord)
265 else writeln( 'ERROR: pt1 and pt2 are the same point',
266 'in compute_t procedure' );
267
268 if DEBUG_COMPUTE_T then writeln('***** In compute_t, t =',
269 , t);
270
271 compute_t := t;
272 end; { compute_t }
273
274 procedure modularize_t ( var t : t_value;
275 fl, f2, ml, m2 : t_value );
276 { actually modularize t here;
277 t' = ml + ( t - fl ) * ( m2 - ml ) / ( f2 - fl ) }
278
279 begin

```

```

281 t := ml + ( t - fl ) * ( m2 - ml ) / ( f2 - fl );
282 end; { modularize_t }
283
284 procedure adjust_t_value( var P : principal_axis;
285 fl, f2, ml, m2 : t_value );
286 { This routine use ml and m2 as the factors to adjust the
287 t parameters in the principal axis P; formula :
288 t' = ml + ( t - fl ) * ( m2 - ml ) / ( f2 - fl ) }
289 var
290 s : segment_ptr;
291 l : layer_ptr;
292
293 begin
294
295 { go thru each bound and t-value dependent item. }
296 s := P.segment_head;
297 while ( s <> nil ) do
298 begin
299
300 modularize_t ( s^.lower_bound, fl, f2, ml, m2 );
301 modularize_t ( s^.upper_bound, fl, f2, ml, m2 );
302
303 l := s^.layer_head;
304 while ( l <> nil ) do
305 begin
306 if l^.inner_curve^.curve_type = arc then
307 modularize_t (
308 l^.inner_curve^.center.x_coord,
309 fl, f2, ml, m2 );
310 if l^.outer_curve^.curve_type = arc then
311 modularize_t (
312 l^.outer_curve^.center.x_coord,
313 fl, f2, ml, m2 );
314 l := l^.next_layer;
315 end; { inner while }
316 s := s^.next_segment;
317 end; { while }
318 end; { adjust_t_value }
319
320 function check_direction( P1, P2 : principal_axis ) : direction;
321 { check to see if P1 and P2 are in the same, opposite or
322 non-co-linear direction }
323 const
324 EPS = 0.0001;
325 var
326 xl, yl, z1, x2, y2, z2, dl, d2 : real;
327
328 begin
329 { get three vector components of P1 and P2 }
330 xl := P1.start_point.x_coord - P1.end_point.x_coord;
331 yl := P1.start_point.y_coord - P1.end_point.y_coord;
332 z1 := P1.start_point.z_coord - P1.end_point.z_coord;
333 dl := sqrt( xl * xl + yl * yl + z1 * z1 );

```

```

337 x1 := x1 / dl;
338 y1 := y1 / dl;
339 z1 := z1 / dl;
340
341 x2 := P2.start_point.x_coord - P2.end_point.x_coord;
342 y2 := P2.start_point.y_coord - P2.end_point.y_coord;
343 z2 := P2.start_point.z_coord - P2.end_point.z_coord;
344 d2 := sqrt( x2 * x2 + y2 * y2 + z2 * z2 );
345 x2 := x2 / d2;
346 y2 := y2 / d2;
347 z2 := z2 / d2;
348
349 { check linearity }
350 if ( abs(x1 - x2) <= EPS) and ( abs(y1 - y2) <= EPS) and
351 ( abs(z1 - z2) <= EPS) then
352   check_direction := same_dir
353 else if ( abs(x1 + x2) <= EPS) and ( abs(y1 + y2) <= EPS)
354   and ( abs(z1 + z2) <= EPS) then
355   check_direction := opposite_dir
356 else
357   check_direction := non_collinear;
358 end; { check_direction }
359
360 procedure reverse_direction( var P: principal_axis );
361 { adjust the t-parameters in P; make all the t-value reverse }
362 var
363   s, ps, ns : segment_ptr;
364   l : layer_ptr;
365   pt : D3_point;
366   t : t_value;
367   y : Y_value;
368 begin
369
370   pt := P.start_point;
371   P.start_point := P.end_point;
372   P.end_point := pt;
373
374   s := P.segment_head;
375   ps := nil;
376   while ( s <> nil ) do
377     begin
378       t := s.lower_bound ;
379       s.lower_bound := 1.0 - s.upper_bound;
380       s.upper_bound := 1.0 - t;
381
382       l := s.layer_head;
383       while ( l <> nil ) do
384         begin
385           y := l.inner_curve.curve_start;
386           l.inner_curve.curve_start :=
387             l.inner_curve.curve_end;
388           l := l.inner_curve;
389           l.inner_curve.curve_end := y;
390           l := l.inner_curve;
391         end
392       if l.inner_curve.curve_type = arc then

```

```

393   l.inner_curve.center.x_coord := 1.0 -
394     l.inner_curve.center.x_coord;
395
396   y := l.outer_curve.curve_start;
397   l.outer_curve.curve_start :=
398     l.outer_curve.curve_end;
399   l.outer_curve.curve_end := y;
400
401   if l.outer_curve.curve_type = arc then
402     l.outer_curve.center.x_coord := 1.0 -
403       l.outer_curve.center.x_coord;
404
405   l := l.next_layer;
406 end; { inner while }
407
408 ns := s.next_segment;
409 s.next_segment := ps;
410 ps := s;
411 s := ns;
412 end; { while }
413
414 P.segment_head := ps;
415 end; { reverse_direction }
416
417 procedure normalization ( var P1, P2 : principal_axis;
418   var pt1, pt2 : D3_point );
419 { normalize both P1 and P2 to have the same t-parameter
420   basis; note: after difference operation, the principal
421   axis should re-normalized. }
422 var
423   t1, t2 : t_value;
424   m1, m2 : real;
425   dir : direction;
426 begin
427   if DEBUG NORMALIZATION then writeln(
428     '*** before normalization, P1.start and end point=',
429     P1.start_point.x_coord, P1.start_point.y_coord,
430     P1.start_point.z_coord, P1.end_point.x_coord,
431     P1.end_point.y_coord, P1.end_point.z_coord);
432   if DEBUG NORMALIZATION then writeln(
433     '*** before normalization, P2.start and end point=',
434     P2.start_point.x_coord, P2.start_point.y_coord,
435     P2.start_point.z_coord, P2.end_point.x_coord,
436     P2.end_point.y_coord, P2.end_point.z_coord);
437   if DEBUG NORMALIZATION then writeln(
438     '*** before normalization, P2.low, up =',
439     P2.segment_head.lower_bound,
440     P2.segment_head.upper_bound);
441   { check the direction of two axes, ok if the same,
442     inverse one if in reverse direction, reject it

```

```

449   if not co_linear ]
450   dir := check_direction( P1, P2 );
451   case dir of
452     same_dir : { ok; go thru it }
453     if DEBUG_NORMALIZATION then writeln(
454       '##### In normalization, '
455       'two axes same direction' );
456     opposite_dir : if (P1.start_point.z_coord >
457       P1.end_point.z_coord) then
458       reverse_direction( P1 )
459     else
460       reverse_direction( P2 );
461     non_co_linear : writeln('----> error, '
462       ' two axes do not co-linear');
463   end; { case }
464
465   { compute the new starting and ending points }
466   t1 := compute_t( P1.start_point, P1.end_point,
467     P2.start_point );
468   t2 := compute_t( P1.start_point, P1.end_point,
469     P2.end_point );
470
471   if DEBUG_NORMALIZATION then writeln(
472     '##### IN normalization, t1, t2 =', t1, t2 );
473
474   if 0.0 <= t1 then pt1 := P1.start_point
475     else pt1 := P2.start_point;
476   if 1.0 >= t2 then pt2 := P1.end_point
477     else pt2 := P2.end_point;
478
479   { adjust P1 and P2 using the new starting and
480     ending points }
481   m1 := compute_t( pt1, pt2, P1.start_point );
482   m2 := compute_t( pt1, pt2, P1.end_point );
483   adjust_t_value( P1, 0.0, 1.0, m1, m2 );
484
485   m1 := compute_t( pt1, pt2, P2.start_point );
486   m2 := compute_t( pt1, pt2, P2.end_point );
487   adjust_t_value( P2, 0.0, 1.0, m1, m2 );
488
489   if DEBUG_NORMALIZATION then writeln(
490     '##### after normalization, P2.low, up =',
491     P2.segment_head.lower_bound,
492     P2.segment_head.upper_bound);
493
494   if DEBUG_NORMALIZATION then writeln(
495     '##### result of normalization, P1.start and end point=',
496     P1.start_point.x_coord, P1.start_point.y_coord,
497     P1.start_point.z_coord, P1.end_point.x_coord,
498     P1.end_point.y_coord, P1.end_point.z_coord);
499   end; { normalization }
500
501   function get_curve( var p : layer_ptr;
502     var flag : lower_or_upper ) : curve_ptr;
503
504   { get next curve from the layers; if none, then return nil.

```

```

505   this routine is somewhat like get_next function, but it
506   will be used by compute_intersection and compute_layer. }
507   begin
508
509     case flag of
510       upper : if p = nil then
511         get_curve := nil
512       else begin
513         flag := lower;
514         get_curve := p.outer_curve ;
515       end;
516     lower : begin
517       flag := upper;
518       get_curve := p.inner_curve ;
519       p := p.next_layer;
520     end;
521   end; { case }
522
523   end; { get_curve }
524
525   procedure line_line( C1, C2 : curve_ptr;
526     S1, S2 : segment_ptr;
527     a, b : t_value;
528     var lptr : intersect_rec_ptr );
529   { compute the intersection of two lines C1 and C2;
530     results are appended to lptr. }
531   label
532   R;
533   const
534     SLOPE_EPSILON = 0.001;
535   var
536     x1, x2, x3, x4, dx, t1, x : t_value;
537     y1, y2, y3, y4, dy : real;
538     p : intersect_rec_ptr;
539   begin
540
541     { get the first line segment }
542     x1 := S1.lower_bound;
543     x2 := S1.upper_bound;
544     y1 := C1.curve_start;
545     y2 := C1.curve_end;
546
547     { get the second line segment }
548     x3 := S2.lower_bound;
549     x4 := S2.upper_bound;
550     y3 := C2.curve_start;
551     y4 := C2.curve_end;
552
553     if DEBUG_CURVE_CURVE then
554       writeln('$$$$$ Enter In line_line, '
555         'x1,y1,x2,y2,x3,y3,x4,y4 ',
556         x1,y1,x2,y2,x3,y3,x4,y4 );
557
558     if ( ( x1 = x3 ) and ( y1 = y3 ) ) and
559       ( ( x2 = x4 ) and ( y2 = y4 ) ) then

```

```

561   { lines coincidence }
562   begin
563   { return }
564   end
565   else if (x1 = x2) and (x3 = x4) then
566     begin
567     {return}
568     end
569   else begin
570     if (x1 < x2) and (x3 < x4) then
571       begin
572         if abs( (y2 - y1)/(x2 - x1) -
573              (y4 - y3)/(x4 - x3) ) <= SLOPE_EPSILON
574           then begin
575             goto R;
576           end;
577         end;
578         dy := y4 - y3;
579         dx := x4 - x3;
580       if DEBUG_CURVE_CURVE then writeln(
581         '$$$$$$ In line_line, dx, dy, y1, y2 =',
582         dx, dy, y1, y2 );
583       end;
584       t1 := dy * (x3 - x1) - dx * (y3 - y1);
585       t2 := t1 / ( dy * (x2 - x1) - dx * (y2 - y1) );
586       x := x1 + t1 * ( x2 - x1 );
587       if ( a < x ) and ( x < b ) then
588         begin
589           new( p );
590           p^.intersection := x;
591           p^.next_intersect := lptr;
592           lptr := p;
593         if DEBUG_CURVE_CURVE then writeln(
594           '$$$$$$ In line_line, intersection t =',
595           , x );
596         end;
597       end;
598     R :
599     end;
600   end;
601   end; { line_line }
602 end;
603 procedure line_arc( Cl, C2 : curve_ptr;
604                   Sl, S2 : segment_ptr; a, b : t_value;
605                   var lptr : intersect_rec_ptr );
606 { compute the intersection of a line Cl and an arc C2;
607 results are appended to lptr.
608 }
609 var
610   x0, x1, x2, x3, x4, dx, x, length : real;
611   y0, y1, y2, y3, y4, dy, t1, t2, AA, d : real;
612   p : intersect_rec_ptr;
613   theta, ans : array[1..2] of real;
614   num, i : integer;
615   begin

```

```

617   length := sqrt( (max_point.x_coord - min_point.x_coord)
618                 *(max_point.x_coord - min_point.x_coord)
619                 + (max_point.y_coord - min_point.y_coord)
620                 *(max_point.y_coord - min_point.y_coord)
621                 + (max_point.z_coord - min_point.z_coord)
622                 *(max_point.z_coord - min_point.z_coord) );
623
624   { get the first line segment }
625   x1 := Sl^.lower_bound * length;
626   x2 := Sl^.upper_bound * length;
627   y1 := Cl^.curve_start;
628   y2 := Cl^.curve_end;
629
630   { get the second arc segment }
631   x3 := S2^.lower_bound * length;
632   x4 := S2^.upper_bound * length;
633   y3 := C2^.curve_start;
634   y4 := C2^.curve_end;
635   x0 := C2^.center.x_coord * length;
636   y0 := C2^.center.y_coord;
637
638   if DEBUG_CURVE_CURVE then
639     writeln( '$$$$$$ In line_arc, x1, y1, x2, y2',
640             x1, y1, x2, y2 );
641   if DEBUG_CURVE_CURVE then
642     writeln( '$$$$$$ In line_arc, x0, y0, x3, y3, x4, y4',
643             x0, y0, x3, y3, x4, y4 );
644
645   if ( ( x1 = x3 ) and ( y1 = y3 ) ) and
646       ( ( x2 = x4 ) and ( y2 = y4 ) ) then
647     {curves intersect at end point}
648     begin
649       { return }
650     end
651   else begin { compute intersection of line and arc }
652     dx := x2 - x1;
653     dy := y2 - y1;
654     AA := ( (y0 - y1) * dx - (x0 - x1) * dy ) /
655           C2^.radius;
656     d := dx * dx + dy * dy - AA * AA;
657
658     if DEBUG_CURVE_CURVE then writeln(
659       '$$$$$$ In line_arc, dx, dy, AA, d',
660       dx, dy, AA, d );
661     if d >= 0.0 then
662       begin
663         if abs( AA + dy ) <= 0.000001 then
664           begin
665             theta[1] := PI;
666             theta[2] := -PI;
667           end
668         else begin
669           t1 := (-dx + sqrt(d)) / ( AA + dy );

```

```

673 t2 := (-dx - sqrt(d)) / ( AA + dy );
674 theta[1] := 2.0 * arctan( t1 );
675 theta[2] := 2.0 * arctan( t2 );
676 end;
677
678 num := 0;
679 case C2^.which_half of
680   down : { take negative angles }
681     for i := 1 to 2 do
682       if theta[i] <= 0.0 then
683         begin
684           num := num + 1;
685           ans[num] := theta[i];
686         end;
687       up : { take positive angles }
688         for i := 1 to 2 do
689           if theta[i] >= 0.0 then
690             begin
691               num := num + 1;
692               ans[num] := theta[i];
693             end;
694           end; { case }
695         for i := 1 to num do
696           begin
697             x := x0 + cos( ans[i] ) * C2^.radius;
698             if ( a < (x/length) ) and
699               ( x < (b/length) ) then
700               begin
701                 new( p );
702                 p^.intersection := x / length;
703                 p^.next_intersect := lptr;
704                 lptr := p;
705               end;
706             end; { for }
707           end; { if }
708         end; { else }
709       end; { line_arc }
710     end;
711   procedure arc_arc( C1, C2 : curve_ptr;
712     S1, S2 : segment_ptr; a, b : t_value;
713     var lptr : intersect_rec_ptr );
714 { compute the intersection of an arc C1 and an arc C2;
715   results are appended to lptr.
716 }
717 var
718   xcl, xc2, xl, x2, x3, x4, dx, x, length : real;
719   ycl, yc2, yl, y2, y3, y4, dy, tl, t2, AA, d : real;
720   p : intersect_rec_ptr;
721   theta, ans : array[1..2] of real;
722   num, i : integer;
723
724 length := sqrt( (max_point.x_coord - min_point.x_coord)
725   *(max_point.x_coord - min_point.x_coord)

```

```

729   *(max_point.y_coord - min_point.y_coord)
730   *(max_point.y_coord - min_point.y_coord)
731   *(max_point.z_coord - min_point.z_coord)
732   *(max_point.z_coord - min_point.z_coord) );
733
734 { get the first arc segment }
735 xl := S1^.lower_bound * length;
736 x2 := S1^.upper_bound * length;
737 yl := C1^.curve_start;
738 y2 := C1^.curve_end;
739 xcl := C1^.center.x_coord * length;
740 ycl := C1^.center.y_coord;
741
742 { get the second arc segment }
743 x3 := S2^.lower_bound * length;
744 x4 := S2^.upper_bound * length;
745 y3 := C2^.curve_start;
746 y4 := C2^.curve_end;
747 xc2 := C2^.center.x_coord * length;
748 yc2 := C2^.center.y_coord;
749
750 if ( ( xl = x3 ) and ( yl = y3 ) ) and
751   ( ( x2 = x4 ) and ( y2 = y4 ) ) then
752   {curves intersect at end point}
753   begin
754     { return }
755   end
756 else begin { compute intersection of line and arc }
757   dx := xc2 - xcl;
758   dy := yc2 - ycl;
759   if DEBUG_EVALUATE_CSG then writeln(
760     '$$$$ In arc-arc, C2 radius =', C2^.radius );
761   AA := ( C1^.radius * C1^.radius - C2^.radius *
762     C2^.radius - dx * dx - dy * dy ) /
763     ( 2.0 * C2^.radius );
764   d := dx * dx + dy * dy - AA * AA ;
765   if d >= 0.0 then
766     begin
767       d := sqrt( d );
768       if abs( AA + dy ) <= 0.000001 then
769         begin
770           theta[1] := PI;
771           theta[2] := -PI;
772         end
773       else begin
774         t1 := ( dy + d ) / ( AA + dx );
775         t2 := ( dy - d ) / ( AA + dx );
776         theta[1] := 2.0 * arctan( t1 );
777         theta[2] := 2.0 * arctan( t2 );
778       end;
779       num := 0;
780       case C2^.which_half of
781
782
783
784

```

```

785 down : { take negative angles }
786 for i := 1 to 2 do
787   if theta[i] <= 0.0 then
788     begin
789       num := num + 1;
790       ans[num] := theta[i];
791     end;
792   up : { take positive angles }
793   for i := 1 to 2 do
794     if theta[i] >= 0.0 then
795       begin
796         num := num + 1;
797         ans[num] := theta[i];
798       end; { case }
799     for i := 1 to num do
800       begin
801         x := xc2 + cos( ans[i] ) *
802           C2.radius;
803         if (( a < (x/length) ) and
804            ( x < (b/length) )) and
805            (( x1 < x ) and ( x < x2 )) and
806            (( x3 < x ) and ( x < x4 )) then
807           begin
808             new( p );
809             p^.intersection := x / length;
810             p^.next_intersect := Iptr;
811             Iptr := p;
812           end; { for }
813         end; { if }
814       end; { else }
815     end; { arc_arc }
816   procedure compute_intersection( S1, S2 : segment_ptr;
817     var Iptr : intersect_rec_ptr;
818     a, b : t_value);
819   { compute the intersection points of curves in S1 and S2;
820     results are stored in the list pointed by Iptr. }
821   var
822     p1, p2 : layer_ptr;
823     flag1, flag2 : lower_or_upper;
824     curvel, curve2 : curve_ptr;
825   begin
826     if ( S1 = nil ) or ( S2 = nil ) then
827       Iptr := nil
828     else begin
829       p1 := sl^.layer_head;
830       flag1 := upper;
831       curvel := get_curve( p1, flag1 );
832       p2 := sl^.layer_head;
833       flag2 := lower_or_upper;
834       curve2 := get_curve( p2, flag2 );
835       while ( curvel <> nil ) do
836         begin
837           p2 := S2^.layer_head;
838           flag2 := upper;
839           curve2 := get_curve( p2, flag2 );
840           while ( curve2 <> nil ) do
841             begin
842               case curve2^.curve_type of
843                 line : if curve2^.curve_type = line
844                   then
845                     line_line( curvel, curve2,
846                               sl, S2, a, b, Iptr )
847                   else
848                     line_arc( curvel, curve2,
849                               sl, S2, a, b, Iptr );
850                 arc : if curve2^.curve_type = line
851                   then
852                     line_arc( curve2, curvel,
853                               sl, S2, a, b, Iptr )
854                   else
855                     arc_arc( curvel, curve2,
856                               sl, S2, a, b, Iptr );
857                 end; { case }
858               curve2 := get_curve( p2, flag2 );
859             end; { curve2 }
860             curvel := get_curve( p1, flag1 );
861           end; { curvel }
862         end; { compute_intersection }
863       procedure sort_intersection( var Iptr : intersect_rec_ptr );
864       { sort the intersection points is ascending order;
865         linear sort is implemented here. }
866       var
867         p, q, l1, l2 : intersect_rec_ptr;
868         done : boolean;
869       begin
870         if DEBUG_INTERSECT then
871           begin
872             writeln('#### In sort_intersection, begin dump Iptr list');
873             p := Iptr;
874             while ( p <> nil ) do
875               begin
876                 writeln('#### In sort_intersection, intersection = '
877                   & p^.next_intersect;
878                   end;
879                 p := p^.next_intersect;
880             end;
881             writeln('#### In sort_intersection, end dump Iptr list');
882           end;
883         done := false;
884         while ( Iptr <> nil ) do
885           begin
886             if DEBUG_INTERSECT then
887               begin
888                 writeln('#### In sort_intersection, begin dump Iptr list');
889                 p := Iptr;
890                 while ( p <> nil ) do
891                   begin
892                     writeln('#### In sort_intersection, intersection = '
893                       & p^.next_intersect;
894                     end;
895                     p := p^.next_intersect;
896                   end;
897                 writeln('#### In sort_intersection, end dump Iptr list');

```

```

841 while ( curvel <> nil ) do
842   begin
843     p2 := S2^.layer_head;
844     flag2 := upper;
845     curve2 := get_curve( p2, flag2 );
846   while ( curve2 <> nil ) do
847     begin
848       case curve2^.curve_type of
849         line : if curve2^.curve_type = line
850           then
851             line_line( curvel, curve2,
852                       sl, S2, a, b, Iptr )
853           else
854             line_arc( curvel, curve2,
855                       sl, S2, a, b, Iptr );
856         arc : if curve2^.curve_type = line
857           then
858             line_arc( curve2, curvel,
859                       sl, S2, a, b, Iptr )
860           else
861             arc_arc( curvel, curve2,
862                       sl, S2, a, b, Iptr );
863         end; { case }
864       curve2 := get_curve( p2, flag2 );
865     end; { curve2 }
866     curvel := get_curve( p1, flag1 );
867   end; { curvel }
868 end; { compute_intersection }
869 procedure sort_intersection( var Iptr : intersect_rec_ptr );
870 { sort the intersection points is ascending order;
871 linear sort is implemented here. }
872 var
873   p, q, l1, l2 : intersect_rec_ptr;
874   done : boolean;
875 begin
876   if DEBUG_INTERSECT then
877     begin
878       writeln('#### In sort_intersection, begin dump Iptr list');
879       p := Iptr;
880       while ( p <> nil ) do
881         begin
882           writeln('#### In sort_intersection, intersection = '
883             & p^.next_intersect;
884           end;
885           p := p^.next_intersect;
886         end;
887       writeln('#### In sort_intersection, end dump Iptr list');

```



```

897 end;
898
899 if Iptr <> nil then
900   begin
901     p := Iptr;
902     p := p^.next_intersect;
903     while ( p <> nil ) do
904       begin
905         q := p;
906         if q^.intersection <= Iptr^.intersection then
907           begin
908             p := p^.next_intersect;
909             q^.next_intersect := Iptr;
910             Iptr := q;
911           end
912         else begin
913             I2 := Iptr;
914             I1 := I2^.next_intersect;
915             done := false;
916             while ( I1 <> nil ) and ( not done ) do
917               begin
918                 if I1^.intersection >=
919                   q^.intersection then
920                   done := true
921                 else begin
922                     I2 := I1;
923                     I1 := I1^.next_intersect;
924                   end;
925                 end; { while }
926             end;
927             p := p^.next_intersect;
928             q^.next_intersect := I1;
929             I2^.next_intersect := q;
930           end;
931         end; { while }
932       end; { if }
933     end; { sort_intersection }
934 end; { screen_intersection }
935
936 procedure screen_intersection( var Iptr : intersect_rec_ptr;
937   a, b : t_value );
938 { insert a and b into the intersection list and make sure that
939 this list is bounded by a and b, and there are no duplicate
940 intersection points there. }
941 var
942   P, q, I1, I2 : intersect_rec_ptr;
943 begin
944   new( I1 );
945   I1^.intersection := a;
946   I1^.next_intersect := Iptr;
947   new( I2 );
948   I2^.intersection := b;
949   I2^.next_intersect := nil;
950   p := Iptr;
951   Iptr := I1;

```

```

953   q := I1;
954   while ( p <> nil ) do
955     begin
956       if ( p^.intersection <= I1^.intersection ) or
957         ( p^.intersection = q^.intersection ) then
958         begin
959           p := p^.next_intersect;
960           q^.next_intersect := p;
961         end
962       else if p^.intersection >= I2^.intersection then
963         begin
964           p := nil;
965         end
966       else begin
967           q := p;
968           p := p^.next_intersect;
969         end; { while }
970       q^.next_intersect := I2;
971     end; { screen_intersection }
972 end;
973
974 function compute_point( t : t_value; C : curve_ptr;
975   x1, x2 : t_value ) : y_value;
976 { given a parameter t, compute the value associated with t
977 on the curve C; i.e. C(t) }
978 var
979   m, length : real;
980   y : y_value;
981 begin
982   case C^.curve_type of
983     line : begin
984       m := ( C^.curve_end - C^.curve_start ) /
985         ( x2 - x1 );
986       y := C^.curve_start + m * ( t - x1 );
987       compute_point := y ;
988     end;
989     arc : begin
990       length := sqrt(
991         (max_point.x_coord - min_point.x_coord)
992         * (max_point.x_coord - min_point.x_coord)
993         + (max_point.y_coord - min_point.y_coord)
994         * (max_point.y_coord - min_point.y_coord)
995         + (max_point.z_coord - min_point.z_coord)
996         * (max_point.z_coord - min_point.z_coord));
997       m := ( t - C^.center.x_coord ) *
998         ( t - C^.center.x_coord ) *
999         length * length;
1000       m := C^.radius * C^.radius - m;
1001       m := sqrt(abs( m ));
1002       if C^.which_half = up then
1003         y := C^.center.y_coord + m

```

```

1009     else
1010         y := C^.center.y_coord - m;
1011         compute_point := y ;
1012     end;
1013     end; { case }
1014
1015 end; { compute_point }
1016
1017 procedure copy( C, D : curve_ptr );
1018 { copy the curve pointed by C to D; }
1019 begin
1020
1021     D^.curve_type := C^.curve_type;
1022     case C^.curve_type of
1023     line : begin
1024         D^.curve_start := C^.curve_start;
1025         D^.curve_end := C^.curve_end;
1026     end;
1027     arc : begin
1028         D^.curve_start := C^.curve_start;
1029         D^.curve_end := C^.curve_end;
1030         D^.center.x_coord := C^.center.x_coord;
1031         D^.center.y_coord := C^.center.y_coord;
1032         D^.radius := C^.radius;
1033         D^.which_half := C^.which_half;
1034     end;
1035     end; { case }
1036
1037 end; { copy }
1038
1039 function fix_curve( C : curve_ptr; low, upp : t_value;
1040 S : segment_ptr ) : curve_ptr;
1041 { fix the start and end points of the curve C }
1042 var
1043     x1, x2 : t_value;
1044     D : curve_ptr;
1045 begin
1046     new( D );
1047     copy( C, D );
1048     x1 := S^.lower_bound;
1049     x2 := S^.upper_bound;
1050     D^.curve_start := compute_point( low, C, x1, x2 );
1051     D^.curve_end := compute_point( upp, C, x1, x2 );
1052
1053     fix_curve := D ;
1054
1055 end; { fix_curve }
1056
1057 procedure fix_bound( var q : layer_ptr; low, upp : t_value;
1058 S : segment_ptr );
1059 { This is for direct copy of intervals ( segments S ).
1060 go thru the q chain to fix the strat and end points
1061 of the curves. }
1062 var
1063     p, r : layer_ptr;
1064
1065 begin
1066     { must save q (i.e. copy list) before operation }
1067     p := q;
1068     r := nil;
1069     while p <> nil do
1070     begin
1071         if r = nil then
1072             new( q );
1073             r := q;
1074         else begin
1075             new( r^.next_layer );
1076             r := r^.next_layer;
1077         end
1078     end;
1079
1080     r^.inner_curve := fix_curve( p^.inner_curve,
1081 low, upp, S );
1082     r^.outer_curve := fix_curve( p^.outer_curve,
1083 low, upp, S );
1084     r^.next_layer := nil;
1085
1086     p := p^.next_layer;
1087 end; { while }
1088
1089 end; { fix_bound }
1090
1091 function distance_to_axis( C : curve_ptr; low, upp : t_value;
1092 S : segment_ptr ) : real;
1093 { compute the distance from the mid-point of the curve C
1094 to principal axis; if C = nil, (means no data) then a
1095 minimal distance is assumed. }
1096 const
1097     min_distance = -1.0;
1098 begin
1099     if C = nil then
1100         distance_to_axis := min_distance
1101     else begin
1102         distance_to_axis := compute_point(
1103             (low + upp) / 2.0, C,
1104             S^.lower_bound, S^.upper_bound ) ;
1105     end;
1106 end; { distance_to_axis }
1107
1108 function null_layer( a, b : curve_ptr ) : boolean;
1109 { to test if curve a and b form a null layer }
1110 const
1111     EPS = 0.0001;
1112 var
1113     result : boolean;
1114 begin
1115
1116
1117
1118
1119
1120

```

```

1121 result := false;
1122 if ( abs (a^.curve_start - b^.curve_start) <= EPS ) and
1123 ( abs (a^.curve_end - b^.curve_end) <= EPS ) and
1124 ( a^.curve_type = b^.curve_type ) then
1125 begin
1126   case a^.curve_type of
1127     line : result := true;
1128     arc : result := (a^.radius = b^.radius) and
1129           (a^.which_half = b^.which_half) and
1130           (abs(a^.center.x_coord -
1131              b^.center.x_coord) <= EPS) and
1132           (abs(a^.center.y_coord -
1133              b^.center.y_coord) <= EPS);
1134   end; { case }
1135 end;
1136 null_layer := result;
1137 { null_layer }
1138 end; { null_layer }
1139
1140 procedure add_layer( upper_range, lower_range : curve_ptr;
1141                    var q, r : layer_ptr );
1142 { add a new layer bounded by [upper_range lower_range]
1143   to the q chain ; r always points to the head of the q chain. }
1144 begin
1145
1146   if not null_layer( upper_range, lower_range ) then
1147     { null_layer will be removed }
1148     begin
1149
1150       if DEBUG ADD_LAYER then writeln(
1151         '$$$ In add_layer, uppercurve(start,end), ',
1152         'lowercurve(start,end)=' ,
1153         upper_range^.curve_start, upper_range^.curve_end,
1154         lower_range^.curve_start, lower_range^.curve_end );
1155
1156       if q = nil then { the first one }
1157         begin
1158           q^.next_layer := nil;
1159           q^.outer_curve := upper_range;
1160           q^.inner_curve := lower_range;
1161           r := q;
1162         end
1163       else
1164         begin { normal process }
1165           new( r^.next_layer );
1166           r := r^.next_layer;
1167           r^.outer_curve := upper_range;
1168           r^.inner_curve := lower_range;
1169           r^.next_layer := nil;
1170         end;
1171       end;
1172     end;
1173
1174   end; { add_layer }
1175
1176 procedure layer_union ( var q : layer_ptr; low, upp : t_value;

```

```

1177   sl, s2 : segment_ptr );
1178 { actually union of layers specified in sl and s2; both sl and
1179   s2 are not nil; new layers are pointed by q and bounded by
1180   low and upp; }
1181 var
1182   flag1, flag2 : lower_or_upper;
1183   pl, p2, r : layer_ptr;
1184   cl, c2 : curve_ptr;
1185   next_state : state;
1186   dl, d2 : real;
1187   a, b : curve_ptr;
1188 begin
1189   flag1 := upper;
1190   flag2 := upper;
1191   pl := sl^.layer_head;
1192   p2 := s2^.layer_head;
1193   cl := get_curve( pl, flag1 );
1194   c2 := get_curve( p2, flag2 );
1195   dl := distance_to_axis( cl, low, upp, sl );
1196   d2 := distance_to_axis( c2, low, upp, s2 );
1197   q := nil;
1198   next_state := minuszero;
1199
1200 repeat
1201   case next_state of
1202     minuszero :
1203       if dl >= d2 then { + }
1204         begin
1205           a := fix_curve( cl, low, upp, sl );
1206           next_state := pluszero;
1207           cl := get_curve( pl, flag1 );
1208           dl := distance_to_axis( cl, low, upp, sl );
1209         end
1210       else
1211         begin { l }
1212           a := fix_curve( c2, low, upp, s2 );
1213           next_state := minusone;
1214           c2 := get_curve( p2, flag2 );
1215           d2 := distance_to_axis( c2, low, upp, s2 );
1216         end;
1217       pluszero :
1218         if dl >= d2 then { - }
1219           begin
1220             b := fix_curve( cl, low, upp, sl );
1221             add_layer( a, b, q, r );
1222             next_state := minuszero;
1223             cl := get_curve( pl, flag1 );
1224             dl := distance_to_axis( cl, low, upp, sl );
1225           end
1226         else
1227           begin { l }
1228             next_state := plusone;
1229             c2 := get_curve( p2, flag2 );
1230             d2 := distance_to_axis( c2, low, upp, s2 );
1231           end;
1232

```

```

1233 plusone :
1234   if dl >= d2 then [ - ]
1235   begin
1236     next_state := minusone;
1237     cl := get_curve( pl, flag1 );
1238     dl := distance_to_axis( cl, low, upp, sl );
1239   end
1240   else
1241     begin [ 0 ]
1242     next_state := pluszero;
1243     c2 := get_curve( p2, flag2 );
1244     d2 := distance_to_axis( c2, low, upp, s2 );
1245   end;
1246   minusone :
1247   if dl >= d2 then [ + ]
1248   begin
1249     next_state := plusone;
1250     cl := get_curve( pl, flag1 );
1251     dl := distance_to_axis( cl, low, upp, sl );
1252   end
1253   else
1254     begin [ 0 ]
1255     b := fix_curve( c2, low, upp, s2 );
1256     add_layer( a, b, q, r );
1257     next_state := minuszero;
1258     c2 := get_curve( p2, flag2 );
1259     d2 := distance_to_axis( c2, low, upp, s2 );
1260   end;
1261   end; [ case ]
1262   until ( cl = nil ) and ( c2 = nil ) ;
1263 end; [ layer_union ]
1264
1265 procedure layer_difference ( var q : layer_ptr;
1266   low, upp : t_value;
1267   sl, s2 : segment_ptr );
1268 { actually difference of layers specified in sl and s2;
1269 both sl and s2 are not nil; new layers are pointed by q
1270 and bounded by low and upp; }
1271 var
1272   flag1, flag2 : lower_or_upper;
1273   pl, p2, r : layer_ptr;
1274   cl, c2 : curve_ptr;
1275   next_state : state;
1276   dl, d2 : real;
1277   a, b : curve_ptr;
1278 begin
1279   flag1 := upper;
1280   flag2 := upper;
1281   pl := sl^layer_head;
1282   p2 := s2^layer_head;
1283   cl := get_curve( pl, flag1 );
1284   c2 := get_curve( p2, flag2 );
1285   dl := distance_to_axis( cl, low, upp, sl );
1286   d2 := distance_to_axis( c2, low, upp, s2 );

```

```

1289   q := nil;
1290   next_state := minuszero;
1291 repeat
1292   if DEBUG_LAYER_DIFFERENCE then writeln(
1293     '##### In layer difference, dl, d2=', dl, d2 );
1294   case next_state of
1295     minuszero :
1296       if dl >= d2 then [ + ]
1297       begin
1298         a := fix_curve( cl, low, upp, sl );
1299         next_state := pluszero;
1300         cl := get_curve( pl, flag1 );
1301         dl := distance_to_axis( cl, low, upp, sl );
1302       end
1303     else
1304       begin [ 1 ]
1305         next_state := minusone;
1306         c2 := get_curve( p2, flag2 );
1307         d2 := distance_to_axis( c2, low, upp, s2 );
1308       end;
1309     pluszero :
1310       if dl >= d2 then [ - ]
1311       begin
1312         b := fix_curve( cl, low, upp, sl );
1313         add_layer( a, b, q, r );
1314         next_state := minuszero;
1315         cl := get_curve( pl, flag1 );
1316         dl := distance_to_axis( cl, low, upp, sl );
1317       end
1318     else
1319       begin [ 1 ]
1320         b := fix_curve( c2, low, upp, s2 );
1321         add_layer( a, b, q, r );
1322         next_state := plusone;
1323         c2 := get_curve( p2, flag2 );
1324         d2 := distance_to_axis( c2, low, upp, s2 );
1325       end;
1326     plusone :
1327       if dl >= d2 then [ - ]
1328       begin
1329         next_state := minusone;
1330         cl := get_curve( pl, flag1 );
1331         dl := distance_to_axis( cl, low, upp, sl );
1332       end
1333     else
1334       begin [ 0 ]
1335         a := fix_curve( c2, low, upp, s2 );
1336         next_state := pluszero;
1337         c2 := get_curve( p2, flag2 );
1338         d2 := distance_to_axis( c2, low, upp, s2 );
1339       end;
1340     minusone :
1341       if dl >= d2 then [ + ]

```

```

1345 begin
1346   next_state := plusone;
1347   C1 := get_curve( p1, flag1 );
1348   d1 := distance_to_axis(C1,low,upp,S1);
1349   else
1350     begin { 0 }
1351       next_state := minuszero;
1352       C2 := get_curve( p2, flag2 );
1353       d2 := distance_to_axis(C2,low,upp,S2 );
1354     end;
1355     until ( C1 = nil ) and ( C2 = nil ) ;
1356   end; { case }
1357   if DEBUG_LAYER_DIFFERENCE then writeln(
1358     '**** In layer_difference, before fix_bound ****' );
1359   end; { layer_differenc }
1360
1361 procedure compute_layer( operation : axis_op_kind;
1362   var q : layer_ptr; low, upp : t_value;
1363   S1, S2 : segment_ptr );
1364 { compute the new layers by S1 operation S2 and bounded by
1365   low and upp; results should be sorted according to depth.
1366   S1 or S2 may be nil, in this case direct copy the curve is
1367   possible which depends on what operation is applied. New
1368   layer is pointed by q. }
1369 var
1370   S : segment_ptr;
1371 begin
1372   case operation of
1373     union : begin
1374       if DEBUG_COMPUTE_LAYER then writeln(
1375         '$$$ compute_layer, union op, low, upp=
1376         ',low,upp);
1377       if ( S1 <> nil ) and ( S2 <> nil ) then
1378         layer_union( q, low, upp, S1, S2 )
1379       else if ( S1 = nil ) and ( S2 = nil ) then
1380         q := nil
1381       else begin
1382         if S1 = nil then
1383           S := S2
1384         else if S2 = nil then
1385           S := S1;
1386         q := S^.layer_head;
1387         fix_bound( q, low, upp, S );
1388       end;
1389     difference :
1390       begin
1391         if DEBUG_COMPUTE_LAYER then writeln(
1392           '$$$ compute_layer, difference op,

```

```

1401   ',low,upp= ',low,upp);
1402   if S1 = nil then
1403     q := nil
1404   else if S2 <> nil then
1405     layer_difference( q, low, upp, S1, S2 )
1406   else begin
1407     q := S1^.layer_head;
1408     fix_bound( q, low, upp, S1 );
1409   end;
1410 end; { case }
1411 end; { compute_layer }
1412
1413 procedure add_segment( operation : axis_op_kind;
1414   a, b : t_value; var p : segment_ptr;
1415   S1, S2 : segment_ptr );
1416 { add a newly created segment [a b] into the principal axis P;
1417   where p is the segment pointer. This routine should include
1418   interval intersection computation and layer merge operation.}
1419 var
1420   Iptr : intersect_rec_ptr;
1421   low, upp : t_value;
1422   if a >= b then
1423     begin
1424       { return } { null interval; do nothing }
1425     end
1426   else begin
1427     compute_intersection( S1, S2, Iptr, a, b );
1428     sort_intersection( Iptr );
1429     screen_intersection( Iptr, a, b );
1430     low := Iptr^.intersection;
1431     Iptr := Iptr^.next_intersect;
1432   repeat
1433     upp := Iptr^.intersection;
1434     p^.lower_bound := low;
1435     p^.upper_bound := upp;
1436     compute_layer(operation, p^.layer_head,
1437       low, upp, S1, S2);
1438     if p^.layer_head <> nil then
1439       begin
1440         new( p^.next_segment );
1441         p := p^.next_segment;
1442         p^.next_segment := nil;
1443       end;
1444     until Iptr = nil;
1445   end;
1446   low := upp;
1447   Iptr := Iptr^.next_intersect;
1448   until Iptr = nil;
1449 end;

```

```

1457 end;
1458
1459 end; { add_segment }
1460
1461 function compute_D3_point( pt1, pt2 : D3_point;
1462 t : t_value ) : D3_point;
1463 { given two end points pt1 and pt2, and a parameter t,
1464 compute its coordinate }
1465 begin
1466
1467 compute_D3_point.x_coord := pt1.x_coord +
1468 t * (pt2.x_coord-pt1.x_coord);
1469 compute_D3_point.y_coord := pt1.y_coord +
1470 t * (pt2.y_coord-pt1.y_coord);
1471 compute_D3_point.z_coord := pt1.z_coord +
1472 t * (pt2.z_coord-pt1.z_coord);
1473
1474 end; { compute_D3_point }
1475
1476 function get_next( var flag : lower_or_upper;
1477 var done : boolean;
1478 var S, NS : segment_ptr ) : t_value;
1479 { get next interval value for use, done is set while no more
1480 can be obtained. S should always point to the interval
1481 currently dealt with. }
1482 const
1483 max_t_value = 2.0;
1484 begin
1485
1486 if done then
1487 get_next := max_t_value
1488 else begin
1489 case flag of
1490 lower : begin
1491 S := NS;
1492 flag := upper;
1493 if S <> nil then
1494 get_next := S.lower_bound
1495 else begin
1496 done := true;
1497 get_next := max_t_value;
1498 end;
1499 upper : begin
1500 flag := lower;
1501 NS := S.next_segment;
1502 get_next := S.upper_bound
1503 end;
1504 end; { case }
1505 end;
1506
1507 end; { get_next }
1508
1509 function partition ( operation : axis_op_kind;
1510 P1, P2 : principal_axis ) : principal_axis;
1511 { compute new intervals for union/difference operation;

```

```

1513 basically it performs the merge step of merge_sort algorithm.)
1514 var
1515 PX : principal_axis;
1516 S1, S2, NS1, NS2 : segment_ptr;
1517 flag1, flag2 : lower_or_upper;
1518 done1, done2 : boolean;
1519 a, b, al, a2, actual_t, first_t, last_t : t_value;
1520 pt1, pt2 : D3_point;
1521 p : segment_ptr;
1522 next_state : state;
1523
1524 begin
1525 { need code to initialize PX.start_point and PX.end_point }
1526 PX.start_point := min_point;
1527 PX.end_point := max_point;
1528
1529 new( PX.segment_head );
1530 p := PX.segment_head;
1531 p.next_segment := nil;
1532 NS1 := P1.segment_head;
1533 NS2 := P2.segment_head;
1534 flag1 := lower;
1535 flag2 := lower;
1536 done1 := false;
1537 done2 := false;
1538
1539 al := get_next( flag1, done1, S1, NS1 );
1540 a2 := get_next( flag2, done2, S2, NS2 );
1541 next_state := minuszero;
1542
1543 repeat
1544 case next_state of
1545 minuszero :
1546 if al <= a2 then
1547 begin
1548 a := al;
1549 next_state := pluszero;
1550 al := get_next( flag1, done1, S1, NS1 );
1551 end
1552 else begin
1553 a := a2;
1554 next_state := minusone;
1555 a2 := get_next( flag2, done2, S2, NS2 );
1556 end;
1557 pluszero :
1558 if al <= a2 then
1559 begin
1560 b := al;
1561 next_state := minuszero;
1562 add_segment(operation, a, b, p, S1, nil);
1563 al := get_next( flag1, done1, S1, NS1 );
1564 end
1565 else begin
1566 b := a2;
1567 next_state := plusone;
1568 add_segment(operation, a, b, p, S1, nil);
1569

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1569 a := a2;
1570 a2 := get_next( flag2, done2, S2, NS2 );
1571 end;
1572 plusone :
1573   if a1 <= a2 then
1574     begin
1575       b := a1;
1576       next_state := minusone;
1577       add_segment( operation, a, b, p, S1, S2 );
1578       a := a1;
1579       a1 := get_next( flag1, done1, S1, NS1 );
1580     end
1581   else begin
1582     b := a2;
1583     next_state := pluszero;
1584     add_segment( operation, a, b, p, S1, S2 );
1585     a := a2;
1586     a2 := get_next( flag2, done2, S2, NS2 );
1587   end;
1588 minusone :
1589   if a1 <= a2 then
1590     begin
1591       b := a1;
1592       next_state := plusone;
1593       if operation = union then
1594         add_segment( operation, a, b,
1595                     p, nil, S2 );
1596       a := a1;
1597       a1 := get_next( flag1, done1, S1, NS1 );
1598     end
1599   else begin
1600     b := a2;
1601     next_state := minuszero;
1602     if operation = union then
1603       add_segment( operation, a, b, p,
1604                   nil, S2 );
1605     a := a2;
1606     a2 := get_next( flag2, done2, S2, NS2 );
1607   end;
1608 end; { case }
1609 until done1 and done2;
1610
1611 { need code to delete the last element of the segment
1612   ( null segment ) }
1613 p := PX.segment_head;
1614 if p^.next_segment = nil then
1615   begin
1616     PX.segment_head := nil;
1617     actual_t := 0.0;
1618   end;
1619 while ( p <> nil ) do
1620   begin
1621     if p^.next_segment^.next_segment = nil then
1622       begin
1623         p^.next_segment := nil;
1624

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1625     actual_t := p^.upper_bound;
1626   end;
1627   p := p^.next_segment;
1628 end;
1629 first_t := PX.segment_head^.lower_bound;
1630 last_t := actual_t;
1631
1632 { if operation = difference, should reshape the t
1633   parameters because the first and last several segments
1634   might be removed. recalculate the end point and the t
1635   parameters. }
1636 if ( operation = difference ) and
1637   (( last_t < 1.0 ) or ( first_t > 0.0 )) then
1638   begin
1639     pt1 := PX.start_point;
1640     pt2 := PX.end_point;
1641     PX.start_point := compute_D3_point(
1642       pt1, pt2, first_t );
1643     PX.end_point := compute_D3_point(
1644       pt1, pt2, last_t );
1645     adjust_t_value( PX, first_t, last_t, 0.0, 1.0 );
1646   end;
1647
1648 partition := PX ;
1649
1650 if DEBUG_PARTITION then writeln(
1651   '*** In partition, PX.start and end point=',
1652   PX.start_point.x_coord, PX.start_point.y_coord,
1653   PX.start_point.z_coord, PX.end_point.x_coord,
1654   PX.end_point.y_coord, PX.end_point.z_coord );
1655 end; { partition }
1656
1657 function smooth_curve( x1, x2 : t_value; C1 : curve_ptr ;
1658                       x3, x4 : t_value; C2 : curve_ptr ) : boolean;
1659 { check to see if the curve C1 and C2 are smoothly connected.
1660   note that C1 and C2 have the same curve type here. }
1661 const
1662   EPSILON = 0.00001;
1663   INFINITY = 99999.99999;
1664 var
1665   m1, m2 : real;
1666 begin
1667   case C1^.curve_type of
1668     line : begin { check if they have the same slope }
1669       if abs(x1 - x2) <= EPSILON then
1670         m1 := INFINITY
1671       else m1 := ( C1^.curve_start -
1672                   C1^.curve_end ) / ( x1 - x2 );
1673       if abs(x3 - x4) <= EPSILON then
1674         m2 := INFINITY
1675       else m2 := ( C2^.curve_start -
1676                   C2^.curve_end ) / ( x3 - x4 );
1677       if abs(m1 - m2) <= EPSILON then
1678

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1681      smooth_curve := true
1682    else smooth_curve := false ;
1683  end;
1684  arc : begin
1685    if Cl^.radius <> C2^.radius then
1686      smooth_curve := false
1687    else if Cl^.which_half <> C2^.which_half
1688      then smooth_curve := false
1689    else if ( Cl^.center.x_coord <>
1690      C2^.center.x_coord ) or
1691      ( Cl^.center.y_coord <>
1692      C2^.center.y_coord ) then
1693      smooth_curve := false
1694    else
1695      smooth_curve := true ;
1696  end;
1697  end; { case }
1698  end; { smooth_curve }
1699
1700  function check_smooth( q, r : segment_ptr;
1701    a, b : layer_ptr ) : boolean;
1702  { within segments q and r, check if layer a and b
1703  are smoothly connected. }
1704  var
1705    flag : boolean;
1706    x1, x2, x3, x4 : t_value;
1707  begin
1708
1709    x1 := q^.lower_bound;
1710    x2 := q^.upper_bound;
1711    x3 := r^.lower_bound;
1712    x4 := r^.upper_bound;
1713
1714    flag := smooth_curve( x1, x2, a^.inner_curve,
1715      x3, x4, b^.inner_curve );
1716
1717    if flag then
1718      flag := smooth_curve( x1, x2, a^.outer_curve,
1719        x3, x4, b^.outer_curve );
1720
1721    check_smooth := flag ;
1722  end; { check_smooth }
1723
1724  function check_connect( q, r : segment_ptr ) : boolean;
1725  { To check if q and r are connected; check the following
1726  condition :
1727  the same connecting point, the same curve type,
1728  smooth at the connecting point,
1729  the same number of layers. }
1730  var
1731    a, b : layer_ptr;
1732    done : boolean;
1733  begin
1734    if q^.upper_bound <> r^.lower_bound then
1735      check_connect := false
1736    else

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1737  begin
1738    a := q^.layer_head;
1739    b := r^.layer_head;
1740    done := false;
1741    while not done do
1742      begin
1743        if ( a^.inner_curve^.curve_type <>
1744          b^.inner_curve^.curve_type ) or
1745          ( a^.outer_curve^.curve_type <>
1746          b^.outer_curve^.curve_type )
1747        then begin
1748          done := true;
1749          check_connect := false ;
1750        end
1751        else if ( a^.inner_curve^.curve_end <>
1752          b^.inner_curve^.curve_start ) or
1753          ( a^.outer_curve^.curve_end <>
1754          b^.outer_curve^.curve_start ) then
1755          begin
1756            done := true;
1757            check_connect := false ;
1758          end
1759          else if not check_smooth( q, r, a, b ) then
1760            begin
1761              done := true;
1762              check_connect := false ;
1763            end
1764            else { ok for this layer, let's try next }
1765              begin
1766                a := a^.next_layer;
1767                b := b^.next_layer;
1768                if ( a < nil ) and ( b <> nil ) or
1769                  ( b = nil ) and ( a <> nil ) then
1770                  begin
1771                    done := true;
1772                    check_connect := false ;
1773                  end
1774                  else if ( a = nil ) and ( b = nil ) then
1775                    begin
1776                      done := true;
1777                      check_connect := true ;
1778                    end;
1779                  end;
1780                end; { while }
1781              end; { else }
1782            end; { check_connect }
1783
1784          procedure merge_curve( C1, C2 : curve_ptr );
1785          { merge the curve C1 and C2; result is in C1; }
1786          begin
1787            case C1^.curve_type of
1788              line : C1^.curve_end := C2^.curve_end;
1789              arc : C1^.curve_end := C2^.curve_end;
1790            end; { case }
1791          end; { merge }
1792        end; { case }

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1793 end; { merge_curve }
1794
1795 procedure connect( q, r : segment_ptr );
1796 { merge q and r into a new interval pointed by q.
1797 }
1798 var
1799   p, s : layer_ptr;
1800 begin
1801   { merge segment }
1802   q^.upper_bound := r^.upper_bound;
1803   q^.next_segment := r^.next_segment;
1804   { merge each layer }
1805   p := q^.layer_head;
1806   s := r^.layer_head;
1807   while ( p <> nil ) do
1808     begin
1809       merge_curve( p^.inner_curve, s^.inner_curve );
1810       merge_curve( p^.outer_curve, s^.outer_curve );
1811     end;
1812     p := p^.next_layer;
1813     s := s^.next_layer;
1814   end; { while }
1815 end; { connect }
1816
1817 procedure merge_interval ( var P : principal_axis );
1818 { try to merge consecutive intervals into a larger one;
1819 this step is important to enforce the uniqueness of
1820 this principal-axis representation. }
1821 var
1822   q, r : segment_ptr;
1823 begin
1824   q := P.segment_head;
1825   r := q^.next_segment;
1826   while ( r <> nil ) do
1827     begin
1828       if check_connect( q, r ) then
1829         begin
1830           connect( q, r );
1831           r := q^.next_segment;
1832         end
1833       else
1834         begin
1835           q := r;
1836           r := q^.next_segment;
1837         end;
1838       { end }
1839     end;
1840   end; { end }
1841
1842 if DEBUG_MERGE_INTERVAL then writeln(
1843   '**** In merge_interval, P.start and end point=',
1844   P.start_point.x_coord, P.start_point.y_coord,
1845   P.start_point.z_coord, P.end_point.x_coord,
1846   P.end_point.y_coord, P.end_point.z_coord );

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1849   P.end_point.y_coord, P.end_point.z_coord);
1850 end; { merge_interval }
1851
1852 function axis_operation ( operation : axis_op_kind,
1853   union/difference P1, P2 : principal_axis ) : principal_axis;
1854 { this operation performs the actual union/difference of
1855 two CSG trees; it first calls normalization procedure
1856 to make two trees have t-parameters on the same basis,
1857 then applies the merge routine of merge-sort algorithm
1858 to union/difference them together. }
1859 var
1860   P : principal_axis;
1861 begin
1862   if DEBUG_AXIS_OP then writeln('**** enter normalization');
1863   normalization ( P1, P2, min_point, max_point );
1864   if DEBUG_AXIS_OP then writeln('**** leave normalization');
1865   if DEBUG_AXIS_OP then writeln('**** enter partition');
1866   P := partition ( operation, P1, P2 );
1867   if DEBUG_AXIS_OP then writeln('**** leave partition');
1868   if DEBUG_AXIS_OP then writeln('**** leave merge_interval');
1869   merge_interval ( P );
1870   if TRACE_AXIS_OP then dump_axis( P );
1871   if DEBUG_AXIS_OP then writeln('**** enter merge_interval');
1872   merge_interval ( P );
1873   if TRACE_AXIS_OP then dump_axis( P );
1874   if DEBUG_AXIS_OP then writeln('**** leave merge_interval');
1875 end; { axis_operation := P };
1876
1877 procedure rot_x( angle : real; var p : D3_point );
1878 { rotate the point p about x-axis by amount of angle;
1879 since the rotation is relative to local coordinate system
1880 of its own, translation must be done before and after
1881 rotation. }
1882 var
1883   Y, z : real;
1884 begin
1885   { translation to the origin of local coordinate }
1886   p.x_coord := p.x_coord - p0.x_coord ;
1887   p.y_coord := p.y_coord - p0.y_coord ;
1888   p.z_coord := p.z_coord - p0.z_coord ;

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  ( rotate )
  y := p.y_coord;
  z := p.z_coord;
  angle := angle * PI / 180.0 ; { convert to radian degree }
  p.y_coord := y * cos( angle ) - z * sin( angle );
  p.z_coord := y * sin( angle ) + z * cos( angle );
  { translation back to the global coordinate }
  p.x_coord := p.x_coord + p0.x_coord ;
  p.y_coord := p.y_coord + p0.y_coord ;
  p.z_coord := p.z_coord + p0.z_coord ;
end; { rot_x }

procedure rot_y( angle : real; var p : D3_point;
  p0 : D3_point );
{ rotate the point p about y-axis by amount of angle,
  since the rotation is relative to local coordinate system
  of its own, translation must be done before and after
  rotation. }
var
  x, z : real;
begin
  { translation to the origin of local coordinate }
  p.x_coord := p.x_coord - p0.x_coord ;
  p.y_coord := p.y_coord - p0.y_coord ;
  p.z_coord := p.z_coord - p0.z_coord ;
  { rotate }
  x := p.x_coord;
  z := p.z_coord;
  angle := angle * PI / 180.0 ; { convert to radian degree }
  p.x_coord := x * cos( angle ) + z * sin( angle );
  p.z_coord := - x * sin( angle ) + z * cos( angle );
  { translation back to the global coordinate }
  p.x_coord := p.x_coord + p0.x_coord ;
  p.y_coord := p.y_coord + p0.y_coord ;
  p.z_coord := p.z_coord + p0.z_coord ;
end; { rot_y }

procedure rot_z( angle : real; var p : D3_point;
  p0 : D3_point );
{ rotate the point p about z-axis by amount of angle,
  since the rotation is relative to local coordinate system
  of its own, translation must be done before and after
  rotation. }
var
  x, y : real;
begin
  { translation to the origin of local coordinate }
  p.x_coord := p.x_coord - p0.x_coord ;

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  p.y_coord := p.y_coord - p0.y_coord ;
  p.z_coord := p.z_coord - p0.z_coord ;
  { rotate }
  x := p.x_coord;
  y := p.y_coord;
  angle := angle * PI / 180.0 ; { convert to radian degree }
  p.y_coord := x * cos( angle ) - y * sin( angle );
  p.y_coord := x * sin( angle ) + y * cos( angle );
  { translation back to the global coordinate }
  p.x_coord := p.x_coord + p0.x_coord ;
  p.y_coord := p.y_coord + p0.y_coord ;
  p.z_coord := p.z_coord + p0.z_coord ;
end; { rot_z }

procedure translate( xform : xform_matrix ; var p : D3_point );
{ translate the point p by the amount of the translation part
  of xform }
begin
  p.x_coord := p.x_coord + xform.translate_x;
  p.y_coord := p.y_coord + xform.translate_y;
  p.z_coord := p.z_coord + xform.translate_z;
end; { translate }

procedure transform ( xform : xform_matrix;
  var p : principal_axis );
{ transform the tree by a transformation matrix;
  actually the principal axis is transformed,
  more specifically only the two end points are
  transformed. }
begin
  { P.start_point := p.start_point * transform_matrix; }
  { P.end_point := p.end_point * transform_matrix; }
  { note other variables are transformation invariant }

  { let's do the rotation first }
  if xform.rotate_x <> 0.0 then
    rot_x( xform.rotate_x, P.end_point, P.start_point );
  if xform.rotate_y <> 0.0 then
    rot_y( xform.rotate_y, P.end_point, P.start_point );
  if xform.rotate_z <> 0.0 then
    rot_z( xform.rotate_z, P.end_point, P.start_point );
  { and then do the translation }
  if ( xform.translate_x <> 0.0 ) or
    ( xform.translate_y <> 0.0 ) or
    ( xform.translate_z <> 0.0 ) then
    begin

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2017 translate( xform, P.start_point );
2018 translate( xform, P.end_point );
2019 end;
2020
2021 end; { transform }
2022
2023 function build_axis ( T : CSG_tree ) : principal_axis;
2024 { convert the primitive solid T to the principal_axis
2025 type representation }
2026 var
2027   P : principal_axis;
2028   s : segment_ptr;
2029   l : layer_ptr;
2030   height : real;
2031 begin
2032
2033   P.start_point := T.prim_solid.base;
2034   P.end_point := T.prim_solid.top;
2035
2036   if DEBUG_BUILD_AXIS then writeln(
2037     '**** In build_axis,P.start and end point=',
2038     P.start_point.x_coord, P.start_point.y_coord,
2039     P.start_point.z_coord, P.end_point.x_coord,
2040     P.end_point.y_coord, P.end_point.z_coord);
2041
2042   if DEBUG_BUILD_AXIS then writeln(
2043     '**** In build_axis,Prim_solid.base and top point=',
2044     T.prim_solid.base.x_coord,T.prim_solid.base.y_coord,
2045     T.prim_solid.base.z_coord, T.prim_solid.top.x_coord,
2046     T.prim_solid.top.y_coord,T.prim_solid.top.z_coord);
2047
2048   new( P.segment_head );
2049
2050   s := P.segment_head;
2051   s.lower_bound := 0.0;
2052   s.upper_bound := 1.0;
2053   s.next_segment := nil;
2054   new( s.layer_head );
2055
2056   l := s.layer_head;
2057   l.next_layer := nil;
2058   new( l.inner_curve );
2059   new( l.outer_curve );
2060
2061   case T.prim_solid.solid_type of
2062     cylinder :
2063       begin
2064         l.outer_curve.curve_type := line;
2065         l.outer_curve.curve_start :=
2066           T.prim_solid.radius;
2067         l.outer_curve.curve_end :=
2068           T.prim_solid.radius;
2069
2070         l.inner_curve.curve_type := line;
2071         l.inner_curve.curve_start := 0.0;

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2073   l.inner_curve.curve_end := 0.0;
2074 end;
2075
2076 cone :
2077   begin
2078     l.outer_curve.curve_type := line;
2079     l.outer_curve.curve_start :=
2080       T.prim_solid.radius_b;
2081     l.outer_curve.curve_end :=
2082       T.prim_solid.radius_t;
2083
2084     l.inner_curve.curve_type := line;
2085     l.inner_curve.curve_start := 0.0;
2086     l.inner_curve.curve_end := 0.0;
2087 end;
2088
2089 torus :
2090   begin
2091     l.outer_curve.curve_type := arc;
2092     height := ( T.prim_solid.inner_radius +
2093       T.prim_solid.outer_radius ) / 2.0;
2094     l.outer_curve.curve_start := height;
2095     l.outer_curve.curve_end := height;
2096     l.outer_curve.center.x_coord := 0.5;
2097     l.outer_curve.center.y_coord := height;
2098     l.outer_curve.radius :=
2099       T.prim_solid.outer_radius - height;
2100     l.outer_curve.which_half := up;
2101
2102     l.inner_curve.curve_type := arc;
2103     l.inner_curve.curve_start := height;
2104     l.inner_curve.curve_end := height;
2105     l.inner_curve.center.x_coord := 0.5;
2106     l.inner_curve.center.y_coord := height;
2107     l.inner_curve.radius := height -
2108       T.prim_solid.inner_radius;
2109     l.inner_curve.which_half := down;
2110 end;
2111
2112 end; { case }
2113
2114 build_axis := P ;
2115
2116 end; { build_axis }
2117
2118 function evaluate_CSG ( T : CSG_tree_ptr ) : principal_axis;
2119 { recursive call itself to evaluate the CSG tree T;
2120 it is basically a tree traversal algorithm.
2121 Our assumption is that the two operands at this point must
2122 be coaxial although they might not be coaxial at the lower
2123 part of the tree; }
2124 var
2125   P, P1, P2 : principal_axis;
2126 begin
2127   if DEBUG_EVALUATE_CSG then writeln(
2128

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2129     '*** in evaluate_csg , partno =', T^.partno );
2130
2131 case T^.node_type of
2132 movement :
2133 begin
2134 P := evaluate_CSG ( T^.child );
2135 transform( T^.move, P );
2136 evaluate_CSG := P ;
2137
2138 if DEBUG_EVALUATE_CSG then writeln(
2139 '*** finish evaluate_csg , partno =',
2140 T^.partno );
2141
2142 end;
2143
2144 union_op :
2145 begin
2146 P1 := evaluate_CSG ( T^.left_child );
2147 P2 := evaluate_CSG ( T^.right_child );
2148
2149 if DEBUG_EVALUATE_CSG then writeln(
2150 '### axis operation for partno =',
2151 T^.partno );
2152
2153 P := axis_operation ( union, P1, P2 );
2154 evaluate_CSG := P ;
2155
2156 if DEBUG_EVALUATE_CSG then writeln(
2157 '*** finish evaluate_csg , partno =',
2158 T^.partno );
2159
2160 if DEBUG_EVALUATE_CSG then writeln(
2161 '!!! result of union,P.start and end point=',
2162 P.start_point.x_coord,P.start_point.y_coord,
2163 P.start_point.z_coord, P.end_point.x_coord,
2164 P.end_point.y_coord,P.end_point.z_coord);
2165
2166 end;
2167 diff_op :
2168 begin
2169 P1 := evaluate_CSG ( T^.left_child );
2170 P2 := evaluate_CSG ( T^.right_child );
2171
2172 if DEBUG_EVALUATE_CSG then writeln(
2173 '### axis operation for partno =',
2174 T^.partno );
2175
2176 P := axis_operation ( difference, P1, P2 );
2177 evaluate_CSG := P ;
2178
2179 if DEBUG_EVALUATE_CSG then writeln(
2180 '*** finish evaluate_csg , partno =',
2181 T^.partno );
2182
2183 end;
2184 primitive:
2185 begin

```

```

2185 P := build_axis ( T );
2186 evaluate_CSG := P ;
2187
2188 if DEBUG_EVALUATE_CSG then writeln(
2189 '*** finish evaluate_csg , partno =',
2190 T^.partno );
2191
2192 end;
2193
2194 end; { case }
2195
2196 end; { evaluate_CSG }
2197
2198 procedure get_input_data( var CSG : CSG_tree_ptr );
2199 { get the input CSG tree from outside world; it may be
2200 from the output of an interactive graphical geometrical
2201 modeller; EEC5487 term project, for example. }
2202 const
2203 obj_path = 'object';
2204 mov_path = 'movement';
2205 prim_path = 'primitives';
2206
2207 var
2208 openstatus : status_st;
2209
2210 procedure read_tree( T : CSG_tree_ptr );
2211 { recursively read data from object, primitives, and
2212 movement files to construct a CSG tree T }
2213 var
2214 part1, part2, part3 : integer;
2215 ch : char;
2216 op : array[1..9] of char;
2217 prm : array[1..8] of char;
2218 ax, ay, az, x, y, z, xl, yl, zl : real;
2219
2220 begin
2221 readln( obj_file, part1, ch, ch, ch, ch, ch, ch, ch,
2222 op, part2, part3 );
2223
2224 if DEBUG_INPUT then writeln(
2225 '**** read object, objno =',part1, op);
2226
2227 if (T^.partno <> CSG_ROOT) and
2228 ( T^.partno <> part1 ) then
2229 writeln('----> read_tree error, ',
2230 ' input tree structure incorrect!');
2231
2232 if op = ' union' then T^.node_type := union_op
2233 else if op = ' differ' then T^.node_type := diff_op
2234 else if op = ' primitive' then T^.node_type := primitive
2235 else if op = ' moved_obj' then T^.node_type := movement
2236 else writeln('illegal CSG node type!');
2237
2238
2239 case T^.node_type of
2240

```

```

2241 union_op,
2242 diff_op :
2243 begin
2244     new( T^.left_child );
2245     T^.left_child^.partno := part2;
2246     read_tree( T^.left_child );
2247
2248     new( T^.right_child );
2249     T^.right_child^.partno := part3;
2250     read_tree( T^.right_child );
2251
2252 end;
2253 primitive :
2254 begin
2255     readln( prim_file, part1, ch,
2256             ch, ch, ch, ch, ch, prn,
2257             xl, yl, zl, x, y, z,
2258             ax, ay, az );
2259
2260     if DEBUG_INPUT then writeln(
2261         '**** read primitive, no =',
2262         part1, ' ', prn );
2263
2264     if ( part1 <> T^.partno ) then
2265         writeln('----> read tree error, ',
2266             'input tree structure incorrect!!');
2267
2268         [ must treat cone specially ]
2269     if prn = 'cylinder' then
2270         T^.prim_solid.solid_type := cylinder
2271     else if prn = 'cone' then
2272         T^.prim_solid.solid_type := cone
2273     else if prn = 'torus' then
2274         T^.prim_solid.solid_type := torus
2275     else writeln( 'illegal CSG ',
2276         'primitive node type!!');
2277
2278     case T^.prim_solid.solid_type of
2279         cylinder :
2280             begin
2281                 T^.prim_solid.base.x_coord := x;
2282                 T^.prim_solid.base.y_coord := y;
2283                 T^.prim_solid.base.z_coord := z;
2284                 T^.prim_solid.top.x_coord := x;
2285                 T^.prim_solid.top.y_coord := y;
2286                 T^.prim_solid.top.z_coord := z+xl;
2287                 T^.prim_solid.radius := yl;
2288             end;
2289         cone :
2290             begin
2291                 T^.prim_solid.base.x_coord := x;
2292                 T^.prim_solid.base.y_coord := y;
2293                 T^.prim_solid.base.z_coord := z;
2294                 T^.prim_solid.top.x_coord := x;
2295                 T^.prim_solid.top.y_coord := y;
2296

```

```

2297     T^.prim_solid.top.z_coord := z+xl;
2298     T^.prim_solid.radius_b := yl;
2299     T^.prim_solid.radius_t := 0.0;
2300 end;
2301 torus :
2302 begin
2303     T^.prim_solid.base.x_coord := x;
2304     T^.prim_solid.base.y_coord := y;
2305     T^.prim_solid.base.z_coord := z-yl;
2306     T^.prim_solid.top.x_coord := x;
2307     T^.prim_solid.top.y_coord := y;
2308     T^.prim_solid.top.z_coord := z+yl;
2309     T^.prim_solid.inner_radius := xl-yl;
2310     T^.prim_solid.outer_radius := xl+yl;
2311 end;
2312 end; { case }
2313
2314 if ax <> 0.0 then
2315     rot_x( ax, T^.prim_solid.top,
2316           T^.prim_solid.base );
2317
2318 if ay <> 0.0 then
2319     rot_y( ay, T^.prim_solid.top,
2320           T^.prim_solid.base );
2321
2322 if az <> 0.0 then
2323     rot_z( az, T^.prim_solid.top,
2324           T^.prim_solid.base );
2325
2326 end;
2327 movement :
2328 begin
2329     readln( mov_file, part1,
2330             T^.move.translate_x,
2331             T^.move.translate_y,
2332             T^.move.translate_z,
2333             T^.move.rotate_x,
2334             T^.move.rotate_y,
2335             T^.move.rotate_z );
2336
2337 if DEBUG_INPUT then writeln(
2338     '**** read movement, no =', part1 );
2339
2340 if ( part1 <> part3 ) then
2341     writeln('----> read tree error, ',
2342         'input tree structure incorrect!!');
2343
2344     new( T^.child );
2345     T^.child^.partno := part2;
2346     read_tree( T^.child );
2347
2348 end; { case }
2349
2350 end; { read_tree }
2351
2352 begin

```

```

2353 { read data and construct them as a CSG tree;
2354 there are three data files; one is the movement,
2355 another is primitives, the third one is the
2356 construction }
2357 open( obj_file, obj_path, 'OLD', openstatus.all );
2358 open( mov_file, mov_path, 'OLD', openstatus.all );
2359 open( prim_file, prim_path, 'OLD', openstatus.all );
2360
2361 reset( obj_file );
2362 reset( mov_file );
2363 reset( prim_file );
2364
2365 if DEBUG INPUT then writeIn(
2366 '*** open files in get_input_data');
2367
2368 new( CSG );
2369 CSG.parno := CSG_ROOT;
2370 read_tree( CSG );
2371
2372 if DEBUG INPUT then writeIn(
2373 '*** finish get_input_data');
2374
2375 end; { get_input_data }
2376
2377 procedure compute_length_diameter( PX : principal_axis;
2378 var L, D : real );
2379 { compute the maximal length L of PX; note:
2380 PX might be broken into several segments. Compute the
2381 extremal boundary (diameter D) from the axis }
2382 var
2383 p, q : segment_ptr;
2384 len, maxlen : t_value;
2385 xl, yl, zl : real;
2386 maxht, height, local_height : real;
2387 C : curve_ptr;
2388
2389 begin
2390 maxlen := 0.0;
2391 len := 0.0;
2392 maxht := 0.0;
2393 local_height := 0.0;
2394 p := PX.segment_head;
2395
2396 while p <> nil do
2397 begin
2398 C := p^.layer_head^.outer_curve;
2399 case C^.curve_type of
2400 line : if C^.curve_start > C^.curve_end then
2401 height := C^.curve_start
2402 else height := C^.curve_end;
2403 arc : if C^.which_half = up then
2404 begin
2405 if ( C^.center.x_coord >=
2406 p^.lower_bound ) and
2407 ( C^.center.x_coord <=

```

```

2409 p^.upper_bound ) then
2410 height := C^.center.y_coord +
2411 C^.radius
2412 else if C^.center.x_coord >=
2413 p^.upper_bound then
2414 height := C^.curve_end
2415 else
2416 height := C^.curve_start;
2417
2418 end { down }
2419 begin
2420 if C^.curve_start > C^.curve_end then
2421 height := C^.curve_start
2422 else
2423 height := C^.curve_end;
2424 end; { case }
2425
2426 if height > local_height then
2427 local_height := height;
2428
2429 len := len + p^.upper_bound - p^.lower_bound;
2430 q := p^.next_segment;
2431 if q <> nil then
2432 begin
2433 if ( q^.lower_bound - p^.upper_bound ) > 0.0001
2434 then
2435 begin { non-continuous segments }
2436 if len > maxlen then
2437 begin
2438 maxlen := len;
2439 maxht := local_height;
2440 end;
2441 len := 0.0;
2442 local_height := 0.0;
2443 end;
2444 p := q;
2445 end;
2446
2447 if len > maxlen then
2448 begin
2449 maxlen := len;
2450 maxht := local_height;
2451 end;
2452
2453 xl := PX.start_point.x_coord - PX.end_point.x_coord;
2454 yl := PX.start_point.y_coord - PX.end_point.y_coord;
2455 zl := PX.start_point.z_coord - PX.end_point.z_coord;
2456
2457 if DEBUG ID then writeIn(
2458 '*** Compute L&D, xl, yl, zl, maxlen =',
2459 xl, yl, zl, maxlen );
2460
2461 L := maxlen * sqrt( xl * xl + yl * yl + zl * zl );
2462
2463 D := 2.0 * maxht;
2464

```

```

2577 file_name := 'output.data';
2578 name_size := 11;
2579
2580 { open the disk file for external storage }
2581
2582 gpr_$open_bitmap_file( gpr_$create, file_name,
2583   name_size, version, window_size,
2584   groups, header, attribs, filebm, created, status);
2585 if status.all <> status.$ok then writeln(
2586   '-----> In produce_$ok then writeln(
2587   'bitmap file for the specified file name');
2588
2589 { set current bitmap for block transfer }
2590
2591 gpr_$set_bitmap( filebm, status);
2592 if status.all <> status.$ok then writeln(
2593   '-----> In produce_bitmap, Can not set ',
2594   'the bitmap file');
2595
2596 { blok transfer data from current bitmap to disk }
2597
2598 gpr_$pixel_blt( init_bitmap, window,
2599   window.window_base, status);
2600 if status.all <> status.$ok then writeln(
2601   '-----> In produce_bitmap, Can not Block ',
2602   'transfer bitmap file');
2603
2604 { set back the original display bitmap }
2605
2606 gpr_$set_bitmap( init_bitmap, status);
2607
2608 end; { produce_bitmap_file }
2609
2610 procedure terminate_graph;
2611 { terminate the graphic environment ;
2612   probably make a hard copy. }
2613 begin
2614   gpr_$terminate( false, status );
2615
2616 end; { terminate_graph }
2617
2618 procedure draw_line( xl, yl, x2, y2 : real );
2619 { draw a line from ( xl, yl ) to ( x2, y2 );
2620   this procedure should interface graphic routine }
2621 var
2622   xx1, yy1, xx2, yy2 : integer;
2623   begin
2624     if DEBUG_DRAW then writeln(
2625       '&&&& draw_line,x1,y1,x2,y2', xl,yl,x2,y2);
2626
2627     procedure draw_arc( xc, yc, radii : real;
2628       half : up_or_down;
2629       xl, yl, x2, y2 : real );
2630     { draw an arc of a circle where its center is
2631       at ( xc, yc ) of radius radii; the arc starts from
2632       ( xl, yl ) to ( x2, y2 ) and half specifies either
2633       the upper part or the lower part in the circle
2634       this arc belongs to. This routine should interface
2635       the graphic routine to draw arc, circle or lines
2636       if approximation is used. }
2637     var
2638       x3, y3, dy : real;
2639       first, middle, last : gpr_$position_t;
2640     begin
2641       if DEBUG_DRAW then writeln(
2642         '&&&& draw_arc,xc,yc,xl,y1,x2,y2',
2643         xc,yc,xl,y1,x2,y2);
2644
2645       { compute the middle point of the arc }
2646       x3 := ( xl + x2 ) / 2.0;
2647       dy := sqrt( radii * radii -
2648         ( x3 - xc ) * ( x3 - xc ) );
2649       if half = up then y3 := yc + dy
2650         else y3 := yc - dy;
2651
2652       { convert real coordinates to integer coordinates }
2653       first.x_coord := trunc( xl + 0.5 ) + x_origin ;
2654       first.y_coord := trunc( yl + 0.5 ) + y_origin ;
2655       middle.x_coord := trunc( x3 + 0.5 ) + x_origin;
2656       middle.y_coord := trunc( y3 + 0.5 ) + y_origin;
2657       last.x_coord := trunc( x2 + 0.5 ) + x_origin;
2658       last.y_coord := trunc( y2 + 0.5 ) + y_origin;
2659
2660       { call graphic routine , check Apollo Domain
2661         graphic routine }
2662       gpr_$move( first.x_coord , first.y_coord , status );
2663       gpr_$arc_3p( middle, last, status );
2664
2665     end; { draw_arc }
2666
2667 procedure draw_curve( C : curve_ptr; xl, x2 : real );
2668
2669 end; { draw_curve }

```

```

2465
2466      writeln('-----> length of the part = ', L);
2467      writeln('-----> diameter of the part = ', D);
2468
2469 end; { compute_length_diameter }
2470
2471 procedure compute_profile ( PX : principal_axis );
2472 { obtain profile of part from PX, intend to support
2473   NC application }
2474 const
2475   { to adjust the output on the center of the screen }
2476   x_origin = 24;
2477   Y_origin = 512;
2478
2479   var
2480     p : segment_ptr;
2481     ratio, xl, Yl, zl, x_position : real;
2482     aa, bb : real;
2483     disp_bm_size : gpr_$offset_t;
2484     init_bitmap : gpr_$bitmap_desc_t;
2485     window : gpr_$window_t;
2486
2487   procedure init_graph;
2488   { initialize the Apollo Domain graphic environment.
2489     Clipping to window is implemented, be careful !! }
2490   begin
2491     disp_bm_size.x_size := 1024;
2492     disp_bm_size.y_size := 1024;
2493     gpr_$init( gpr_$borrow, 1, disp_bm_size, 0,
2494               init_bitmap, status );
2495
2496     { set window clipping }
2497     window.window_base.x_coord := 0;
2498     window.window_base.y_coord := 0;
2499     window.window_size.x_size := 1024;
2500     window.window_size.y_size := 1024;
2501
2502     { set 'exclusive or' raster operation }
2503     gpr_$set_raster_op( 0, 6, status );
2504
2505     gpr_$set_clip_window( window, status );
2506     gpr_$set_clipping_active( true, status );
2507     gpr_$clear( -2, status );
2508
2509   end; { init_graph }
2510
2511   procedure hold_screen;
2512   { hold screen for a moment, say 30 seconds }
2513   const
2514     one_second = 250000;
2515   var
2516     wait_time : time_$clock_t;
2517     status : status_$t;
2518   begin
2519
2520

```

```

2521     wait_time.high16 := 0;
2522     wait_time.low32 := 30 * one_second;
2523     time_$wait( time_$relative, wait_time, status);
2524
2525   end; { hold_screen }
2526
2527   procedure produce_bitmap_file;
2528   { store the bitmap of the whole screen into disk,
2529     open the disk file for it
2530     if file exists or other errors, write error messages }
2531   var
2532     window : gpr_$window_t ;
2533     { window of the display bitmap }
2534     hi_plane : gpr_$plane_t ;
2535     groups : integer ;
2536     { the number of groups in external bitmap }
2537     status : status_$t;
2538     { returned status }
2539     version : gpr_$version_t;
2540     { version number of bitmap file }
2541     header : gpr_$bmf_group_header_array_t;
2542     { descriptor of external group bitmap header }
2543     attrs : gpr_$attribute_desc_t;
2544     { attributes which the bitmap will use }
2545     filebm : gpr_$bitmap_desc_t;
2546     { descriptor of bitmap }
2547     created : boolean;
2548     { specifies whether the bitmap file was created }
2549     file_name : name_$pname_t;
2550     { name of the external file for bitmap }
2551     name_size : integer;
2552     { length of the file name }
2553
2554   begin
2555     { set parameters for the window of operation }
2556     hi_plane := 0;
2557     groups := 1;
2558     window.window_base.x_coord := 0;
2559     window.window_base.y_coord := 0;
2560     window.window_size.x_size := 1000;
2561     window.window_size.y_size := 800;
2562     with header[0] do
2563       begin
2564         n_sects := hi_plane + 1;
2565         pixel_size := 1;
2566         allocated_size := 1;
2567         bytes_per_line := 0;
2568         bytes_per_sect := 0;
2569       end; { with }
2570     gpr_$allocate_attribute_block( attrs, status);
2571
2572     { get_file_name( file_name, name_size); }
2573
2574
2575
2576

```



```

2689 { draw a curve specified by C;
2690   xl ans x2 are horizontal bounds. }
2691 var
2692   xc, yc : real;
2693   begin
2694
2695     case C^.curve_type of
2696       line : draw_line( xl, C^.curve_start,
2697         x2, C^.curve_end );
2698       arc : begin
2699         xc := ratio * C^.center.x_coord;
2700         yc := C^.center.y_coord;
2701         draw_arc( xc, yc, C^.radius,
2702           C^.which_half, xl,
2703           C^.curve_start, x2,
2704           C^.curve_end );
2705       end;
2706     end; { case }
2707
2708   end; { draw_curve }
2709
2710 procedure draw_curve_buddy( C : curve_ptr; xl, x2 : real );
2711 { draw the counterpart of a curve specified by C;
2712   xl ans x2 are horizontal bounds. }
2713 var
2714   D : curve_ptr;
2715   begin
2716
2717     new( D );
2718     D^.curve_type := C^.curve_type;
2719     case C^.curve_type of
2720       line : begin
2721         D^.curve_start := - C^.curve_start;
2722         D^.curve_end := - C^.curve_end;
2723       end;
2724       arc : begin
2725         D^.curve_start := - C^.curve_start;
2726         D^.curve_end := - C^.curve_end;
2727         D^.center.x_coord := C^.center.x_coord;
2728         D^.center.y_coord := - C^.center.y_coord;
2729         D^.radius := C^.radius;
2730         if C^.which_half = up then
2731           D^.which_half := down
2732         else
2733           D^.which_half := up;
2734       end;
2735     end; { case }
2736
2737     draw_curve( D, xl, x2 );
2738
2739   end; { draw_curve_buddy }
2740
2741 procedure draw_layer( q : layer_ptr; xl, x2 : real );
2742 { to draw the layer curves by going thru q chain;
2743   each counterpart of the curve in q must also be drawn;
2744

```

```

2745   }
2746   begin
2747
2748     while q <> nil do
2749       begin
2750         draw_curve( q^.outer_curve, xl, x2 );
2751         draw_curve( q^.inner_curve, xl, x2 );
2752         draw_line( xl, q^.outer_curve^.curve_start,
2753           xl, q^.inner_curve^.curve_start );
2754         draw_line( x2, q^.outer_curve^.curve_end,
2755           x2, q^.inner_curve^.curve_end );
2756
2757         draw_curve_buddy( q^.inner_curve, xl, x2 );
2758         draw_curve_buddy( q^.outer_curve, xl, x2 );
2759         draw_line( xl, q^.outer_curve^.curve_start,
2760           xl, q^.inner_curve^.curve_start );
2761         draw_line( x2, q^.outer_curve^.curve_end,
2762           x2, q^.inner_curve^.curve_end );
2763
2764         q := q^.next_layer;
2765       end;
2766     end; { draw_layer }
2767
2768   end;
2769
2770 begin
2771
2772   { initialize the Apollo graphic environment }
2773   init_graph;
2774
2775   xl := PX.end_point.x_coord - PX.start_point.x_coord;
2776   yl := PX.end_point.y_coord - PX.start_point.y_coord;
2777   zl := PX.end_point.z_coord - PX.start_point.z_coord;
2778   ratio := sqrt( xl * xl + yl * yl + zl * zl );
2779   x_position := 0.0;
2780
2781   p := PX.segment_head;
2782
2783   while p <> nil do
2784     begin
2785       x_position := ratio * p^.lower_bound;
2786       xl := ratio * p^.upper_bound;
2787       draw_layer( p^.layer_head, x_position, xl );
2788
2789       p := p^.next_segment;
2790     end; { while }
2791
2792   { hold the screen for a movement }
2793   hold_screen;
2794
2795   { generate the bitmap file for dump }
2796   produce_bitmap_file;
2797
2798   { terminate the graphic environment }
2799   terminate_graph;
2800

```

```

2801
2802 end; { compute_profile }
2803
2804 { MAIN PROCEDURE }
2805 { show the applications }
2806 { P is the principal axis; CSG is the CSG tree to be evaluated;
2807 L is the length of this part and D is the diameter of the part;
2808 other properties for various code assignment can also be derived.
2809 }
2810 begin
2811
2812 { debugging switches }
2813 DEBUG_INPUT := false;
2814 DEBUG_EVALUATE_CSG := false;
2815 DEBUG_BUILD_AXIS := false;
2816 DEBUG_COMPUTE_T := false;
2817 DEBUG_CURVE_CURVE := false;
2818 DEBUG_INTERSECT := false;
2819 DEBUG_AXIS_OP := false;
2820 DEBUG_NORMALIZATION := false;
2821 DEBUG_PARTITION := false;
2822 DEBUG_ADD_LAYER := false;
2823 DEBUG_COMPUTE_LAYER := false;
2824 DEBUG_LAYER_DIFFERENCE := false;
2825 DEBUG_MERGE_INTERVAL := false;
2826 DEBUG_ID := false;
2827 DEBUG_DRAW := false;
2828 TRACE_AXIS := false;
2829
2830 { get input CSG tree }
2831 get_input_data ( CSG );
2832
2833 { transform the CSG tree into axis representation }
2834 P := eval
2835
2836 { find applications }
2837 { support part/shape classification and
2838 process planning }
2839 compute_length_diameter( P, L, D );
2840
2841 { support NC }
2842 compute_profile( P );
2843
2844 end. { main }
2845

```

***Appendix 2. Input CSG Data for "Bottle"***

\*\*\*\*\*  
 BOTTLE  
 \*\*\*\*\*

OBJECT

O#	OP	L_O#	R_O#	MV#/dummy
-1	differ	99	98	
99	union	44	45	
44	union	36	37	
36	differ	34	35	
34	differ	33	30	
33	union	27	29	
27	union	17	18	
17	union	11	12	
11	union	5	6	
5	differ	3	4	
3	union	1	2	
1	primitive	1	1	
2	primitive	2	2	
4	primitive	4	4	
6	differ	9	8	
9	differ	7	10	
7	primitive	7	7	
10	primitive	10	10	
8	primitive	8	8	
12	differ	13	14	
13	union	15	16	
15	primitive	15	15	
16	primitive	16	16	
14	primitive	14	14	
18	differ	19	20	
19	union	19	22	
21	differ	23	24	
23	primitive	23	23	
24	primitive	24	24	
22	differ	25	26	
25	primitive	25	25	
26	primitive	26	26	
20	primitive	20	20	
29	differ	31	32	
31	primitive	31	31	
32	primitive	32	32	
30	primitive	30	30	
35	primitive	35	35	
37	differ	40	42	
40	differ	38	39	
38	primitive	38	38	
39	primitive	39	39	
42	differ	41	43	
41	primitive	41	41	
43	primitive	43	43	
45	differ	52	53	
52	union	50	51	
50	union	48	49	

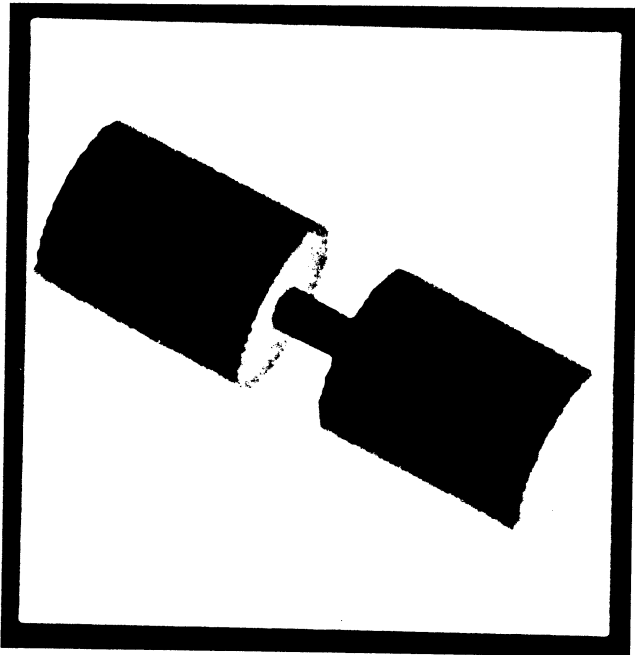
48 differ  
 46 primitive  
 47 primitive  
 49 primitive  
 51 primitive  
 53 primitive  
 98 primitive

O#	GC	XL	YL	ZL	X	Y	Z	ax	ay	az
1	torus	0.5000	0.2000	0.0000	0.0000	0.0000	0.2000	0.0000	0.0000	0.0000
2	cylinder	0.4000	0.5000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000
4	cylinder	0.4000	0.3000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000
7	cone	1.0000	1.2500	0.0000	0.0000	0.0000	1.0000	180.00	0.0000	0.0000
10	cone	0.8400	1.0500	0.0000	0.0000	0.0000	1.0000	180.00	0.0000	0.0000
8	cylinder	0.4000	0.5000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000
15	torus	1.2500	0.3000	0.0000	0.0000	0.0000	1.3000	0.0000	0.0000	0.0000
16	cylinder	0.6000	1.2500	0.0000	0.0000	0.0000	1.0000	0.0000	0.0000	0.0000
14	cylinder	0.6000	1.0500	0.0000	0.0000	0.0000	1.0000	0.0000	0.0000	0.0000
23	cone	1.9672	1.9672	0.0000	0.0000	0.0000	2.3172	180.00	0.0000	0.0000
24	cylinder	1.6000	1.2500	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000
25	torus	1.6844	0.4000	0.0000	0.0000	0.0000	2.6000	0.0000	0.0000	0.0000
26	torus	1.6844	0.2500	0.0000	0.0000	0.0000	2.6000	0.0000	0.0000	0.0000
20	cone	1.9086	1.9086	0.0000	0.0000	0.0000	2.4586	180.00	0.0000	0.0000
31	cone	1.9672	1.9672	0.0000	0.0000	0.0000	2.8828	0.0000	0.0000	0.0000
32	cylinder	1.6000	2.0000	0.0000	0.0000	0.0000	3.3172	0.0000	0.0000	0.0000
30	cone	1.9086	1.9086	0.0000	0.0000	0.0000	2.7414	0.0000	0.0000	0.0000
35	cylinder	0.6000	1.6844	0.0000	0.0000	0.0000	2.3000	0.0000	0.0000	0.0000
38	torus	1.6844	0.4000	0.0000	0.0000	0.0000	3.5400	0.0000	0.0000	0.0000
39	torus	1.6844	0.2600	0.0000	0.0000	0.0000	3.5400	0.0000	0.0000	0.0000
41	cylinder	0.8800	2.2500	0.0000	0.0000	0.0000	3.1000	0.0000	0.0000	0.0000
43	cylinder	0.4412	1.6000	0.0000	0.0000	0.0000	3.3172	0.0000	0.0000	0.0000
46	cone	3.2500	2.5000	0.0000	0.0000	0.0000	5.0000	180.00	0.0000	0.0000
47	cylinder	2.0500	1.6000	0.0000	0.0000	0.0000	1.7000	0.0000	0.0000	0.0000
49	torus	2.1525	0.3000	0.0000	0.0000	0.0000	5.0000	0.0000	0.0000	0.0000
51	cylinder	0.3000	2.1525	0.0000	0.0000	0.0000	5.0000	0.0000	0.0000	0.0000
53	cone	3.2200	2.1525	0.0000	0.0000	0.0000	5.0100	180.00	0.0000	0.0000
98	cube	3.0000	3.0000	5.5000	0.0000	-3.0000	0.0000	0.0000	0.0000	0.0000

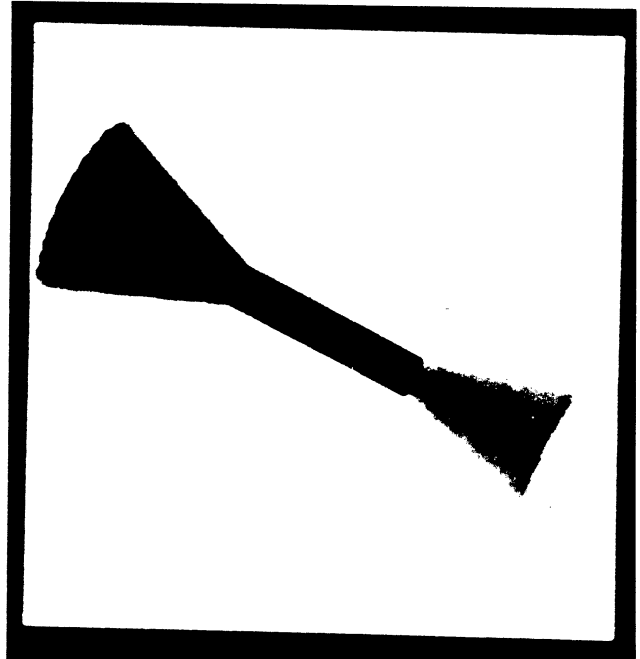
PRIMITIVE

### ***Appendix 3. Shaded Pictures by Ray Casting***

- (a) corresponds to Figure 6.1
- (b) corresponds to Figure 6.2
- (c) corresponds to Figure 6.3
- (d) corresponds to Figure 6.4
- (e) corresponds to Figure 7, the "bottle"



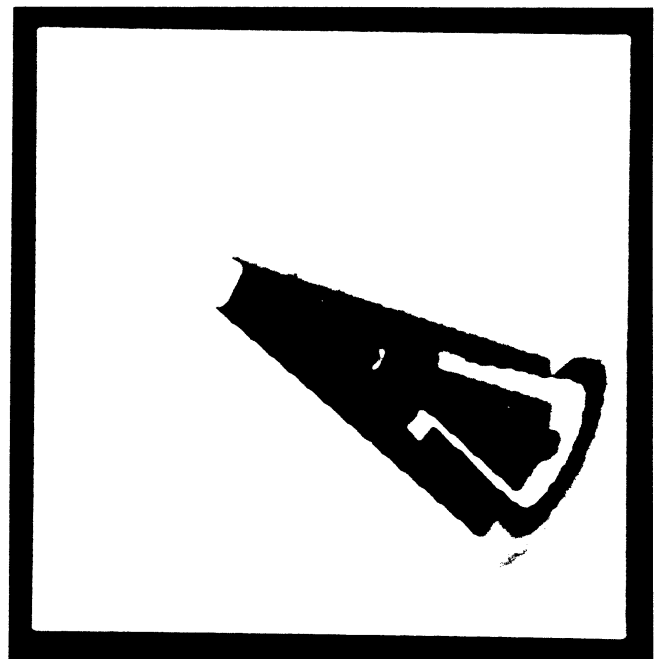
(a)



(b)



(c)

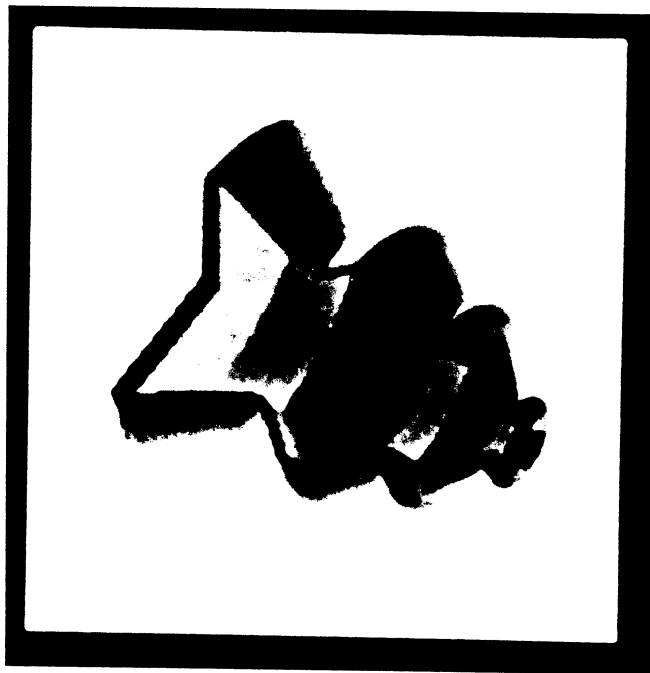


(d)

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(e)