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RQA-1—THE RECURSIVE QUEUE ANALYZER

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ABSTRACT

The Recursive Queue Analyzer, RQA-1, is a computer program designed to evaluate the equilibrium joint probability distributions of queue lengths and system conditions in very large, finite Markovian queueing systems. It will rapidly treat discrete-time and continuous-time systems having as many as 5000 states through numerical solution of equations of balance. The algorithm is generally very much faster and more accurate than equivalent Monte-Carlo simulation. This report is both a description of the organization of the program and a manual for program use.

1. INTRODUCTION

The Recursive Queue Analyzer, RQA-1, is a computer program designed to evaluate the equilibrium joint probability distributions of queue lengths and system conditions in very large, finite Markovian queueing systems. This report is both a description of the organization of the program and a manual for program use.

In the study of networks of queues and servers, the Markovian models represent an exceedingly important class of models. This program has been designed to facilitate the analysis of both discrete-time and continuous-time Markov chains having as many as 5,000 states. The primary design goal has been to provide a computation fast enough to encourage experimentation with the models in the study of system design. This has been achieved through efficient use of available (32K) high-speed storage in the computer and through careful program design. The program was written in the MAD language¹ with selective use of the UMAP assembly language. Although this work has been motivated by the need for it in the design of large computer systems, the potential area of application is quite broad.

For a continuous-time Markov chain with a finite state space, one can always write the Kolmogoroff differential equation

$$\dot{p}_t = p_t Q \quad (1)$$

where p_t is a vector whose i -th element is the probability that the system is in the i -th state at time t , \dot{p} represents the time derivative of p , and Q is a matrix of constants called the transition intensity matrix of the chain. The matrix Q is assumed to have the properties that

$$\sum_j q_{ij} = 0 \quad (2a)$$

for all states i , and

$$q_{ij} \geq 0 \quad (2b)$$

for all $i \neq j$. An equilibrium probability distribution is, by definition, a solution to the equation

$$p Q = 0 . \quad (3)$$

However, any solution to Eq. (3) is also a solution to

$$p = p(\Delta Q + I) \quad (4)$$

where Δ is a scalar constant and I is the identity matrix. The purpose of RQA-1 is to solve Eq. (4).

For a discrete-time Markov chain with a finite state space, one can always write the Chapman-Kolmogoroff equation

$$p_{n+1} = p_n A \quad (5)$$

where p_n is the vector whose i -th element is the probability that the system is in the i -th state at time n , and A is a matrix of constants called the transition matrix of the chain. The matrix A is assumed to have the properties that

$$\sum_j a_{ij} = 1 \quad (6a)$$

for all states i , and

$$1 \geq a_{ij} \geq 0 \quad (6b)$$

for all i and j . A real matrix satisfying these properties is called a stochastic matrix. An equilibrium probability distribution in this case is a solution to the equation

$$p = p A . \quad (7)$$

Because of the obvious similarity between Eqs. (4) and (7), RQA-1 can conveniently solve either equation through a suitable specification of the input. The most common application, however, is expected to be the continuous-time case, and, hence, a major part of the theoretical discussion and examples is devoted to that case.

The state space will usually originate as a multidimensional vector space with components corresponding to lengths of queues, numbers of occupied servers of a particular type, and similar quantities. However, since the number of such states is assumed finite, it is always possible to represent the state as a single variable i whose values are in correspondence to distinct values of the vector. The variable i is the state variable described in the preceding paragraphs. Nevertheless, the elements p_i of the solution to Eqs. (4) and (7) can be interpreted as joint probabilities of the variables represented in the original multidimensional state space, and appropriate marginal and conditional distributions can be calculated from them as desired.

This report will first treat, in Section 2, the numerical methods which are applied to the solution of Eqs. (3) and (7). This will be followed in Section 3, by a treatment of the technique used for the storage of the transition (intensity) matrix information in the form of what we call a "transition table" and several illustrations of the technique needed for the modeling of systems and the determination of transition tables. Section 4 and Appendix A constitute a description of the structure of the program and the function of its component routines, while Section 5 explains in detail how the program is applied. A person whose only purpose in using RQA-1 is to obtain a set of equilibrium joint probability distributions should be able to skip Sections 2 and 4 without major loss, while a person desiring to alter or to transliterate the program for other computers or purposes will, of course, be most interested in those sections.

The conclusion, Section 6, will discuss the performance of the program and some of the problems associated with its use.

2. NUMERICAL ANALYSIS FOR RQA-1

The RQA-1 employs an iterative procedure to determine the equilibrium probabilities for homogeneous continuous-time and discrete-time Markov chains. An equilibrium probability is found even when the chain is cyclic or when the chain does not have a unique equilibrium probability distribution.

This section is concerned with the conditions for convergence and for uniqueness of solutions. It details the mathematical justification for the procedures used.

2.1 CONDITIONS FOR CONVERGENCE

If G is an arbitrary m -th degree complex matrix, we will say that the iteration process

$$p^{(n+1)} = p^{(n)} G, \quad (8)$$

where $p^{(n)}$ is an m -dimensional row vector, converges if a nonzero vector

$$\lim_{n \rightarrow \infty} p^{(n)} = \pi \quad (9)$$

exists. We will say that the iteration process converges uniquely if it converges and if the vector, π , is unique (i.e., independent of the choice of $p^{(0)}$, the initial iterate). Clearly, convergence of the process to π implies that $\pi = \pi G$.

A set of sufficient conditions for the iteration process (Eq. (8)) to converge is that there be an eigenvalue of G at unity, that all other eigenvalues of G which are not equal to unity have modulus less than unity, and that all the unity eigenvalues of G have elementary divisors² of first degree. For, from Eq. (8) we may clearly write that

$$p^{(n)} = p^{(0)} G^n. \quad (10)$$

Then from the confluent form of Sylvester's Theorem,³ we may also write (making use of the above conditions) that

$$p^{(0)} G^n = \lambda_1^n (a_1 y_1 + a_2 y_2 + \dots + a_r y_r) + o(\lambda_{r+1}^n) \quad (11)$$

where λ_i represents the i -th eigenvalue (assumed ordered so that $\lambda_1 = \lambda_2 = \dots = \lambda_r > |\lambda_{r+1}| \geq \dots \geq |\lambda_m|$), r is the multiplicity of the dominant eigenvalue λ_1 , the a_i are scalar constants, and the y_i are m -dimensional row-vectors. If $\lambda_1 = 1$, Eq. (11) evidently converges to $(a_1 y_1 + \dots + a_r y_r)$, a row-vector. Note also that the rate of convergence is determined by λ_{r+1} , the subdominant eigenvalue.

2.2 CONVERGENCE FOR CONTINUOUS-TIME MARKOV CHAINS

For the continuous-time Markov chains, we are concerned with the convergence of the process

$$p^{(n+1)} = p^{(n)} (\Delta Q + I) \quad (12)$$

where Q is a transition intensity matrix.

Let $R = \max_i |q_{ii}|$. Then, if $\Delta \leq 1/R$, the matrix $(\Delta Q + I)$ is found to be a stochastic matrix, and, hence, it can be shown⁴ that every elementary divisor corresponding to a unity eigenvalue is of the first degree. Since the degree of these elementary divisors is not dependent upon the value of Δ , this implies that they are of first degree regardless of Δ as long as $\Delta \leq \frac{1}{R}$.

Moreover, it is a direct consequence of a theorem of Gershgorin⁵ that the eigenvalues of a transition intensity matrix Q will lie on the disc of complex values γ such that

$$|\gamma - R| \leq R \quad (13)$$

where $R = \max_i |q_{ii}|$. The eigenvalues of $\Delta Q + I$ must, therefore, lie on a **disc** of complex values such that

$$|\lambda - (\Delta R + 1)| \leq \Delta R. \quad (14)$$

Thus, if Δ is chosen so that $\Delta < 1/R$, then there can be no eigenvalues of $(\Delta Q + I)$ having unit modulus which are not of unity value, and the iteration process (Eq. (12)) is convergent to a solution of Eq. (4).

On the other hand, while a choice of $\Delta < 1/R$ is sufficient for convergence, it is not always necessary. Indeed, a choice of Δ slightly greater than $1/R$ can often lead to faster convergence than one slightly less, but one cannot know by any simple test when this is so. Hence, in RQA-1, Δ is chosen equal to $.99/R$ to guarantee convergence for continuous-time models.

2.3 CONVERGENCE FOR DISCRETE-TIME MARKOV CHAINS

For discrete-time Markov chains, we are concerned with the convergence of

$$p^{(n+1)} = p^{(n)} A \quad (15)$$

where A is a stochastic matrix. As a result of the theorems cited above, A is known to have only elementary divisors of first degree associated with its unity eigenvalues and to have eigenvalues lying on the unit disc. In order to guarantee convergence, one may transform Eq. (15) into the iteration process

$$p^{(n+1)} = p^{(n)} (.99A + .01I), \quad (16)$$

noting that

$$p = p(.99A + .01I) \quad (17)$$

has the same set of solutions as Eq. (7). As a result, the eigenvalues of $(.99A + .01I)$ lie on the disc of complex values λ' such that

$$|\lambda' - .01| \leq .99 \quad (18)$$

which does not admit values of unit modulus not equal to unity. The iteration process (Eq. (16)) is used in RQA-1 for discrete-time models.

2.4 UNIQUE CONVERGENCE

Under certain circumstances, a Markov chain may not have a unique equilibrium probability distribution. These are the cases in which the limiting probability distribution ($\lim_{t \rightarrow \infty} p_t$ or $\lim_{n \rightarrow \infty} p_n$) is dependent upon the initial distribution, p_0 . It is in those circumstances that the RQA-1 algorithm will not con-

verge uniquely. It converges, rather, to an equilibrium distribution which is consistent with the initial iterate chosen.

This should cause no difficulty since such cases are rare in practice, and the user of RQA-1 will usually be aware of the dependence upon initial conditions through knowledge of his model. In general, a sufficient condition for the convergent iteration processes above to converge uniquely is that it be possible (with nonzero probability) to reach each state from every other state in a finite time. A transition matrix having this property (known as irreducibility) is called a primitive matrix. Necessary conditions for uniqueness are somewhat stronger.

If nonuniqueness of a solution is suspected, a simple technique can be used to test the model. The solution can be found by starting with an initial iterate having unit value for some (arbitrary) state and zero for all others. If the resulting solution has no states with a calculated zero probability, the solution is unique. If it has states with a calculated zero probability, then the solution process should be repeated using as an initial iterate a probability distribution which is nonzero wherever the previous solution was zero (within computational error) and zero wherever the previous solution was nonzero. If the new solution is within the calculational error of the previous solution, it is the unique solution. If not, there is no unique solution.

For more detail on the conditions leading to nonunique equilibrium distributions, Parzen,⁶ Gantmacher,⁴ or Chung⁷ are useful references.

2.5 ESTIMATION OF ERROR

It can be shown that, if the elementary divisors of the subdominant eigenvalues λ_{r+1} of a stochastic matrix G are of a degree m which is small compared to n , and if n is sufficiently large, then the n -th iterate can be approximated roughly by

$$p^{(n)} \triangleq G^n p^{(0)} \approx \pi + \lambda_{r+1}^n n^{m-1} x \quad (19)$$

where x is a constant vector and π is the convergent solution. Also, it can be shown that the difference between successive iterates can be approximated roughly by

$$p^{(n+1)} - p^{(n)} \approx \lambda_{r+1}^n (\lambda_{r+1} - 1) n^{m-1} x. \quad (20)$$

A fairly crude, but useful, estimate $\hat{\lambda}_{r+1}$ of $|\lambda_{r+1}|$ can be obtained from

$$\hat{\lambda}_{r+1} \triangleq |p^{(n+1)} - p^{(n)}|^{\frac{1}{n}} \approx |\lambda_{r+1}| |\lambda_{r+1} - 1|^{\frac{1}{n}} \frac{m-1}{n} x^{\frac{1}{n}} \quad (21)$$

since, for n sufficiently large, all but the first multiplicative term (λ_{r+1}) approach unity. Using this estimate, it can be argued from Eq. (20) that

$$\frac{p^{(n+1)} - p^{(n)}}{1 - \hat{\lambda}_{r+1}} \approx \lambda_{r+1}^n n^{m-1} x \quad (22)$$

which is an approximation to the error term in Eq. (19).

Provided that the $\hat{\lambda}_{r+1}$ calculated is less than unity and greater than about .5, $(\lambda_{r+1} - 1)^{1/R}$ can be assumed to be near enough to unity, and the estimate so provided is, at the very least, a more realistic estimate of error than $(p^{(n+1)} - p^{(n)})$.

2.6 REPRESENTATION OF MATRICES

In order to store information completely describing matrices A or Q , with A or Q having the requisite high degree, a special approach is required. Moreover, because it is desirable to have the matrix elements expressed as algebraic functions of a set of parameters which may assume different values at different executions of the program, this approach must allow expression of the matrix elements in a literal form in the program. More will be said about the latter requirement in Section 3. For now, we will consider the problem of compressing the information needed to describe A or Q in such a way that computations are not appreciably slowed by the decoding required.

Clearly, a 5,000-degree matrix, when stored as a two-index array, requires 25,000,000 locations of storage which is unreasonable for a "fast" program. However, both A and Q are generally sparse matrices (have mostly zero-valued elements) and will usually have a high degree of repetition of equal element values. Hence, a scheme of storage which lists location information along with value information is a plausible starting point.

We construct a set of four vectors, together called a transition table, which implicitly define the matrix. Let us call them α , β , γ , and B^* and denote their i -th elements by α_i , β_i , γ_i , and B_i , respectively. The quadruple $(\alpha_i, \beta_i, \gamma_i, B_i)$ specifies one or more elements of a matrix in the following manner.

*In the program listing, these correspond to the variables named "TRANS," "FIRST," "DEST," and "LAST," respectively.

The value of the element is α_i , and its matrix co-ordinates are (β_i, γ_i) . Due to the repetitive nature of these matrices, the element α_i may occur in other locations of the matrix with co-ordinates $(\beta_i + r\delta, \gamma_i + r\delta)$ where δ is a constant (fixed throughout the transition table) and r takes values $0, 1, 2, \dots, (B_i - \beta_i)/\delta$. In other words, the quadruple $(\alpha_i, \beta_i, \gamma_i, B_i)$ specifies the occurrence in the matrix of an element with value α_i at co-ordinates $(\beta_i, \gamma_i), (\beta_i + \delta, \gamma_i + \delta), \dots, (B_i, \gamma_i + (B_i - \beta_i))$. This set of co-ordinates will be called the range of the quadruple. The set $\{\beta_i, \beta_i + \delta, \dots, B_i\}$ will be called the β -range of the quadruple. The constant δ will be called the period of the table.

The choice of a value for δ will be discussed in detail in Section 3. Also in Section 3, we will carry out the construction of transition tables for three systems of varying degrees of complexity. These systems will be characterized as consisting of a single dominant queueing facility and a complex service facility. The degree of the Q-matrix for such a system will depend on the maximum number of jobs allowed in the queue. This number will be a parameter in the transition table, and, in this way, a transition table of fixed size can represent a Q matrix of arbitrary order.

2.7 TIMING

Each iteration (Eqs. (12) or (16)) is carried out in the following sequence:

- (1) The vector to contain p^{n+1} is initially zeroed.
- (2) α_i is multiplied by $p_{\beta_i}^{(n)}$ and accumulated into $p_{\gamma_i}^{(n+1)}$.
- (3) α_i is multiplied by $p_{\beta_i+\delta}^{(n)}$ and accumulated into $p_{\gamma_i+\delta}^{(n+1)}$
if $\beta_i + \delta \leq B_i$.
- (4) Step (3) is repeated until the range of the quadruple is exhausted.
- (5) Steps (2) through (4) are repeated using $\alpha_i, \beta_i, \gamma_i, B$ for
 $i = 2, 3, \dots$ until all quadruples have been treated.

As a result of this procedure, the number of multiplications per iteration is equal to the number of nonzero elements in the transition matrix. In general, if the average number of different events which can occur at any time is E and the number of states is S , then the number of multiplications per iteration is ES . The convergence error after n iterates is $O(\lambda_{r+1}^n)$. Thus, if ϵ is the desired error, the number of iterations needed will be $I \approx \log \epsilon / \log |\lambda_{r+1}|$, and the total number of multiplications required in the main loop is

$$M = ES \log \epsilon / \log |\lambda_{r+1}| . \quad (23)$$

In contrast, a rule of thumb estimate for the standard deviation of a measurement of a probability p in a continuous-time Markov chain using simulation is

$$S \approx \frac{2p}{|\gamma|T} \quad (24)$$

where T is the length of the simulation (in system-time) and γ is the nonzero eigenvalue of the matrix Q having smallest modulus. For ϵ to be equal to two standard deviations, this means that T should be

$$T = \frac{8p}{|\gamma|\epsilon^2} , \quad (25)$$

and the number of random numbers generated would need to be

$$K = \frac{8 v p}{|\gamma|\epsilon^2} \quad (26)$$

where v is the mean rate of occurrence of events.

It can be shown that $|\gamma|/v$ is usually of the order of $-\log |\lambda_{r+1}|$ which, in turn, is normally much smaller than unity. Hence, a relatively crude measure of the computation time of simulation and RQA-1 calculation is obtained from Eqs. (23) and (26) by approximating the ratio K/M by

$$\frac{K}{M} = 0 \left(\frac{8p}{\epsilon^2 \log \epsilon} \cdot \frac{1}{ES} \right) . \quad (27)$$

Since $\epsilon^2 \log \epsilon$ is very small for $\epsilon < .01$,

ϵ	$\frac{1}{\epsilon^2 \log \epsilon}$
.1	43.5
.01	2,170
.001	145,000 ,

we conclude that RQA-1 has a large advantage over simulation if desired accuracy is moderately high. This conclusion is strengthened considerably by the fact that a much greater amount of "overhead" calculation is required in simulations than with the RQA-1 algorithm.

3. EXAMPLES OF TRANSITION TABLES

The preparation of the transition table is the most technical aspect of RQA-1 use. As stated in the previous section, the aim is to code a representation for a transition matrix which has the flexibility to include a wide class of problems and which makes efficient use of storage. The general format of the transition table has already been discussed. The purpose of this section is to illustrate the procedure by which the set of quadruples can be developed when a statement of the control rules and the structure of a Markovian queueing system is given. The examples are obviously not exhaustive, but it is hoped that they are sufficiently representative and that a user of RQA-1 familiar with them will have a minimum of trouble deriving transition tables for a wide variety of problems. Future reports will expand upon the process of modeling general, large queueing systems by Markov chains and by RQA-1 transition tables.

The procedure illustrated here is capable of being mechanized by computer programs to some degree once the principles are understood. In general, however, each queueing model represents a distinct programming problem, and it appears difficult to generalize to automatic techniques. For this reason, the present report will limit itself to essentials. More sophisticated general methods are the object of current research by the Laboratory.

All three examples discussed in this section represent continuous-time Markov chains containing one main queue of jobs to be served by a service facility. In each succeeding example, the service facility has an increasingly complex structure. The derived transition table in each case represents a transition intensity matrix Q .

The pattern to be followed in transcribing a Markovian model for RQA-1 is:

(1) Determine the state space and a suitable mapping of the states (which are almost always vectors) to a set of integers.

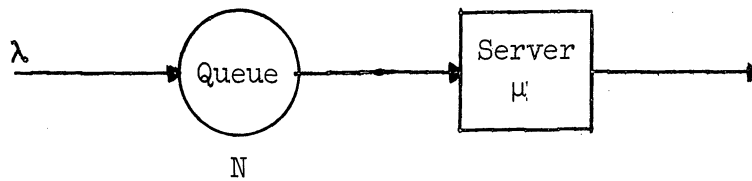
(2) Determine the most useful value of δ . This will generally be the value of δ which allows the transitions to be described with the least number of quadruples. The choice of state-space mapping strongly influences what is possible here, and steps (1) and (2) may have to be attempted several times for a complex model. Computation time does not depend upon the number of quadruples used. Thus, the primary necessity here is to remain within available storage where space for only 2,000 quadruples is allotted. A secondary necessity is to keep the transition table preparation from becoming unnecessarily unwieldy.

(3) Choose a sequence for generating the quadruples. This sequence must insure that every possible event of the model is considered for every

possible state of the model and that no event is considered more than once for any state. It must also insure that, when an event-state pair is considered, it is incorporated by extending the range of an already prepared quadruple whenever possible rather than creating a new quadruple.

3.1 SIMPLE SERVER SYSTEM

The first and simplest example is illustrated in the diagram below:



Jobs arrive as a Poisson process with mean rate λ . The duration of job processing time is independent of arrival and is exponentially distributed with mean time $1/\mu$. Jobs are executed on a first-come first-served basis. If more than N jobs are present in the queue, arrivals are simply rejected.

The first step is to define the state space for the system. Because of the simplicity of the example, the definition of state is artificial in this instance but provides a pattern for future use. The system state is defined by the number of jobs in the waiting line and by whether or not the service facility is occupied. Thus, a state is designated by the ordered pair (i, j) where i takes values in the set $I = \{0, 1, 2, \dots, N\}$ and j takes values in the set $J = \{0, 1\}$. In this case, the cardinality C of the set J is 2.

It is apparent that not all ordered pairs (i, j) in the Cartesian product $I \times J$ represent states. In fact, the set $\{(i, 0) \mid i > 0\}$ is not a set of states, since, physically, a queue can exist only when there is a job in service. In general, the Cartesian product $I \times J$ can be partitioned into two sets: the set of states of the system, and the set of ordered pairs which are not states which we will call points. Thus, for this example, the set of states of the system is $\{(0, 0), (0, 1), (1, 1), (2, 1), \dots, (N, 1)\}$, and the set of points is $\{(1, 0), (2, 0), \dots, (N, 0)\}$.

As stated previously, the state description within the program is one-dimensional. A natural mapping, T , from the two-dimensional description (i, j) to the single dimension, k , is

$$k = T(i, j) = Ci + j + 1. \quad (28)$$

Thus, the one-dimensional state description k runs from $k = 1$ corresponding to $(i, j) = (0, 0)$ to $k = C(N + 1)$ corresponding to $(i, j) = (N, C - 1)$ which is equal to $(N, 1)$, in this case. In this example, the set of states is $\{1, 2, 4, 6, \dots, 2N + 2\}$.

At any time, the system can change state due to a job arrival or a service completion or can remain in the same state if neither of these occurs. The transition table can be partitioned into three classes of quadruples: those associated with job arrivals, those associated with service completions, and those associated with no change, i.e.,

- (1) Job arrivals,
- (2) service completions, and
- (3) no change.

We can subdivide each of the above classes by the first dimension

- (a) $i = 0$;
- (b) $i = 1$;
- (c) $i = 2$, etc.,

and further subdivide them by the second dimension

- (i) $j = 0$
- (ii) $j = 1$.

A subclass is interpreted as a collection of transitions from a particular state caused by a particular event. For example, subclass (1) (b) (ii) represents the transitions from state $(1, 1)$ caused by job arrivals.

One can then create quadruples of the form $(\alpha, \beta, \gamma, B)$ by sequencing through each subclass, making certain that the transition rate associated with each subclass corresponds to the α of a quadruple, and that the state associated with the subclass is in the β -range of the same quadruple (see Section 2.4). The state corresponding to the destination of the transition must be the γ corresponding to that β . The following detailed construction should clarify the above discussion.

In this example, δ , the period of the transition table, is chosen to be equal to 2, since, as will be seen, the states to which equal rates apply differ most often by 2. It will usually be the case, when the model contains a main queue, that δ is most conveniently chosen to be equal to C , the cardinality of the set J of all possible combinations of values of all state variables other than the queue length. However, in general, δ may be any other value (constant throughout the transition table) which allows the matrix to be described with a smaller number of quadruples.

Beginning with the subclass (1) (a) (i), the first two parts of a quadruple can immediately be written,

$$\alpha = \lambda$$

$$\beta = 1 .$$

The fact that $\alpha = \lambda$ is obvious since class (1) is concerned with arrivals and jobs arrive at the system with a rate λ . $\beta = 1$ is the result of applying the mapping T to the state designated by (a) (i), namely (0,0). To determine γ , we note that an arrival in state (0,0) produces a transition to state (0,1) and, therefore, $\gamma = T(0,1) = 2$. To determine B, assume $B = r\delta + \beta$ for $r \neq 0$; then if λ applies for the transition

$$1 \rightarrow 2 ,$$

it should next apply for the transition

$$1 + \delta \rightarrow 2 + \delta$$

$$\text{i.e.,} \quad 3 \rightarrow 4 .$$

But in the listing of states, 3 does not occur and hence $B = \beta = 1$. Thus, the first quadruple $(\lambda, 1, 2, 1)$ applies only for the transition from state 1 to state 2 and, therefore, refers only to state 1, or subclass (1) (a) (i).

The next subclass is (1) (a) (ii). We can immediately write

$$\alpha = \lambda$$

$$\beta = T(0,1) = 2 .$$

An arrival in state (0,1) produces a transition to state (1,1) and, therefore, $\gamma = T(1,1) = 4$. Again, assume $B = r\delta + \beta$ for $r \neq 0$; then λ will apply for the transition

$$2 + \delta \rightarrow 4 + \delta$$

$$\text{i.e.,} \quad 4 \rightarrow 6 .$$

To determine B, it is necessary to find the maximum r such that

$$r \delta + \gamma \leq \delta(N + 1)$$

and such that the corresponding transition is a correct description of an arrival, i.e., in this case, $r = N + 1 - 4/2 = N - 1$. In general, the maximum r is

desired for which the transition intensity α for states $\beta + r\delta$ to $\gamma + r\delta$ is still valid. For all the examples to be treated, this will correspond to the maximum r for which $\gamma + r\delta$ is in the state space. Therefore, $B = 2(N-1)+2 = 2N$. Thus, the second quadruple $(\lambda, 2, 4, 2N)$ has the range $\{(2,4), (4,6), \dots, (2N, 2N+2)\}$ and, therefore, refers to the subclasses (1) (a) (ii), (1) (b) (ii), (1) (c) (ii) etc. Since (b) (i), (c) (i), etc., are not states, all the subclasses of class (1) have been referred to. A similar procedure now begins for class (2).

The β for the quadruple associated with subdivision (2) (a) (i) would be 1, but, since there can be no service completions in the 1 state, the corresponding α is zero and requires no quadruple. We next turn to (2) (a) (ii) for which $\beta = 2$ and $\alpha = \mu$, the service rate. A service completion in state 2 produces a transition to state 1, i.e., $\gamma = 1$. Assuming $B = r\delta + \beta$, for $r \neq 0$, the next transition should be

$$4 \rightarrow 3 .$$

But this is impossible since a completion does not take state 4 into state 3 (which in fact, is not a state). Hence, the resulting quadruple is $(\mu, 2, 1, 2)$ and refers only to subclass (2) (a) (ii).

Since (b) (i), (c) (i), etc., are not states, the next subclass is (2) (b) (ii) for which $\alpha = \mu$ and $\beta = 4$. The corresponding γ is 2. B will be determined from the maximum r such that

$$r\delta + \gamma \leq \delta (N + 1)$$

$$\text{i.e., } r = N + 1 - 2/2 = N .$$

Therefore, $B = r\delta + \beta = 2N + 2$. The quadruple $(\mu, 4, 2, 2N + 2)$ has the range $\{(4,2), (6,4), \dots, (2N + 2, 2N)\}$ which refers to subclasses (2) (b) (ii), (2) (c) (ii), etc., and thereby completes class (2).

The quadruples for class (3) can easily be constructed by using the property of Q that the row sums are zero. For all these quadruples, it is also the case that $\beta = \gamma$. For class (3) (a) (i), it can be readily determined that

$$\alpha = -\lambda$$

$$\beta = 1$$

$$\gamma = 1$$

since the only way to leave state 1 is through an arrival with transition intensity λ . Thus, the practical rule for determining the α 's of quadruples for

class (3) is that α equals minus the sum of all transition intensities associated with leaving that particular state. For the above quadruple, it is apparent that $B = 1$ since state 1 is the only state for which there will be no change if there are no arrivals.

The next subclass is (3) (a) (ii) for which we write

$$\alpha = -(\lambda + \mu)$$

$$\beta = 2$$

$$\gamma = 2.$$

The fact that $\alpha = -(\lambda + \mu)$ is an illustration of the above rule where the sum of the transition intensities is $\lambda + \mu$. This quadruple will refer to all states which have the same properties as state 2, i.e., states 4, 6, ..., 2N. Hence, B is 2N, and the quadruple is $(-(\lambda + \mu), 2, 2, 2N)$.

For the final state, $2N + 2$, the only transition intensity associated with leaving is μ ; therefore, the necessary quadruple is $(-\mu, 2N + 2, 2N + 2, 2N + 2)$.

This completes class (3), thereby completing references to all subclasses. The transition table is listed in its entirety in Table I. We will now go on to consider a more complex problem in considerably less detail.

TABLE I

TRANSITION TABLE—"SIMPLE SERVER" MODEL

	α	β	γ	B
Job Arrivals	λ	1	2	1
	λ	2	4	2N
Service Completions	μ	2	1	2
	μ	4	2	2N + 2
No Changes	$-\lambda$	1	1	1
	$-\lambda - \mu$	2	2	2N
	$-\mu$	2N + 2	2N + 2	2N + 2

3.2 FIXED-SIZE MEMORY-SHARED COMPUTER MODEL ("FIXMEM" MODEL)

This model was discussed in the paper by Fife and Rosenberg.⁸ It is depicted in Fig. 1.

Jobs arrive at the loading queue as a Poisson process with mean rate λ . Each job finds some number n of jobs in queue ahead of it. The queue is cleared on a first-come first-served basis, and the job enters the first stage of "service" where the mean rate of loading of jobs is μ_1 and the number of jobs being loaded is n_1 ($n_1 = 0$ or 1). If more than N jobs are in queue, the arriving job is rejected.

After loading is completed, a job enters the second state of service, the processing phase. There are 5 "channels" in this stage, and n_2 is the number which are occupied. The mean rate at which jobs are processed in a single block or channel is μ_2 . A job leaves the system when its processing time ends.

The principal complexity of the model is introduced by the following constraint:

$$n_1 + n_2 \leq 5 .$$

This implies that loading for no job can be begin while all five blocks are processing jobs.

The state description consists of the ordered triple (n, n_1, n_2) where n takes values $0, 1, \dots, N$, n_1 takes values zero or 1, and n_2 takes values $0, 1, 2, 3, 4, 5$, with the restriction that $n_1 + n_2 \leq 5$. Table II lists all the 11 possible occupancy configurations (n_1, n_2) and associates a value j to each.

TABLE II
INTERNAL-STATE DESIGNATION—"FIXMEM" MODEL

n_1	n_2	Internal State j
0	0	1
0	1	2
0	2	3
0	3	4
0	4	5
0	5	6
1	0	7
1	1	8
1	2	9
1	3	10
1	4	11

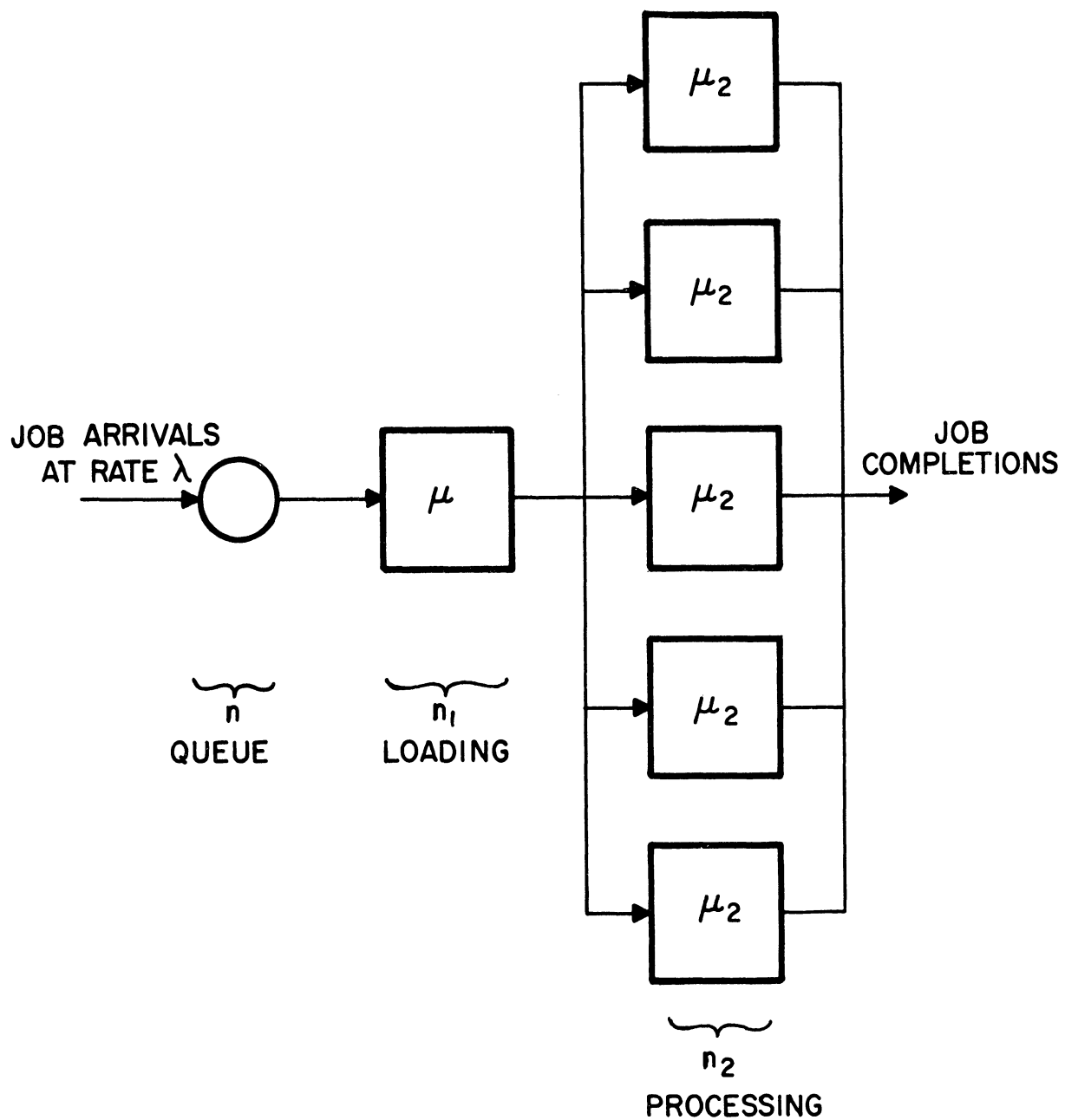


Fig. 1. Fixed-size memory-shared computer model.

The values of j are here called "internal states." Thus, for this example, we have the set $I = \{1,2,\dots,N\}$ and the set $J = \{1,2,\dots,11\}$ and, again, we take $\delta = C = 11$ where C is the cardinality of J . The mapping T for this example reduces to

$$k = T(i,j) = 11i + j .$$

The set of "points" is $\{(i,j) | i > 0, j = 1,2,3,4,5\}$.

Table III lists the 47 quadruples which are necessary to refer to all the subclasses of this system. There are more subclasses for the example than for the simple server system since there are 11 internal states rather than 2. We will not carry out the complete derivation of this table but, rather, we will consider a few illustrative cases.

We will derive the quadruples which refer to state $(0,6)$ for all three classes. For class (1), we can write

$$\alpha = \lambda$$

$$\beta = 6 .$$

An arrival at state 6 will remain in the queue resulting in state $(1,6)$, i.e., state 17. Thus $\gamma = 17$. It is obvious that later arrivals will also remain in the queue. Thus, we are looking for a $B = w\delta + \beta$ such that $w\delta + \gamma \leq \delta(N+1)$ where w is the largest integer for which this inequality holds. The above yields

$$w \leq N + 1 - 17/11 .$$

Therefore,

$$w = N - 1 .$$

This results in $B = 11(N-1) + 6$. The quadruple $(\lambda, 6, 17, 11(N-1) + 6)$, which is number (6) in the listing, has the range $\{(6,17), (17,28), \dots, (11(N-1) + 6, 11N + 6)\}$.

For class (2), we write

$$\alpha = 5\mu_2$$

$$\beta = 6$$

since state 6 represents the configuration in which all 5 processing blocks are occupied, and jobs are processed at a mean rate of μ_1 per block, the rate for

TABLE III
"FIXMEM" MODEL QUADRUPLES

No.	α	β	γ	B
(a) <u>Job Arrivals</u>				
1	λ	1	7	1
2	λ	2	8	2
3	λ	3	9	3
4	λ	4	10	4
5	λ	5	11	5
6	λ	6	17	$11(N-1) + 6$
7	λ	7	18	$11(N-1) + 7$
8	λ	8	19	$11(N-1) + 8$
9	λ	9	20	$11(N-1) + 9$
10	λ	10	21	$11(N-1) + 10$
11	λ	11	22	$11(N-1) + 11$
(b) <u>Service Completions</u>				
12	μ_2	2	1	2
13	$2\mu_2$	3	2	3
14	$3\mu_2$	4	3	4
15	$4\mu_2$	5	4	5
16	$5\mu_2$	6	5	6
17	$5\mu_2$	17	11	$11N + 6$
18	μ_1	7	2	7
19	μ_1	18	8	$11N + 7$
20	μ_2	8	7	$11N + 8$
21	μ_1	8	3	8
22	μ_1	19	9	$11N + 8$
23	$2\mu_2$	9	8	$11N + 9$
24	μ_1	9	4	9
25	μ_1	20	10	$11N + 9$
26	$3\mu_2$	10	9	$11N + 10$
27	μ_1	10	5	10
28	μ_1	21	11	$11N + 10$
29	$4\mu_2$	11	10	$11N + 11$
30	μ_1	11	6	$11N + 11$
(c) <u>No Changes</u>				
31	$-\lambda$	1	1	1
32	$-\lambda - \mu_2$	2	2	2
33	$-\lambda - 2\mu_2$	3	3	3
34	$-\lambda - 3\mu_2$	4	4	4
35	$-\lambda - 4\mu_2$	5	5	5
36	$-\lambda - 5\mu_2$	6	6	$11(N-1) + 6$
37	$-5\mu_2$	$11N + 6$	$11N + 6$	$11N + 6$
38	$-\lambda - \mu_1$	7	7	$11(N-1) + 7$
39	$-\mu_1$	$11N + 7$	$11N + 7$	$11N + 7$
40	$-\lambda - \mu_1 - \mu_2$	8	8	$11(N-1) + 8$
41	$-\mu_1 - \mu_2$	$11N + 8$	$11N + 8$	$11N + 8$
42	$-\lambda - \mu_1 - 2\mu_2$	9	9	$11(N-1) + 9$
43	$-\mu_1 - 2\mu_2$	$11N + 9$	$11N + 9$	$11N + 9$
44	$-\lambda - \mu_1 - 3\mu_2$	10	10	$11(N-1) + 10$
45	$-\mu_1 - 3\mu_2$	$11N + 10$	$11N + 10$	$11N + 10$
46	$-\lambda - \mu_1 - 4\mu_2$	11	11	$11(N-1) + 11$
47	$-\mu_1 - 4\mu_2$	$11N + 11$	$11N + 11$	$11N + 11$

l blocks in use, where $l \leq 5$ is $l\mu_1$. In this instance, $l = 5$, and, therefore, $\alpha = 5\mu_2$. A single service completion will result in a state in which $l = 4$, and, therefore, $\gamma = 5$. If there is a queue, then a service completion in the processing block will result in the job at the head of the queue moving into the loading phase. Thus, two quadruples, namely $(5\mu_2, 6, 5, 6)$ and $(5\mu_2, 17, 11, 11N+6)$, will be necessary to form the range $\{(6,5), (17,11), (28,22), \dots, (11N+6, 11N)\}$.

Class (3) will require two quadruples for the same range: for the first $\alpha = -(\lambda + 5\mu_2)$, and for the second $\alpha = -5\mu_2$. These quadruples are numbered (36) and (37), respectively. For the states $6, 17, 28, \dots, 11(N-1) + 6$, transitions are possible in two ways: an arrival with transition intensity λ , and a service completion with transition intensity $5\mu_2$. For state $11N + 6$, a transition is only possible through a service completion.

For a final case, we will derive the quadruples which refer to state (0,9) for all three classes. State 9 is the service configuration consisting of 2 jobs being processed and 1 job being loaded. Thus, there are two ways for service completions to occur.

The quadruple associated with class (1) is number (9) in the listing, and, at this point, its derivation should be straightforward.

For class (2), the relevant quadruples are numbered (23), (24) and (25). The α for quadruple (23) is $2\mu_2$ which is the transition intensity associated with a service completion in the processing blocks. The range of this quadruple is $\{(9,8), (20,19), \dots, (11N+9, 11N+8)\}$. The service configuration for state 8 is shown in Table II. There are two quadruples necessary to describe service completions in the loading block. When there are no jobs in the queue, a service completion in the loading block results in an additional processing block being employed, i.e., state 4 results. When there are jobs in the queue, the job at the head of the queue moves into the recently vacated loading block and state 10 results.

The class (3) quadruples are numbered (42) and (43). The range of quadruple (42) is $\{(9,9), (20,20), \dots, (11(N-1) + 9, 11(N-1) + 9)\}$, and its α is $-(\lambda + \mu_1 + 2\mu_2)$, as expected. Quadruple (42) refers to state $11N + 9$ with an α of $-(\mu_1 + 2\mu_2)$.

These few cases should be sufficient to indicate the structure of the entire transition table. The next example will introduce an additional complexity in that service completions will depend more closely upon the nature of job arrivals.

3.3 RANDOM-SIZE MEMORY-SHARED COMPUTER MODEL ("RANMEM" MODEL)*

The model depicted in Fig. 2 bears a resemblance to the previous model discussed, but it is only a superficial one. As before, the queue will be cleared on a first-come first-served basis. There will be three types of jobs arriving, distinguishable by the number of service blocks or storage units they require. The jobs will arrive as a Poisson process with mean arrival rates $\lambda_1, \lambda_2, \lambda_3$ for one, two, and three storage unit jobs, respectively. These rates are independent of time and state of the system. For convenience, we set $\lambda = \lambda_1 + \lambda_2 + \lambda_3$.

We assume that the system is nonpreemptive, i.e., a job once begun is completed. Each type of job will have a basic mean service rate when it is being processed alone. These service rates are μ_1, μ_2, μ_3 for one, two, and three storage unit jobs, respectively. However, when two jobs are being processed simultaneously, the mean service rate at which each job is completed is assumed to be one-half the basic mean service rate. For example, if a one-unit and a two-unit job are being processed, the mean service rate for the one-unit job is $(1/2)\mu_1$ and for the two-unit job $(1/2)\mu_2$. Similarly, for three one-unit jobs, each will have a mean service rate of $(1/3)\mu_1$. Upon completion of any type of job, all blocks being used for that job become free simultaneously.

This model is particularly interesting because it requires the evaluation of conditional probabilities (conditional upon the state of the system) in order to determine the correct transition intensities. This is true because jobs having different properties are allowed to be mixed in the queue, and the nature of the first job in the queue is dependent statistically upon the condition of the system.

The state description consists of the ordered quadruple (n, n_1, n_2, n_3) where n is the number of jobs in the waiting line, and n_1, n_2, n_3 are the numbers of one-, two-, and three-unit jobs in service, respectively. Table IV lists all

TABLE IV
INTERNAL STATE DESIGNATION—"RANMEM" MODEL

n_1	n_2	n_3	Internal State, j
0	0	0	1
0	0	1	2
0	1	0	3
1	1	0	4
1	0	0	5
2	0	0	6
3	0	0	7

*This model corresponds to one suggested by Dr. R. V. Evans, now of UCLA, for multiprogrammed time-shared computer operation.⁹

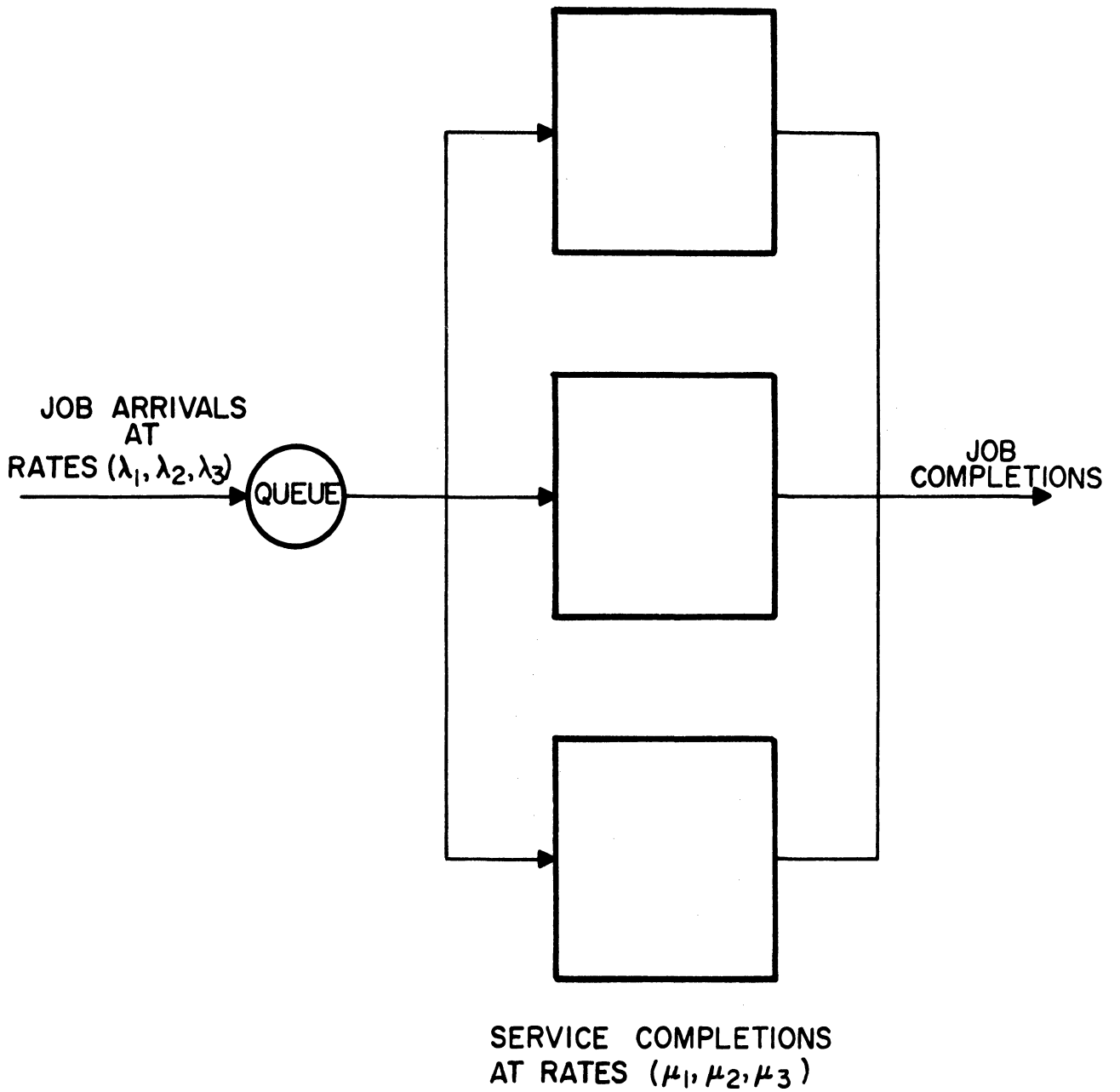


Fig. 2. Random-size memory-shared computer model (RANMEM)

the possible occupancy configurations of which there are 7. For this example, the set J is $\{1,2,\dots,7\}$ with $\delta = C = 7$, and the set I is $\{0,1,\dots,N\}$. The mapping T reduces to $k = T(i,j) = 7i + j$. We will again call the value of j the "internal state" of the system. The set of "points" is $\{(i,j) | i > 0, j = 1\}$.

Table V lists the 60 quadruples which are necessary to refer to all the subclasses of the state space. We will consider only one illustrative case.

We will derive the quadruples which refer to state $(n,3), n = 0,1,\dots,N$, for all three classes. State 3 represents a two-unit job in service. The arrival of a one-unit job with transition intensity λ_1 produces a transition to state 4. The arrivals of either a two-unit or a three-unit job with intensity $\lambda_2 + \lambda_3$ do not change the internal state but do increase the queue length. Both these quadruples (5) and (6) refer only to the single state 3, since for state 10, the arrival of any type of job (with intensity $\lambda = \lambda_1 + \lambda_2 + \lambda_3$) produces a transition to state 17. Then quadruple (14) will have the range $\{(10,17), (17,24), \dots, (7(N-1) + 3, 7N + 3)\}$. Thus, three quadruples are necessary for the class (1), subclasses $(n,3), n = 0,1,\dots,(N)$.

In state 3, a service completion will produce a transition to state 1 with intensity μ_2 . Quadruple (18) will have only the single element (3,1) in the range, because the transitions resulting from service completion of a two-unit job are quite different for state 10. In state 10, the single job in queue can only be a two-unit or a three-unit job (a one-unit job cannot remain in the queue when there is space for it in service. If, for example, only a two-unit job in service, the one-unit job cannot remain in the queue).

We have here the situation in which the identity of the first job in the queue depends partially on the type of job in service. For example, if there is a two-unit job in service, the first job in the queue cannot be a one-unit job. If there is a one-unit job in service, the first in queue cannot be a one-unit or a two-unit job but must be a three-unit job. This dependency will introduce complexities into the form of the transition intensities. If there is a two-unit job in the queue, a service completion will result in a transition with intensity $\mu_2\lambda_2/(\lambda_2 + \lambda_3)$; similarly, for a three-unit job in the queue, the transition intensity will be $\mu_2\lambda_3/(\lambda_2 + \lambda_3)$. In the first case, the $\lambda_2/(\lambda_2 + \lambda_3)$ factor is the probability of a two-unit job being in queue given that there is a two-unit or a three-unit job in queue.

Quadruples (27) and (28) account for these cases. Quadruple (2) has a complete range because, given a three-unit job at the head of the queue, the remaining composition of the queue is unimportant. But quadruple (27) has only the single element (10,3) in its range, since transitions from state 17 depend on the nature of the first two jobs in the queue if the first in queue is a two-unit job.

The only two cases to consider for state 17 are that the first in queue is a two-unit job and the second in queue is not a one-unit job, and that the

TABLE V

"RANMEN" MODEL QUADRUPLES

No.	α	β	γ	B
(a) <u>Job Arrivals</u>				
1	λ_1	1	5	1
2	λ_2	1	3	1
3	λ_3	1	2	1
4	λ	2	9	$7(N-1) + 2$
5	λ_1	3	4	3
6	$\lambda_2 + \lambda_3$	3	10	3
7	λ	4	11	$7(N-1) + 4$
8	λ_1	5	6	5
9	λ_2	5	4	5
10	λ_3	5	12	5
11	λ_1	6	7	6
12	$\lambda_2 + \lambda_3$	6	13	6
13	λ	7	14	$7(N-1) + 7$
14	λ	10	17	$7(N-1) + 3$
15	λ	12	19	$7(N-1) + 5$
16	λ	13	20	$7(N-1) + 6$
(b) <u>Service Completions</u>				
17	μ_3	2	1	2
18	μ_2	3	1	3
19	$\mu_1/2$	4	3	4
20	$\mu_2/2$	4	5	4
21	μ_1	5	1	5
22	μ_1	6	5	6
23	μ_1	7	6	7
24	$\mu_3\lambda_3/\lambda$	9	2	$7N + 2$
25	$\mu_3\lambda_2/\lambda$	9	3	9
26	$\mu_3\lambda_1/\lambda$	9	5	9
27	$\mu_2\lambda_2/(\lambda_2 + \lambda_3)$	10	3	10
28	$\mu_2\lambda_3/(\lambda_2 + \lambda_3)$	10	2	$7N + 3$
29	$\mu_1(\lambda_3 + \lambda_2)/2\lambda$	11	10	$7N + 4$
30	$\mu_2\lambda_3/2\lambda$	11	12	$7N + 4$
31	$(\mu_1\lambda_1 + \mu_2\lambda_2)/2\lambda$	11	4	$7N + 4$
32	$\mu_2\lambda_1/2\lambda$	11	6	11
33	$\mu_1\lambda_3/(\lambda_2 + \lambda_3)$	13	12	$7N + 6$
34	$\mu_1\lambda_2/(\lambda_2 + \lambda_3)$	13	4	$7N + 6$
35	μ_1	12	2	$7N + 5$
36	$\mu_1(\lambda_2 + \lambda_3)/\lambda$	14	13	$7N + 7$
37	$\mu_1\lambda_1/\lambda$	14	7	$7N + 7$
38	$\mu_3\lambda_2(\lambda_2 + \lambda_3)/\lambda^2$	16	10	$7N + 2$
39	$2\mu_3\lambda_2\lambda_1/\lambda^2$	16	4	$7N + 2$
40	$\mu_3\lambda_1\lambda_3/\lambda^2$	16	12	$7N + 2$
41	$\mu_3\lambda_1^2/\lambda$	16	16	16
42	$\mu_2\lambda_2/\lambda$	17	10	$7N + 3$
43	$\mu_2\lambda_2\lambda_1/(\lambda(\lambda_2 + \lambda_3))$	17	4	$7N + 3$
44	$\mu_2\lambda_1(\lambda_2 + \lambda_3)/2\lambda^2$	18	13	$7N + 4$
45	$\mu_2\lambda_1^2/2\lambda^2$	18	7	$7N + 4$
46	$\mu_3\lambda_1^2(\lambda_2 + \lambda_3)/\lambda_3$	23	13	$7N + 2$
47	$\mu_3\lambda_1^3/\lambda^3$	23	7	$7N + 2$
(c) <u>No Changes</u>				
48	$-\lambda$	1	1	1
49	$-\mu_3$	$7N + 2$	$7N + 2$	$7N + 2$
50	$-\lambda - \mu_3$	2	2	$7(N-1) + 2$
51	$-\mu_2$	$7N + 3$	$7N + 3$	$7N + 3$
52	$-\lambda - \mu_2$	3	3	$7(N-1) + 3$
53	$-\mu_1/2 - \mu_2/2$	$7N + 4$	$7N + 4$	$7N + 4$
54	$-\lambda - \mu_1/2 - \mu_2/2$	4	4	$7(N-1) + 4$
55	$-\mu_1$	$7N + 5$	$7N + 5$	$7N + 5$
56	$-\lambda - \mu_1$	5	5	$7(N-1) + 5$
57	$-\mu_1$	$7N + 6$	$7N + 6$	$7N + 6$
58	$-\lambda - \mu_1$	6	6	$7(N-1) + 6$
59	$-\mu_1$	$7N + 7$	$7N + 7$	$7N + 7$
60	$-\lambda - \mu_1$	7	7	$7(N-1) + 7$

second in queue is a one-unit job. In the first case, the transition intensity will be

$$\mu_2 \lambda_2 / \lambda = \mu_2 \frac{\lambda_2}{\lambda_2 + \lambda_3} \frac{\lambda_2 + \lambda_3}{\lambda} ,$$

i.e., μ_2 multiplied by the probability that the first job in queue is a two-unit job and the second is either a two-unit or a three-unit job. The transition will be made to state 10. The complete quadruple is number (42). In the second case, the transition intensity will be

$$\mu_2 \frac{\lambda_2}{\lambda_2 + \lambda_3} \frac{\lambda_1}{\lambda} ,$$

i.e., μ_2 multiplied by the probability that the first job in queue is a two-unit job and the second is a one-unit job. Thus, after the service completion two jobs move into service resulting in state 4. Number (43) in Table Vb is the complete quadruple.

For this problem, it will be advisable to determine the class (3) quadruples by examination of those classes (1) and (2). The sum of the transition intensities referring to state 3 is $\mu_2 + \lambda_1 + \lambda_2 + \lambda_3$ resulting in the transition intensity $-(\lambda + \mu_2)$. To determine the range of the quadruple with this α , we will consider α 's for states 10 and 17. For state 10, the following is the sum of the transition intensities.

$$\lambda + \mu_2 \lambda_2 / (\lambda_2 + \lambda_3) + \mu_2 \lambda_3 / (\lambda_2 + \lambda_3) .$$

This also reduces to $\lambda + \mu_2$. For state 17, the following is the sum of the transition intensities:

$$\lambda + \mu_2 \lambda_3 / (\lambda_2 + \lambda_3) + \mu_2 \lambda_2 / \lambda + \mu_2 \lambda_2 \lambda_1 / (\lambda (\lambda_2 + \lambda_3)) .$$

This again reduces to $\lambda + \mu_2$. Thus, quadruple (52) will refer to all states except the final state. For the final state, the α will be simply $-\mu_2$ since quadruple (14) with $\alpha = \lambda$ does not refer to the final state.

A few additional observations may be useful. In quadruples (19) and (20), the transition intensities μ_1 and μ_2 are each divided by 2. This is a result of an assumption made earlier in which it was stated that, if two jobs were run simultaneously, the mean rate at which each would be completed would be 1/2 the basic mean rate. When two like jobs (i.e., one-unit jobs) are being run, the mean rate for each is $(1/2)\mu_1$ but the mean rate for a service completion will be the sum, namely μ_1 . Quadruple (39) has a factor 2 associated with the transition intensity. This arises from the fact that two like transition intensities have been added in order to arrive at this result.

3.4 SUMMARY

It should be noted that all three transition tables have a common feature. Each represents a system with the size of the state space limited by the maximum number of allowed jobs in the queue. As we stated previously, this is only one kind of system for which a transition table may be usefully employed. Additional complexities can be introduced by having elements of the quadruples expressed as functions of some or all of the parameters of the queueing model.

As illustrated, the method for deriving a transition table provides an orderly procedure for determining a complete set of quadruples. What is actually done is to enumerate all events which can occur for each state. The division of events between arrivals and service completions in the examples is artificial and done only for convenience in ordering tasks. For problems in which some events cannot be identified as either an arrival or a service completion, some other division will be suitable. In any case, one will continue to enumerate the transitions for each event and each state. Of course, sequences other than those suggested in the examples could be followed in systematic evaluation of the quadruples.

In some instances, the construction of a partial state diagram will greatly aid the process of determining the transition table.⁸ For complex systems, the state diagram may prove very difficult to draw and may not provide much insight into the regularities of the state space.

Once again, we stress the point that examples have been presented to suggest possible forms for transition tables. For most problems, it will be the user's responsibility to devise transition tables with those properties most useful to his special purposes.

4. PROGRAM DESCRIPTION

The purpose of this section is to discuss programming details of the RQA-1 program. For the most part, the block diagram (Fig. 3) and the MAD listings (Appendix A) of the program provide full documentation of the program. However, a few remarks will be made here to aid in their interpretation. We shall begin these remarks by discussing the roles of the major subroutines.

More detail of the procedures to be followed by a user, of definitions of user specified parameters, and of diagnostic printouts is provided in Section 5 which acts as a sort of "user's manual" for RQA-1.

4.1 THE DISCIPLINE SUBROUTINE (DISCPL)

The function of the Discipline Subroutine is to prepare the transition table which describes the transition probabilities or intensities of the problem to be solved. This subroutine is unique to each application or model solved by RQA-1, and hence, must be written by the user.

The form of the transition table internal to RQA-1 is a set of four vectors called TRANS, FIRST, DEST, and LAST corresponding to the α , β , γ , and B respectively, of the quadruples discussed in Section 3. The i -th element of each of the four vectors together form the i -th quadruple. The DISCPL subroutine is basically a program for the evaluation of these vectors in terms of parameters which may be altered at execution time. These parameters are the counterparts of the λ , μ , N, etc., of the examples in Section 3 and appear only in DISCPL.

In addition to preparing the transition table, DISCPL has the responsibility of reading all parameters supplied by the user at execution and of providing essential control information descriptive of the queueing model being analyzed.

Sample DISCPL routines for the three models used as examples in Section 3 can be found in Section 5 (Table VI, VII, and VIII).

4.2 STOCHASTIC CONVERSION SUBROUTINE (STOCAL)

The role of this subroutine is to determine whether the transition table in memory represents a transition intensity matrix Q (hence, a continuous-time model) or a stochastic matrix A (hence, a discrete-time model). If it

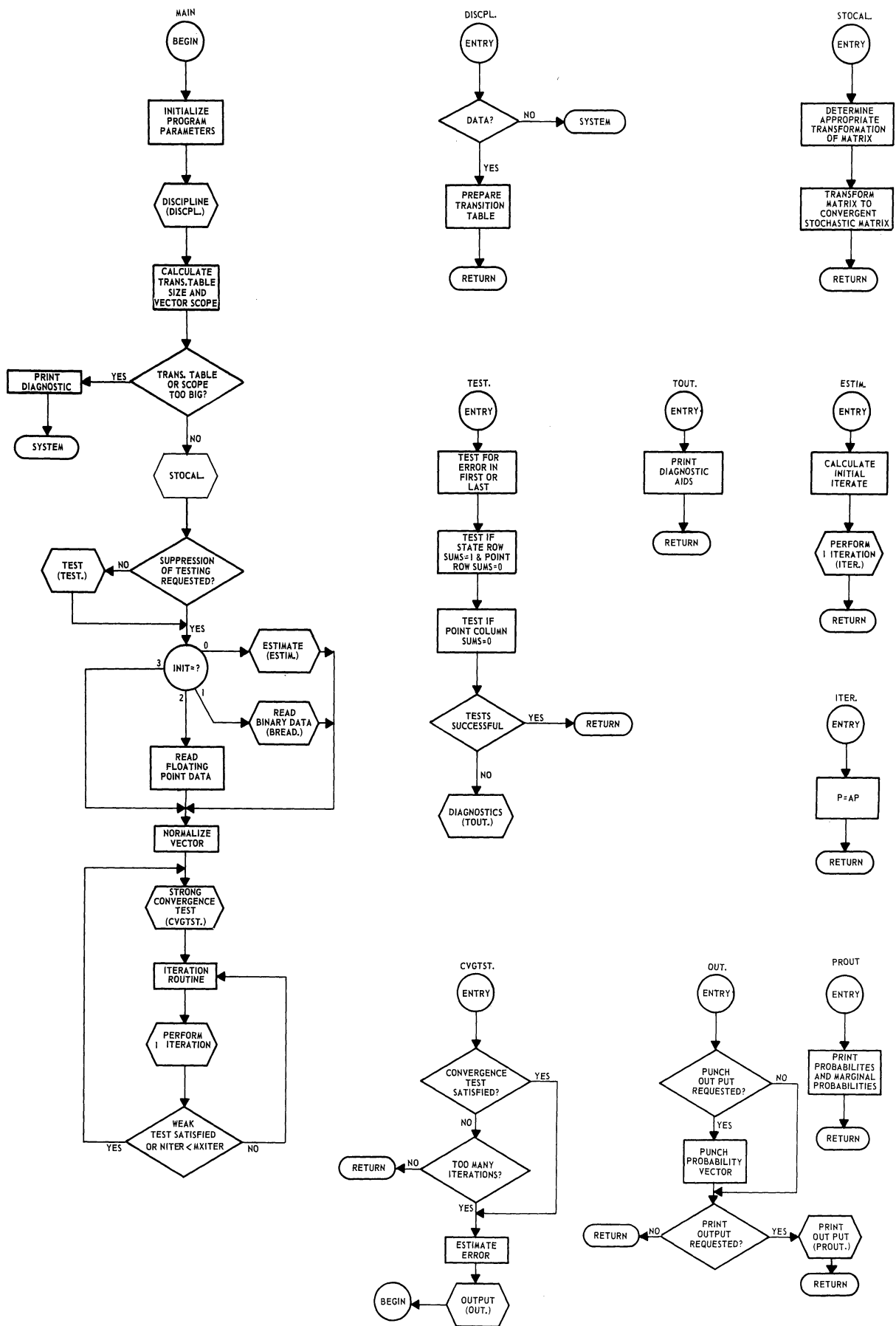


Fig. 3. Flow diagram for RQA-1.

is the former, STOCAL converts the transition table to represent the stochastic matrix $(.99/R)Q + I$, where $R = \max |q_{i,i}|$. If it is the latter, STOCAL converts the transition table to represent the new stochastic matrix $(.99A + .01I)$. These conversions assure the convergence of the resulting iteration processes to an equilibrium distribution of the specified model (cf., Sections 2.2 and 2.3).

The determination of the type of matrix represented is accomplished by determination of the element of the TRANS vector having least value. If it is equal to or greater than zero, the transition table represents a stochastic matrix. If, on the other hand, it is negative, the transition table represents a transition intensity matrix, and the value of the element having least value is equal to $-R$.

4.3 THE TEST SUBROUTINES TEST AND TOUT

The purpose of TEST is to determine whether the current transition table does indeed represent a stochastic matrix, and whether it adheres to the necessary properties of a valid transition table.

Using the transition table, the elements corresponding to each row of the matrix are summed and successively stored in a prezeroed vector. If the matrix is stochastic, the rows corresponding to states will have a one in the associated position of the vector and these corresponding to points will have a zero. In addition, all elements in columns corresponding to points must be zero. In the TEST subroutine, it is determined whether or not the above conditions exist. If they do not exist, the subroutine TOUT is called, and various diagnostics are printed to help determine the error(s) in the transition table.

For each quadruple in a valid transition table, the following equation must hold:

$$B - \beta = r\delta \quad \text{for } r \in \{0,1,2,\dots\} . \quad (29)$$

If this condition is not met for all the quadruples of the transition table, the subroutine TOUT is called, and appropriate diagnostics are printed.

4.4 THE ESTIMATION SUBROUTINE ESTIM

Before beginning the iterations, an initial iterate must be selected. The RQA-1 main program allows this iterate to be selected in four different ways by the user: it can be read in as input data either in binary or floating point form, it can be taken to be the final iterate of the previous run (which has remained in storage), or it can be calculated by the subroutine called ESTIM.

As its name suggests, ESTIM is a subroutine for estimating a satisfactory initial iterate and is based on a generalization of the simple server system discussed in detail in Section 3.1. For a simple server model, the probability of queue length i is proportional to ρ^i , where $\rho = \lambda/\mu$ (e.g., ref. 10). ESTIM prepares an initial iterate so that the k -th element of the iterate is proportional to $(\rho^*)^k$, where ρ^* is a constant. If $G = \{g_{ij}\}$ is the matrix represented by the transition table, then ρ^* is calculated as

$$\rho^* = \left(\frac{\lambda^*}{\mu^*} \right)^{\frac{1}{v}} \quad (30)$$

where

$$\lambda^* = \sum_{i,j:i>j} g_{ij} \quad (31)$$

$$\mu^* = \sum_{i,j:i>j} g_{ij} \quad (32)$$

$$v = \frac{\sum_{i,j: g_{ij} \neq 0} |i - j|}{\eta} \quad (33)$$

and η is the total number of nonzero, off-diagonal elements of the matrix. For the simple queue, these expressions reduce to $\rho^* = \lambda/\mu$. The only real justification for applying this calculation to more complex models is that it has been found to be quite effective, particularly when the model being solved originally had a multivariate state space and the mapping to the linear state space had a major queue of the system as the most significant variable in a Cartesian mapping.

ESTIM calls the iteration subroutine once before returning to the calling program, producing a vector having nonzero values only in states. The vector supplied to the calling program will not have a unit sum. Hence, subsequent normalization by the calling program is required.

4.5 THE ITERATION SUBROUTINE, ITER

Whereas the main program and all the other subroutines are written in the procedure-oriented language, MAD, the ITER subroutine is written in the machine-oriented language, UMAP.* The purpose of this is to make this subroutine as efficient as possible.

The only argument of this subroutine is the number of iterations to be carried out. Thus, a typical call of the subroutine would be EXECUTE ITER.(N) where N is the number of iterations. One iteration consists of the multiplication of the current iterate by the matrix. The subroutine requires two vectors in erasable store: one for the last iterate, and one for the iterate being formed. A switch is set which indicates which is the current iterate, thus removing the necessity for continually transposing newly computed iterates.

Because of the importance of this subroutine, we would like to discuss its operation in more detail. The two vectors mentioned above are actually stored in one vector which is represented in the program by the symbol V. This vector has 10,000 locations, thereby allotting 5,000 to each vector. To indicate which half of V contains the current iterate and which half will be used to store the newly computed iterate, two symbols, VFLAG and CVFLAG, are defined in the program. These symbols can only assume the values 0 or 1 and are constrained to have complementary values, i.e., when VFLAG = 1, CVFLAG = 0.

The current iterate is always referred to by calling V(VFLAG,I). When VFLAG = 0, the first half of V contains the current iterate. When VFLAG = 1, the second half of V contains the current iterate. To illustrate how a typical iteration would proceed, assume that VFLAG = 1. This means that the current iterate occupies the second half of V. During the process of multiplying this vector by the matrix, values of the newly calculated iterate are stored in the first half of V which had been zeroed previously. When this single iteration is complete, VFLAG is set equal to 0, and the second half of V is zeroed. Thus, the current iterate occupies the first half of V and will be correctly referenced by calling V(VFLAG,I).

As mentioned above, the multiplication process is carried out by sequencing through the transition table. For a particular quadruple $(\alpha_i, \beta_i, \gamma_i, B_i)$, α_i is multiplied by the value of the current iterate at location β_i , and the result is added to the value at location γ_i of the iterate being formed. This process is continued for the range of the quadruple, i.e., until α_i is multiplied by the value of the current iterate at location B_i with the result being added to the value at location $\gamma_i + B_i - \beta_i$ of the iterate being

* For users who do not have a UMAP assembler available, a MAD version of ITER is also included in the listings of Appendix A. Its use in place of the UMAP version will approximately quadruple the execution time of RQA-1.

formed. The program then proceeds to $(\alpha_{i+1}, \beta_{i+1}, \gamma_{i+1}, B_{i+1})$ and repeats the process until all quadruples are exhausted.

This type of vector-matrix multiplication is quite efficient because only nonzero elements of the matrix are referenced. Also, no significant penalty in computing time results from the manner of describing the matrix.

4.6 THE ITERATION STRATEGY (MAIN AND CVGTST)

Upon completion of the calculation and normalization of the initial iterate, the main program produces a first call to the convergence testing subroutine CVGTST. This first call of CVGTST simply sets the underflow traps so that underflows produce normal zeros, and returns to the calling program (MAIN) with an index corresponding to a nonzero element of the initial iterate.

The main program continues to iterate until successive values of the specified element have an absolute difference which is less than the convergence criterion ϵ , or until the specified maximum number of iterations has been taken. At that point CVGTST is called again.

If we define $p^{(n)}$ to be the n -th iterate and $p_i^{(n)}$ to be the value of the i -th element of the n -th iterate, the maximum difference between successive iterates $e^{(n)}$ is given by

$$e^{(n)} = \max_i | p_i^{(n)} - p_i^{(n-1)} | .$$

The iteration process is said to have converged after n iterations for error criterion ϵ if $e^{(n)} \leq \epsilon$. For every call of CVGTST after the first, CVGTST calculates $e^{(n)}$ and tests for convergence. If the process has not converged, this subroutine returns to the main program with the integer i_M corresponding to the state for which the absolute value between successive iterates is largest. The idea is not to calculate i_M for each iteration but to use an initially calculated i_M as a coarse test and to recalculate only when this coarse test is satisfied. Then the main program iterates until the absolute difference of the i_M -th element is less than ϵ or the specified number of iterations has been taken. Another call of CVGTST then results, and the process is repeated.

On the other hand, whenever the tests within CVGTST are satisfied, CVGTST prints appropriate comments and error estimates, resets the underflow error trap, calls the output subroutine OUT2, and transfers out (via an "error return") to the beginning of the main program for a calculation of another model (if data are available).

The purpose of this two-level testing procedure is to avoid the expense of calculating i_M for each iteration by using a value of i_M in a coarse test for convergence and recalculating i_M (using CVGTST) only when the coarse test has been satisfied.

4.7 MISCELLANEOUS SYSTEM ROUTINES

For the sake of completeness, the library routines explicitly called by RQA-1 will be listed here with a brief mention of their function.

(a) BNBCD.(X) converts its argument X from binary to binary coded decimal (Hollerith) form suitable for punching or printing characters.

(b) BPUNCH.(\mathcal{L}) punches the values of the list \mathcal{L} in column binary format.*

(c) BREAD.(\mathcal{L}) reads the values of the list \mathcal{L} from column binary cards.

(d) ELOG.(X) computes the natural logarithm of X.

(e) ERROR. produces a call to the monitor system with a dump of memory if requested.

(f) FTRAP. sets the underflow trapping procedure so that underflows are replaced by normal zeros until NTRAP is executed.

(g) NSDBCD.(X,Y) provides for punching the BCD values X,Y in columns 72-80 of the next card punched. (Sequencing of the last characters of this field is accomplished by The University of Michigan Executive System.)

(h) NTRAP. resets the underflow trapping procedure to its normal condition so that an underflow produces an exit to the monitor system (as in ERROR).

(i) SYSTEM. produces a call to the monitor system.

(j) TODAY.(DATE, YEAR) gives to DATE, and YEAR BCD values corresponding to the day, month, and year that the program is executed.

(k) ZERO.(\mathcal{L}) sprays normal zeros into all locations specified by the list \mathcal{L} .

* Actually, in The University of Michigan system, it writes binary card images on the "punch output" tape, for offline punching.

4.8 CONCLUSION

We would like to make some remarks on the program as a whole. Except for the two subroutines DISCPL and PROUT,* the program operates on a state space which is one-dimensional. Thus part of the program needs no information with respect to the structure of the states, i.e., their multidimensional nature. The transition table created by DISCPL describes a one-dimensional process, but DISCPL also contains some information about the structure of the states which is passed on to PROUT. For example, if the state designation were two-dimensional, DISCPL would contain values for the upper bounds of each dimension. In the examples of Section 3, the upper bound for the internal state variable is C which is fixed, and the upper bound for the queue state variable is N which is put in as data. The upper bounds are needed only in the subroutine PROUT to prepare a suitable output format. PROUT was written specifically for a two-dimensional state description and presents the final iterate in terms of such a state description.

Thus, the two subroutines DISCPL and PROUT can be considered user-controlled. DISCPL must be supplied by the user, while PROUT may be replaced by the user if he desires output formats other than those built into RQA-1 or if he wishes to perform some additional calculation before printing results.

*PROUT is the program, called by OUT2, which prepares and prints the printed state probabilities and marginal probabilities.

5. PROCEDURES FOR THE USE OF THE PROGRAM

This section describes procedures for the use of RQA-1. It is concerned primarily with the data which must be supplied by the user and with the interpretation of the output produced by the program. It is assumed here that the program is used with The University of Michigan Executive System.

In normal use, the object deck of RQA-1 is followed by the DISCPL subroutine, a "\$DATA" card, and the input data. Since the DISCPL subroutine must be prepared by the user, we will give a detailed discussion of its preparation. The data are in five classes, divided into two distinct groups of cards. In order of their appearance following the \$DATA card, these are

GROUP 1

- (1) Program control parameters
- (2) Structural parameters
- (3) Discipline parameters
- (4) Transition table parameter

GROUP 2

- (1) Initial Iterate.

Following the description of the programming of a DISCPL subroutine will be a discussion of the input data, the form of both normal and abnormal output, the structure of the program card deck, and the format of binary cards occurring as punched output.

5.1 THE DISCIPLINE SUBROUTINE, DISCPL

The chief input to RQA-1 provided by the user is the discipline subroutine. This user-generated program provides RQA-1 with a complete description of the model to be solved in the form of code which generates the transition table. In addition to correctly describing his model, the user must insure, while writing this subroutine, that the number of quadruples generated does not exceed 2,000 and that the linear index of every state does not exceed 5,000.*

The discipline subroutines must carry out certain necessary functions for RQA-1:

* Failure to do the former will usually result in a violation of low-core storage protection in the University of Michigan system and probably worse difficulties in other systems. Failure to do the latter will simply result in an error return to system with suitable diagnostics printed out.

- (1) Specifying the complete transition table
- (2) Reading or specifying the Group 1 parameters
- (3) Specifying a name for the model.

The transition table consists of four vectors called TRANS, FIRST, DEST, and LAST corresponding respectively to the α , β , γ , and B of Section 3. The values of elements of these vectors will usually be dependent upon the values of discipline parameters which are supplied at execution time and, hence, cannot be treated as simple numerical vectors to be read as data. Rather, they are computed by the execution of program segments contained in the discipline subroutine. This facilitates examination of the behavior of a model as parameters are varied, since the program may be simply rerun with different numerical values of the parameters supplied.

Normally, the discipline subroutine must contain all the statements within the dashed blocks in Fig. 4. The first block contains the instruction to read the Group 1 data, and in this section, it will be assumed that this is done by the simplified input/output statements in MAD. The second block defines the model name ("name" should be replaced by any desired twelve characters, including spaces) and sets up necessary communication to the remainder of RQA-1 through PROGRAM COMMON.

Group 1 parameters which are not read as data may be declared or calculated in DISCPL. The Group 1 parameters will be defined later.

The vectors TRANS, FIRST, DEST, and LAST of the transition table can be set up in several ways, only one of which is exhibited in the listings of the DISCPL subroutines shown in Tables VI, VII, and VIII for the examples of Section 3. This method involves the use of two list-manipulating statements in MAD, namely: SET LIST TO and SAVE DATA.¹ SET LIST TO TRANS sets up a vector named TRANS, and SAVE DATA $\alpha_1, \alpha_2, \dots, \alpha_n$ causes $\alpha_1, \alpha_2, \dots, \alpha_n$ to be the n entries of this vector where the α 's can be any arithmetic expression acceptable to MAD.

While the illustrated technique is quite useful, other methods will also be found suitable to particular problems. For example, the quadruples may be generated in a series of THROUGH loops when some systematic evaluation technique is available or the number of quadruples to be generated depends on data. Appendix B contains a discussion of a group of subroutines, contributed by J. H. Jackson, which have been found to be useful in the generation of quadruples under quite general circumstances. On the whole, the degree of innovation which is possible is limited chiefly by the user's desires and needs.

In determining the quadruples of the transition table, the period DELTA (δ) can be treated as a constant or a parameter. If it is a constant, as is the case in the examples of Section 3, it is set in DISCPL, but, if it is to be treated as a parameter, it can be defined in DISCPL or read in as data.

TABLE VI

"SIMPLE SERVER" DISCIPLINE SUBROUTINE

\$COMPILE MAD

DISC0000

000787 09/09/65

2 09 11.6 PM

MAD (09 AUG 1965 VERSION) PROGRAM LISTING

```

EXTERNAL FUNCTION                                *001
ENTRY TO DISCIPL.                                *002
INPUT PARAMETERS ARE L, MU, EPSI, NQ
PRINT COMMENT $8                                *003
1 S $                                           *003
READ AND PRINT DATA                            *004
MAXJ = 2                                        *005
DELTA = 2                                       *006
SET LIST TO TRANS                               *007
SAVE DATA L,L,MU,MU,-L,-L-MU,-MU             *008
SET LIST TO FIRST                              *009
SAVE DATA 1,2,2,4,1,2,2*NQ+2                *010
SET LIST TO DEST                               *011
SAVE DATA 2,4,1,2,1,2,2*NQ+2                *012
SET LIST TO LAST                              *013
SAVE DATA 1,2*NQ ,2,2*NQ+2,1,2*NQ,2*NQ+2    *014
VECTOR VALUES NAME = $SIMPLE MODEL$         *015
INTEGER FIRST,DELT,LAST,INIT,PCH,PRSW,DELTA,MAXJ,MAXITER,NQ *016
PROGRAM COMMON TRANS(2001),FIRST(2000),DEST(2000),LAST(2000), *017
1 PCH,PRSW,MAXITER,EPSI,DELTA,INIT,RUN,TSM,NAME(1),MAXJ *017
FUNCTION RETURN                                 *018
END OF FUNCTION                                *019
DISC0010
DISC0020
DISC0030
DISC0040
DISC0050
DISC0062
DISC0063
DISC0080
DISC0090
DISC0100
DISC0110
DISC0120
DISC0130
DISC0140
DISC0150
DISC0160
DISC0200
DISC0220
DISC0230
DISC0270
DISC0280

```


TABLE VIII

"RANMEM" DISCIPLINE SUBROUTINE

\$COMPILE MAD DISC0000 001195 09/13/65 12 58 6.7 PM

MAD (09 AUG 1965 VERSION) PROGRAM LISTING

```

EXTERNAL FUNCTION
ENTRY TO DISCPL.
INPUT PARAMETERS ARE L1, L2, L3, L, NU1, NU2, NU3, EPSI, NQ
PRINT COMMENT $8
1 S $
READ AND PRINT DATA
MAXJ = 7
DELTA = 7

SET LIST TO TRANS
SAVE DATA L1,L2,L3,L,L1,L2+L3,L,L1,L2,L3,L,L1,L2+L3,L,L,L,L,
NU3,NU2,NU1/2.,NU2/2.,NU1,NU1,
2 NU1,NU3*L3/L,NU3*L2/L,NU3*L1/L,NU2*L2/(L2+L3),NU2*L3/(L
2+L3),NU1*(L3+L2)/(2.*L),NU2*L3/(L2+L3),NU1*L1/(2.*L)+NU2*L2/(
4 2.*L),NU2*L1/(2.*L),NU1*L3/(L2+L3),NU1*L2/(L2+L3),NU1,NU1*(L2
5 +L3)/L,NU1*L1/L,NU3*L2*(L2+L3)/(L*L),2.*NU3*L2*L1/(L*L),NU3*L
6 1*L3/(L*L),NU3*L1*L1/(L*L),NU2*L2/L,NU2*L2*L1/(L2+L3)*L),NU2
7 *L1*(L2+L3)/(2.*L*L),NU2*L1*L1/(2.*L*L),NU3*L1*L1*(L3+L2)/(L*
8 L*L),NU3*L1*L1*(L*L*L)
SAVE DATA
-L,
1 -NU3,-L-NU3,-NU2,-L-NU2,-NU1/2.,-L-NU1/2.,-NU2/2.,-NU1,
2 -L-NU1,-NU1,-L-NU1,-NU1,-L-NU1

```

```

SET LIST TO FIRST
SAVE DATA 1,1,1,2,3,3,4,5,5,6,6,7,10,12,13,2,3,4,4,5,6,7,9,
1 9,9,13,16,11,11,11,13,13,12,14,14,16,16,16,16,17,17,18,18,
2 23,23,1,7*NQ+2,2,7*NQ+3,3,7*NQ+4,4,7*NQ+5,5,7*NQ+6,6,7*NQ+7,7

SET LIST TO DEST
SAVE DATA 5,3,2,9,4,10,11,6,4,12,7,13,14,17,19,20,1,1,3,5,1,5
1 ,6,2,3,5,3,2,16,12,4,6,12,4,2,13,7,10,4,12,6,10,4,13,7,13,7,1
2 ,7*NQ+2,2,7*NQ+3,3,7*NQ+4,4,7*NQ+5,5,7*NQ+6,6,7*NQ+7,7

```

```

DISC0000 *001
DISC0010 *002
DISC0020 *003
DISC0040 *003
DISC0040 *004
DISC0062 *005
DISC0063 *006
DISC0070 *007
DISC0080 *008
DISC0090 *008
DISC0100 *008
DISC0110 *008
DISC0120 *008
DISC0130 *008
DISC0140 *008
DISC0150 *008
DISC0160 *008
DISC0170 *009
DISC0180 *009
DISC0190 *009
DISC0200 *010
DISC0210 *011
DISC0220 *011
DISC0230 *011
DISC0240 *012
DISC0250 *013
DISC0260 *013
DISC0280 *013

```

TABLE VIII (Concluded)

SET LIST TO LAST		
SAVE DATA 1,1,1,7*(NQ-1)+2,3,3,7*(NQ-1)+4,5,5,5,5,6,6,7*(NQ-1)+	DISC0290	*014
1 7,7*(NQ-1)+3,7*(NQ-1)+5,7*(NQ-1)+6,2,3,4,4,5,6,7*(NQ)+2,9	DISC0300	*015
2 ,9,10,7*(NQ)+3,7*(NQ)+4,7*(NQ)+4,7*(NQ)+4,11,7*(NQ)	DISC0310	*015
3 +6,7*(NQ)+6,7*(NQ)+5,7*(NQ)+7,7*(NQ)+7,7*(NQ)+2,	DISC0330	*015
4 7*(NQ)+2,7*(NQ)+2, 16,7*(NQ)+3,7*(NQ)+3,	DISC0340	*015
5 7*(NQ)+4,7*(NQ)+4,7*(NQ)+2,7*(NQ)+2	DISC0345	*015
SAVE DATA 1,7*NQ+2,7*(NQ-1)+2,7*NQ+3,7*(NQ-1)+3,7*NQ+4,7*(NQ-	DISC0350	*015
1 1)+4,7*NQ+5,7*(NQ-1)+5,7*NQ+6,7*(NQ-1)+6,7*NQ+7,7*(NQ-1)+7	DISC0351	*016
	DISC0352	*016
VECTOR VALUES NAME = 3RANMEM MODELS	DISC0360	*017
INTEGER FIRST,DEST,LAST,INIT,PCH,PRSW,DELTA,MAXJ,MAXITER,NQ	DISC0400	*018
PROGRAM COMMON TRANS(2001),FIRST(2000),DEST(2000),LAST(2000),	DISC0420	*019
1 PCH,PRSW,MAXITER,EPSI,DELTA,INIT,RUN,ISW,NAME(1),MAXJ	DISC0430	*019
FUNCTION RETURN	DISC0470	*020
END CF FUNCTION	DISC0480	*021

```

EXTERNAL FUNCTION
ENTRY TO DISCPL
PRINT COMMENT $8          INPUT PARAMETERS $
READ AND PRINT DATA

```

.....

Portion of program defining Group 1 parameters which are not read as data

.....

Portion of program defining transition table in terms of
DELTA, TRANS, FIRST, DEST, LAST

.....

```

VECTOR VALUES NAME = $ "Name"$
INTEGER FIRST, DEST, LAST, INIT, PCH, PRSW, DELTA, MAXJ, MXITER
PROGRAM COMMON TRANS(2001), FIRST(2000), DEST(2000), LAST(2000), PCH, PRSW,
    MXITER, ERSI, DELTA, INIT, RUN, TSW, NAME(1), MAXJ
FUNCTION RETURN
END OF FUNCTION

```

Fig. 4 DISCPL subroutine.

We will complete this discussion with a brief reference to the three DISCPL subroutines listed in Tables VI, VII, and VIII. In all these cases, the SET LIST TO and SAVE DATA statements have been used. All these examples have been formulated in terms of a single queue and a complex server. The parameter NQ is used in the transition tables to represent the maximum queue length. Of course, any integer-valued variables introduced by the user must be declared. For these examples, DELTA was chosen equal to C (the cardinality of the set J) and so the following pair of statements is used:

```

DELTA = x
MAXJ  = x

```

where x is 2,7,11 for the SIMPLE QUEUE, RANMEM MODEL and FIXMEM MODEL respectively. (MAXJ is a structural parameter (Group 1) which, in these cases, is also equal to the cardinality of J.)

Group 1 parameters are defined in the next three subsections.

5.2 PROGRAM CONTROL PARAMETERS

The program control parameters are named PCH, PRSW, MXITER, EPSI, INIT, and TSW. The mode, preset value, and range of permitted values of each of these variables are given in Table IX. If these parameters are not specified on any particular run, they have, for that run, the value last supplied on a previous run. If no previous run has specified their value, they have the "preset" value shown.

TABLE IX

PROGRAM CONTROL DATA			
Name	Mode	Preset Value	Other Values
PCH	INTEGER	1	0,1
PRSW	INTEGER	1	0,1
MXITER	INTEGER	100	$\geq 0, < 10^5$
EPSI	FLOATING POINT	10^{-4}	$\geq 0, \leq 1.$
INIT	INTEGER	0	0,1,2,3
TSW	INTEGER	1	0,1

EPSI is the value of the convergence test criterion. Thus, when the maximum absolute difference between corresponding elements of successive iterates is less than or equal to EPSI, the convergence test is satisfied. MXITER is the maximum number of iterations to be allowed in the calculation. If this number of iterations has not been satisfied, the calculation is terminated and output supplied.

Table X lists the significance of each of the values of the variables PCH, PRSW, INIT, and TSW.

TABLE X
THE FUNCTIONS OF PCH, PRSW, INIT, and TSW

Name	Value	Function
PCH	1	- Final iterate is punched in binary format.
	0	- Final iterate is not punched.
PRSW	1	- Final iterate is printed by execution of PROUT.
	0	- Final iterate is not printed. PROUT is not called.
INIT	0	- The initial iterate is calculated by the subroutine ESTIM.
	1	- The library subroutine BREAD is called and an initial iterate in binary format is read in. This input is the punched output of a previous run.
	2	- The initial iterate is read in by READ DATA statement and is in floating point format.
	3	- The final iterate from the previous run is used as the initial iterate for the present run.
TSW	1	- The transition table is tested to assure its stochastic character and other properties. (TEST is called).
	0	- The transition table is not tested.

5.3 STRUCTURAL PARAMETER

Structural parameters convey dimensionality information to be output printing subroutines. For the output printing subroutine provided (PROUT), there is only one such parameter, MAXJ. This parameter indicates the maximum extent of the least significant coordinate J of the two-dimensional output array of proba-

bilities. In PROUT, a state with index K will be treated as a state with vector value (I,J) as determined by the mapping

$$K = \text{MAXJ} * I + J + 1$$

corresponding to Eq. (28). Thus, MAXJ corresponds to the cardinality of the set {J} described in Section 3.

The preset value and the range of MAXJ are given in Table XI. If the preset value is used (by not specifying MAXJ) or if a value greater than 2500 is specified, the printed output will supply the results ten-to-a-row and fifty-to-a-block.

TABLE XI

STRUCTURAL PARAMETER

Name	Mode	Preset Value	Restriction
MAXJ	INTEGER	10	MAXJ \leq 2500

5.4 DISCIPLINE PARAMETERS

These data consist of the required parameters for the particular queue discipline being run. They are defined by DISCPL only. If they are not floating point constants, their mode must be declared. For the examples discussed in Section 3, such parameters are the L's and NU's. For the FIXMEM model (Table VII), a data card with discipline parameters would be

$$L = 1.0, \text{ NU1} = 1.5, \text{ NU2} = .25, \dots$$

In general, these parameters are the independent variables involved in the quadruples of the transition table.

5.5 TRANSITION TABLE PARAMETER

This parameter is DELTA, which is an integer. Its use was discussed in Section 5.1. It is treated essentially like a structural parameter, but, while structural parameters refer to the state-space structure, DELTA is actually a part of the description of the transition table. It has a preset value of 1.

5.6 INITIAL ITERATE

As seen in Table X, when INIT is 1 or 2, the initial iterate is read in as data. For INIT = 1, the initial iterate must be presented in column binary format. It is usually the set of cards containing a final iterate which was punched in binary mode in a previous run as the result of setting PCH = 1.

The form of this iterate is described in Section 5.11. The entire set of output cards as described there must be used as input for the initial iterate, i.e., the header card followed by the binary data.

For INIT = 2, the initial iterate must be supplied in floating point form to be read by a READ DATA type statement. Unspecified values are automatically zero. Before being used, the initial iterate will be normalized to a probability vector. Therefore, at least one element must be specified. Typical data cards would be

$$V(1) = .25, \quad V(3) = .50, \quad V(10) = .25^*$$

or

$$V(1) = .20, .15, .10, .05, \quad V(50) = .20, .15, .10, .05^* \quad .$$

Values can be given for any $V(I)$ where I runs from 1 to the largest integer denoting a state.

These data are Group 2 data and must appear after the Group 1 data of the run.

5.7 FORMAT OF INPUT DATA

There is no specified order for the Group 1 data. These data are in the format of cards to be read by a READ DATA statement in MAD and must be followed (according to that format) by "*" .

Following the Group 1 data will be the Group 2 data, if required (i.e., if INIT = 1 or INIT = 2). Following this will be the data for the next run, if further runs of the model with different data values are desired. Any number of sets of data may be stacked in this way.

The run number sequence which identifies all output of RQA-1 can be started (restarted) at a user-specified value on any run by insertion of a card just ahead of the Group 1 data for the run with the format

THE INITIAL RUN NUMBER IS XXX .

This statement must start in the first column of the card, and words must be separated by a single space. The XXX may be any non-negative integer of three or less digits. If this is not done, runs will be numbered starting with 1 for the first run after loading.

The following may be a typical set of data cards for one run:

```
EPSI = .0005, MXITER = 200, INIT = 2
L = 1,
MU1 = 1.5, MU2 = .25*
V(1) = .5, .3, .1, .1*
```

Alternatively, another set may be:

```
THE INITIAL RUN NUMBER IS 75
INIT = 1
L = 1, MXITER = 50, MU1 = 1.5, MU2 = .25*
.....
(Binary card deck as described in Section 5.11)
.....
```

Note that, in this case, EPSI will be the same value as it was on the previous run (or it will be .0001 if this is the first set of data) and that, in both cases, PCH, PRSW, and TSW are the same as they were in previous runs (or their preset values if this is the first set of data).

5.8 NORMAL OUTPUT

The normal output of RQA-1, illustrated in Table XII, should be self-explanatory.

However, it should be noted that the probabilities are printed in terms of a two-dimensional state designation where PROB(I,J) is the limiting state

TABLE XII

TYPICAL OUTPUT OF RQA-1

RQA RESULTS FOR FIXMEM MODEL DISCIPLINE SEPT 1, 1965 VERSION OF RQA-1

RUN 1, 16 SEP 1965

INPUT PARAMETERS

L=1, NU1=1.5, NU2=0.5,

NQ=15*

THE INITIAL ITERATE WAS DETERMINED BY THE SUBROUTINE ESTIM.....

***** SUCCESSFUL RUN *****

CONVERGENCE TEST CRITERION (EPSI) = .00010000

MAXIMUM ALLOWED ITERATIONS (MXITER) = 100

ITERATIONS TAKEN = 89

MAXIMUM DIFFERENCE BETWEEN LAST TWO ITERATES = 1.5E-04

ESTIMATED ORDER OF ERROR = .1E-03 (APPROXIMATE CONVERGENCE FACTOR = .9016573)

THE SOLUTION VECTOR HAS BEEN PUNCHED

TABLE XII (Continued)

THE FOLLOWING ARE THE LIMITING STATE PROBABILITIES, $PROB(I, J)$,
 $PROB(I)$ IS THE SUM OF $PROB(I, J)$ OVER ALL STATES WITH THE VALUE I FOR THE FIRST VARIABLE

I	PROB(I, J)									
	J = 0	1	2	3	4	5	6	7	8	9
0	.52348530	.08872072	.08792652	.05749297	.02716861	.00869784	.02980799	.05949457	.05947242	.03980388
1	.15267809	.00000000	.00000000	.00000000	.00000000	.00889547	.02000391	.04019641	.04069529	.02792383
2	.10608881	.00000000	.00000000	.00000000	.00000000	.00712090	.01348743	.02730851	.02796113	.01950914
3	.07300190	.00000000	.00000000	.00000000	.00000000	.00521510	.00912379	.01858519	.01916128	.01347094
4	.04970614	.00000000	.00000000	.00000000	.00000000	.00364714	.00617654	.01261922	.01303470	.00916726
5	.03345351	.00000000	.00000000	.00000000	.00000000	.00247474	.00417181	.00851817	.00877507	.00614247
6	.02223266	.00000000	.00000000	.00000000	.00000000	.00164044	.00280297	.00569988	.00583454	.00405082
7	.01457771	.00000000	.00000000	.00000000	.00000000	.00106571	.00186844	.00377258	.00382645	.00262885
8	.00942488	.00000000	.00000000	.00000000	.00000000	.00067974	.00123296	.00246581	.00247318	.00167895
9	.00600726	.00000000	.00000000	.00000000	.00000000	.00042634	.00080406	.00158983	.00157486	.00105586
10	.00377696	.00000000	.00000000	.00000000	.00000000	.00026350	.00051756	.00101066	.00098847	.00065495
11	.00234679	.00000000	.00000000	.00000000	.00000000	.00016089	.00032868	.00063387	.00061297	.00040237
12	.00144560	.00000000	.00000000	.00000000	.00000000	.00009656	.00020615	.00039372	.00037850	.00024682
13	.00088689	.00000000	.00000000	.00000000	.00000000	.00005556	.00012857	.00024633	.00023900	.00014833
14	.00054715	.00000000	.00000000	.00000000	.00000000	.00002964	.00008311	.00017104	.00014868	.00008031
15	.00033959	.00000000	.00000000	.00000000	.00000000	.00001674	.00008897	.00011723	.00007530	.00003186

TABLE XII (Concluded)

PROB(J) IS THE SUM OF PROB(I,J) OVER ALL STATES WITH THE VALUE J FOR THE SECOND VARIABLE

CHECKSUM	PROB(J)									
	J = 0	1	2	3	4	5	6	7	8	9
.99999920	.04460908	.08872072	.08792652	.05749297	.02716861	.04748632	.09083293	.18282300	.18525183	.12699661
	.06769266									

probability of state (I,J). The marginal probabilities PROB(I) and PROB(J), which are the sums of PROB(I,J) over J and I respectively, are also listed. Thus, for a fixed value of I, PROB(I) is printed and followed on the same line by the values of PROB(I,J) as J runs over its allowed values. CHECKSUM, the sum of all the elements of the probability vector, may not be exactly equal to 1 due to roundoff.

5.9 DIAGNOSTICS OF ABNORMAL OUTPUT

A run can terminate without supplying answers to the given problem for several reasons; there may be no input data, the number of allowable states or entries in the transition table may be exceeded, or there may be an error in the transition table. In each of these instances, appropriate statements and diagnostic aids are printed. This termination is caused by execution of the MAD instruction

EXECUTE ERROR.

which causes a return to the operating system preceded by a dump of storage if requested.

The following is a further explanation of a few of these diagnostics which may not be self-explanatory:

- (1) THE FOLLOWING QUADRUPLES HAVE AN ERROR IN FIRST OR LAST.

This is printed if $LAST(I) - FIRST(I)$ is not an integer multiple of DELTA. This is inconsistent with the rules for forming a transition table.

- (2) THE FOLLOWING QUADRUPLES HAVE AN ERROR IN DEST.

This is printed if there is a transition specified into a point, and there is no other transition with that point as a source.

- (3) THE FOLLOWING ROW SUMS, $SUM(I)$, OF THE STOCHASTIC MATRIX ARE INCORRECT.

The appearance of this statement indicates that the matrix represented by the transition table (at this point, it should be stochastic) is not stochastic. If a row sum corresponding to a point is nonzero, then a transition is being defined for something other than a state. If a row sum corresponding to a state is less than one, then a possible transition is being overlooked. If it is greater than one, an extraneous transition has been included.

5.10 FORMAT OF PROGRAM DECK

Two formats will be described; one for which the DISCPL subroutine is to be compiled, and the other for which it has been compiled and, similar to the rest of the program, is now on binary cards. For the first case, the following is the deck format:

```
2 yellow ID cards
"$ COMPILER MAD, EXECUTE, (etc.)" card
DISCPL MAD deck
"$ BINARY" card
RQA-1 binary deck
"$DATA" card
DATA
```

For the second case, the following is the deck format:

```
2 yellow ID cards
"$ BINARY, EXECUTE (etc.)" card
RQA-1 binary deck including DISCPL binary deck
"$DATA" card
Data
```

5.11 FORMAT OF THE PUNCHED FINAL ITERATE

As discussed in Section 5.2, when PCH = 1, the final iterate will be punched on binary cards. To illustrate the form of this output, we will assume that the name of the model is SIMPLE QUEUE, the date is MAY 25, 1965, and this is the 15th run.

The first card punched, called the header card, will have the following form:

```
SIMPLE QUEUE RUN 15 25 MAY 1965                SIMPLE Q.
```

Since, in the University of Michigan system, binary cards have only columns 73-80 interpreted (printed on the top of card), this card will be identified by the first 8 characters in the model name appearing in columns 73-80. The next card will contain the linear index corresponding to the maximum state, called SCOPE, punched in binary mode. Columns 73-80 of this card will contain

```
25 MAY 01
```

i.e., the day, the month, and the number 01. The succeeding cards will contain the values of the final iterate in binary mode. The number of such values will be SCOPE. Columns 73-75 of these cards will contain

(i.e., the run number). Columns 78-80 will contain the numbers 001, 002, etc. (i.e., a sequential numbering of the cards).

Each final iterate punched by RQA-1 will automatically have this output card format.

6. CONCLUSION

The purpose of the RQA-1 program is to evaluate equilibrium joint-probability distributions of queue lengths and system conditions for very large, finite Markov queueing systems. Although only these probabilities are obtained by the program, it is possible to estimate from them a wide variety of other performance statistics. The expected waiting and throughput times are easily determined as well as the probability distributions of any of the state variables (e.g., queue lengths, channel occupancies) by taking appropriate sums of state probabilities. Special programs can be written to determine the above statistics which will have as their input the probability distribution calculated by the RQA-1 program.

As part of this section, we will present some timing data for computer runs of the queueing models discussed previously.

6.1 TIMING EXPERIENCE

Results are shown in Table XIII(a) for the FIXMEM model with λ and μ_1 fixed and μ_2 varied. Table XIII(b) gives the results for one run of the model with 611 states. It is interesting to note that for each model the number of iterations per second varies inversely as the number of states. Table XIV presents some results for the RANMEM model for which the discipline parameters are fixed but EPSI and the number of states are varied. For the problem with 77 states and $\text{EPSI} = 10^{-4}$, the number of iterations per second is 48.3. For the problem with 357 states and $\text{EPSI} = 10^{-4}$, this number is 10. The ratio of states is $357/77 = 4.6$, and the ratio of iterations per second is $48.3/10 = 4.8$.

It should be noted that these results have been achieved for very sparse matrices. In fact, for the FIXMEM model, there is an average of four non-zero elements per row. Thus, for a 611 state model, the ratio of the number of nonzero elements to the total number of elements is $(4 \times 611)/(611 \times 611) = .0065$.

Using the equation developed in Section 2, we can compare the number of multiplications required for a solution of a typical FIXMEM model by simulation with the number of multiplications required by the numerical methods employed in RQA-1. For the FIXMEM model with $\mu_2 = .3333$, the EPSI satisfied was $.4058 \times 10^{-4}$ which corresponds to an ϵ , in the equation, of about 10^{-3} . A simulation for a model with these parameters would require at least 100 times as many multiplications as the numerical procedure or about one minute as compared to one-half second. For the FIXMEM model with $\mu_2 = .5$ and $\text{EPSI} = .6429 \times 10^{-4}$, a simulation would require about 150 times as many multiplications or about 4-1/2 minutes as compared with 2 seconds for an $\epsilon = 10^{-3}$.

TABLE XIII

"FIXMEM" MODEL

(The initial vector is estimated for each run)

 $\lambda = 1.0, \mu_1 = 1.5$

μ_2	EPFI ($\times 10^{-4}$)	Number of Iterations	Iteration of (Seconds)	Number of Iterations per Second	Total Time (Seconds)	Maximum Probability
(a) <u>Number of States = 101</u>						
1.0	0.4884	73	2.4	30.4	3.6	0.1251
0.5	0.6429	64	1.8	35.5	3.6	0.0801
0.3333	0.4058	23	0.6	38.4	2.4	0.0525
0.2857	0.9508	63	1.8	35.0	3.6	0.0411
TOTALS		223	6.6	33.8	13.2	
(b) <u>Number of States = 611</u>						
0.5	1.0	54	8.4	6.4	17.4	0.0783

TABLE XIV

"RANMEM" MODEL

(The initial vector is estimated for each run)

$\lambda_1 = 1.5, \lambda_2 = .3, \lambda_3 = .2;$
 $\mu_1 = 2, \mu_2 = 1.5, \mu_3 = 1.2$

Number of States	EPSI x 10 ⁻⁴	Number of Iterations	Iteration of (Seconds)	Number of Iterations per Second	Total Time (Seconds)	Maximum Probability
77	1.0	58	1.2	48.3	3.6	0.3828
77	0.1	102	2.4	42.5	4.8	0.3844
77	0.0478	115	3.0	38.3	5.4	0.3845
TOTALS		275	6.6	41.7	13.8	
357	1.0	60	6.0	10.0	13.8	0.3812

6.2 FINAL REMARKS

This work represents the first stage in exploiting numerical solution of stochastic models for the analysis of large systems. The major disadvantage of the technique is currently the relatively great difficulty of preparing correct discipline subroutines for complex models. It is hoped that this disadvantage can be overcome by the development of automatic programs to do this work from a very formalized description of the model to be solved (i.e., a translator for a problem-oriented language). With that accomplished, we believe the RQA-1 and its successors represent a very powerful tool indeed for system analysis.

The programs listed have used the May 1, 1965, version of MAD and the May 5, 1965, version of UMAP. These programs are known as the September 1, 1965, version of RQA-1. They are designed to operate with The University of Michigan Executive System, although an effort was made to make the program as independent as possible. We believe that with the information presented here a potential user could adapt the program to any comparable system.

6.3 ACKNOWLEDGMENT

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APPENDIX A

LISTING OF RQA-1

This program was written in the MAD and UMAP languages and processed in The University of Michigan Executive System. The listings are complete for the September 1, 1965, version of the RQA-1 program. The listings for the DISCPL subroutines of the examples of Section 3 are given in Section 5. Two versions of ITER are supplied, one in UMAP and the other in MAD.


```

EXECUTE BREAD,(SCOPE)                                MAIN0202 *048 01
EXECUTE BREAD,(V(VFLAG,1)...V(VFLAG,SCOPE))          MAIN0203 *049 01

OR WHENEVER INIT.E.2
EXECUTE ZERO,(V(VFLAG,1)...V(VFLAG,SCOPE),V(CWFLAG,1)...V(CVF
1 LAG,SCOPE))
VFLAG = 0
CWFLAG = 1
PRINT FORMAT $I18,$I4,H+THE INITIAL ITERATE IS GIVEN BELOW+/I1--$
READ AND PRINT DATA                                MAIN0210 *050 01
                                                    MAIN0220 *051 01
                                                    MAIN0221 *051 01
                                                    MAIN0222 *052 01
                                                    MAIN0223 *053 01
                                                    MAIN0224 *054 01
                                                    MAIN0230 *055 01

OTHERWISE
PRINT COMMENT $8 THE INITIAL ITERATE WAS THE RES    MAIN0231 *056 01
1 ULT OF THE PREVIOUS RUN$                          MAIN0234 *057 01
END OF CONDITIONAL                                MAIN0240 *058 01

*****
INITIAL VECTOR NORMALIZED
*****
VECSUM=0.
THROUGH VEC1, FOR I=1,1,I.G.SCOPE
VECSUM=VECSUM+V(VFLAG,I)
THROUGH VEC2, FOR I=1,1,I.G.SCOPE
V(VFLAG,I)=V(VFLAG,I)/VECSUM

*****
ITERATION ROUTINE AND CONVERGENCE TESTING
*****
NIITER=0
MAXIX=CVGTST.(BEGIN)
SAVE1=V(VFLAG,MAXIX)
SAVE2=SAVE1
EXECUTE ITER,(I)
SAVE1=V(VFLAG,MAXIX)
WHENEVER.ABS.(SAVE2-SAVE1).LE.EPSI.OR.MXITER.LE.NIITER, TRANSFER TO MAIN1
TRANSFER TO MAIN2

INTEGER INIT,MXITER,I,SCOPE,VFLAG,NITER,MAXIX,MAXJ,
1 CVFLAG,DELTA,FIRST,DESI,LAST,EXIT1,EXIT2,J,CVGTST, RUN,DATE,
2 YEAR,SCOPE,NAME,TEMP,SCOPE1,NTRANS,CHK,PCH,TSW,PRSW,RUNP(1)
EQUIVALENCE (TRANS,NTRANS),(TRANS(2001),CHK)
PROGRAM COMMON TRANS(2001),FIRST(2000),DEST(2000),LAST(2000),
1 PCH,PRSW,MXITER,EPST,DELTA,INIT,RUN,TSW,NAME(1),MAXJ
ERASABLE DUMMY(100),VFLAG,CVFLAG,SCOPE,NITER,V(0...1)*5000,
1 ARATEM,VECSUM,EXIT1,EXIT2,EXIT3,I,J,DATE,YEAR,DIADD
END OF PROGRAM

```

```

MAIN0202 *048 01
MAIN0203 *049 01
MAIN0210 *050 01
MAIN0220 *051 01
MAIN0221 *051 01
MAIN0222 *052 01
MAIN0223 *053 01
MAIN0224 *054 01
MAIN0230 *055 01
MAIN0231 *056 01
MAIN0234 *057 01
MAIN0240 *058 01
MAIN0260 *059
MAIN0270 *060
MAIN0280 *061
MAIN0290 *062
MAIN0300 *063
MAIN0440 *064
MAIN0450 *065
MAIN0460 *066
MAIN0470 *067
MAIN0480 *068
MAIN0490 *069
MAIN0500 *070
MAIN0510 *071
MAIN0520 *072
MAIN0530 *072
MAIN0535 *072
MAIN0545 *073
MAIN0550 *074
MAIN0560 *074
MAIN0580 *075
MAIN0585 *075
MAIN0600 *076

```

01
01

MAD (09 AUG 1965 VERSION) PROGRAM LISTING

REFERENCES ON

EXTERNAL FUNCTION
ENTRY TO STOCAL.

RATEM=0.
THROUGH STOC1, FOR I=1,I,I.G.NTRANS
WHENEVER RATEM.G.TRANS(I), RATEM=TRANS(I)
CONTINUE

STOC1

WHENEVER RATEM.GE.0.
ARATEM =1./0.99
DIADD=.01
TRANSFER TO STOC3
END OF CONDITIONAL

ARATEM=(ABS.RATEM)/.99

DIADD=1.0
THROUGH STOC2, FOR I=1,I,I.G.NTRANS
WHENEVER FIRST(I).NE.DEST(I)
TRANS(I)=TRANS(I)/ARATEM
OTHERWISE

STOC3

TRANS(I)=(TRANS(I)/ARATEM)+DIADD
END OF CONDITIONAL
CONTINUE

STOC2

INTEGER I, NTRANS, VFLAG, CVFLAG, DELTA, SCOPE, NITER, FIRST, DEST, LA
1 ST, MXITER, INIT, PCH, MAXJ, SCOPET
EQUIVALENCE (TRANS, NTRANS)
PROGRAM COMMON TRANS(2001), FIRST(2000), DEST(2000), LAST(2000),
1 PCH, PRSN, MXITER, EPSI, DE-FA, INIT, RUN, TSM, NAME(1), MAXJ
ERASABLE DUMMY(100), VFLAG, CVFLAG, SCOPE, NITER, V(0...1)*5000),
1 ARATEM, VECSUM, EXIT1, EXIT2, EXIT3, I, J, DATE, YEAR, DIADD
FUNCTION RETURN
END OF FUNCTION

*001

STOC0010
STOC0020
STOC0030
STOC0040
STOC0050
STOC0060

*001
*002
*003
*004
*005
*006

01
01

STOC0070
STOC0090
STOC0100
STOC0105
STOC0106

*007
*008
*009
*010
*011

01
01
01
01

STOC0110
STOC0115
STOC0120
STOC0130
STOC0140
STOC0150
STOC0160
STOC0170
STOC0180

*012
*013
*014
*015
*016
*017
*018
*019
*020

01
01
01
01
01
01
01
01

STOC0190
STOC0200
STOC0210
STOC0220
STOC0230
STOC0245
STOC0246
STOC0260
STOC0270

*021
*022
*023
*024
*025
*026

REFERENCES ON

EXTERNAL FUNCTION
ENTRY TO TEST.

TEST FOR LAST(I)-FIRST(I)=R*DELTA, WHERE R IS AN INTEGER

ZERO.(V(CVFLAG,1)..V(CVFLAG,SCOPE))
EXIT1=0

THROUGH TES1, FOR I=1,1,I.G.NTRANS
THROUGH TES2, FOR J=FIRST(I),DELTA,J.G.LAST(I)
V(CVFLAG,J)=V(CVFLAG,J)+TRANS(I)
WHENEVER J.E.LAST(I)+DELTA.AND.LAST(I).LE.SCOPE.AND.FIRST(I)
I .G.0, TRANSFER TO TES1

EXIT1=EXIT1+1
WHENEVER EXIT1.E.1,EXECUTE TOUTT.
EXECUTE TOUT1.
CONTINUE

TES1

TEST FOR STATE ROW SUMS = 1. AND POINT ROW SUMS = 0.

EXIT2=0
THROUGH TES3, FOR I=1,1,I.G.SCOPE
WHENEVER V(CVFLAG,I).E.0.

TRANSFER TO TES3
OR WHENEVER .ABS.(1-V(CVFLAG,I)).LE.1.E-7
TRANSFER TO TES3

OTHERWISE
EXIT2=EXIT2+1
WHENEVER EXIT2.E.1,EXECUTE TOUTC.
EXECUTE TOUT2.
END OF CONDITIONAL
CONTINUE

TES3

TEST FOR POINT COLUMNS = 0

EXIT3=0
THROUGH TES5, FOR I=1,1,I.G.NTRANS
THROUGH TES4, FOR J=DEST(I),DELTA,J.G.LAST(I)-FIRST(I)+DEST(I)
1)

*001

TEST0010
TEST0020

*001
*002

TEST0030 *003
TEST0040 *004
TEST0050 *005
TEST0060 *006
TEST0070 *007
TEST0090 *008
TEST0095 *008
TEST0096 *009
TEST0100 *010
TEST0110 *011
TEST0120 *012

01
02
01
01
01
01

TEST0130 *013
TEST0140 *014
TEST0150 *015
TEST0160 *016
TEST0170 *017
TEST0180 *018
TEST0190 *019
TEST0200 *020
TEST0210 *021
TEST0220 *022
TEST0230 *023
TEST0240 *024

01
01
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01
01
01
01

TEST0251 *025
TEST0252 *026
TEST0254 *027

01

WHENEVER V(CVFLAG,J).NE.0,TRANSFER TO TES4
EXIT3 = EXIT3+1
WHENEVER EXIT3.E.1,EXECUTE TOUTM.
EXECUTE TOUT1.
TRANSFER TO TES5
CONTINUE
CONTINUE

TES4
TES5

WHENEVER EXIT1.NE.0.OR.EXIT2.NE.0.OR.EXIT3.NE.0, EXECUTE
1 TOUT3.
FUNCTION RETURN

INTEGER CVFLAG,EXIT1,I,NTRANS,J,FIRST,DELTA,LAST,EXIT2,SCOPE,
1 VFLAG,NITER,DEST,MXITER,INIT,PCH,MAXJ,EXIT3
EQUIVALENCE (TRANS,NTRANS)
PROGRAM COMMON TRANS(2001),FIRST(2000),DEST(2000),LAST(2000),
1 PCH,PRSM,MXITER,EPSI,DELTA,INIT,RUN,TSM,NAME(1),MAXJ
ERASABLE DUMMY(100),VFLAG,CVFLAG,SCOPE,NITER,V((0...1)*5000),
1 ARATEM,VECSUM,EXIT1,EXIT2,EXIT3,I,J,DATE,YEAR,DIADD
END OF FJUNCTION

TEST0255 *028
TEST0256 *029
TEST0257 *030
TEST0258 *031
TEST0259 *032
TEST0260 *033
TEST0261 *034

TEST0262 *035
TEST0263 *035
TEST0264 *036

TEST0280 *037
TEST0290 *037
TEST0300 *038
TEST0310 *039
TEST0320 *039
TEST0340 *040
TEST0345 *040
TEST0360 *041

MAD (09 AUG 1965 VERSION) PROGRAM LISTING

REFERENCES ON

EXTERNAL FUNCTION

ENTRY TO TOUTT.

PRINT COMMENT \$1 THE FOLLOWING QUADRUPLES HAVE

1 AN ERROR IN FIRST OR LAST. J IS FIRST + N*DELTA, WHERE N IS

2 \$ AN INTEGER.\$

PRINT COMMENT \$

PRINT FORMAT TITLE

FUNCTION RETURN

ENTRY TO TOUTM.

PRINT COMMENT \$1 THE FOLLOWING QUADRUPLES HAVE

1 AN ERROR IN DEST. J IS A POINT INTO WHICH TRANSITION IS INDI

2 CALLED.\$

PRINT FORMAT TITLE

FUNCTION RETURN

ENTRY TO TOUTL.

PRINT FORMAT OUT1,I,TRANS(I),FIRST(I),DEST(I),LAST(I),J

FUNCTION RETURN

ENTRY TO TOUTC.

PRINT COMMENT \$1 THE FOLLOWING ROW SUMS, SUM (I)

1 , OF THE STOCHASTIC MATRIX ARE INCORRECT.\$

PRINT FORMAT TITLEC

FUNCTION RETURN

ENTRY TO TOUT2.

PRINT FORMAT OUT2,I,(CVFLAG,I)

FUNCTION RETURN

ENTRY TO TOUT3.

PRINT COMMENT \$1 THE ABOVE IS A FATAL ERROR IN T

HE TRANSITION TABLE

PRINT FORMAT \$IHC,\$I4,H+IF (IU) IS THE MATRIX DESCRIBED BY D

1 ISCP. THEN THIS TABLE DESCRIBES+I4.7,H*(IU) +,E14.7,

2 H*(I)**\$,I./ARATEM,DIADD

PRINT FORMAT TITLE

THROUGH TU, FOR I=1,I.G.NTRANS

PRINT FORMAT OUT1,I,TRANS(I),FIRST(I),DEST(I),LAST(I)

WHENEVER SCOPE .G.500,TRANSFER TO TUC

TU

PRINT COMMENT \$1

1 OF THE STOCHASTIC MATRIX \$

PRINT COMMENT \$8\$

THROUGH TUT, FOR I=1,8,I+7,G.SCOPE

PRINT FORMAT OUT4,I,I+7,(CVFLAG,I)...V(CVFLAG,I+7)

WHENEVER I.LE.SCOPE,PRINT FORMAT OUT4,I,SCOPE,V(CVFLAG,I)...V

1 (CVFLAG,SCOPE)

EXECUTE ERROR.

TUO

*001

TOUT0010 *001

TOUT0020 *002

TOUT0030 *003

TOUT0040 *003

TOUT0045 *003

TOUT0046 *004

TOUT0050 *005

TOUT0060 *006

TOUT0061 *007

TOUT0062 *008

TOUT0063 *008

TOUT0064 *008

TOUT0065 *009

TOUT0066 *010

TOUT0070 *011

TOUT0080 *012

TOUT0090 *013

TOUT0100 *014

TOUT0110 *015

TOUT0120 *015

TOUT0130 *016

TOUT0140 *017

TOUT0150 *018

TOUT0160 *019

TOUT0170 *020

TOUT0180 *021

TOUT0190 *022

TOUT0195 *022

TOUT0200 *023

TOUT0201 *023

TOUT0202 *023

TOUT0220 *024

TOUT0230 *025

TOUT0240 *026

TOUT0245 *027

TOUT0250 *028

TOUT0260 *028

TOUT0261 *029

TOUT0270 *030

TOUT0280 *031

TOUT0285 *032

TOUT0286 *032

TOUT0290 *033

01

01

```

VECTOR VALUES TITLE=$1H8,S15,1H1,S8,8HTRANS(I),S11,8HFIRST(I
1 ),S8,7HDEST(I),S9,7HLAST(I),S11,1HJ/1H *$
VECTOR VALUES OUT1=$117,F15,8,4116*$
VECTOR VALUES TITLE=$1H8,S15,1H1,S9,6HSUM(I)/1H *$
VECTOR VALUES OUT2=$14,I4,S4,F11,8*$
VECTOR VALUES TITLE=$1H8,S15,1H1,S8,8HTRANS(I),S11,8HFIRST(I
1 ),S8,7HDEST(I),S9,7HLAST(I)/1H *$
VECTOR VALUES OUT4=$14,I4,3H -,I4,S3,8(E8,6,S3)/1H *$

INTEGER FIRST,DEST,LAST,I,J,CVFLAG,NTRANS,VFLAG,DELTA,SCOPE,
1 NITER,MXITER,INIT,PCH,MAXJ,EXIT1,EXIT2,EXIT3
EQUIVALENCE (NTRANS,NTRANS)
PROGRAM COMMON TRANS(2001),FIRST(2000),DEST(2000),LAST(2000),
1 PCH,PRSW,MXITER,EPST,DELTA,INIT,RUN,ISM,NAME(I),MAXJ
1 ARATEM,VECSUM,EXIT1,EXIT2,EXIT3,I,J,DATE,YEAR,DIADD
FUNCTION RETURN
END OF FUNCTION

```

```

TOUT0380 *034
TOUT0390 *034
TOUT0400 *035
TOUT0410 *036
TOUT0420 *037
TOUT0430 *038
TOUT0440 *038
TOUT0460 *039

TOUT0490 *040
TOUT0500 *040
TOUT0510 *041
TOUT0520 *042
TOUT0530 *042
TOUT0540 *043
TOUT0550 *043
TOUT0590 *044
TOUT0600 *045

```

MAD (03 JAN 1966 VERSION) PROGRAM LISTING

REFERENCES ON

EXTERNAL FUNCTION
ENTRY TO ESTIM.

LAMSTR=0.
MUSTR=0.
T2=0
T3=0

THROUGH ES1, FOR I=1,1,I.G.NTRANS
DIFSTR=FIRST(I)-DEST(I)
T1=((LAST(I)-FIRST(I))/DELTA)+1
WHENEVER DIFSTR.E.0,TRANSFER TO ES1
WHENEVER DIFSTR.G.0,MUSTR=MUSTR+TRANS(I)*T1
WHENEVER DIFSTR.L.0,LAMSTR=LAMSTR+TRANS(I)*T1
T2=T2+.ABS.DIFSTR*T1
T3=T3+T1
CONTINUE
RHOSTR=(LAMSTR/MUSTR).P.((1.*T3)/T2)

ES1

***** ESTIMATE ONLY VALUES WHICH ARE LARGER THAN DESIRED ACCURACY WHEN NORMALIZED. *****

***** ESTIMATE ONLY VALUES WHICH ARE LARGER THAN DESIRED ACCURACY WHEN NORMALIZED. *****

***** ESTIMATE ONLY VALUES WHICH ARE LARGER THAN DESIRED ACCURACY WHEN NORMALIZED. *****

T1=1
T2=SCOPE
SCOPE=.ABS.(ELOG.(EPSI)/ELOG.(RHOSTR))
WHENEVER SCOPE.GE.SCOPE,TRANSFER TO ES2
WHENEVER RHOSTR.GE.1.,T1=SCOPE-SCOPE
WHENEVER RHOSTR.L.1.,T2=SCOPE
EXECUTE ZERO.(V(1)..V(SCOPE))
RHON=1.
CVFLAG=1
VFLAG=0

ES2

THROUGH ES3, FOR I=1,1,I.G.T2
RHON=RHON#RHOSTR
V(I)=RHON

ES3

***** ESTIMATE ONLY VALUES WHICH ARE LARGER THAN DESIRED ACCURACY WHEN NORMALIZED. *****

ZERO OUT VALUES AT POINTS BY PERFORMING AN ITERATION

***** ESTIMATE ONLY VALUES WHICH ARE LARGER THAN DESIRED ACCURACY WHEN NORMALIZED. *****

EXECUTE ITER.(1)
FUNCTION RETURN

1 INTEGER T1,T2,T3,SCOPE,DIFSTR,I,NTRANS,FIRST,DEST,LAST,
DELTA,K,J,CVFLAG,VFLAG,SCOPE,NITER,MXITER,INIT,PCH,MAXJ
EQUIVALENCE (TRANS,NTRANS)

*C01
ESTM0010
ESTM0020
ESTM0030
ESTM0040
ESTM0050
ESTM0060
ESTM0070
ESTM0080
ESTM0090
ESTM0100
ESTM0110
ESTM0120
ESTM0130
ESTM0140
ESTM0150
ESTM0160

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01
01

*017
ESTM0170
*018
ESTM0180
*019
ESTM0190
*020
ESTM0200
*021
ESTM0210
*022
ESTM0220
*023
ESTM0230
*024
ESTM0240
*025
ESTM0250
*026
ESTM0260
*027
ESTM0270
*028
ESTM0280
*029
ESTM0290

01
01

*030
ESTM0300
*031
ESTM0310
*032
ESTM0320
*033
ESTM0330

01
01

```
PROGRAM COMMON TRANS(2001),FIRST(2000),DEST(2000),LAST(2000),
PCH,PRSW,MXITER,EPSI,DELTA,INIT,RUN,TSM,NAME(1),MAXJ
ERASABLE DUMMY(100),VFLAG,CVFLAG,SCOPE,NITER,(0...1)*5000),
ARATEM,VECSUM,EXIT1,EXIT2,EXIT3,I,J,DATE,YEAR,DIADC
END CF FUNCTION
```

```
ESTM0380
ESTM0390
ESTM0410
ESTM0420
ESTM0430
```

```
*034
*034
*035
*C35
*036
```


PAGE 1 (MOD 05 MAY 1965) UNIVERSITY OF MICHIGAN 13 SEP 1965

PROGRAM LENGTH (OCTAL) 00066 PROGRAM COMMON LENGTH 17520 LOWEST ERASABLE DEFINED 54206 ID ITER0000
 INITIAL TRANSFER VECTOR 00000 TRANSFER VECTOR LENGTH 00000

PROGRAM ENTRY POINT LOCATIONS AND NAMES
 00000 ITER

PROGRAM LISTING (NUMBERS IN OCTAL)

	ENTRY	ITER	START	LXA	VFLAG,1
00000	0634 00 4 00060	010	ITER	SXA	XR4,4
00001	0500 60 4 00001	00		CLA*	1,4
00002	0601 00 0 00064	010		STO	N
00003	0500 00 0 27511	00		CLA	DELTA
00004	0767 00 0 00022	00		ALS	18
00005	0622 00 0 00044	010		STD	D1
00006	0622 00 0 00045	010		STD	D2
00007	0534 00 1 77632	010	*	START	
00010	0500 00 1 00063	010		CLA	AD2,1
00011	0621 00 0 00042	010		STA	RV1
00012	0621 00 0 00043	010		STA	RV2
00013	0621 00 0 00025	010		STA	RV3
00014	0500 00 0 77632	010		CLA	VFLAG
00015	0322 00 0 00065	010		ERA	=1
00016	0734 00 1 00000	00		PAX	.1
00017	0500 00 1 00063	010		CLA	AD2,1
00020	0621 00 0 00040	010		STA	OPV1
00021	0500 00 0 77627	010		CLA	NITER
00022	0400 00 0 00065	010		ADD	=1
00023	0601 00 0 77627	010		STO	NITER
00024	0534 00 1 77630	010		LXA	SCOPE,1
00025	0600 00 1 00000	00	RV3	STZ	**1
00026	2 00001 1 00025	010	*	TIX	**1,1,1
00027	0534 00 4 13721	00		LXA	TRANS,4
00030	0500 00 4 27504	00	KI	CLA	LAST,4
00031	0767 00 0 00022	00		ALS	18
00032	0622 00 0 00037	010		STD	L1
00033	0500 00 4 17642	00		CLA	FIRST,4
00034	0734 00 2 00000	00		PAX	,2
00035	0500 00 4 23563	00		CLA	DEST,4
00036	0734 00 1 00000	00	*	PAX	,1
00037	3 00000 2 00046	010	*	L1	TXH
00040	0560 00 2 00000	00	OPV1	LDQ	L2,2,**
00041	0260 00 4 13721	00	F1	FMP	**2
00042	0300 00 1 00000	00	RV1	FAD	TRANS,4
00043	0601 00 1 00000	00	RV2	STO	**1
00044	1 00000 1 00045	010	D1	TXI	**1,1,**
00045	1 00000 2 00037	010	D2	TXI	L1,2,**
00046	2 00001 4 00030	010	L2	TIX	KI,4,1
00047	0500 00 0 77632	010	*	CLA	VFLAG
00050	0601 00 0 77631	010		STO	CVFLAG
00051	0322 00 0 00065	010		ERA	=1

ID ITER0001

ID ITER0002

ID ITER0003

ADDRESS	DATA	SYMBOL	VALUE	TYPE	LOCATION	SYMBOL	VALUE	TYPE	LOCATION
00052	0601 00 0	77632	010		STO	VFLAG			
00053	0500 00 0	00064	010		CLA	N			
00054	0402 00 0	00065	010		SUB	=1			
00055	0601 00 0	00064	010		STO	N			
00056	0520 00 0	00064	010		ZET	N			
00057	0020 00 0	00007	010		TRA	START			
00060	0774 00 4	00000	00		AXT	**4			
00061	0020 00 4	00002	00		TRA	2,4			
00062	0 0000 0	77626	010		PZE	V			
00063	0 0000 0	66016	010		PZE	V-5000			
00064	0 0000 0	00000	00		PZE				
		13721			PGMCOM	2001			
		17642			PGMCOM	2000			
		23563			PGMCOM	2000			
		27504			PGMCOM	2000			
		27510			PGMCOM	3			
		27511			PCLIST	DELTA			
		27517			PGMCOM	5			
		77632			ERAS	100			
		77632			ERLIST	VFLAG,CVFLAG,SCOPE,NITER			
		77631			VFLAG				
		77630			CVFLAG				
		77627			SCOPE				
		77626			NITER				
					V				
					ERAS	10000			
					END				

PROGRAM LITERALS

00065 000000000001 00

*** 00066 IS THE FIRST LOCATION NOT USED BY THIS PROGRAM.

REFERENCES TO SYMBOLS DEFINED IN PROGRAM.

TYPE SYMBOL VALUE REFERENCES TO SYMBOL (SYMBOL TYPES -- R - RELOCATABLE, A - ABSOLUTE, E - ERASABLE, T - TRANSFER VECTOR)

R	AD1	00064	00062
R	AD2	00063	00010,00017,00063
E	CVFLAG	77631	00059,00065
R	O1	00044	00005,00044
R	O2	00045	00006,00045
A	DELTA	27511	00003,00065
A	DEST	23563	00035,00065
R	FI	00041	00041
A	FIRST	17642	00033,00065
R	ITER	00000	00000,00000,00000
R	K1	00030	00030,00046
R	L1	00037	00032,00037,00045
R	L2	00046	00037,00046
A	LAST	27504	00030,00065
R	N	00064	00002,00053,00055,00056,00064
E	NITER	77627	00021,00023,00065
R	OPV1	00040	00020,00040
R	RV1	00042	00011,00042
R	RV2	00043	00012,00043
R	RV3	00025	00013,00025
E	SCOPE	77630	00024,00065
R	START	00007	00037,00057

A TRANS 13721 00027,00041,00065
 E V 77626 00062,00063,00065
 E VFLAG 77632 00007,00014,00047,00052,00065
 R XR4 00060 00000,00060

REFERENCES TO PROGRAM LITERALS

LOC REFERENCES
 00065 00015,00022,00051,00054

FLAGGED LINES	FATAL 0	NON-FATAL 0	****	DEFINITIONS	SYMBOL 26	OPERATION 0	MACRO 0	RMT 0
MACRO CALLS	CORRECT 0	INCORRECT 0	****	CARD COUNTS	INPUT 71	OUTPUT 4		

MAD (09 AUG 1965 VERSION) PROGRAM LISTING

REFERENCES ON *001
 EXTERNAL FUNCTION *001
 ENTRY TO CVGTST. *002

 IF NITER = 0, SET TRAP PROCEDURE SO THAT UNDERFLOWS
 PRODUCE ZEROS. ALSO FIND SUITABLE INDEX FOR WEAK TEST.

 VECTOR VALUES FTSM=0
 WHENEVER FTSM.E.0
 WHENEVER NITER .GE. MXITER
 MAXDIF=0.
 TRANSFER TO CV3
 END OF CONDITIONAL
 FTSM=1
 EXECUTE FTRAP.
 THROUGH CV0, FOR I=1,1,VVFLAG,I).G.0..OR.I.E.SCOPE
 FUNCTION RETURN I
 END OF CONDITIONAL

CV0

FIND MAXIMUM DIFFERENCE AND TEST AGAINST EPSI.

 MAXDIF=0.
 THROUGH CV2, FOR I=1,1,I.G.SCOPE
 DIF=.ABS.(V(VFLAG,I)-V(CVFLAG,I))
 WHENEVER MAXDIF.G.DIF, TRANSFER TO CV2
 MAXDIF=DIF
 MIX=1
 CONTINUE

CV2

 WHENEVER MAXDIF .LE. EPSI
 PRINT COMMENT \$8 ***** SUCC
 1 LESSFUL RUN *****\$
 TRANSFER TO CV1

 OR WHENEVER NITER .GE. MXITER
 PRINT COMMENT \$8***** THE SPECIFIED NUMBER OF ITERATIONS
 1 HAS BEEN SATISFIED ***** \$

CV3

ESTIMATE SUBDOMINANT EIGENVALUE AND ERROR.

CVGT0003 *003
 CVGT0004 *004
 CVGT0005 *005
 CVGT0006 *006
 CVGT0007 *007
 CVGT0008 *008
 CVGT0014 *009
 CVGT0015 *010
 CVGT0016 *011
 CVGT0017 *012
 CVGT0018 *013

CVGT0030 *014
 CVGT0040 *015
 CVGT0050 *016
 CVGT0060 *017
 CVGT0070 *018
 CVGT0080 *019
 CVGT0090 *020
 CVGT0100 *021
 CVGT0105 *022
 CVGT0110 *022
 CVGT0115 *023
 CVGT0120 *024
 CVGT0125 *025
 CVGT0130 *025

CVGT0030 *014
 CVGT0040 *015
 CVGT0050 *016
 CVGT0060 *017
 CVGT0070 *018
 CVGT0080 *019
 CVGT0090 *020
 CVGT0100 *021
 CVGT0105 *022
 CVGT0110 *022
 CVGT0115 *023
 CVGT0120 *024
 CVGT0125 *025
 CVGT0130 *025

```

CVGT0131 *026 01
CVGT0132 *027 01
CVGT0133 *028 02
CVGT0134 *029 02
CVGT0135 *030 02
CVGT0136 *031 02
CVGT0137 *032 02
CVGT0138 *033 02

```

```

CVGT0139 *034 01
CVGT0140 *035 01
CVGT0145 *036 01

```

```

CVGT0150 *037 01
CVGT0155 *038 01
CVGT0160 *039 01
CVGT0165 *040 01

```

```

CVGT0170 *041
CVGT0180 *041
CVGT0182 *041
CVGT0184 *041
CVGT0186 *041
CVGT0180 *041

```

```

CVGT0185 *042
CVGT0186 *042
CVGT0190 *043
CVGT0200 *043
CVGT0205 *044
CVGT0210 *044
CVGT0290 *045

```

```

CW1 LDIF=MAXDIF.P.(1./NITER)
WHENEVER LDIF.GE.1..OR.LDIF.L..5
PRINT FORMAT CVFL,EPSI,MXITER,NITER,MAXDIF,MAXDIF,0.
OR WHENEVER (1.-LDIF).P.(1./NITER).L.LDIF
PRINT FORMAT CVFL,EPSI,MXITER,NITER,MAXDIF,0.,LDIF
OTHERWISE
PRINT FORMAT CVFL,EPSI,MXITER,NITER,MAXDIF,MAXDIF/LDIF,LDIF
END OF CONDITIONAL

```

```

*****

```

```

RESET TRAP PROCEDURE SO THAT UNDERFLOWS AGAIN CAUSE
ERROR RETURNS.

```

```

*****

```

```

EXECUTE NTRAP,
FTSM=0

```

```

EXECUTE OUT2.

```

```

*****

```

```

THE ERROR RETURN TRANSFERS TO 'BEGIN' IN THE MAIN PROGRAM

```

```

*****

```

```

ERROR RETURN
OTHERWISE
FUNCTION RETURN MIX
END OF CONDITIONAL

```

```

VECTOR VALUES CVFL=I1H-,S14,35HCONVERGENCE TEST CRITERION (EP
1 SI) =F11.8/1H0,S12,37HMAXIMUM ALLOWED ITERATIONS (MXITER) =I5
2 /1H0,S31,18HITERATIONS TAKEN =I5/1H0, S3,48HMAXIMUM DIFFERENC
3 E BETWEEN LAST TWO ITERATES =E8.1/1H0, S23,26HESTIMATED ORDER
4 OF ERROR =E8.1,S3,33H(APPROXIMATE CONVERGENCE FACTOR =F9.7,
5 IH)*$

```

```

INTEGER I,SCOPE,VFLAG,CVFLAG,MIX,DELTA,NITER,FIRST,DEST,LAST,
1 MXITER,INIT
PROGRAM COMMON TRANS(2001),FIRST(2000),DEST(2000),LAST(2000),
1 PCH,PRSM,MXITER,EPSI,DELTA,INIT,RUN,ISM,NAME(I),MAXJ
ERASABLE DUMMY(100),VFLAG,CVFLAG,SCOPE,NITER,V((0..1)*5000),
1 ARATEM,VECSUM,EXIT1,EXIT2,EXIT3,I,J,DATE,YEAR,DIADD
END OF FUNCTION

```

\$COMPILE MAD,EXECUTE,DUMP,I/O DUMP,PRINT OBJECT,PUNCH OBJECT QUIT0000

MAD (09 AUG 1965 VERSION) PROGRAM LISTING

REFERENCES ON

EXTERNAL FUNCTION

ENTRY TO OUT2.

WHENEVER PCH=NE.1, TRANSFER TO UT9

PRINT COMMENT \$- THE SOLUTION VECTOR HAS BEEN PUNCHED\$

OUTPUT VECTOR CARDS WILL HAVE THE FOLLOWING ID IN
COLUMNS 73-80 .

(1) MODEL NAME

(2) DATE

(3) RUN NO. XX 001

ALL SUCCEEDING CARDS WILL HAVE RUN NO. XX VECTOR CARD NO.

PUNCH FORMAT \$IH , S14,2C6,S1,4HRUN ,I3,2H, ,2C6,S23,C6,C2*\$,

1 NAME,NAME(1),RUN,DATE,YEAR,NAME,NAME(1)

EXECUTE NSDBC.D.(DATE,\$ \$)

EXECUTE BPUNCH.(SCOPE)

EXECUTE NSDBC.D.((BNBCD.(RUN).LS.18).V,\$000 0\$, \$01\$)

EXECUTE BPUNCH.(V(VFLAG,1)..V(VFLAG,SCOPE))

UT9 WHENEVER PRSW.E.1,EXECUTE PROUT.

RUN = RUN + 1

INTEGER NITER,MXITER,I,SCOPE,J,K,DELTA,VFLAG,CVFLAG,FIRST,DES

1 I, LAST, INIT, MAXJ, S, RUN, DATE, YEAR, NAME, BNBCD., PCH, PRSW

PROGRAM COMMON TRANSI(2001), FIRST(2000), DEST(2000), LAST(2000),

1 PCH, PRSW, MXITER, EPSI, DELTA, INIT, RUN, TSM, NAME(1), MAXJ

ERASABLE DUMMY(100), VFLAG, CVFLAG, SCOPE, NITER, V(0...1)*5000),

1 ARATEM, VECSUM, EXIT1, EXIT2, EXIT3, I, J, DATE, YEAR, DIADD

FUNCTION RETURN

END OF FUNCTION

001632 09/15/65 6 18 16.2 PM

*001

QJTT0010 *001

QJTT0020 *002

QJTT0300 *003

QJTT0320 *004

QJTT0351 *005

QJTT0352 *005

QJTT0354 *006

QJTT0355 *C07

QJTT0356 *008

QJTT0357 *009

QJTT0370 *010

QJTT0380 *011

QJTT0430 *012

QJTT0440 *012

QJTT0470 *013

QJTT0480 *013

QJTT0495 *014

QJTT0496 *014

QJTT0510 *015

QJTT0520 *016

MAD (09 AUG 1965 VERSION) PROGRAM LISTING

REFERENCES ON

EXTERNAL FUNCTION
ENTRY TO PROUT.

*001
*001
*002

PJUT0010
PJUT0020

PROBABILITY VECTOR PRINTED OUT

PRINT FORMAT \$IH1,S20,H+THE FOLLOWING ARE THE LIMITING STATE
1 PROBABILITIES, PROB(I,J).+/S20, H+ PROB(I) IS THE SUM OF PR
2 OB(I,J) OVER ALL STATES WITH THE VALUE I FOR THE FIRST VARIAB
3 LE+/IH-*\$
WHENEVER (SCOPE/MAXJ)+1.G.2500.OR.MAXJ.G.2500
PRINT COMMENT \$ *** MAXJ, THE MAXIMUM COLUMN INDEX HAS B
1 EEN SET TO FIFTY TO ALLOW MORE EFFICIENT PRINTING OF THE PROB
2 ABILITY VECTOR.\$

MAXJ=50
END OF CONDITIONAL
PRINT FORMAT \$IH-,S3,IH1,S4,8H PROB(I),S40,9HPR0B(I,J)/IH-,
1 S19,3HJ =,S2,10(I1,S10)*\$, (I=0,1,I.G.9.OR.1.G.(MAXJ-1),I)
VECSUM=0.
EXECUTE ZERO.(MARG1...MARG1((SCOPE/MAXJ)+1),MARG2...MARG2(MAX
1 J))

THROUGH UTTA, FOR K=1,MAXJ,K.G.SCOPE
THROUGH UTTB, FOR S=K,1,S.G.(K+MAXJ-1).OR.S.G.SCOPE
I=(S-1)/MAXJ
I=S-1-I*MAXJ

UTTB
MARG1(I)=MARG1(I)+V(VFLAG,S)
MARG2(J)=MARG2(J)+V(VFLAG,S)
VECSUM=VECSUM+MARG1(I)
WHENEVER (K+MAXJ-1).G.SCOPE
PRINT FORMAT LINE,I,MARG1(I),V(VFLAG,K)...V(VFLAG,SCOPE)
OTHERWISE

UTTA
PRINT FORMAT LINE,I,MARG1(I),V(VFLAG,K)...V(VFLAG,K+MAXJ-1)
END OF CONDITIONAL
VECTOR VALUES LINE=#IH0,I4,S2,F10.8,S2,10(F10.8,S1)/(
1 S19,10(F10.8,S1))*\$

1 ALL STATES WITH THE VALUE J FOR THE SECOND VARIABLE+/IH-,
2 S8,8HCHECKSUM,S40,7HPR0B(J)/IH0,S19,3HJ =,S2,10(I1,S10)*\$,
3 (I=0,1,I.G.9.OR.1.G.(MAXJ-1),I)
PRINT FORMAT \$IH0,S6,F10.8,S2,10(F10.8,S1)/(S19,10(F10.8,
1 S1))*\$,VECSUM,MARG2(I0)...MARG2(MAXJ-1)

INTEGER NITER,MXITER,MAXJ,SCOPE,K,I,J,VFLAG,CVFLAG,S
DIMENSION MARG1(2500),MARG2(2500)
EQUIVALENCE (MARG1,FIRST),(MARG1(2001),TRANS),(MARG2,LAST)
1 ,(MARG2(2001),DEST)
PROGRAM COMMON TRANS(2001),FIRST(2000),DEST(2000),LAST(2000),

POUT0060 *003
POUT0070 *003
POUT0080 *003
POUT0090 *003
POUT0095 *004
POUT0100 *005
POUT0105 *005
POUT0106 *005
POUT0110 *006
POUT0115 *007
POUT0120 *008
POUT0121 *008
POUT0122 *009
POUT0127 *010
POUT0128 *010
POUT0130 *011
PJUT0140 *012
POUT0150 *013
POUT0160 *014
PJUT0170 *015
POUT0180 *016
POUT0190 *017
PJUT0200 *018
PJUT0210 *019
POUT0220 *020
POUT0230 *021
POUT0240 *022
PJUT0250 *023
POUT0260 *023
POUT0270 *024
POUT0280 *024
PJUT0290 *024
POUT0291 *024
POUT0310 *025
POUT0320 *025

POUT0330 *026
POUT0340 *027
POUT0345 *028
PJUT0346 *028
POUT0350 *029

```
1 PCH, PRSW, MXITER, EPSI, DELTA, INIT, RUN, TSW, NAME(1), MAXJ
  ERASABLE DUMMY(100), VFLAG, CVFLAG, SCOPE, NITER, V((0...1)*5000),
1 ARATEM, VECSUM, EXIT1, EXIT2, EXIT3, I, J, DATE, YEAR, DIADD
  FUNCTION RETURN
  END OF FUNCTION
```

```
POUT0360
POUT0380
POUT0390
POUT0400
POUT0410
```

```
*029
*030
*030
*031
*032
```


MAD (09 AUG 1965 VERSION) PROGRAM LISTING

REFERENCES ON

```

EXTERNAL FUNCTION (N)
ENTRY TO ITER.
NORMAL MODE IS INTEGER
FLOATING POINT V, TRANS
THROUGH END, FOR CNTR=1,1,CNTR.G.N
EXECUTE ZERO.(V(CVFLAG,1)...V(CVFLAG,SCOPE))
NITER = NITER+1
THROUGH KLOOP, FOR K=1,1,K.G.NTRANS
J=FIRST(K)
I=DEST(K)
V(CVFLAG,I)=V(VFLAG,J)*TRANS(K)+V(CVFLAG,I)
I=I+DELTA
J=J+DELTA
WHENEVER J.LE.LAST(K),TRANSFER TO COMPLP
VFLAG=VFLAG.EV.1
CVFLAG=CVFLAG.EV.1
PROGRAM COMMON TRANS(2001),FIRST(2000),DEST(2000),LAST(2000),
1 PCH,PRSW,MXITER,EPSI,DELTA,INIT,RUN,TSW,NAME(1),MAXJ
\ ERASABLE DUMMY(100),VFLAG,CVFLAG,SCOPE,NITER,V(10...1)*5000),
1 ARATEM,VECSUM,EXIT1,EXIT2,EXIT3,I,J,DATE,YEAR,DIADD
EQUIVALENCE (NTRANS,TRANS)
FUNCTION RETURN
END OF FUNCTION
    
```

COMPLP

KLOOP

END

```

ITER0010 *001
ITER0020 *001
ITER0030 *002
ITER0040 *003
ITER0050 *004
ITER0060 *005
ITER0070 *006
ITER0080 *007
ITER0090 *008
ITER0100 *009
ITER0110 *010
ITER0120 *011
ITER0130 *012
ITER0140 *013
ITER0150 *014
ITER0160 *015
ITER0170 *016
ITER0180 *017
ITER0190 *017
ITER0210 *018
ITER0220 *018
ITER0230 *019
ITER0240 *020
ITER0250 *021
    
```

```

01
01
01
02
02
02
02
02
02
02
01
01
01
    
```

APPENDIX B

SOME SUBROUTINES FOR USE IN GENERATING QUADRUPLES IN RQA-1

James H. Jackson

B.1 PURPOSE

A three-entry subroutine has been written in an attempt to simplify the programming of discipline subroutines for use with RQA-1. The entries and their functions are as follows:

SETUP - Initializes parameters internal to the subroutine

QUAD - Produces the quadruples which represent transitions between states (the off-diagonal elements of the Q matrix)

DIAG - Produces the quadruples which represent the diagonal elements of the Q matrix.

B.2 CALLING SEQUENCES (MAD)

SETUP. (STATE.)

STATE is the name of an integer-valued internal function whose value is the state number represented by the current values of the state variables. Its values may range from 1 through 5,000.

QUAD. (RATE, CHANGE, REPEAT)

RATE is a floating point expression whose value is the transition rate for the generated quadruple.

CHANGE is an integer expression whose value is the amount by which the state number is changed when a transition indicated by the generated quadruple occurs.

REPEAT is an integer expression chosen such that the generated quadruple will apply to STATE., STATE. +DELTA, ..., STATE.+REPEAT*DELTA.

DIAG.

This entry has no arguments.

B.3 RESTRICTIONS

SETUP must be called before either QUAD or DIAG is called, otherwise an error comment will be printed and control will be returned to the system.

MAXJ must be initialized before SETUP is called.

QUAD must be called as many times as needed to produce all of the quadruples which represent transitions between states. The last call to QUAD must occur before DIAG is called.

After DIAG is called, no more calls should be made to the subroutine until RQA-1 has processed the generated data.

B.4 DISCIPLINE PROGRAM STRUCTURE

The discipline program which calls SETUP, QUAD, and DIAG may have the following structure:

```
EXTERNAL FUNCTION
ENTRY TO DISCIPL.
READ AND PRINT DATA
MAXJ = ----
DELTA = ----
SETUP.(STATE.)
THROUGH S, FOR ----
----
THROUGH S, FOR ----
WHENEVER ----, TRANSFER TO S
WHENEVER ----, QUAD.(----,----,----)
----
WHENEVER ----, QUAD.(----,----,----)
CONTINUE
DIAG.
FUNCTION RETURN
FLOATING POINT EPSI, ----
NORMAL MODE IS INTEGER
PROGRAM COMMON DUMMY (8004), PCH, PRSW, MXITER, EPSI,
lDELTA, INIT, RUN, TSW, NAME(1), MAXJ
VECTOR VALUES NAME = ----
INTERNAL FUNCTION STATE. (X) = ----
END OF FUNCTION
```

When the discipline subroutine is called, it reads a set of data. This data should include values of the parameters which describe the model as well as the values of PCH, PRSW, MXITER, EPSI, INIT, RUN, and TSW which are to be different from the preset values determined by RQA-1.

After the data is read, MAXJ and DELTA are initialized. Then SETUP is called to initialize internal parameters to the subroutine.

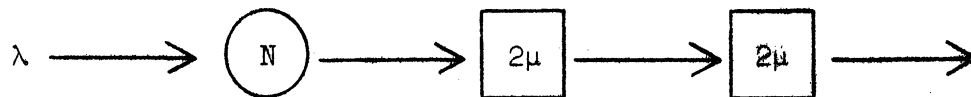
A set of nested iterations is then begun. Each iteration variable is a state variable. Hence, within the innermost loop, all possible combinations of state variables will be produced.

Since impossible combinations of state variables may not always be eliminated by constraining the set of values each variable may assume, the first conditional statement shown specifies a transfer to the end of the loop when any constraint of the model is violated.

The remaining conditional statements generate quadruples which represent transitions between states whenever transitions may occur from the state represented by the current values of the state variables. When all of these quadruples have been generated, the diagonal elements of the Q matrix are generated, and control is returned to RQA-1.

B.5 EXAMPLE

As an illustration of the programming of a discipline routine using SETUP, QUAD, and DIAG, one might consider a second-order Erlang channel with a finite queue:



- N = Maximum length of queue
- λ = Arrival rate
- μ = Processing rate

He may choose to define the state variables as follows:

- A = Number of Erlang channels busy
- B = Number of Erlang channels in second phase
- C = Queue Length

The ranges of these state variables are the following:

$$\begin{aligned}
 0 &\leq A \leq 1 \\
 0 &\leq B \leq A \\
 0 &\leq C \leq N .
 \end{aligned}$$

Since RQA-1 will print a two-dimensional probability array, the state space must be described in terms of two state variables to facilitate interpretation of the results. One may choose one of these two variables to be the state of

the service and the other to be the queue length. Since the state of the service is completely defined by A and B, it will be denoted by S(A,B), whereas the queue length is simply C.

S(A,B) may be defined by the following table:

A	B	S(A,B)
0	0	0
1	0	1
1	1	2

from this table, one may derive the following algebraic expression for S(A,B):

$$S(A,B) = A + B .$$

Since S(A,B) has three possible values, and since the minimum values of the function STATE must be at least one, the value of STATE, S(A,B,C), may be expressed as follows:

$$S(A,B,C) = S(A,B) + 3C + 1 ,$$

or

$$S(A,B,C) = A + B + 3C + 1 .$$

We may now define all possible transitions for the model by the following table:

<u>Condition</u>	<u>Event</u>	<u>Change</u>	<u>Rate</u>
$0 \leq C \leq N$ $A = 1$ $B = 0$	B ↑	1	2μ
$1 \leq C \leq N$ $B = 1$	C ↓ B ↓	-4	2μ
$C = 0$ $B = 1$	A ↓ B ↓	-2	2μ
$C = 0$ $A = 0$	A ↑	1	λ
$0 \leq C \leq N-1$ $A = 1$	C ↑	3	λ

From this table, one can infer that many of the transitions are repeated for the various values of the queue length C. Consequently, C should be chosen to be the first index for the probability array, and DELTA should be chosen to be the number of values of S(A,B), or 3. MAXJ is then 3.

By applying the above table of possible transitions and the program form of Section 4, one may obtain the following discipline program for the model:

```

EXTERNAL FUNCTION
ENTRY TO DISCPL.
READ AND PRINT DATA
MAXJ = 3
DELTA = 3
SETUP.(STATE.)
THROUGH S, FOR A = 0, 1, A. G. 1
THROUGH S, FOR B = 0, 1, B. G. A.
THROUGH S, FOR C = 0, 1, C. G. N. OR. C. G. 1
WHENEVER C. G. 0. AND. A. E. 0, TRANSFER TO S
WHENEVER C. E. 0. AND. A. E. 1. AND. B. E. 0, QUAD.(2.*MU,1,N)
WHENEVER C. E. 1. AND. B. E. 1, QUAD.(2.*MU,-4,N-1)
WHENEVER C. E. 0. AND. B. E. 1, QUAD.(2.*MU,-2,0)
WHENEVER C. E. 0. AND. A. E. 0, QUAD.(LAMBDA,1,0)
WHENEVER C. E. 0. AND. A. E. 1, QUAD.(LAMBDA,3,N-1)
S CONTINUE
DIAG.
FUNCTION RETURN
FLOATING POINT EPSI, MU, LAMBDA
NORMAL MODE IS INTEGER
PROGRAM COMMON DUMMY (8004), PCH, PRSW, MXITER, EPSI,
1DELTA, INIT, RUN, TSW, NAME(1),MAXJ
VECTOR VALUES NAME = $EXAMPLES$
INTERNAL FUNCTION STATE.(X) = A + B + 3* C + 1
END OF FUNCTION

```

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PROGRAM LENGTH (OCTAL) 00352 PROGRAM COMMON LENGTH 17520
 INITIAL TRANSFER VECTOR 00000 TRANSFER VECTOR LENGTH 00002
 LOWEST ERASABLE DEFINED 77626 ID QUAD 000

PROGRAM ENTRY POINT LOCATIONS AND NAMES
 00002 SETUP
 00031 QUAD
 00105 DIAG

EXTERNAL ROUTINES CALLED BY THIS PROGRAM
 00000 334751314563 00 .PRINT
 00001 255151465160 00 ERROR

PROGRAM LISTING (NUMBERS IN OCTAL)

	ENTRY	SETUP	QUAD	DIAG
00002	0634 00 1 00025 010	SETUP	XRA	XRA,1
00003	0534 00 1 77631 010	LXA	LXA	Z,1
00004	0500 00 1 00250 010	CLA	CLA	ADR,1
00005	0621 00 0 00023 010	STA	STA	VA
00006	0621 00 0 00075 010	STA	STA	VB
00007	0621 00 0 00076 010	STA	STA	VC
00010	0621 00 0 00122 010	STA	STA	COND
00011	0621 00 0 00131 010	STA	STA	VD
00012	0621 00 0 00132 010	STA	STA	VE
00013	0621 00 0 00140 010	STA	STA	HIGH
00014	0621 00 0 00142 010	STA	STA	VF
00015	0400 00 0 00351 010	ADD	ADD	=1
00016	0621 00 0 00160 010	STA	STA	VG
00017	0774 00 1 00001 00	AXT	AXT	1,1
00020	0634 00 1 13721 00	SXA	SXA	TRANS,1
00021	0634 00 1 00251 010	SXA	SXA	SW,1
00022	0774 00 1 11610 00	AXT	AXT	5000,1
00023	0600 00 1 00000 00	VA	VA	**1
00024	2 00001 1 00023 010	TIX	TIX	VA,1,1
00025	0774 00 1 00000 00	XRA	AXT	**1
00026	0500 60 4 00001 00	CLA*	CLA*	1,4
00027	0621 00 0 00044 010	STA	STA	FN
00030	0020 00 4 00001 00	TRA	TRA	1,4
00031	0500 00 0 00251 010	QUAD	CLA	SW
00032	0402 00 0 00351 010	SUB	SUB	=1
00033	-0100 00 0 00171 010	TNZ	TNZ	ERSETA
00034	0634 00 4 00045 010	SXA	SXA	XRB,4
00035	0634 00 2 00102 010	SXA	SXA	XRC,2
00036	0634 00 1 00103 010	SXA	SXA	XRD,1
00037	0534 00 2 13721 00	LXA	LXA	TRANS,2
00040	3 03720 2 00201 010	TXH	TXH	ERC,2,2000
00041	0500 60 4 00001 00	CLA*	CLA*	1,4
00042	0601 00 2 13721 00	STO	STO	TRANS,2
00043	0634 00 2 00046 010	SXA	SXA	XRB,1,2
00044	0074 00 4 00000 00	FN	TSX	**4
00045	0774 00 4 00000 00	XRB	AXT	**4
00046	0774 00 2 00000 00	AXT	AXT	**2
00047	0340 00 0 00350 010	CAS	CAS	=5000
00050	0020 00 0 00212 010	TRA	TRA	ERD
00051	0020 00 0 00053 010	TRA	TRA	**2
00052	-0120 00 0 00220 010	TMI	TMI	ERE

ID QUAD 002

ID QUAD 003

00053	0601	00	2	17642	00	STO	FIRST,2
00054	0400	60	4	00002	00	ADD*	2,4
00055	0601	00	2	23563	00	STO	DEST,2
00056	0402	60	4	00002	00	SUB*	2,4
00057	0601	00	0	77777	010	STO	TMP
00060	0560	60	4	00003	00	LDQ*	3,4
00061	0200	00	0	27511	00	MPY	DELTA
00062	0131	00	0	00000	00	XCA	
00063	0400	00	0	77777	010	ADD	TMP
00064	0601	00	2	27504	00	STO	LAST,2
00065	0734	00	1	00000	00	PAX	,1
00066	-0634	00	1	00073	010	SXD	THRU,1
00067	0534	00	1	27511	00	LXA	DELTA,1
00070	-0634	00	1	00077	010	SXD	INC,1
00071	0500	00	2	17642	00	CLA	FIRST,2
00072	0734	00	1	00000	00	PAX	,1
00073	3	00000	1	00100	010	THRU	BQ,1,**
00074	0502	00	2	13721	00	CLS	TRANS,2
00075	0300	00	1	00000	00	VB	FAD **,*1
00076	0601	00	1	00000	00	VC	STO **,*1
00077	1	00000	1	00073	010	INC	THRU,1,**
00100	1	00001	2	00101	010	BQ	TXI **1,2,1
00101	0634	00	2	13721	00	SXA	TRANS,2
00102	0774	00	2	00000	00	XRC	AXT **,*2
00103	0774	00	1	00000	00	XRD	AXT **,*1
00104	0020	00	4	00001	00	TRA	1,4
00105	0500	00	0	00251	010	DIAG	CLA SW
00106	0402	00	0	00351	010	SUR	=1
00107	-0100	00	0	00175	010	TNZ	ERSEIB
00110	0634	00	1	00164	010	SXA	XRE,1
00111	0634	00	2	00165	010	SXA	XRF,2
00112	0634	00	4	00166	010	SXA	XRG,4
00113	0774	00	1	00001	00	AXT	1,1
00114	0534	00	4	27511	00	LXA	DELTA,4
00115	-0634	00	4	00127	010	SXD	BKD,4
00116	-0634	00	4	00156	010	SXD	BKZ,4
00117	0534	00	4	13721	00	LXA	TRANS,4
00120	3	03720	4	00201	010	TXH	ERC,4,2000
00121	0622	00	0	00130	010	STD	BKD+1
00122	-0520	00	1	00000	00	COND	**,*1
00123	1	00001	1	00122	010	TXI	*-1,1,1
00124	0754	00	1	00000	00	PXA	,1
00125	0734	00	2	00000	00	PAX	,2
00126	3	11610	1	00164	010	TXH	XRE,1,5000
00127	1	00000	2	00130	010	TXI	**1,2,**
00130	3	00000	2	00140	010	TXH	HIGH,2,**
00131	0500	00	1	00000	00	VD	CLA **,*1
00132	0302	00	2	00000	00	VE	FSB **,*2
00133	0760	00	0	00003	00	SSP	
00134	0340	00	0	27510	00	CAS	EPSI
00135	0020	00	0	00140	010	TRA	HIGH
00136	0020	00	0	00127	010	TRA	RKD
00137	0020	00	0	00127	010	TRA	BKD
00140	0500	00	1	00000	00	HIGH	CLA **,*1
00141	0601	00	4	13721	00	STO	TRANS,4
00142	0600	00	1	00000	00	VF	STZ **,*1
00143	0754	00	1	00000	00	PXA	,1
00144	0601	00	4	17642	00	STO	FIRST,4

ID QUAD 004

ID QUAD 005

ID QUAD 006

00145	0601.00	4	23563	00	STO	DEST,4
00146	0754.00	2	00000	00	PXA	,2
00147	0402.00	0	27511	00	SUB	DELTA
00150	0601.00	4	27504	00	STO	LAST,4
00151	1.00001	4	00152	010	TXI	**1,4,1
00152	3.03720	4	00201	010	TXH	ERC,4,2000
00153	1.00001	1	00154	010	TXI	**1,1,1
00154	-0634.00	2	00157	010	SXD	STPZ,2
00155	0634.00	1	00162	010	SXA	ENDZ,1
00156	1.00000	1	00157	010	BKZ	TXI STPZ,1,**
00157	3.00000	1	00162	010	STPZ	TXH ENDZ,1,**
00160	0600.00	1	00000	00	VG	STZ **,1
00161	0020.00	0	00156	010	TRA	BKZ
00162	0774.00	1	00000	00	ENDZ	AXT **,1
00163	0020.00	0	00122	010	TRA	COND
00164	0774.00	1	00000	00	XRE	AXT **,1
00165	0774.00	2	00000	00	XRE	AXT **,2
00166	0774.00	4	00000	00	XRG	AXT **,4
00167	0600.00	0	00251	010	STZ	SW
00170	0020.00	4	00001	00	TRA	1,4
00171	0074.00	4	00000	010	ERSEIA	PRINT FA,0
00172	-1.00000	0	00252	010		
00173	-1.00000	0	00000	00		
00174	0074.00	4	00001	010	CALL	ERROR
00175	0074.00	4	00000	010	ERSEIB	PRINT FB,0
00176	-1.00000	0	00261	010		
00177	-1.00000	0	00000	00		
00200	0074.00	4	00001	010	CALL	ERROR
00201	0774.00	1	03720	00	ERC	AXT 2000,1
00202	0074.00	4	00000	010	PRINT	FC
00203	-1.00000	0	00270	010		
00204	-1.00000	1	13721	00	IOP	TRANS,1
00205	-1.00000	1	17642	00	IOP	FIRST,1
00206	-1.00000	1	23563	00	IOP	DEST,1
00207	-1.00000	1	27504	00	IOP	LAST,1
00210	-1.00000	0	00000	00	ENDIO	
00211	0074.00	4	00001	010	CALL	ERROR
00212	0601.00	0	77777	010	STO	TMP
00213	0074.00	4	00000	010	PRINT	FD,IMP,0
00214	-1.00000	0	00306	010		
00215	-1.00000	0	77777	010		
00216	-1.00000	0	00000	00		
00217	0020.00	0	00225	010	TRA	TBL
00220	0601.00	0	77777	010	STO	TMP
00221	0074.00	4	00000	010	PRINT	FE,IMP,0
00222	-1.00000	0	00315	010		
00223	-1.00000	0	77777	010		
00224	-1.00000	0	00000	00		
00225	0534.00	1	13721	00	TBL	LXA TRANS,1
00226	-3.00001	1	00246	010	TXL	END+1,1,1
00227	1.77777	1	00230	010	TXI	**1,1,-1
00230	-0634.00	1	00235	010	SXD	MAX,1
00231	0074.00	4	00000	010	PRINT	FIBL
00232	-1.00000	0	00324	010		
00233	0600.00	0	77777	010	STZ	TMP
00234	0774.00	1	00001	00	AXT	1,1
00235	3.00000	1	00245	010	MAX	END,1,**
00236	-1.00000	1	13721	00	IOP	TRANS,1

ID QUAD 007

ID QUAD 008


```

R FN 00044 00027,00044
R FTBL 00324 00232,00324
R HIGH 00140 00013,00130,00135,00140
R INC 00077 00070,00077
R .PRINT 00000 00171,00175,00202,00213,00221,00231
R LAST 27504 00064,00150,00207,00241,00350
R MAX 00235 00230,00235,00244
R MAKJ 27517 00350
R QVAD 00031 00000,00002,00031
R SETUP 00002 00000,00002,00002
R STPZ 00157 00154,00156,00157
R SW 00251 00021,00031,00105,00167,00251
R TBL 00225 00217,00225
R THRU 00073 00066,00073,00077
R TMP 77777 00057,00063,00212,00215,00220,00223,00242,00243,00350
R TRANS 13721 00020,00037,00042,00074,00101,00117,00141,00204,00225,00236,00350
R V 77626 00247,00250,00350
R VA 00023 00005,00023,00024
R VB 00075 00006,00075
R VO 00076 00007,00076
R VD 00131 00011,00131
R VE 00132 00012,00132
R VF 00142 00014,00142
R VB 00160 00016,00160
R KRA 00025 00002,00025
R XRB 00045 00034,00043,00045
R XRC 00102 00035,00102
R XRD 00103 00036,00103
R XRE 00164 00110,00126,00164
R XRF 00165 00111,00165
R XRG 00166 00112,00166
R Z 77631 00003,00350

```

REFERENCES TO PROGRAM LITERALS

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RGC REFERENCES
00350 00047
00251 00015,00032,00106

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FLAGGED LINES FATAL 0 NON-FATAL 0 **** DEFINITIONS/ SYMBOL 55 OPERATION 0 MACRO 0 RMT 0
MACRO CALLS CORRECT 0 INCORRECT 0 **** CARD COUNTS INPUT 177 OUTPUT 13

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