

Improving Power System Voltage Stability by Using Demand Response to Maximize the Distance to the Closest Saddle-Node Bifurcation

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Abstract—Temporal load shifting via demand response can be used to improve power system frequency stability. Recent work has shown that spatio-temporal load shifting can also be used to improve power system static voltage stability. However, the best static voltage stability metric is an open question. In this paper, we propose a method to improve power system static voltage stability by maximizing the distance to the closest saddle-node bifurcation of the power flow. Specifically, we formulate a nonlinear nonconvex optimization problem in which we choose loading patterns that maximize this distance while also constraining the total system loading to remain constant so that the actions do not affect frequency stability. We derive the KKT conditions and solve the resulting nonlinear system of equations using the Newton-Raphson method and check if the solution is a local minimum. Using a 4-bus system and the IEEE 9-bus system as our test cases, we explore the performance of the algorithm and the accuracy of the obtained solutions. We compare the solution to those obtained using other voltage stability metrics including the smallest singular value of the power flow Jacobian and the loading margin, finding that all approaches produce different solutions. Using Kundur’s two area system, we also explore some algorithm convergence issues.

I. NOTATION

Sets

S_{PV} Set of all PV buses
 S_{PQ} Set of all PQ buses
 S_{DR} Set of buses with demand responsive loads

Variables & Parameters

θ_i Voltage angle at bus i
 V_i Voltage magnitude at bus i
 P_i Real power injection at bus i
 Q_i Reactive power injection at bus i
 d Distance to the closest Saddle-Node Bifurcation
 x System state vector
 λ System parameter vector (power injections)
 Λ Feasible set of λ
 n_{dr} Number of buses with demand responsive loads
 n_e Number of engineering limits
 m Length of system state and parameter vectors
 w Left eigenvector corresponding to zero eigenvalue
 α_i Ratio between real and reactive demand at bus i
 β Weighting matrix
 μ, γ Lagrange multipliers
 ζ Constant

Functions

$\mathcal{F}(\cdot): \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}^m$ Standard power flow
 $g_1(\cdot): \mathbb{R}^{n_{dr}} \rightarrow \mathbb{R}^{2n_{dr}}$ Demand response limits
 $g_2(\cdot): \mathbb{R}^m \rightarrow \mathbb{R}^{n_e}$ Engineering limits
 $h(\cdot): \mathbb{R}^{2n_{dr}} \rightarrow \mathbb{R}^{n_{dr}+1}$ Demand response assumptions

We use the superscript ‘c’ to denote states/parameters at the closest Saddle-Node Bifurcation (SNB), the superscript ‘*’ to denote states/parameters at the solution, and the superscript ‘0’ to denote states/parameters at the initial operating point. For notational simplicity, we assume that each bus has at most one generator or one load. The notation $X \succ 0$ means that X is a positive definite matrix.

II. INTRODUCTION

Demand response can be used to improve power system economics and reliability [1]. There has been a significant amount of recent research into the development of strategies that enable demand response resources to provide frequency regulation via temporal load shifting, e.g., [2], [3]. However, demand response can also be used to improve other types of power system stability, for example, static voltage stability via load shedding [4] and spatio-temporal load shifting [5], [6]. Spatio-temporal load shifting refers to increasing/decreasing the load at various points in the network while forcing the total loading to remain constant and then paying back the load changes in future time intervals. In this way, frequency stability is unaffected because the total loading is unchanged in every time period. Additionally, each load receives the same amount of energy over the entire horizon as it would have received without demand response.

The best static voltage stability metric is an open question. Our previous research investigated use of the loading margin [7] and the smallest singular value (SSV) of the power flow Jacobian [8] within the spatio-temporal load shifting problem [5], [6]. However, the loading margin specifies the direction of the changes to power injections precipitating an instability and the SSV gives only indirect information about the distance to instability [9].

In this paper, we explore the use of the distance to the closest saddle-node bifurcation of the power flow as the stability metric we would like to maximize by spatially shifting load across a network within a single time step (we leave the full spatio-temporal problem to future work). The distance to the closest saddle-node bifurcation (SNB) is a well-known stability metric [10]. Past work [11] showed that the optimal control direction to move the system away from instability is antiparallel to the normal vector at the closest SNB. The idea is generalized in [12] for computing the optimal design of system parameters (i.e., shunt and

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series compensation) to improve this distance. The benefit of this approach is that the resulting optimization problem can be solved by formulating the Karush-Kuhn-Tucker (KKT) conditions, solving the nonlinear system of equations using the Newton-Raphson method, and checking if the solution is a local minimum by using the iterative method proposed in [13]. By reinitializing the nonlinear system solver and repeating this process many times we may find the global minimum, though we have no guarantee. We note that, in practice, limit-induced bifurcations (LIB) may occur before SNBs. We do not consider LIBs here; in future work we will explore algorithmic approaches to maximize the distance to the closest SNB or LIB.

Our contributions are as follows. 1) We formulate the optimization problem and derive its KKT conditions. 2) We conduct case studies using a 4-bus system and the IEEE 9-bus system and explore the performance of the algorithm and the accuracy of the solution. In particular, we find that our algorithm is able to maximize the distance to the globally closest SNB for the 4-bus system but does not find the globally closest SNB for the 9-bus system, instead maximizing the distance to a locally closest SNB. However, the globally closest SNB of the 9-bus system is unrealistic. 3) We compare our solution to those obtained by formulations that use other stability metrics. We find that all approaches produce different results and we discuss the implications of this finding. 4) Using Kundur's two area system, we explore algorithm convergence issues.

The remainder of this paper is organized as follows. Section III describes the problem. Section IV reviews the optimization formulation used for finding the closest SNB. Section V formulates our problem and presents our solution approach. Section VI shows the results of our case studies and Section VII concludes the paper.

III. PROBLEM DESCRIPTION

A conceptual illustration of the problem is shown in Fig. 1. The power flow solvability boundary (black curve) is defined by a set of SNBs, where λ denotes power injections. Suppose the initial operating point with injections equal to λ^0 is not sufficiently far from its closest SNB. The system operator would like to increase this distance, which is a measure of static voltage stability. It could do so through generator redispatch, load shedding, and/or spatial load shifting. Here, we investigate the impact of spatial load shifting.

While generators take time to respond to dispatch commands, demand responsive loads can respond quickly if coordinated via low-latency communications systems. Load shedding reduces quality of service to consumers and requires an equivalent decrease in generation to maintain system frequency. In contrast, spatial load shifting decreases and increases loads at various points in the network while maintaining the total loading so as not to affect system frequency. Aggregations of loads such as residential and commercial air conditioning systems can both decrease and increase their power consumption for short periods of time. So long as the energy is "paid back" within a short period of time, quality of service can be maintained. While it would likely be uneconomical to purpose-build demand response

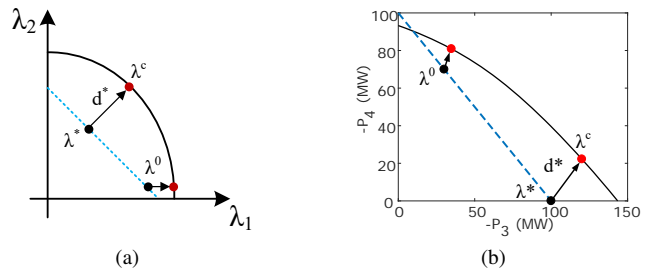


Fig. 1. Illustration of the problem. The black curve is the power flow solvability boundary and the dashed line shows the feasible solutions. λ^0 is the initial operating point; λ^* is the optimal operating point with the maximum shortest distance to the boundary and λ^c is the corresponding closest SNB. (a) Conceptual illustration. (b) 4-bus system example.

capability for this application, it could be one of many services that demand responsive loads could provide in future power networks.

In Fig. 1a, the blue dashed line is the feasible range of the injections, including the requirement that the total loading is constant. Our goal is to determine injections λ^* corresponding to the optimal operating point along the blue dashed line that maximize the distance d^* to the closest SNB λ^c . Figure 1b shows an example using a simple four bus system that will be discussed in detail in Section VI-A. The optimization problem is:

$$\max_{\lambda^* \in \Lambda^*} \left(\min_{\lambda^c \in \Lambda^c} \|\lambda^c - \lambda^*\|_2 \right), \quad (1)$$

where Λ defines the feasible set of λ .

In our formulation, we assume that the generator real power outputs do not change with the exception of that of the slack bus, which changes its output to compensate for the change in system losses that occurs when the load is spatially shifted. (Alternatively, we could have assumed that the total load plus losses remain constant, i.e., that the loads manage the change in system losses and none of the generator real power outputs change.) Additionally, we assume that PV bus voltages are fixed. Therefore, we choose only the real and reactive power consumption of each demand responsive load, which is modeled as constant power with constant power factor. In practice, the system operator could simultaneously redispatch generators and demand responsive loads to improve the stability margin, though the generators may be ramp limited. However, here we focus on characterizing the response of demand responsive loads alone.

IV. CLOSEST SADDLE-NODE BIFURCATIONS

We first review the approach for computing the closest SNB to a given operating point. The standard power flow equations [14] can be expressed as:

$$\mathcal{F}(x, \lambda) = f(x) - \lambda = 0, \quad (2)$$

where $x \in \mathbb{R}^m$ is the system state vector, $\lambda \in \mathbb{R}^m$ is the system parameter vector and $\mathcal{F} : \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}^m$. In this paper, we assume $x = [\theta_{i \in \mathcal{S}_{PV}}; \theta_{i \in \mathcal{S}_{PQ}}; V_{i \in \mathcal{S}_{PQ}}]$ and $\lambda = [P_{i \in \mathcal{S}_{PV}}; P_{i \in \mathcal{S}_{PQ}}; Q_{i \in \mathcal{S}_{PQ}}]$ (unless otherwise stated). The SNB is reached when the power flow Jacobian becomes

singular:

$$\frac{\partial f^T}{\partial x} w = 0, \quad (3)$$

where $w \in \mathbb{R}^m$ is a left eigenvector corresponding to the zero eigenvalue of the power flow Jacobian matrix. To obtain a unique solution of w , we normalized the left eigenvector such that $w^T w - 1 = 0$.

As discussed in [10], for a given operating point (x^0, λ^0) , if the distance to bifurcation is defined as Euclidean distance $d = \|\lambda^c - \lambda^0\|_2$, then the closest SNB can be found by solving the following optimization problem:

$$\min_{x^c, \lambda^c, w} \frac{1}{2} \|\lambda^c - \lambda^0\|_2^2 \quad (4a)$$

$$\text{subject to } \mathcal{F}(x^c, \lambda^c) = 0 \quad (4b)$$

$$\frac{\partial f^T}{\partial x} \Big|_{(x=x^c)} w = 0 \quad (4c)$$

$$w^T w - 1 = 0. \quad (4d)$$

To solve (4), we derive the KKT conditions. The Lagrange function is:

$$\mathcal{L} = \frac{1}{2} \|\lambda^c - \lambda^0\|_2^2 + \mu_1^T \mathcal{F}(x^c, \lambda^c) + \mu_2^T \frac{\partial f^T}{\partial x} \Big|_{(x=x^c)} w + \mu_3 (w^T w - 1), \quad (5)$$

where $\mu_1 \in \mathbb{R}^m$, $\mu_2 \in \mathbb{R}^m$, $\mu_3 \in \mathbb{R}$ are Lagrange multipliers. Therefore, the KKT conditions are:

$$\frac{\partial \mathcal{L}}{\partial x^c} = \mu_1^T \frac{\partial \mathcal{F}}{\partial x} \Big|_{(x=x^c)} + \mu_2^T \frac{\partial}{\partial x} \left(\frac{\partial f^T}{\partial x} w \right) \Big|_{(x=x^c)} = 0 \quad (6a)$$

$$\frac{\partial \mathcal{L}}{\partial \lambda^c} = (\lambda^c - \lambda^0)^T + \mu_1^T \frac{\partial \mathcal{F}}{\partial \lambda} = 0 \quad (6b)$$

$$\frac{\partial \mathcal{L}}{\partial w} = \mu_2^T \frac{\partial f^T}{\partial x} \Big|_{(x=x^c)} + 2\mu_3 w^T = 0 \quad (6c)$$

$$(4b) - (4d) \quad (6d)$$

From (6b) we know that $\mu_1^T \partial \mathcal{F} / \partial \lambda \neq 0$. Also, $\partial \mathcal{F} / \partial \lambda = -I$. Therefore, the Lagrange multiplier μ_1 must be nonzero. If we post-multiply (6c) by w , the first term becomes zero and since w is not zero, μ_3 must be zero. Then μ_2 is either zero or a right eigenvector corresponding to the zero eigenvalue of the power flow Jacobian (making the first term of (6c) zero). Assume μ_2 is a right eigenvector. Post-multiplying (6a) by μ_2 results in the first term becoming zero, and therefore the second term, which has quadratic form, must also equal zero. This is only possible if μ_2 lies in the null space of the (symmetric) matrix of that second term. Accordingly, the second term of (6a) must equal zero. Alternatively, if $\mu_2 = 0$ then that second term in (6a) is zero. In either case, the first term of (6a) must equal zero, so μ_1 must be a left eigenvector corresponding to the zero eigenvalue of the power flow Jacobian. Since both μ_1 and w are left eigenvectors corresponding to the zero eigenvalue of the power flow Jacobian, we can set $\mu_1 = \zeta_1 w$, where $\zeta_1 \neq 0$ is a scalar.

Hence, a locally closest SNB must satisfy the following equations:

$$\mathcal{F}(x^c, \lambda^c) = 0 \quad (7a)$$

$$\frac{\partial f^T}{\partial x} \Big|_{(x=x^c)} w = 0 \quad (7b)$$

$$w^T w - 1 = 0 \quad (7c)$$

$$(\lambda^c - \lambda^0) - \zeta_1 w = 0. \quad (7d)$$

Reference [10] proposed a similar set of equations, the only difference being that instead of (7d) they use the more general equation $(\lambda^c - \lambda^0) - (\partial \mathcal{F}^T / \partial \lambda) w = 0$ since they allow λ to be any system parameter whereas we define λ as power injections. Equation (7) is a set of $3m + 1$ nonlinear equations with $3m + 1$ unknowns. Direct methods, for instance, the Newton-Raphson method, or iterative methods such as the one given in [15] can be used to compute the numerical solutions to (7). Note that the KKT conditions are just necessary conditions giving us minima, maxima, and saddle points. Solutions obtained with Newton-Raphson need to be checked to ensure they are minima. In contrast, the iterative method in [15] guarantees that the solution is a local minimum, i.e., a locally closest SNB. The distance to the locally closest SNB is $d = \|\lambda^c - \lambda^0\|_2 = \|\zeta_1 w\|_2 = |\zeta_1|$. We can attempt to find the globally closest SNB by computing all of the locally closest SNBs using different initializations and determining the minimum d . This may be computationally intractable for large systems and we have no guarantee that we will obtain the globally closest SNB.

V. OPTIMIZATION FORMULATION

In our problem, we need to determine both the parameters λ^* corresponding to the optimal operating point and the parameters λ^c corresponding to the closest SNB. Since the real power injections at PV buses and the real and reactive power injections at PQ buses without demand responsive loads are unchanged, we divide λ^* into two parts. The controlled power injections $\lambda_1^* = [P_{i \in \mathcal{S}_{\text{DR}}}; Q_{i \in \mathcal{S}_{\text{DR}}}]$ are limited by the flexibility of the demand responsive loads:

$$g_1(\lambda_1^*) = \begin{bmatrix} P_i - \bar{P}_i, \forall i \in \mathcal{S}_{\text{DR}} \\ -P_i + \underline{P}_i, \forall i \in \mathcal{S}_{\text{DR}} \end{bmatrix} \leq 0, \quad (8)$$

where $g_1 : \mathbb{R}^{2n_{\text{dr}}} \rightarrow \mathbb{R}^{2n_{\text{dr}}}$ and $\underline{P}_i, \bar{P}_i$ are the lower and upper limits of the range of allowed changes to the real power consumption of the demand responsive loads. The uncontrolled power injections are $\lambda_2^* = [P_{i \in \mathcal{S}_{\text{PV}}}; P_{i \in \mathcal{S}_{\text{PQ}} \setminus \mathcal{S}_{\text{DR}}}; Q_{i \in \mathcal{S}_{\text{PQ}} \setminus \mathcal{S}_{\text{DR}}}] = \lambda_2^0$.

Our goal is to determine λ_1^* that maximizes the distance to its closest SNB. Therefore, the decision variables of the optimization problem are the system state vectors x^c, x^* , system parameter vectors λ^c, λ_1^* and the left eigenvector w . The optimization problem is:

$$\min_{x^c, \lambda^c, x^*, \lambda_1^*, w} -\frac{1}{2} (\lambda^c - \lambda^*)^T \beta (\lambda^c - \lambda^*) \quad (9a)$$

$$\text{subject to } \mathcal{F}(x^c, \lambda^c) = 0 \quad (9b)$$

$$\mathcal{F}(x^*, \lambda^*) = 0 \quad (9c)$$

$$\frac{\partial f^T}{\partial x} \Big|_{(x=x^c)} w = 0 \quad (9d)$$

$$w^T w - 1 = 0 \quad (9e)$$

$$h(\lambda_1^*) = 0 \quad (9f)$$

$$g_1(\lambda_1^*) \leq 0 \quad (9g)$$

$$g_2(x^*) \leq 0. \quad (9h)$$

The objective (9a) maximizes a weighted distance instead of the Euclidean distance ($\beta \succ 0$). Constraints (9b) and (9c) are the standard power flow equations for the SNB and the optimal operating point, respectively. Constraint (9d) implies that (x^c, λ^c) is an SNB. The left eigenvector w is normalized in (9e). Equation (9f) ensures our demand response assumptions are enforced at λ_1^* , specifically, 1) the total loading is constant and 2) the load is modeled as constant power with constant power factor:

$$h(\lambda_1^*) = \left[\sum_{P_i^* \in \mathcal{S}_{\text{DR}}} P_i^* - \sum_{Q_i^0 \in \mathcal{S}_{\text{DR}}} Q_i^0 \right] = 0, \quad (10)$$

where $h : \mathbb{R}^{2n_{\text{dr}}} \rightarrow \mathbb{R}^{n_{\text{dr}}+1}$. The inequality constraint (9g) is defined in (8). The inequality constraint (9h) specifies the engineering limits at (x^*, λ^*) . They include limits on the voltage magnitudes at PQ buses, the reactive power injections at PV buses and the slack bus, and the line flows ($g_2 : \mathbb{R}^m \rightarrow \mathbb{R}^{n_e}$). The Lagrange function of (9) is:

$$\begin{aligned} \mathcal{L} = & -\frac{1}{2}(\lambda^c - \lambda^*)^T \beta (\lambda^c - \lambda^*) + \mu_1^T \mathcal{F}(x^c, \lambda^c) \\ & + \mu_4^T \mathcal{F}(x^*, \lambda^*) + \mu_2^T \frac{\partial f^T}{\partial x} \Big|_{(x=x^c)} w + \mu_3(w^T w - 1) \\ & + \mu_5^T h(\lambda_1^*) + \gamma_1^T g_1(\lambda_1^*) + \gamma_2^T g_2(x^*), \end{aligned} \quad (11)$$

where $\mu_1, \mu_2, \mu_4 \in \mathbb{R}^m$, $\mu_3 \in \mathbb{R}$, $\mu_5 \in \mathbb{R}^{n_{\text{dr}}+1}$, $\gamma_1 \in \mathbb{R}^{2n_{\text{dr}}}$ and $\gamma_2 \in \mathbb{R}^{n_e}$ are Lagrange multipliers. The KKT conditions are:

$$\frac{\partial \mathcal{L}}{\partial x^c} = \mu_1^T \frac{\partial f}{\partial x} \Big|_{(x=x^c)} + \mu_2^T \frac{\partial}{\partial x} \left(\frac{\partial f^T}{\partial x} w \right) \Big|_{(x=x^c)} = 0 \quad (12a)$$

$$\frac{\partial \mathcal{L}}{\partial \lambda^c} = -(\lambda^c - \lambda^*)^T \beta^T + \mu_1^T \frac{\partial \mathcal{F}}{\partial \lambda} = 0 \quad (12b)$$

$$\frac{\partial \mathcal{L}}{\partial x^*} = \mu_4^T \frac{\partial f}{\partial x} \Big|_{(x=x^*)} + \gamma_2^T \frac{\partial g_2}{\partial x} \Big|_{(x=x^*)} = 0 \quad (12c)$$

$$\frac{\partial \mathcal{L}}{\partial \lambda_1^*} = (\lambda_1^c - \lambda_1^*)^T \beta_1^T + \mu_4^T \frac{\partial \mathcal{F}}{\partial \lambda_1} + \mu_5^T \frac{\partial h}{\partial \lambda_1^*} + \gamma_1^T \frac{\partial g_1}{\partial \lambda_1^*} = 0 \quad (12d)$$

$$\frac{\partial \mathcal{L}}{\partial w} = \mu_2^T \frac{\partial f^T}{\partial x_1} \Big|_{(x=x^c)} + 2\mu_3 w^T = 0 \quad (12e)$$

$$\text{equality constraints (9b) - (9f)} \quad (12f)$$

$$\gamma_{1,j} g_{1,j}(\lambda_1^*) = 0, \forall j = 1, \dots, 2n_{\text{dr}} \quad (12g)$$

$$\gamma_{2,k} g_{2,k}(x^*) = 0, \forall k = 1, \dots, n_e \quad (12h)$$

$$\gamma_1 \geq 0, \quad \gamma_2 \geq 0 \quad (12i)$$

$$\text{inequality constraints (9g) - (9h)} \quad (12j)$$

As before, μ_1 equals a constant times w , i.e., $\mu_1 = \zeta_2 w$, the second term of (12a) is equal to zero, and $\mu_3 = 0$. Therefore, an optimal solution should satisfy the following equations:

$$\mu_4^T \frac{\partial f}{\partial x} \Big|_{(x=x^*)} + \gamma_2^T \frac{\partial g_2}{\partial x} \Big|_{(x=x^*)} = 0 \quad (13a)$$

$$-\beta(\lambda^c - \lambda^*) - \zeta_2 w = 0 \quad (13b)$$

$$\beta_1(\lambda_1^c - \lambda_1^*) + \frac{\partial \mathcal{F}^T}{\partial \lambda_1} \mu_4 + \frac{\partial h^T}{\partial \lambda_1} \mu_5 + \frac{\partial g_1^T}{\partial \lambda_1} \gamma_1 = 0 \quad (13c)$$

$$\text{equality constraints (12f) - (12h)} \quad (13d)$$

$$\text{inequality constraints (12i) - (12j),} \quad (13e)$$

where β_1 is the partition of β corresponding to λ_1 . There are $5m + 5n_{\text{dr}} + n_e + 2$ equations and unknowns in (13a)-(13d). The solution algorithm is as follows. First, we initialize the Newton-Raphson solver to find the solution to (13a)-(13d). We check to see if the solution also satisfies (13e). If so, we check whether λ^c is a locally closest SNB to λ^* by using the iterative method of [15]. If so, then we check whether λ^c is a globally closest SNB to λ^* by testing different initializations within the iterative method to determine if there is a closer SNB to λ^* than λ^c . If we find that λ^c is the globally closest SNB then λ^* is the desired solution. Otherwise, we reinitialize the Newton-Raphson solver in the direction of the globally closest SNB to find a new λ^* and repeat the process.

In our cases studies, we compare the performance of our method to that of a brute force method. Specifically, for all possible loading patterns within a discrete mesh in which the total loading is constant, we compute the distance to the closest SNB via the method of [15]. The optimal loading pattern is the pattern associated with the maximum distance.

VI. CASE STUDY

All computation is done in MATLAB and with the help of MATPOWER [16] on an Intel(R) i7-4720HQ CPU with 16 GB of RAM. The base MVA for all cases is 100 MVA and we set $\beta = I$. The number of the equality constraints greatly influences the computation time of our method, therefore, we neglect (9h) in our case studies. In each case, our initial operating points satisfy (9h) and we also find that the optimal solutions we obtain also satisfy (9h).

A. Simple 4-bus system results

We first apply our method to the simple 4-bus system as shown in Fig. 2a. Bus 1 is the slack bus at a voltage of 1 pu, bus 2 is a PV bus outputting 10 MW at a voltage of 1 pu, and buses 3 and 4 are PQ buses with demand responsive loads of 30 MW and 70 MW, respectively. The reactance of the lines are $x_{13} = j0.5, x_{23} = x_{34} = j0.25$ p.u.

When λ only includes the real power injections at the PQ buses (i.e., $\lambda = [P_3; P_4]$), the solution is as shown in Fig. 1b. Specifically, the black curve is the power flow solvability boundary; the dashed blue line represents the total loading constraint, i.e., $P_3 + P_4 = -100$ MW; and the optimal loading pattern is $\lambda^* = [-100, 0]$ MW, which maximizes the shortest distance to the boundary.

If we instead define $\lambda = [P_{2-4}; Q_{3-4}]$, the initial distance to the closest SNB is $d = 0.0879$. The optimal solution determined by our method is $P_3^* = -63.74$ MW and $P_4^* = -36.26$ MW, and $d^* = 0.1264$, which is consistent with the optimal loading pattern obtained via the brute force method, as shown in Fig. 2b.

B. IEEE 9-bus system results

We next evaluate our method using the IEEE 9-bus system using the data available in MATPOWER [16]. The system has 1 slack bus (bus 1), 2 PV buses (buses 2 and 3), and 6 PQ

TABLE I
IEEE 9-BUS SYSTEM: INITIAL AND OPTIMAL POWER INJECTIONS (P.U.)

	P_2	P_3	P_4	P_5	P_6	P_7	P_8	P_9	Q_4	Q_5	Q_6	Q_7	Q_8	Q_9
λ^0	1.6300	0.8500	0.0000	-0.9000	0.0000	-1.0000	0.0000	-1.2500	0.0000	-0.3000	0.0000	-0.3500	0.0000	-0.5000
λ^c	1.0629	0.2508	-0.1248	-1.5821	-0.6003	-1.3436	-0.5726	-1.7980	-0.2961	-0.7620	-0.0638	-0.3343	-0.0741	-0.9325
λ^*	1.6300	0.8500	0.0000	-1.0842	0.0000	-0.7386	0.0000	-1.3272	0.0000	-0.3614	0.0000	-0.2585	0.0000	-0.5309

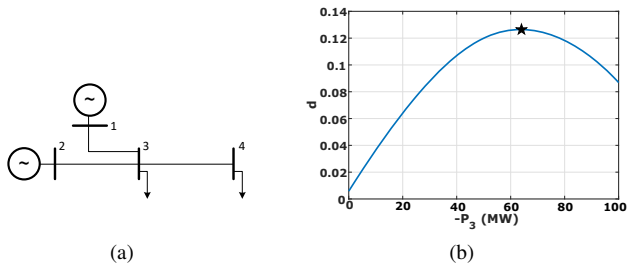


Fig. 2. (a) Single line diagram for the 4-bus system. (b) The distance to the closest SNB as a function of P_3 .

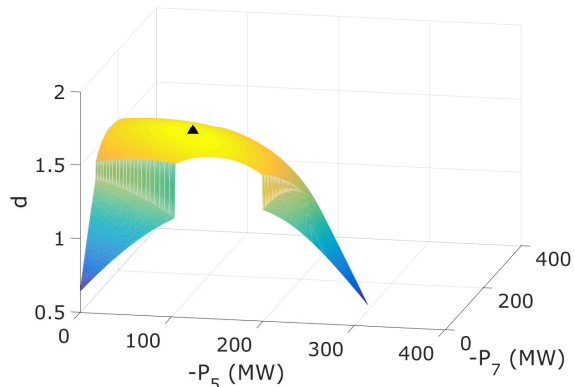


Fig. 3. The distance to the closest SNB as a function of P_5 and P_7 .

buses (buses 4-9). We model the entire load at buses 5, 7 and 9 (315 MW) as demand responsive. Hence, the system parameter vector is $\lambda = [P_{2-9}; Q_{4-9}]$ and the controlled power injections are $\lambda_1^* = [P_5^*; P_7^*; P_9^*; Q_5^*; Q_7^*; Q_9^*]$. We assume the system is initially operating at the operating point given within MATPOWER (see Table I, λ^0).

The optimal solution obtained by our method is given in Table I. The corresponding maximum distance is $d_\beta^* = 1.6263$, the optimal loading pattern is $P_5 = -108.42$ MW, $P_7 = -73.86$ MW, and $P_9 = -132.72$ MW. To verify the results, we compare the solution of our method to that of the brute force method. We use 5000 different directions as initializations of the iterative method of [15] to find locally closest SNB to λ^* and then determine the globally closest SNB. Figure 3 shows the distance to the closest SNB as a function of P_5 and P_7 (where $P_9 = -315 - P_5 - P_7$ since the total loading must be constant). The triangle represents the maximum distance obtained by the brute force method: $P_5 = -108$ MW, $P_7 = -74$ MW, $P_9 = -133$ MW and $d = 1.6263$, which is consistent with the solution of our method. There exist discontinuities on the surface in Fig. 3 because the feasibility boundary is very likely a folded hypersurface, so the distance is not continuous.

We have verified that λ^c is a locally closest SNB to λ^* but we cannot guarantee that this SNB is the globally

TABLE II
VOLTAGE AND REACTIVE POWER (P.U.) AT THE SNBs

	V_4	V_5	V_6	V_7	V_8	V_9	Q_1	Q_2	Q_3
SNB 1	0.5618	0.1593	0.5812	0.0795	0.4969	0.3571	7.8432	8.2343	7.1842
SNB 2	0.7780	0.6907	0.9071	0.8841	0.9009	0.6946	4.9147	1.6245	1.5874

closest SNB since the brute force method only explores 5000 random directions. Recently, [17] proposed a new enumeration search strategy to identify multiple local minima to a related optimization problem. Applying this strategy to (4), we obtain a closer λ^c to our λ^* with a distance $d = 0.1718$. This solution satisfies the KKT conditions (7) and may be the globally closest SNB to λ^* . The voltage magnitudes at the PQ buses and the reactive power injections at the buses with generators corresponding to this SNB (SNB 1) and the SNB that our method finds (SNB 2) are given in Table II. For both, the voltage magnitudes are low and the generator reactive power injections are high; however, SNB 1 is particularly unrealistic. Our method moves the system away from the relatively realistic locally closest SNB (SNB 2) but unfortunately there is a closer SNB (SNB 1), which it does not find. This example points to one of the drawbacks of our approach: we cannot guarantee that we will find the globally closest SNB so we might push the system away from a locally closest SNB and end up closer to the globally closest SNB.

We also compared this optimal solution to those obtained using other voltage stability metrics including the smallest singular value (SSV) of the power flow Jacobian and the loading margin (LM). Table III summarizes the results. The maximum SSV and LM cases are obtained from [5]. The results show that we obtain different loading patterns when maximizing different stability metrics, which is not surprising since the different margins capture different kinds of “distance to instability.” The loading margin describes the distance to voltage instability for power injection changes in a single direction, while the SSV and the distance to the closest SNB do not specify the direction. The SSV of the power flow Jacobian describes the distance to the singularity of power flow Jacobian matrix, which is an indirect measure of distance. In contrast, the distance to the closest SNB is a measure of distance in the parameter (power injection) space.

A disadvantage of using our method is that it relies on good initializations, whereas the iterative linear programming method used to maximize the SSV of the power flow Jacobian does not have this issue. The computation time for the 9-bus system is comparable for both approaches; however, it is not yet clear how the computational time/requirements compare for realistically-sized systems. Another disadvantage of our method is that we have no convergence guarantee, as we will show next.

TABLE III

OPTIMAL LOADING PATTERNS FOR DIFFERENT STABILITY METRICS

	$-P_5$ (MW)	$-P_7$ (MW)	$-P_9$ (MW)	SSV -	LM (MW)	d (p.u.)
max SSV	75	167	73	0.8995	516	1.5819
max LM	97	135	83	0.8984	566	1.6033
max d	108	74	133	0.8898	408	1.6263

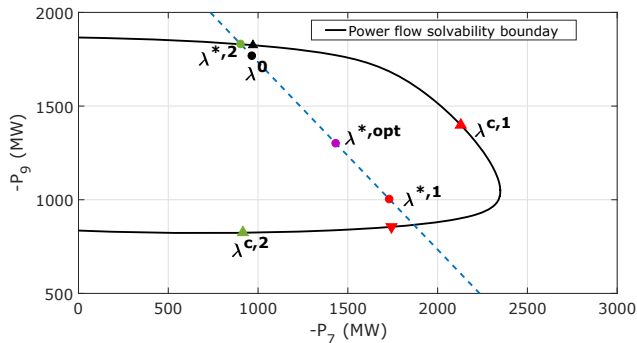


Fig. 4. The power flow solvability boundary of the Kundur system. The blue dashed line represents the total load constant constraint.

C. Convergence issues: Kundur's two area system results

Kundur's two area system [18] has 4 generators and 2 loads. We model the entire load at buses 7 and 9 (2134 MW) as demand responsive and set $\lambda = [P_7; P_9]$. The power flow solvability boundary is shown in Fig. 4. The black dot is the initial operating point $\lambda^0 = [P_7; P_9] = [-967; -1767]$ MW. The shortest distance between the black dot and the boundary (i.e., the distance from the black dot to the black triangle) is $d^0 = 0.5831$. Our method first finds the solution: $\lambda^{*,1}$ (red dot), $\lambda^{c,1}$ (red upper triangle) with $d^{*,1} = 5.615$; however, the globally closest SNB to $\lambda^{*,1}$ is not $\lambda^{c,1}$ but instead the SNB denoted with the red lower triangle with $d = 1.472$. Initializing the Newton-Raphson solver in the direction of the globally closest SNB to $\lambda^{*,1}$, we find another solution $\lambda^{*,2}$ (green dot), $\lambda^{c,2}$ (green upper triangle) with $d^{*,2} = 10.06$. However, $\lambda^{*,2}$ is on the solvability boundary and so we know that it is not the desired solution. In fact, neither solution is the desired solution. The desired solution is $\lambda^{*,opt}$ (pink dot), which has the maximum shortest distance to the boundary; it can not be obtained with our method. Further research is needed to develop approaches to cope with this problem.

VII. CONCLUSION

In this paper, we formulated a problem to spatially shift demand responsive load to improve static voltage stability. Specifically, we wish to increase the distance between the operating point and the point corresponding to the closest saddle-node bifurcation, which is a measure of static voltage stability. The problem was posed as a nonconvex nonlinear optimization problem and solved by formulating the KKT conditions, applying the Newton-Raphson method to solve them, and checking that the solution is a local minimum. Case study results using a simple 4-bus system and the IEEE 9-bus system showed that the distance to the closest SNB is improved by demand response actions, which increase and decrease individual loads while ensuring the total load

is constant. We also noted several issues with our method, specifically, we cannot guarantee that we find the globally closest SNB and, for some systems, we observe convergence issues. In the future, we would like to develop an improved algorithm that addresses these issues, test our method on larger systems, and compare the magnitude of stability margin improvement achievable with demand response to that achievable with generator redispatch.

VIII. ACKNOWLEDGMENT

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