# Latency Measurements of V2X Communication by Uday Kumar Bejawada

A thesis submitted in partial fulfillment of the requirements for the degree of Master of Science in Engineering (Automotive Systems Engineering) in the University of Michigan-Dearborn 2017

#### **Master's Thesis Committee:**

Professor Weidong Xiang, Chair Associate Professor Jinhua Guo Associate Professor Hafiz Malik

#### Acknowledgements

I take immense pleasure to express my profound gratitude and deep regards to my academic supervisor 'Prof. Weidong Xiang' for his exemplary guidance, monitoring and constant encouragement throughout the course of this thesis. I would also like to thank him for offering me thesis in my field of interest and equipping with all basic needs, which helped my work to be interesting and productive.

I would also like to thank 'Prof. Jiabo Zhang' for giving me an opportunity to take his guidance. I am obliged to colleagues of "laboratory" for their friendly sense of humour, creating energetic and enthusiastic working environment.

Additional thanks to the "University of Michigan Dearborn" who have identified ways to help students succeed.

I am grateful for the cooperation of 'Dr. Pankaj Mallick' and 'Ms. Sherry Boyd' for their encouragement and in departmental aspects during the period of my thesis.

Finally, yet importantly, I would like to express my heartfelt thanks to my beloved parents and brother for their blessings and support, my friends for their help and wishes for the successful completion of this project.

# **Table of Contents**

Acknowledgements	ii
List of Tables	V
List of Figures	vi
List of Abbreviations	vii
Abstract	. viii
CHAPTER 1: INTRODUCTION	1
1.1 DESCRIPTION	2
1.2 SECURITY AND SAFETY IN V2X COMMUNICATION	2
1.3 MINIMUM REQUIREMENTS FOR V2X COMMUNICATION	3
1.3.1 CHANNEL REQUIREMENTS	3
1.3.2 Transmit Power Accuracy	3
1.3.3 PATH HISTORY REFERENCE DESIGN	3
CHAPTER 2: DATA INTERPRETATION	5
2.1 NMEA FROM GARMIN GLOBAL POSITIONING SYSTEM SENSOR	5
2.2 DESIRED INFORMATION OF NMEA (National Marine Electronics Association)	6
2.3 ALGORITHM TO INTERPRET THE NMEA DATA	8
2.3.1 SYNCHRONOUS INPUT/ OUTPUT	8
2.3.2 ASYNCHRONOUS INPUT/ OUTPUT	8
2.4 VEHICLE CAN DATA VIA OBD DONGLE	14
2.5 ALGORITHM TO INTERPRET THE CAN PACKETS	15
CHAPTER 3: NETWORK CONFIGURATION	17

3.1	LOCAL NETWORK	17
3.2	REMOTE NETWORK	19
3.3	ALGORITHM FOR NETWORK INTERFACE	20
CHAP	TER 4: LATENCY MEASUREMENTS	32
4.1	SCENARIOS IMPLEMENTED	32
4.2	SCENARIOS OVERVIEW	32
4.3	Conclusion.	33
REFER	RENCES	34

# **List of Tables**

TABLE 1: GARMIN 10X TECHNICAL SPECIFICATIONS	. 5
TABLE 2: GGPGA DEFINED BY NMEA 0183	. 6

# **List of Figures**

FIGURE 1: RESEARCH WORK APPROACH	1
FIGURE 2: WAAS SYSTEM [5]	6
FIGURE 3: NMEA SENTENCES	7
FIGURE 4: VEHICLE CAN BUS INFORMATION	14
FIGURE 5: DESIRED INFORMATION WITH TIME STAMP	16
FIGURE 6: LOCAL NETWORK CONFIGURATION	17
Figure 7: Algorithm Flowchart	18
FIGURE 8: REMOTE NETWORK CONFIGURATION	19

# **List of Abbreviations**

B	L
BSM (Basic Safety Message)2	Long Term Evolution)
C	P
CAN (Controller Area Network)	Path History (PH)3
D	Path Prediction (PP)3
D2D (Device to Device)7	S
DSRC (Dedicated Short Range	SA (selective availability)5
Communications) 7	SAE (Society of Automotive Engineers)
G	J27352
GNSS (Global Navigation Satellite System)	T
2	TCP (Transmission Control Protocol) 1
GPS (Global Positioning System)7	Time of Flight (TOF)
H	U
Host Vehicle (HV)	UDP (User Datagram Protocol)1

#### **Abstract**

V2X communication plays one of the important roles for future autonomous vehicle. Current V2X communication is more directed towards DSRC (Dedicated Short Range Communications). What if the DSRC law is not enforced and some other technology such as D2D (Device to Device) based on LTE (Long Term Evolution) will provide good results and become possible alternative.

The scope of this project is to have a sniffer that has the capability to process transmit and receive the data and could able to interface all of the following devices DSRC, LTE and D2D (LTE based communication without involvement of Base Station) and any other possible medium of communication for V2X communication. And also, to have a database that has the capability to help in the future of autonomous vehicles. This thesis explains the research study, done for one of the performance evaluation of V2X communication. It can be further continued to develop a device and a server with more functionalities for V2X communication.

In V2V communication network, performance evaluation is one of the key factors that need to be considered. To address the evaluation criteria, I have analyzed the elapsed time in reading the GPS (Global Positioning System) sensor and CAN (Controller Area Network) bus information from the vehicle before transmitting the data to the Network (IEEE802.11p or LTE) and also after the transmission. I have developed a c code, which can read the data from the GPS sensor and the CAN bus, and add a time stamp to the real time data to analyze the elapsed time.

#### **CHAPTER 1: INTRODUCTION**

In V2X communication networks, the elapsed time in receiving the CAN bus information from the CAN network of the vehicle should be taken into consideration, prior to the transmission of data to the other vehicle. Once the data is transmitted either by IEEE802.11p or LTE, Time of Flight (TOF) comes into picture. To estimate the overall performance of V2X communication network these two time factors need to be considered. A mobile service can be used to provide the live updates of the traffic situation on the road to the users 1. To get the live updates, the users need to transmit and receive the data from the remote server when they are driving on the road. The objective of this thesis is to determine the performance of IEEE802.11p and LTE as vehicular networks in real time. The overall approach followed for this research work is shown in Figure 1. Two aspects are contributed to this thesis to estimate the real time performance of vehicular networks.

- 1) Collecting the vehicle information in real time and transmitting the information to a remote TCP (Transmission Control Protocol) server or UDP (User Datagram Protocol) server.
- 2) Calculating the time elapsed to receive CAN data from the CAN network and transmitting to the server at different traffic situations and driving speeds. The elapsed time in transmitting the user data to the server is an essential challenge to estimate the overall performance of the vehicular networks.



Figure 1: Research work approach

#### 1.1 DESCRIPTION

The vehicular networks architecture consists of three layers, access network layer, data aggregation layer and application layer. Vehicles are connected to cellular base station and road side unit in network layer. They act as mobile sensors to collect the traffic information in this layer [2]. The BSM (Basic Safety Message), which is defined in SAE (Society of Automotive Engineers) J2735, is the message used for V2V safety communications. Each vehicle broadcasts BSMs to provide neighboring vehicles with trajectory and status information. The BSM consists of all data elements listed in Part I and selected data elements and data frames listed in Part II of the SAE J2735 standard. Part I contains the vehicle position, speed, heading, acceleration, transmission, steeringwheel angle, brake, and vehicle-size information. Many V2X safety applications need projection information with lane-level granularity. The positioning system includes a GNSS (Global Navigation Satellite System) receiver to enable the System to determine its own position and accurate time. The System also maintains its own path history and calculates its path prediction. Each vehicle broadcasts its time-tagged position (10Hz), heading, speed, and acceleration, its path history and path prediction in the BSM. Hence as a part of this study, GPS information and some CAN data packets from the BSM part1 have been taken into consideration to transmit the data via LTE.

#### 1.2 SECURITY AND SAFETY IN V2X COMMUNICATION

According to [11], to support trust in message exchange between vehicles, BSM signing and verification are performed using a public key digital signature algorithm. Each BSM is transmitted with a signature and either a security certificate containing the public key or a certificate. The BSM with a certificate is transmitted approximately every 500ms, and other BSMs are transmitted with a certificate digest to reduce the overall message length. BSM security is compliant with IEEE 1609.2.

LTE security is compliant with 3GPP [12]. However, this thesis work doesn't discuss the security and focuses only on latency.

#### 1.3 MINIMUM REQUIREMENTS FOR V2X COMMUNICATION

V2X communication should meet the requirements mentioned in [11]. Some of the minimum requirements that need to be met by the V2X communication system are outlined in this section.

#### 1.3.1 CHANNEL REQUIREMENTS

- The System shall transmit BSMs on channel with 10 MHz channel spacing.
- The System shall transmit BSMs using an 802.11 data rate of 6 Mbps.
- After the first BSM generated after System device startup, the System shall generate each subsequent BSM within tolerance rate of (-10msec) to (+10msec) of its scheduled generation time.

#### 1.3.2 Transmit Power Accuracy

The DSRC Radio Subsystem shall meet the 802.11 transmitter requirements for 10 MHz channel spacing with QPSK and ½ rate coding.

#### 1.3.3 PATH HISTORY REFERENCE DESIGN

The Path History (PH) module for the V2V communications System uses a history of the past GNSS locations traversed by the Host Vehicle (HV) and computes an adaptable, concise PH representation of recent vehicle movement over a certain distance (200-300 meters). The PH communicated by a vehicle provides other vehicles with information needed for predicting the roadway geometry.

PH plays an important role in target vehicle classification, relative to the HV, regarding the roadway. There are different methods for design and implementation of the PH module. Three different design methods are described here, each with a slightly different approach.

Path Prediction (PP) for the V2V safety communications System utilizes dynamics information provided by the vehicle to estimate the driver's intended future path. The estimate is provided without dependence on future road geometry information obtained from outside sources (e.g., map databases, vehicle probes). In order to determine if the LTE can meet the minimum V2X communication requirement, the first step that has been considered to determine the latency is outlined in this thesis work.

PP carries out the following basic operations:

- Gathers vehicle dynamics information.
- Computes path radius using dynamics information to represent the driver's intended future path:
- Radius =  $1/\text{curvature}(\rho)$ .
- Computes confidence of the predicted path based upon the ra``te of change of the vehicle dynamics to infer transient conditions (i.e., non-steady-state conditions).

#### **CHAPTER 2: DATA INTERPRETATION**

#### 2.1 NMEA FROM GARMIN GLOBAL POSITIONING SYSTEM SENSOR

The GPS sensor used in this project is Garmin 10X, the technical specifications of which are listed in

Table 1.

Table 1: Garmin 10x technical specifications

Electrical Characteristics			
Input Voltage	4.5-5.5V		
Input Current	500mA@ 5V		
Battery	1100mAhr		
Operating temperature range	-30°C to +60°C powered from external supply		
	-10°C to +60°C powered by removable		
	rechargeable battery		
Receiver	High sensitivity SiRFstarIII GPS receiver;		
	WAAS <sup>1</sup> -enabled		
Acquisition Times	On average for a stationary receiver with a clear		
	view of the sky.		
	✓ Warm:<1 sec		
	✓ Cold:<38 sec		
	✓ Factory Reset:<45 sec		
Update Rate	Update Rate   1 record per second, continuous		
Accuracy			
GPS accuracy	Position: <10 meters, typical		
	Velocity: 0.05 meters/sec RMS		
GPS (WAAS) accuracy	Position: <5 meters, typical		
	Velocity: 0.05m/s (RMS)		
Interfaces			
Bluetooth Serial Profile	Range up to 32 feet or 10 meters		
Protocol	NMEA 0183, Version 2.0, default		

5

<sup>&</sup>lt;sup>1</sup> Wide Area Augmentation System corrects the GPS signal errors caused by ionosphere disturbances, timing and satellite orbit errors. It can give a position accuracy of 3m, 95 percent of the time. WAAS satellite coverage is available in North America.

As shown in Figure 2 the accuracy of GPS system subjects to degradation with SA (selective availability) on and improves with SA off and WAAS.

- ➤ 15m: typical GPS position accuracy without SA.
- ➤ 3-5m: typical differential GPS (DGPS) position accuracy.
- > <3m: typical WAAS position accuracy.

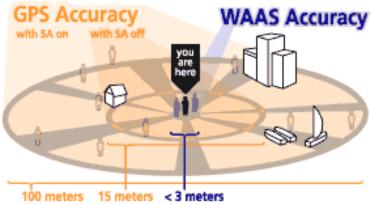


Figure 2: WAAS system [5]

### 2.2 DESIRED INFORMATION OF NMEA (National Marine Electronics Association)

As shown in Figure 3, NMEA 0183 consists of various sentences. The GGPGA consists the global positioning system fix data, which is described in Table 2.

Table 2: GGPGA Defined by NMEA 0183

Name	Example	Description
	Data	
Sentence Identifier	\$GPGGA	Global Positioning System Fix Data
Time	170834	17:08:34 Z
Latitude	4219.2152,	42d 19.2152' N or 42d 19' 12.912" N
	N	
Longitude	08314.0246,	83d 14.0246' W or 83d 14' 1.4760" W
	W	
Fix Quality:	1	Data is from a GPS fix
0=Invalid		
I = GPS fix		
2=DGPS fix		
Number of Satellites	06	6 Satellites are in view
Horizontal Dilution of	1.5	Relative accuracy of horizontal
Precision (HDOP)		position
Altitude	249.7, M	280.2 meters above mean sea level
Height of geoid above WGS84	-34.0, M	-34.0 meters
ellipsoid		

Time since last DGPS update	blank	No last update
DGPS reference station id	blank	No station id
Checksum	*69	Used by program to check for transmission errors

```
$GPGGA,212106.000,4219.2154,N,08314.0246,W,1,07,1.4,249.7,M,-34.0,M,,0000*6A
$GPGSA,A,3,11,07,01,30,09,27,08,,,,,,1.9,1.4,1.3*3C
$GPGSV,3,1,10,07,72,177,39,30,68,291,19,11,56,130,44,08,46,050,14*77
$GPGSV,3,2,10,28,37,290,21,01,34,141,41,13,16,315,25,17,14,227,22*79
$GPGSV,3,3,10,27,14,049,17,09,06,199,15*79
$GPRMC,212106.000,A,4219.2154,N,08314.0246,W,0.27,339.16,290416,,,A*7B
$GPGGA,212107.000,4219.2152,N,08314.0254,W,1,06,2.2,249.4,M,-34.0,M,,0000*69
$GPGSA,A,3,11,07,01,30,27,08,,,,,,3.8,2.2,3.1*33
$GPGSV,3,1,10,07,72,177,39,30,68,291,19,11,56,130,44,08,46,050,14*77
$GPGSV,3,2,10,28,37,290,21,01,34,141,41,13,16,315,25,17,14,227,22*79
$GPGSV,3,3,10,27,14,049,18,09,06,199,15*76
$GPRMC,212107.000,A,4219.2152,N,08314.0254,W,0.28,282.63,290416,,,A*73
$GPGGA,212108.000,4219.2153,N,08314.0255,W,1,06,1.4,249.4,M,-34.0,M,,0000*63
$GPGSA,A,3,11,07,01,30,09,27,,,,,,1.9,1.4,1.3*34
$GPGSV,3,1,10,07,72,177,39,30,68,291,19,11,56,130,44,08,46,050,14*77
$GPGSV,3,2,10,28,37,290,21,01,34,141,41,13,16,315,25,17,14,227,22*79
$GPGSV,3,3,10,27,14,049,18,09,06,199,15*76
$GPRMC,212108.000,A,4219.2153,N,08314.0255,W,0.53,27.70,290416,,,A*4F
$PGRMT,GPS 10x Software v2.20 BT v2.60,,,,,,*7C
$GPGGA,212109.000,4219.2155,N,08314.0253,W,1,07,1.4,249.2,M,-34.0,M,,0000*65
$GPGSA,A,3,11,07,01,30,09,27,08,,,,,,1.9,1.4,1.3*3C
$GPGSV,3,1,10,07,72,177,39,30,68,291,20,11,56,130,44,08,46,050,14*7D
$GPGSV,3,2,10,28,37,290,21,01,34,141,41,13,16,315,25,17,14,227,22*79
$GPGSV,3,3,10,27,14,049,18,09,06,199,14*77
$GPRMC,212109.000,A,4219.2155,N,08314.0253,W,0.95,19.90,290416,,,A*47
$GPGGA,212110.000,4219.2157,N,08314.0253,W,1,06,1.4,249.1,M,-34.0,M,,0000*6D
$GPGSA,A,3,11,07,01,30,09,27,,,,,,1.9,1.4,1.3*34
$GPGSV,3,1,10,07,72,177,39,30,68,291,20,11,56,130,44,08,46,050,14*7D
$GPGSV,3,2,10,28,37,290,21,01,34,141,41,13,16,315,25,17,14,227,22*79
$GPGSV,3,3,10,27,14,049,18,09,06,199,14*77
$GPRMC,212110.000,A,4219.2157,N,08314.0253,W,0.14,10.44,290416,,,A*44
$GPGGA,212111.000,4219.2156,N,08314.0253,W,1,06,1.4,249.3,M,-34.0,M,,0000*6F
$GPGSA,A,3,11,07,01,30,09,27,,,,,,1.9,1.4,1.3*34
$GPGSV,3,1,10,07,72,177,39,30,68,291,20,11,56,130,44,08,46,050,14*7D
$GPGSV,3,2,10,28,37,290,20,01,34,141,41,13,16,315,25,17,14,227,23*79
$GPGSV,3,3,10,27,14,049,18,09,06,199,14*77
```

Figure 3: NMEA sentences

A C program is developed to interpret the latitude and longitude data from GPGGA received. The GPGGA has the following format.

\$GPGGA,hhmmss.ss,llll.ll,a,yyyyy.yy,a,x,xx,x.x,x.x,M,x.x,M,x.x,xxxx\*hh The program is to,

- 1) Read the complete NMEA sentences sent from the GARMIN GPS 10X sensor via Bluetooth.
- 2) Capture GPGGA string and write to a data buffer.
- 3) Extract desired information including timestamp and latitude and longitude.
- 4) Loop for next data.

#### 2.3 ALGORITHM TO INTERPRET THE NMEA DATA

For the serial communication between the GPS module and the windows PC, the Bluetooth interface is established. In windows reading from and writing to communication ports is like the file input and output. Functions for serial I/O uses the same functions as file I/O. Reading and writing to communication ports in windows can be done either synchronously or asynchronously.

#### 2.3.1 SYNCHRONOUS INPUT/ OUTPUT

In synchronous I/O, the calling thread is blocked when an operation is taking place. Once the operation is complete, the function will return and the thread can continue with the next operation. This is useful in multithreaded applications, because when one thread is blocked on an I/O operation, other threads can still perform work. However, the application should serialize the ports correctly. The drawback in this operation is that if one thread is blocked waiting for the I/O operation, all other threads that call a communications API will be blocked until the original operation completes.

#### 2.3.2 ASYNCHRONOUS INPUT/ OUTPUT

In Asynchronous I/O, a port open for overlapped operations allows multiple threads to perform the I/O operations at the same time, thus allowing to do other work while the operations are incomplete. The overlapped operations allow a single thread to issue multiple requests of different type and do work in the background while the operations are incomplete, thus allowing more flexibility and efficiency.

Asynchronous operation is not a good for portability in comparison with synchronous operation, because it is not supported by most of the operating systems. Multithreading is supported by most of the operating systems, therefore synchronous operation might be the best choice for portability. Between issuing requests and processing the results, some synchronization must take place, in both single-threaded and multithreaded applications. One thread should be blocked until the result of an operation is available. Asynchronous I/O allows a thread to do some work between the time of the request and its completion, thus providing the advantage for better responsiveness.

According to [13] for opening and closing a communications resource handle and for performing read and write operations, "CreateFile, closehandle, readfile, and writefile" will provide the basic interface. As the first step to open a handle to a communication resource the process uses the create

file function. For example, specifying COM1 opens a handle to a serial port. The handle returned by 'createfile' to identify the resource in any of the functions that access the resource can be used by any thread of the process. The following attributes need to be specified while calling a create file function, type of read/ write access, whether the handle can be inherited by child process and whether the handle can be used in asynchronous I/O operations. More information about 'createfile' function to open or create an I/O device can be found at [14].

The code developed for reading the GPS NMEA sentences works as follows:

```
Step1:
```

```
// Scans for an available COM port in a specified maximum and minimum range
// The maximum and minimum range must be defined by the user
  for (n = scan max; n \ge scan min; --n)
  {
// Once the COM port is identified, the COM should be opened to read or write the data
    hSerial = CreateFile(device name, GENERIC READ|GENERIC WRITE, 0, 0,
                 OPEN EXISTING, FILE ATTRIBUTE NORMAL, 0);
    if (hSerial!=INVALID HANDLE VALUE)
      // Assign the device number and stop scanning the ports
       dev number = n;
       break; // stop scanning COM ports
    // if it results in invalid handle value display the error code
  }
Step2:
// Once the serial port is successfully opened, set the device parameters
// Baudrate (default 9600), 1 start bit, 1 stop bit, no parity
  dcbSerialParams.DCBlength = sizeof(dcbSerialParams);
  if (GetCommState(hSerial, &dcbSerialParams) == 0)
    exit message("Error getting device state", 1);
  dcbSerialParams.BaudRate = baudrate:
```

```
dcbSerialParams.ByteSize = 8;
dcbSerialParams.StopBits = ONESTOPBIT;
dcbSerialParams.Parity = NOPARITY;
if(SetCommState(hSerial, &dcbSerialParams) == 0)
  exit message("Error setting device parameters", 1);
```

### Step3:

Read the incoming data as characters, to avoid running into infinite loop when no data is received set COM port timeout settings.

```
timeouts.ReadIntervalTimeout = 50;
 timeouts.ReadTotalTimeoutConstant = 50;
 timeouts.ReadTotalTimeoutMultiplier = 10;
 timeouts.WriteTotalTimeoutConstant = 50;
 timeouts.WriteTotalTimeoutMultiplier = 10;
 if(SetCommTimeouts(hSerial, &timeouts) == 0)
    exit message("Error setting timeouts", 1);
Step4:
```

Read the data to a specified buffer of defined length char message buffer[MESSAGE LENGTH]; DWORD bytes read;

#### Step5:

Read the messages in the correct format, by identifying the delimiters that are used to separate the NMEA sentences.

```
while(1)
  // Read next character
  ReadFile(hSerial, &c, 1, &bytes read, NULL);
  if (bytes read != 1) continue;
  if (state == 1)
  {
    // State 1: Waiting for '<' character to start message
    n = 0;
```

```
strcpy(message buffer, "");
          if (c == '<') state = 2;
       else if (state == 2)
          // State 2: Reading id one digit at a time
          if (c \ge 0') \&\& c \le 9' n = 10*n + (c-48);
          else if (c == '@') state = 3;
          else state = 1;
       }
       else if (state == 3)
          // Got a '@' character, so now reading actual message
          if (n != id) state = 1;
          else if (c!='>') strncat(message buffer, &c, 1);
          else if (c == '>') state = 4;
          if (strlen(message_buffer) >= MESSAGE_LENGTH-1) state = 4;
       else if (state == 4)
        {
          // State 4: Display complete message
          fprintf(stdout, "%s\n", message buffer);
          state = 1;
       else state = 1;
  return 0;
Step6:
```

Close the serial port in case of error or when the user presses a specific key defined in the code to terminate program.

#### Step7:

To interpret a specific NMEA sentence out of the buffer data, a string comparison function is required. For this a structure containing the GPS NMEA string info can be created.

```
struct {
  int
       timeZone;
  // flag to be used, when a valid sentence begins
  bool flagRead;
 // valid GPS fix and data available
 bool flagDataReady;
 // hold parsed words for one given NMEA sentence
  char words[20][15];
// hold the received checksum for one given NMEA sentence
char szChecksum[15];
// will be set to true for characters between $ and * only
  bool flagComputedCks;
// used to compute checksum and indicate valid checksum interval (between $ and * in a given
sentence)
// compute numeric checksum for a given sentence
     checksum;
int
// after getting * start cuttings the received checksum
 bool flagReceivedCks;
// Parse the received checksum
     index received checksum;
// Indexing the word
     wordIdx;
  int
       prevIdx;
  int
  int
       nowIdx;
 // globals to store parser results
  bool positionFixIndicator; //GPGGA
```

```
bool dataValid;
                          //GPRMC
                    // GPRMC and GPGGA
  float longitude;
  float latitude;
                    // GPRMC and GPGGA
  unsigned char UTCHour, UTCMin, UTCSec,
                                                       // GPRMC and GPGGA
                     UTCDay, UTCMonth, UTCYear;
                                                       // GPRMC
  int
       satellitesUsed:// GPGGA
  float speed;
                    // GPRMC
}gps;
Step8:
Comparing the string in the buffer to read the data as words in the GPGGA sentence by parsing the
delimiters.
  if (strcmp(gps.words[0], "\$GPGGA") == 0) {
    /*
    // Check GPS Fix: 0=no fix, 1=GPS fix, 2=Dif. GPS fix
    if (gps.words[6][0] == '0') {
       gps.positionFixIndicator = false;
      // clear data
//
        gps.res fLatitude = 0;
//
        gps.res fLongitude = 0;
       gps.flagDataReady = false;
Reading the time information from the GPS
    gps.UTCHour = charToInt(gps.words[1][0]) * 10 + charToInt(gps.words[1][1]);
    gps.UTCHour = gps.UTCHour-4; // Michigan time zone
    gps.UTCMin = charToInt(gps.words[1][2]) * 10 + charToInt(gps.words[1][3]);
    gps.UTCSec = charToInt(gps.words[1][4]) * 10 + charToInt(gps.words[1][5]);
Reading the latitude and longitude information
    gps.latitude = strtof(gps.words[2], NULL);
    gps.longitude = strtof(gps.words[4], NULL);
Reading the number of satellites
```

gps.satellitesUsed = (int)strtof(gps.words[7], NULL);

#### Step9:

Create a data file and write the data to the file. Continue the loop till the user terminates the program or if an error is encountered, and close the file.

#### 2.4 VEHICLE CAN DATA VIA OBD DONGLE

Vehicle CAN data is read from an Open-XC dongle through Bluetooth from Ford. Moreover, Ford provides open sources and data focused API [6]. It works for Ford vehicles and others if you can get the corresponding firmware from their manufacturers. Once the dongle is paired, it starts to transmit the selected CAN packets in the JSON format, shown in Figure 4.

status", "value" false) ("name": transmission gear position", "value". "first") ("name": "anie". "Allowe": neutral] ("name". "accelerator, pedal\_position", value": allow ["indemosited] ("name". "indemosited] ("name". "indemosited

Figure 4: Vehicle CAN bus information

A C program is developed to interpret CAN packet. Its works to,

1) Read JSON strings sent from the OBD dongle via Bluetooth.

- 2) Capture and process interest vehicle data, such as engine speed, vehicle speed, torque, steer angle and brake pedal position as shown in Figure 5.
- 3) Once it encounters the string it writes the data to a buffer.
- 4) In the next step, it divides the string and reads its value.
- 5) Resets the buffer and the process continues.

In each cycle of this process a time stamp is added to calculate the elapsed time for the update rate. Figure 5 shows that the vehicle speed sensor gives output every 100ms and engine speed sensor every 100ms and latitude and longitude information updates approximately every 1sec.

#### 2.5 ALGORITHM TO INTERPRET THE CAN PACKETS

The algorithm to interpret the CAN packets is like the one discussed in section 2.3, except for few changes. Instead of the NMEA sentences the incoming data will be in JSON (Java Script Object Notation) format. JSON is built on two structures, a collection of name value pairs and an ordered list of values. In JSON an object is an unordered set of name/value pairs, that begins with {(left brace) and ends with} (right brace). Each name is followed by: (colon) and the name/value pairs are separated by, (comma). An array is an ordered collection of values, begins with [ (left bracket) and ends with] (right bracket). Values are separated by, (comma)

.

A value can be a string in double quotes, or a number, or true or false or null, or an object or an array. A string is a sequence of zero or more Unicode characters, wrapped in double quotes, using backslash escapes. A character is represented as a single character string. For the number, octal or hexadecimal formats are not used. Therefore, a string comparison to interpret the data with specified delimiters will be used to compared and read the desired information from the NMEA sentence. The string comparison functionality will work on pointer reference.

```
1470617220 s:283 ms, Engine_Speed =1486.000000, elapsed = 85 ms
1470617220:298, Latitude = 42.276459, elapsed = 1001 ms
1470617220:298,Longitude = -83.204651,elapsed = 1001 ms
1470617220:314, Vehicle Speed = 56.899998, elapsed = 100 ms
1470617220 s:414 ms, Engine Speed =1498.000000, elapsed = 131 ms
1470617220:414, Vehicle Speed = 56.789997, elapsed = 100 ms
1470617220 s:483 ms, Engine_Speed =1508.000000, elapsed = 69 ms
1470617220:548, Vehicle Speed = 56.739998, elapsed = 134 ms
1470617220 s:583 ms, Engine Speed =1478.000000, elapsed = 100 ms
1470617220:615, Vehicle Speed = 56.590000, elapsed = 67 ms
1470617220 s:683 ms,Engine_Speed =1476.000000,elapsed = 100 ms
1470617220:715, Vehicle Speed = 56.570000, elapsed = 100 ms
1470617220 s:799 ms, Engine Speed =1468.000000, elapsed = 116 ms
1470617220:815, Vehicle_Speed = 56.449997, elapsed = 100 ms
1470617220 s:884 ms, Engine Speed =1478.000000, elapsed = 85 ms
1470617220:915, Vehicle Speed = 56.340000, elapsed = 100 ms
1470617220 s:984 ms,Engine_Speed =1486.000000,elapsed = 100 ms
1470617221:15, Vehicle Speed = 56.169998, elapsed = 100 ms
1470617221 s:84 ms, Engine Speed =1500.000000, elapsed = 100 ms
1470617221:115, Vehicle Speed = 56.230000, elapsed = 100 ms
1470617221 s:184 ms, Engine_Speed =1506.000000, elapsed = 100 ms
1470617221:216, Vehicle Speed = 56.099998, elapsed = 101 ms
1470617221 s:284 ms, Engine Speed =1488.000000, elapsed = 100 ms
1470617221:300,Latitude = 42.276520,elapsed = 1002 ms
1470617221:300,Longitude = -83.204842,elapsed = 1002 ms
1470617221:316, Vehicle Speed = 55.969997, elapsed = 100 ms
1470617221 s:385 ms, Engine Speed =1484.000000, elapsed = 101 ms
1470617221:416, Vehicle Speed = 55.879997, elapsed = 100 ms
1470617221 s:485 ms, Engine Speed =1480.000000, elapsed = 100 ms
1470617221:516, Vehicle Speed = 55.889999, elapsed = 100 ms
1470617221 s:585 ms, Engine Speed =1478.000000, elapsed = 100 ms
1470617221:617, Vehicle Speed = 55.849998, elapsed = 101 ms
1470617221 s:685 ms, Engine Speed =1462.000000, elapsed = 100 ms
1470617221:717, Vehicle_Speed = 55.739998, elapsed = 100 ms
1470617221 s:785 ms,Engine Speed =1444.000000,elapsed = 100 ms
1470617221:817, Vehicle Speed = 55.610001, elapsed = 100 ms
1470617221 s:886 ms,Engine Speed =1452.000000,elapsed = 101 ms
1470617221:917, Vehicle Speed = 55.549999, elapsed = 100 ms
1470617221 s:986 ms, Engine Speed =1458.000000, elapsed = 100 ms
1470617222:17, Vehicle Speed = 55.480000, elapsed = 100 ms
1470617222 s:86 ms, Engine Speed =1454.000000, elapsed = 100 ms
1470617222:117, Vehicle Speed = 55.309998, elapsed = 100 ms
1470617222 s:186 ms,Engine_Speed =1438.000000,elapsed = 100 ms
1470617222:217, Vehicle_Speed = 55.129997, elapsed = 100 ms
1470617222 s:286 ms,Engine_Speed =1432.000000,elapsed = 100 ms
1470617222:302,Latitude = 42.276573,elapsed = 1002 ms
```

Figure 5: Desired information with time stamp

#### **CHAPTER 3: NETWORK CONFIGURATION**

In order investigate the latency for the LTE, I have categorized the network configuration into two types, local and the remote network. I have done the socket programming to create server and client configuration for both TCP (Transmission Control Protocol) and UDP (User Datagram Protocol).

#### 3.1 LOCAL NETWORK

The local network configuration is as shown in Figure 6. The client is considered to be a car transmitting the CAN packets and GPS information to the server. Once the server receives the information, it will store the data. The client and the server are connected to the local network of University of Michigan Dearborn.

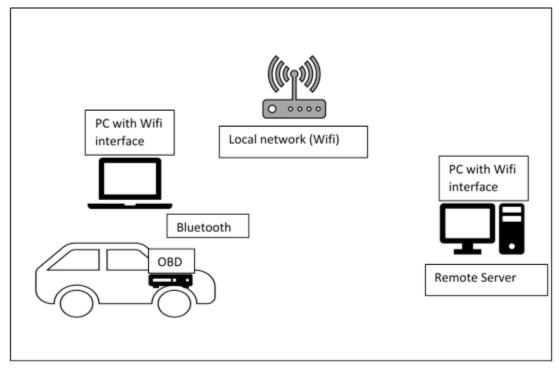


Figure 6: Local network configuration

A code has been developed for the local network configuration as shown in Figure 7. The code functionality and the evaluation is as follows:

#### Step 1:

- ➤ The server is configured with a standard receiving buffer size and IP network listening to a port.
- The client is configured to transmit the data to the server IP address on to the specific port that the server is configured to listen.
- > Dummy packets are transmitted to test the established communication between the server and the client.

#### Step 2:

- ➤ Client is now connected to the OBD dongle to interface the CAN network.
- ➤ Once the client establishes the connection with the server it will start receiving the CAN packets from the OBD dongle.
- ➤ The received packets will be transmitted to the server and the client will not wait for the acknowledgement from the server.
- > The server stores the data sent from the client. It listens to the client until the client closes the connection.

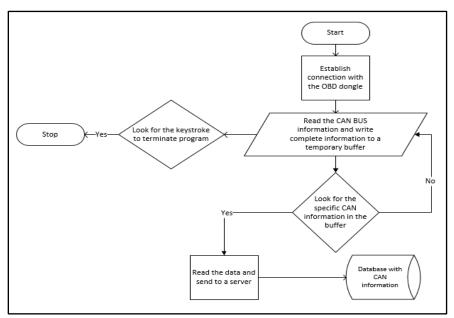


Figure 7: Algorithm Flowchart

The above two steps are tested by connecting the client and server to the local network (Same Wi-Fi network). No data has been lost when the network interface is local. The TCP network connection is like the UDP with one difference that the client transmits the information to the server after receiving the acknowledgement from the server. The TCP network is tested in both local network and through internet.

#### 3.2 REMOTE NETWORK

The remote network configuration is as shown in Figure 8. The client is a car transmitting the CAN (Controller Area Network) packets and GPS (Global Positioning System) information to the server via internet. The client is connected to the LTE (Long Term Evolution) network and the server to the private local area network. Once the server receives the information, it will store the data. Verizon LTE network is used for the network connection.

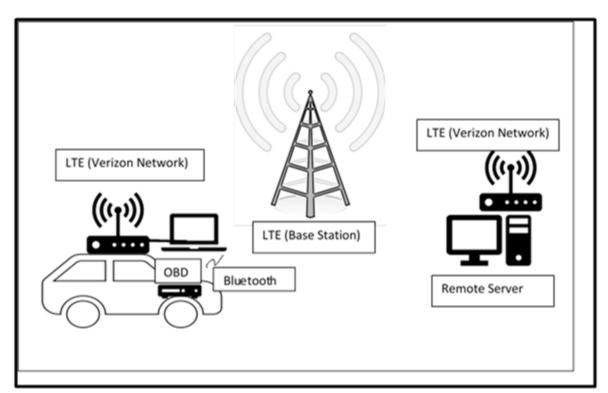


Figure 8: Remote network configuration

For local and the remote network configurations a time stamp is added on the server side and also on the client side to determine the latency. Latency observed for the local network is within the range of 0 to 16ms. Latency measured at different scenarios for the remote network configuration are discussed in chapter 4.

#### 3.3 ALGORITHM FOR NETWORK INTERFACE

Network specific software that communicates with the network specific device driver and the IP layer to provide the IP layer with consistent interface can be considered as the network interface. Each network interface has a network address. Based on the destination address of the packet to be transmitted the IP layer selects the appropriate network interface. A socket will be an endpoint in a two-way communication link between two programs running on the network. It is bounded to a port number so that the TCP layer can identify the application to where data need to be sent. A TCP connection can be uniquely identified by its endpoints. Each endpoint is a combination of an IP address and the port number. Thus, providing the access of having multiple connections between clients and server.

The basic TCP client server connection involves three steps, the client sends a synchronous message, the server sends an acknowledge with the clients synchronous message and contains the servers synchronous message, then the client sends an acknowledge for the servers synchronous message. Synchronous message indicates that the segment is being used to initialize a connection., and acknowledge indicates that the device sending the segment is conveying the acknowledgement for a message it has received. The algorithm works as follows:

#### Step1:

The server opens a passive connection, creating a transmission control block for the connection and prepares for the recipient connection from (Synchronous message) from the client. The client cannot connect to server until the server opens a passive connection.

#### Step2:

The client creates a transmission control block for the connection and sends a synchronous message to the server. The server listens to the client.

#### Step3:

The server receives the synchronous message from the client. It sends a synchronous message and acknowledgement message back to the client that contains an acknowledgment form clients and server synchronization. The client waits to receive an acknowledgement to the synchronous message it has sent, and also for the server synchronous message.

#### Step4:

The client receives the synchronous and acknowledge message from the server and sends an acknowledgement for the server synchronization. The client ready to establish connection. Server waits for the acknowledgement from the client for the synchronous message it has send previously.

#### Step5:

The clients will wait for the server to establish connection. The server receives the acknowledgement to its synchronization and establishes the connection.

#### Step6:

The client and server will be ready for the data transfer.

The client code below will explain the functions performed by the client.

// Attributes to initialize the socket

```
int recvbuflen = DEFAULT BUFLEN;
// Validate the parameters
if (argc != 2) {
   printf("usage: %s server-name\n", argv[0]);
  return 1;
}
// Initialize Winsock
iResult = WSAStartup(MAKEWORD(2,2), &wsaData);
if (iResult != 0) {
   printf("WSAStartup failed with error: %d\n", iResult);
  return 1;
}
ZeroMemory(&hints, sizeof(hints));
hints.ai family = AF UNSPEC;
hints.ai socktype = SOCK STREAM;
hints.ai protocol = IPPROTO TCP;
// Resolve the server address and port
iResult = getaddrinfo(argv[1], DEFAULT PORT, &hints, &result);
if ( iResult != 0 ) {
   printf("getaddrinfo failed with error: %d\n", iResult);
   WSACleanup();
  return 1;
}
// Attempt to connect to an address until one succeeds
for(ptr=result; ptr != NULL ;ptr=ptr->ai next) {
// Create a SOCKET for connecting to server
   ConnectSocket = socket(ptr->ai family, ptr->ai socktype,
     ptr->ai protocol);
   if (ConnectSocket == INVALID SOCKET) {
     printf("socket failed with error: %d\n", WSAGetLastError());
```

```
WSACleanup();
     return 1;
// Connect to server.
   iResult = connect(ConnectSocket, ptr->ai addr, (int)ptr->ai addrlen);
   if (iResult == SOCKET ERROR) {
     closesocket(ConnectSocket);
     ConnectSocket = INVALID SOCKET;
     continue;
   break;
 freeaddrinfo(result);
 if (ConnectSocket == INVALID SOCKET) {
   printf("Unable to connect to server!\n");
   WSACleanup();
   return 1;
 }
// Send an initial buffer
 iResult = send( ConnectSocket, sendbuf, (int)strlen(sendbuf), 0 );
 if (iResult == SOCKET ERROR) {
   printf("send failed with error: %d\n", WSAGetLastError());
   closesocket(ConnectSocket);
   WSACleanup();
   return 1;
 // shutdown the connection since no more data will be sent
 iResult = shutdown(ConnectSocket, SD SEND);
 if (iResult == SOCKET ERROR) {
   printf("shutdown failed with error: %d\n", WSAGetLastError());
   closesocket(ConnectSocket);
```

```
WSACleanup();
    return 1;
  }
  // Receive until the peer closes the connection
  do {
    iResult = recv(ConnectSocket, recvbuf, recvbuflen, 0);
    if ( iResult > 0 )
      printf("Bytes received: %d\n", iResult);
    else if ( iResult == 0 )
      printf("Connection closed\n");
    else
       printf("recv failed with error: %d\n", WSAGetLastError());
  \} while( iResult > 0 );
  // cleanup
  closesocket(ConnectSocket);
  WSACleanup();
The server code below explains the basic functions implemented
  WSADATA wsaData;
  int iResult;
  SOCKET ListenSocket = INVALID SOCKET;
  SOCKET ClientSocket = INVALID SOCKET;
  struct addrinfo *result = NULL;
  struct addrinfo hints;
  int iSendResult;
  char recvbuf[DEFAULT BUFLEN];
  int recvbuflen = DEFAULT BUFLEN;
  // Initialize Winsock
  iResult = WSAStartup(MAKEWORD(2,2), &wsaData);
```

```
if (iResult != 0) {
  printf("WSAStartup failed with error: %d\n", iResult);
  return 1;
}
ZeroMemory(&hints, sizeof(hints));
hints.ai family = AF INET;
hints.ai socktype = SOCK STREAM;
hints.ai protocol = IPPROTO TCP;
hints.ai flags = AI PASSIVE;
// Resolve the server address and port
iResult = getaddrinfo(NULL, DEFAULT PORT, &hints, &result);
if ( iResult != 0 ) {
  printf("getaddrinfo failed with error: %d\n", iResult);
  WSACleanup();
  return 1;
}
// Create a SOCKET for connecting to server
ListenSocket = socket(result->ai family, result->ai socktype, result->ai protocol);
if (ListenSocket == INVALID SOCKET) {
  printf("socket failed with error: %ld\n", WSAGetLastError());
  freeaddrinfo(result);
  WSACleanup();
  return 1;
// Setup the TCP listening socket
iResult = bind( ListenSocket, result->ai addr, (int)result->ai addrlen);
if (iResult == SOCKET ERROR) {
  printf("bind failed with error: %d\n", WSAGetLastError());
  freeaddrinfo(result);
```

```
closesocket(ListenSocket);
  WSACleanup();
  return 1;
freeaddrinfo(result);
iResult = listen(ListenSocket, SOMAXCONN);
if (iResult == SOCKET ERROR) {
  printf("listen failed with error: %d\n", WSAGetLastError());
  closesocket(ListenSocket);
  WSACleanup();
  return 1;
}
// Accept a client socket
ClientSocket = accept(ListenSocket, NULL, NULL);
if (ClientSocket == INVALID_SOCKET) {
  printf("accept failed with error: %d\n", WSAGetLastError());
  closesocket(ListenSocket);
  WSACleanup();
  return 1;
// No longer need server socket
closesocket(ListenSocket);
// Receive until the peer shuts down the connection
do {
  iResult = recv(ClientSocket, recvbuf, recvbuflen, 0);
  if (iResult > 0) {
```

```
printf("Bytes received: %d\n", iResult);
  // Echo the buffer back to the sender
     iSendResult = send( ClientSocket, recvbuf, iResult, 0 );
    if (iSendResult == SOCKET_ERROR) {
       printf("send failed with error: %d\n", WSAGetLastError());
       closesocket(ClientSocket);
       WSACleanup();
       return 1;
     }
     printf("Bytes sent: %d\n", iSendResult);
  else if (iResult == 0)
    printf("Connection closing...\n");
  else {
     printf("recv failed with error: %d\n", WSAGetLastError());
    closesocket(ClientSocket);
     WSACleanup();
    return 1;
} while (iResult \geq 0);
// shutdown the connection since we're done
iResult = shutdown(ClientSocket, SD_SEND);
if (iResult == SOCKET ERROR) {
  printf("shutdown failed with error: %d\n", WSAGetLastError());
  closesocket(ClientSocket);
  WSACleanup();
  return 1;
```

```
// cleanup
closesocket(ClientSocket);
WSACleanup();
```

Now the message buffer filled with either vehicle data or GPS data or both can be transferred to server from the car (Client). However, the above code is the basic functionality that works for only one server and client configuration. To accept multiple clients few more functionalities are developed to establish asynchronous communication (refer to section 2.3). For the local network, the performance is similar to the single server client configuration. However, the performance evaluation is still need to be tested on remote network configuration.

```
The code works as follows:
// Waits for the incoming connections
// Reads the set of socket descriptors
//clear the socket fd set
    FD ZERO(&readfds);
 //add master socket to fd set
     FD SET(master, &readfds);
 //add child sockets to fd set
     for (i = 0; i < max \ clients; i++)
       s = client socket[i];
       if(s > 0)
       {
         FD SET(s, &readfds);
//wait for an activity on any of the sockets, timeout is NULL, so wait indefinitely
     activity = select(0, &readfds, NULL, NULL, NULL);
    if (activity == SOCKET ERROR)
     {
```

```
printf("select call failed with error code : %d" , WSAGetLastError());
       exit(EXIT FAILURE);
//If something happened on the master socket, then its an incoming connection
    if (FD ISSET(master, &readfds))
       if ((new socket = accept(master, (struct sockaddr*)&address, (int *)&addrlen))<0)
         perror("accept");
         exit(EXIT FAILURE);
       //add new socket to array of sockets
       for (i = 0; i < max \ clients; i++)
         if (client socket[i] == 0)
            client socket[i] = new socket;
            printf("Adding to list of client sockets at index %d \n", i);
            break;
s = client socket[i];
//if client present in read sockets
       if (FD ISSET(s, &readfds))
  //get details of the client
         getpeername(s , (struct sockaddr*)&address , (int*)&addrlen);
//Check if it was for closing, and also read the incoming message
         valread = recv( s, buffer, recv buflen, 0);
         if( valread == SOCKET ERROR)
            int error code = WSAGetLastError();
```

```
if(error code == WSAECONNRESET)
              //Somebody disconnected, get his details and close the connection
              //Close the socket and mark as 0 in list for reuse
              closesocket(s);
              client socket[i] = 0;
            }
            else
              printf("recv failed with error code : %d" , error code);
            }
          }
         if (valread == 0)
//Somebody disconnected, get his details and close
//Close the socket and mark as 0 in list for reuse
            closesocket(s);
            client socket[i] = 0;
          }
//Echo back the message that came in
         else
//add null character, if you want to use with printf/puts or other string handling functions
            buffer[valread] = '\0';
      printf("%s:%d - %s \n", inet ntoa(address.sin addr), ntohs(address.sin port), buffer);
            send(s, buffer, valread, 0);
  closesocket(s);
  WSACleanup();
```

The server code developed is a standalone application, waiting for incoming connections, if the user terminates it. Each car connecting as a client will be assigned a unit port address within the application. Whenever the client closes the connection, the port address within the application assigned for that car will be deleted in the array, and will be assigned to new car that connects to the server. Thus, establishing multiple stable connections.

#### **CHAPTER 4: LATENCY MEASUREMENTS**

#### 4.1 SCENARIOS IMPLEMENTED

The scenarios that are considered for the simulation are as follows:

- 1) **Scenario 1:** Vehicle speed data received from the CAN bus is every 100 to 140ms. Projection information received from the GPS sensor is every 1sec. TOF calculated when the car is located in the parking slot within a range of 10miles radius from the server is 0 to 30ms.
- 2) **Scenario 2:** Vehicle speed data received from the CAN bus is every 100 to 140ms. Projection information received from the GPS sensor is every 1sec. TOF calculated when the car is driving at 70mph around 450miles radius from the server in the dense forest area of Pennsylvania with average signal connectivity is 300 to 400ms.
- 3) **Scenario 3:** Vehicle speed data received from the CAN bus is every 100 to 140ms. Projection information received from the GPS sensor is every 1sec. TOF calculated when the car is driving at 40mph around 650miles radius from the server in the heavy traffic of New York city is 150 to 250ms.

#### 4.2 SCENARIOS OVERVIEW

If we consider a worst case scenario that the traffic is moving at 70mph, then for each millisecond each vehicle covers a distance of 0.032 meters or 0.105 foot. If we assume a time latency of 100ms for the traffic update. Then each vehicle could able to transmit or receive a traffic update for every 3.2 meters or 10.5 foot approximately. If we assume that the vehicle speed sensor of vehicle transmits information every 100ms. It adds up to the to a time latency of 100ms resulting to a time lapse of 200ms. Therefore, it takes around 6.4meters or 21foot approximate distance to transmit the information to the server

Similarly, if a GPS sensor could able to update every 100ms and if we assume a time latency of 100ms to transmit to the server. Every 3.2meters or 10.5foot the server gets the projection information of vehicle and it adds to a time latency of network of 100ms in addition to reach the data to the server resulting in 200ms time lapse. Therefore, reduction in the time latency and time lapse plays an essential role in vehicular networks, in receiving the accurate and the reliable information of the traffic update.

#### 4.3 Conclusion

Many possible scenarios such as using IEEE802.11p and cellular networks at different ranges from the server, and the speeds of multiple vehicles connected to the server can still be considered to evaluate the time latency in the vehicular networks. However, the simulations that are outlined in this paper with one vehicle connected to the server at different ranges and speeds, shows how essential it is to determine the time latency to evaluate the overall performance of the vehicular networks.

The simulations have shown that the traffic updates that can be sent to the server or received from the server by the vehicles are affected by factors such as vehicle speed, range of server from the vehicle and the signal connectivity of the network used.

#### REFERENCES

- Ahmed S. Kaseb, Youngsol Koh, Everett Berry, Kyle McNulty, Yung-Hsiang Lu and Edward J. Delp, "Multimedia content creation using global network cameras: The making of CAM2", Signal and Information Processing (GlobalSIP) 2015 IEEE Global Conference on, pp. 15-18, 2015.
- Yunshu Liu, Xuanyu Chen, Cailian Chen, and Xingping Guan," Traffic Big Data Analysis Supporting Vehicular Network Access Recommendation" Communications (ICC), 2016 IEEE International Conference on,pp. 14 July 2016.
- 3. IEEE Trial-Use Standard for Wireless Access in Vehicular environments (WAVE) Resource manager
- 4. Dedicated short range communications (DSRC) message set dictionary SAE J2735\_201603
- 5. http://www8.garmin.com/aboutGPS/waas.html.
- 6. http://openxcplatform.com/
- 7. http://www.nmea.org/content/nmea\_standards/nmea\_0183\_v\_410.asp.
- 8. https://www.garmin.com/support/pdf/NMEA 0183.pdf
- 9. http://www.plaisance-pratique.com/IMG/pdf/NMEA0183-2.pdf
- 10. https://en.wikipedia.org/wiki/World Geodetic System
- 11. http://standards.sae.org/wip/j2945/
- 12. http://www.3gpp.org/DynaReport/33-series.htm
- 13. https://msdn.microsoft.com/en-us/library
- 14. https://msdn.microsoft.com/en-us/library/windows/desktop/aa363858(v=vs.85).aspx