Abstract

While Markov jump systems (MJSs) are more appropriate than LTI systems in terms of modeling abruptly changing dynamics, MJSs (and other switched systems) may suffer from the model complexity brought by the potentially sheer number of switching modes. Much of the existing work on reducing switched systems focuses on the state space where techniques such as discretization and dimension reduction are performed, yet reducing mode complexity receives few attention. In this work, inspired by clustering techniques from unsupervised learning, we propose a reduction method for MJS such that a mode-reduced MJS can be constructed with guaranteed approximation performance. Furthermore, we show how this reduced MJS can be used in designing controllers for the original MJS to reduce the computation cost while maintaining guaranteed suboptimality.

Keywords: Markov Jump Systems, System Reduction, Clustering

1. Introduction

As the control and machine learning communities build tools to model ever more complex dynamical systems, it will become increasingly important to identify redundant aspects of a model and remove them using various unsupervised learning techniques. State dimensionality reduction has long been common in control systems, using PCA and similar techniques. In this paper we consider the setting where switched systems have redundant modes, and we apply clustering – the other most fundamental unsupervised learning technique – to remove redundancies.

Switched systems generalize time-invariant systems and can be used to model abrupt changes in the environment (e.g. weather and road surfaces), controlled plants (e.g. functioning statuses of different components), disturbances, or even control goals (e.g. cost functions in the optimal control). Switched system models are used in a variety of applications including controlling a Mars rover exploring an unknown heterogeneous terrain, solar power generation, investments in financial markets, and communications with packet losses Blackmore et al. (2005); Cajueiro (2002); Loparo and Abdel-Malek (1990); Svensson et al. (2008); Ugrinovskii and Pota (2005); Sinopoli et al. (2005); Truong et al. (2021). Using a pool of modes and allowing them to switch brings versatility but also new challenges: the number of modes can grow easily during modeling. For example, for controlled plants composed with multiple components, if we model each combination of health statuses, e.g. working and faulty, of all components as a mode, then the number of modes grow exponentially with the number of components. Even analysis such as stability verification can become computationally intractable when the number of components is large, thus finding a
systematic and theoretically guaranteed way to reduce the number of modes in switched systems is imperative.

Existing work on (switched) system reduction mainly focuses on reducing the state dimension (Zhang et al., 2003) or constructing finite abstractions for the continuous state space Zamani and Abate (2014). Reducing the mode complexity, however, is still an uncharted field. A closely related line of research is the reduction of hidden Markov models (HMM) (Abate et al., 2014; Lun et al., 2018) and Markov chains (Zhang and Wang, 2019; Bittracher and Schütte, 2021), but HMM and Markov chain models are much simplified versions of switched dynamical models.

In this work, we study how one can perform mode reduction for Markov jump systems (MJS), a class of switched systems with the dynamics of mode switching governed by a Markov chain, and each individual mode characterized by a linear time-invariant (LTI) model. Our main contributions are the following:

• As an extension of Du et al. (2019), by treating the dynamics of each mode as features, we propose a clustering-based method that constructs a mode-reduced MJS.

• The reduced MJS is shown to provably well approximate the original MJS in terms of the trajectory difference.

• We show that LQR controllers designed with the reduced MJS can achieve guaranteed performance on the original MJS while significantly reducing the computational cost.

Our work adds a new dimension, i.e., reduction of modes, to the research of switched system reduction. This framework can be generalized to other problems such as stability analysis, robust and optimal control, invariance analysis, partially observed systems, etc. Other than constructing and analyzing the reduced MJS, the technical tools we develop in this work regarding perturbations can be applied to cases when there are model mismatches, e.g. system estimation errors incurred when dynamics are learned in identification or data-driven adaptive control as in Sattar et al. (2021).

2. Related Work

The work on reduction for stochastic switched systems can be roughly divided into three categories: bisimulation, symbolic abstraction, and order reduction.

**Bisimulation:** To evaluate the equivalency between two stochastic switched systems, notions of (approximate) probabilistic bisimulation are proposed in (Larsen and Skou, 1991; Desharnais et al., 2002, 2004). Approximation metrics from different perspectives (Abate, 2013) are developed to compare two systems, e.g. one(multi)-step transition kernels (Abate et al., 2011) and trajectories (Girard and Pappas, 2007; Tkachev and Abate, 2014; Julius and Pappas, 2009). Unlike existing work which typically defines the exact and approximate bisimulation on the state space, our work takes the mode space into consideration. Existing work on bisimulation is mainly conceptual and more for analysis purposes: they focus on developing bisimulation notions and approximation metrics, and “the majority of the examined approaches assume to be given two similar processes to compare, only a few put forward procedures for model approximation or abstraction with quantified quality” (Abate, 2013). This shortcoming is the work on abstractions address.

**Symbolic Abstraction:** Given a system with continuous state space, abstraction (Alur et al., 2000) considers discretizing the state space and then constructing a finite state symbolic model, which can be used as a surrogate for model verification (Clarke Jr et al., 2018; Kurshan, 2014) or controller synthesis (Maler et al., 1995). The work on abstraction for stochastic hybrid systems starts with the autonomous cases. Under uniform discretization, Abate et al. (2010, 2011) provide approximation
guarantees in terms of the discretization width. An adaptive partition scheme is proposed in Soudjani and Abate (2011), which mitigates the curse of dimensionality suffered by uniform sampling. Since the systems under consideration are autonomous, these work mainly serves verification purposes, but fall short toward controller synthesis goals. Zamani and Abate (2014) addresses this by allowing inputs in the systems. The idea of partitioning the continuous state space is similar to our work except that our partition is performed on the mode space, a.k.a. the discrete state space in hybrid systems, which provides a new yet closely related dimension to existing abstraction work.

**Order Reduction:** Another important line of research on system reduction is order reduction (Gugercin and Antoulas, 2004), where one seeks to reduce the dimension of the state space under certain criterion. With the help of linear matrix inequalities (LMIs), various methods have been applied for MJS, including $H_\infty$ reduction (Zhang et al., 2003), balanced truncation (Kotsalis and Rantzer, 2010), and $H_2$ reduction (Sun and Lam, 2016), etc. Order reduction is also applied to models with adversarial conditions such as time-varying delay (Zhang et al., 2015a) and partial transition probability observation (Zhang et al., 2015b; Shen et al., 2019).

### 3. Preliminaries and Problem Setup

For a matrix $E$, $E(i, :)$ denotes the $i$th row of $E$, and $E(i, j:k)$ denotes the $i$th row preserving only the $j$th to $k$th columns. For any index set $A$, $E(i, A)$ denotes the $i$th row preserving columns given by $A$. Let $\sigma_i(E) (\lambda_i(E))$ denote its $i$th largest singular (eigen) value. For any $s \in \mathbb{N}$, we let $[s] := \{1, 2, \ldots, s\}$. We say $\Omega_{1:r} := \{\Omega_1, \ldots, \Omega_r\}$ is a $r$-cluster partition of $[s]$ if $\bigcup_{i=1}^r \Omega_i = [s]$, $\Omega_i \cap \Omega_j = \emptyset$ for any $i \neq j$, and $\Omega_i \neq \emptyset$. We let $\Omega_{i, j}$ denote the cluster with $i$th largest cardinality. For a sequence of variables $X_0, X_1, \ldots, X_N$, let $X_{0:N} := \{X_i\}_{i=0}^N$.

#### 3.1. Preliminaries

In this work, we consider Markov jump systems (MJSs) with dynamics given by

$$\Sigma := \{x_{t+1} = A_{\omega_t} x_t + B_{\omega_t} u_t, \quad \omega_t \sim \text{Markov Chain}(T)\} \tag{1}$$

where $x_t \in \mathbb{R}^n$ and $u_t \in \mathbb{R}^p$ denote the state and input at time $t$. The switching nature of the dynamics is characterized by $s$ modes $\{A_i, B_i\}_{i=1}^s$ where $A_i \in \mathbb{R}^{n \times n}$ and $B_i \in \mathbb{R}^{n \times p}$ are state and input matrices for mode $i$. The active mode at time $t$ is indexed by $\omega_t \in [s]$, and the mode switching sequence $\omega_{0:t}$ follows a Markov chain with Markov matrix $T \in \mathbb{R}^{s \times s}$, i.e., $P(\omega_{t+1} = j | \omega_t = i) = T(i, j)$. We assume the Markov chain $T$ is ergodic. By properties of ergodicity, $T$ has a unique stationary distribution $\pi \in \mathbb{R}^s$, and we let $\pi_{\max}$ and $\pi_{\min}$ denote the largest and smallest element in $\pi$. In the remaining of the paper, we use $\Sigma := \text{MJS}(A_{1:s}, B_{1:s}, T)$ to denote the ground-truth MJS in (1) that we want to study, and similarly use notation $\text{MJS}(\cdot, \cdot, \cdot)$ to parameterize any MJS with expressions given in (1). We introduce the following two special types of Markov chain, which are closely tied to the main focus of this work.

**Definition 1 (Lumpability and Aggregatability (Buchholz, 1994))** Markov matrix $T$ is lumpable wrt. partition $\Omega_{1:r}$ on $[s]$ if for any $k, l \in [r]$, and $i, i' \in \Omega_k$, $\sum_{j \in \Omega_l} T(i, j) = \sum_{j \in \Omega_l} T(i', j)$.

Lumpability of a Markov chain coincides with the definition of probabilistic bisimulation in Desharnais et al. (2002), which describes an equivalence relation on $[s]$, i.e., two members are equivalent if they belong to the same cluster. For a Markov chain $T$ that is lumpable with respect to partition $\Omega_{1:s}$, we use $\zeta_t \in [r]$ to index the active cluster at time $t$, i.e., $\zeta_t = i$ if and only if $\omega_t \in \Omega_i$, and use $\zeta_{0:t}$ to denote the active cluster switching sequence.
CLUSTERING-BASED MODE REDUCTION FOR MARKOV JUMP SYSTEMS

3.2. Mode-reducibility Conditions

We first define an equivalence relation between two MJSs with different number of modes via a surjection from each mode of the larger MJS to the smaller one, which is characterized by a partition. This extends the bijection idea in Julius and Pappas (2009); Zhang et al. (2003), which can only be used to compare two switched systems with the same number of modes.

Definition 2 (Equivalence between MJSs) Consider two MJSs $\Sigma_1$ and $\Sigma_2$ with the same state and input dimensions $n, p$, but different number of modes $s_1$ and $s_2$ respectively. WLOG, assume $s_1 > s_2$. Let $\{x_t^{(1)}(i), u_t^{(1)}(i), \omega_t^{(1)}\}$ and $\{x_t^{(2)}(i), u_t^{(2)}(i), \omega_t^{(2)}\}$ denote their respective state, input, and mode index. $\Sigma_1$ and $\Sigma_2$ are equivalent if there exists a partition $\Omega_{1:s_2}$ on $[s_1]$ such that $\Sigma_1$ and $\Sigma_2$ have the same transition kernels, i.e., for any $t$, any $k, k' \in [s_2]$, any $x, x' \in \mathbb{R}^n$, and any $u \in \mathbb{R}^p$

\[
P(\omega_{t+1}^{(1)} = k', x_{t+1}^{(1)} = x' | \omega_t^{(1)} = k, x_t^{(1)} = x, u_t^{(1)} = u) = P(\omega_{t+1}^{(2)} = k', x_{t+1}^{(2)} = x' | \omega_t^{(2)} = k, x_t^{(2)} = x, u_t^{(2)} = u). \tag{2}
\]

If one views the discrete mode and the continuous state $\{\omega_t^{(1)}(i), x_t^{(1)}(i)\}$ as a hybrid state (Abate et al., 2011), then Definition 2 generalizes Definition 1 from Markov chains to MJSs. A sufficient condition in terms of the dynamics parameters, which guarantees that an MJS can be reduced to a smaller MJS with equivalency between them, is given by the following.

Definition 3 (Mode-reducibility Condition) $\Sigma$ is mode-reducible with respect to a partition $\Omega_{1:r}$, if its Markov chain $T$ is lumpable with respect to $\Omega_{1,r}$, and modes within the same cluster have the same dynamics, i.e., for any $k \in [r]$, any $i, i' \in \Omega_k$, we have $A_i = A_{i'}$, $B_i = B_{i'}$.

If this condition holds for $\Sigma$, we construct a mode-reduced MJS given by $\tilde{\Sigma} := \text{MJS}(\tilde{A}_{1:r}, \tilde{B}_{1:r}, \tilde{T})$ such that for any $k, l \in [r]$, any $i \in \Omega_k$, $\tilde{A}_k = A_i$ and $\tilde{B}_k = B_i$, $\tilde{T} \in \mathbb{R}^{ns'}$ with $T(k, l) = \sum_{j \in \Omega_l} T(i, j)$, which is illustrated in Fig. 1. Let $\{\tilde{x}_t, \tilde{u}_t, \tilde{\omega}_t\}$ denote the state, input, and mode index for the reduced $\tilde{\Sigma}$. Then, the following result shows that $\tilde{\Sigma}$ and $\Sigma$ are equivalent according to Definition 2. This is more of a teaser result for the ideal mode-reducible case, and formal problems regarding practical cases when $\Sigma$ is not mode-reducible will be formulated later.

Proposition 4 Suppose $\Sigma$ is mode-reducible and $\tilde{\Sigma}$ is constructed as above. Consider the case when the two MJSs have (i) initial mode distributions satisfy $P(\omega_0 \in \Omega_k) = P(\tilde{\omega}_0 = k)$ for all $k \in [r]$, (ii) the same initial states ($x_0 = \tilde{x}_0$), and (iii) the same input sequences ($u_{0:t-1} = \tilde{u}_{0:t-1}$). Then, these two MJSs have the same mode and state transition kernels, i.e., $P(\omega_t \in \Omega_k, x_t = x) = P(\tilde{\omega}_t = k, \tilde{x}_t = x)$ for all $t$, all $k \in [r]$ and $x \in \mathbb{R}^n$. Particularly, there exists a special type of reduced $\tilde{\Sigma}$ such that the modes are synchronized: for all $t$, $\tilde{\omega}_t = \zeta_t$. In this case, $\tilde{x}_t = x_t$ for all $t$. 

Figure 1: Illustration of reduction under mode-reducibility condition.
Proposition 4 first shows the equivalency between $\Sigma$ and $\bar{\Sigma}$ for the transition kernels, then if certain synchrony, exists between $\omega_{0,t}$ and $\bar{\omega}_{0,t}$, the equivalency also exists for the trajectories. The condition $\bar{\omega}_t = \omega_t$ in Proposition 4 essentially establishes a coupling between the Markov chains $\omega_{0,t}$ and $\bar{\omega}_{0,t}$. With $P(\omega_t \in \Omega_k, \omega_t = k) = P(\omega_t \in \bar{\Omega}_k) = P(\bar{\omega}_t = k)$. Establishing coupling between the stochastic systems usually allows for stronger equivalency and approximation metrics result. Similar coupling scheme is implicitly used in Julius and Pappas (2009); Zhang et al. (2003); an optimal coupling by minimizing Wasserstein distance is discussed in Tkachev and Abate (2014); and a weaker coupling using the idea of HMM is discussed in Shen et al. (2019).

In Definition 3 and Proposition 4, one can view $\{\omega_t, x_t\} \in [s] \times \mathbb{R}^n$ as a hybrid state, then $T$ being lumpable guarantees the existence of an equivalence relation in the discrete domain $[s]$, while state/input matrices being the same guarantees this in the continuous domain $\mathbb{R}^n$.

3.3. Problem Formulation

The mode-reducibility condition in Definition 3 and Proposition 4 provide principles under which one could construct a reduced MJS that is equivalent to the original one. Though seemingly strong, the mode-reducibility condition may hold approximately in many practical scenarios that allow for mode reduction. For example, in hybrid power systems, solar panels that are geographically close tend to have similar levels of solar radiation thus similar power dynamics. Hence, we would like to study the following: suppose the mode-reducibility condition holds approximately, how can we construct a reduced MJS? From here, we formulate the following two concrete problems.

**Problem P1 (Lumpable Case)** Assume the dynamics of $\Sigma=MJS(A_{1:s}, B_{1:s}, T)$ are known. Suppose there exists a partition $\Omega_{1:r}$ on $[s]$ such that: for any $\Omega_k, \Omega_l$ and any $i, i' \in \Omega_k$,

$$\|A_i - A_{i'}\|_F \leq \epsilon_A, \quad \|B_i - B_{i'}\|_F \leq \epsilon_B,$$

$$\frac{1}{2} \sum_{l \in [r]} \sum_{j \in \Omega_l} T(i, j) - \sum_{j \in \Omega_l} T(i', j) \leq \epsilon_T. $$

Then, we seek to estimate the partition $\Omega_{1:r}$, construct a reduced MJS, and provide guarantees on the behavior difference incurred by the reduction.

Throughout this work, we will refer to $\epsilon_A, \epsilon_B, \text{ and } \epsilon_T$ as perturbations (that lead to violation of mode-reducibility condition in Definition 3 and prevents the existence of $\bar{\Sigma}$ as in Proposition 4). For the special aggregatable case, we separately formulate a similar problem in Problem P2. Even though solutions to P1 automatically solve P2, as we shall see in the next section, algorithm theoretical guarantees for P2 require milder and more interpretable assumptions than P1.

**Problem P2 (Aggregatable Case)** In Problem P1, replace (4) with $\|T(i, :)^T - T(i', :)^T\|_1 \leq \epsilon_T$. In P2, $A_i, B_i, \text{ and } T(i, :)$ can be thought of as features and provide the position for mode $i$ in the feature space, and pairwise distances between all the modes reflect the partition $\Omega_{1:r}$. The recovery of the partition $\Omega_{1:r}$ resembles the classical clustering problems studied in the machine learning community. However, this resemblance does not directly apply to the weak lumpability condition (4) in P1, as the similarity in the transition probabilities is encoded by the partition $\Omega_{1:r}$. This makes the already non-convex clustering problem even more challenging.

In the aforementioned hybrid power system example where the geographical closeness serves as expert knowledge, one may be tempted to naively attempt construction of a small scale system from the beginning, but this makes it impossible to know how well it approximates the true dynamics.
Our approach would be to instead approximate the dynamics as well as possible with many modes and then use clustering to reduce the overall system. Hence, other than constructing a reduced MJS, we also seek to provide guarantees on the behavior difference and performance degradation incurred by the reduction. These guarantees play a critical role in safety-critical environments.

4. Clustering-based Mode Reduction for MJS

### Algorithm 1: System Reduction for MJS

<table>
<thead>
<tr>
<th>Step</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Construct feature matrix $\Phi$</td>
</tr>
<tr>
<td>2</td>
<td><strong>Case P2</strong></td>
</tr>
<tr>
<td>3</td>
<td>Solve k-means problem: $\hat{\Omega}<em>{1:r}, \hat{c}</em>{1:r} = \arg\min_{\Omega_{1:r}, \epsilon_{1:r}} \sum_{k=1}^{r} \sum_{i \in \Omega_k} | U_r(i,:) - \hat{c}_k |^2 $</td>
</tr>
<tr>
<td>4</td>
<td>Construct $\hat{\Sigma}$, $\forall k, l \in [r]$</td>
</tr>
<tr>
<td>5</td>
<td>$\hat{A}_k = \frac{1}{</td>
</tr>
</tbody>
</table>

Output: $\hat{\Sigma}$ : MJS($\hat{A}_{1:r}, \hat{B}_{1:r}, \hat{T}$)

The clustering-based MJS reduction method is presented in Algorithm 1. We treat the estimation of partition $\Omega_{1:r}$ essentially as a mode clustering problem with dynamics matrices $A_i$, $B_i$ and transition distribution $T(i,:)$ serving as features for mode $i$. We first construct feature matrix $\Phi$ from Line 2 to Line 8, with $\Phi(i,:)$ denoting the features of mode $i$. For the aggregatable case in Problem P2, we simply stack the vectorized $A_i$, $B_i$ and $T(i,:)$, and use $\alpha_A, \alpha_B, \alpha_T$ to denote their weights respectively. One way to choose these weights is to use a normalization, e.g. $\alpha_A = \frac{1}{\max_j \| A_j \|}$, so that these three features would have the same scales. By the definition of aggregatability, similarities of $T(i,:)$ among different modes are direct indicators for the groundtruth partition $\Omega_{1:r}$. For the lumpable case, however, this is not an option since two modes belonging to the same cluster can still have different transition probabilities $T(i,:)$ even if $\epsilon_T = 0$. According to (4), the groundtruth partition $\Omega_{1:r}$ is only embodied in the mode-to-cluster transition probabilities $\sum_{j \in \Omega_l} T(i,j)$ constructed using the groundtruth partition itself. This leaves us in a “chicken-and-egg” dilemma. To deal with this, from Line 5 to 8, we compute the first $r$ left singular vectors $W_r$ of matrix $\text{diag}(\pi)^{\frac{1}{2}} T \text{diag}(\pi)^{-\frac{1}{2}}$, and then weight it by $\text{diag}(\pi)^{-\frac{1}{2}}$ to obtain matrix $S_r \in \mathbb{R}^{s \times r}$, which is used to construct features in $\Phi$ for the lumpable case. We will later justify using $S_r$ as features by showing row similarities in $S_r$ reflect the partition under certain assumptions.

With the feature matrix $\Phi$, to recover the partition, we resort to k-means: in Line 10, k-means is applied to the first $r$ left singular vector $U_r$ of $\Phi$. Extracting the low-rank structure $U_r$ de-noises the impact of the perturbations. Based on the k-means solution $\hat{\Omega}_{1:r}$, we construct the reduced $\hat{\Sigma}$ by averaging modes within the same estimated cluster. When one picks $\alpha_A = \alpha_B = 0$, i.e. only the Markov matrix $T$ is used to cluster the modes, our clustering scheme under the aggregatable case P2 is equivalent to the one in Zhang and Wang (2019). The lumpable case P1, on the other hand, is based on the preliminary analysis in Meiù and Shi (2001).
4.1. Theoretical Guarantees for Clustering

We assume a \((1 + \epsilon)\) solution (Lei et al., 2015) to the k-moons problem in Algorithm 1 can be obtained, i.e., \(\sum_{k \in [r], i \in \Omega_k} \| U_r(i, :) - \hat{c}_k \|^2 \leq (1 + \epsilon) \min_{\epsilon'_k, c_k} \sum_{k \in [r], i \in \Omega_k} \| U_r(i, :) - c_k' \|^2\). We later show how \(\epsilon\) affects the overall clustering performance. To evaluate the performance of partition estimation, we define misclustering rate (MR) as \(\text{MR}(\hat{\Omega}_{1:r}) = \min_{ \mathcal{H} } \sum_{k \in [r]} \sum_{i \in \Omega_k} \| U_r(i, :) - \hat{c}_k \|^2 \), where \(\mathcal{H}\) is the set of all bijections from \([r]\) to \([r]\) so that the comparison finds the best cluster label matching. We define the following averaged feature matrix \(\Phi\) based on the underlying partition \(\Omega_{1:r}\): for all \(i \in [s]\) (suppose \(i \in \Omega_k\) for some \(k \in [r]\)), \(\Phi(i,:) = \frac{1}{|\Omega_k|} \sum_{i' \in \Omega_k} \Phi(i',:)\). By construction, there are up to \(r\) unique rows in \(\Phi\), hence \(\text{rank}(\hat{\Phi}) \leq r\). We first present the clustering guarantee for Problem P2, i.e., the aggregatable case.

**Theorem 5** Consider Problem P2 and Algorithm 1. Suppose \(\hat{\Omega}_{1:r}\) is a \((1 + \epsilon)\)-k-moons solution. Let \(\epsilon_{\text{Agg}} := \alpha A e_A + \alpha B e_B + \alpha T e_T\). Then, if \(\text{rank}(\hat{\Phi}) = r\) and \(\epsilon_{\text{Agg}} \leq \frac{\sigma_r(\hat{\Phi})}{8 \sqrt{s} (2 + \epsilon) |\Omega_{1:r}|}\), we have

\[
\text{MR}(\hat{\Omega}_{1:r}) \leq \frac{64(2 + \epsilon)s}{\sigma_r^2(\hat{\Phi})} \epsilon_{\text{Agg}}^2 \tag{5}
\]

Additionally, if \(\epsilon_{\text{Agg}} \leq \frac{\sigma_r(\Phi)}{8 \sqrt{s} (2 + \epsilon) |\Omega_{1:r}|}\), then \(\text{MR}(\hat{\Omega}_{1:r}) = 0\).

We see the misclustering rate can be bounded by the inner-cluster spread \(\epsilon_{\text{Agg}}\) when it is small enough. As for the inter-cluster distance, i.e., dissimilarity of \(A_i, B_j\) and \(T(i,:)\) of modes in different clusters, it affects the guarantee through the term \(\sigma_r(\hat{\Phi})\). This is because when two modes belonging different clusters have similar features, their rows in the averaged feature matrix \(\Phi\) will also be similar, which could lead to small \(\sigma_r(\Phi)\) and larger error bound.

The clustering guarantee for the lumpable case in Problem P1 is more involved than the aggregatable case. We first provide a few more background notions and definitions that can help the exposition. We say a Markov matrix \(T\) is reversible if there exists a distribution \(\pi \in \mathbb{R}^s\) such that \(\pi(i)T(i,j) = \pi(j)T(j,i)\) for all \(i, j \in [s]\). This condition for reversibility translates to \(\text{diag}(\pi)T = T^\top \text{diag}(\pi)\) when \(T\) is ergodic with stationary distribution \(\pi\). For a reversible Markov matrix that is also lumpable, we have the following property.

**Lemma 6** (Appendix A in Meilă and Shi (2001)) For a reversible Markov matrix \(T\) that is also lumpable with respect to partition \(\Omega_{1:r}\), it is diagonalizable with real eigenvalues. Let \(S \in \mathbb{R}^{s \times s}\) denote an arbitrary eigenvector matrix of \(T\). Then, there exists an index set \(A \subseteq [s]\) with \(|A| = r\) such that for all \(k \in [r]\), for all \(i, j' \in \Omega_k\), we have \(S(i,A) = S(j',A)\).

We say \(T\) in Lemma 6 has informative spectrum if \(A = [r]\) and \(|\lambda_r(T)| > |\lambda_{r+1}(T)|\), which implies that the uniquely defined \(r\) leading eigenvalues have eigenvectors that carry partition information in Lemma 6. For lumpable Markov matrices, we define the \(\epsilon_T\)-neighborhood of \(T\):

\[
\mathcal{L}(T, \Omega_{1:r}, \epsilon_T) := \left\{ T_0 \in \mathbb{R}^{s \times s} : T_0 \text{ is Markovian,} \sum_{k \in [r], l \in \Omega_k} T_0(i_j, j) = \frac{1}{|\Omega_k|} \sum_{i' \in \Omega_k} T(i',j), \|T_0 - T\|_\infty \leq \epsilon_T \right\}. \tag{6}
\]

Then we provide the clustering guarantee for the lumpable case.

**Theorem 7** Consider Problem P1 and Algorithm 1. Let \(\gamma_1 := \sum_{i=2}^s \frac{1}{1 - \lambda_i(T)}, \gamma_2 = \min \{\sigma_r(H) - \sigma_{r+1}(H), 1\}, \gamma_3 = \frac{16\gamma_1^2 \sqrt{\frac{s}{\max \{\text{trace}\(T\)\}}} |\Omega|}{\gamma_2 \sigma_{\text{min}}}, \text{ and } \epsilon_{\text{Lmp}} := \alpha A e_A + \alpha B e_B + \gamma_3 \alpha T e_T\). Assume there exists an
ergodic and reversible \( T_0 \in \mathcal{L}(T, \Omega_{1:r}, \epsilon_T) \) with informative spectrum. Suppose \( \hat{\Omega}_{1:r} \) is a \((1 + \epsilon)\) \textit{k}-means solution. Then, if \( \text{rank}(\hat{\Phi}) = r \), \( \epsilon_T \leq \frac{\epsilon_{\text{min}}}{\tau_{\text{max}}} \), and \( \epsilon_{\text{Lmp}} \leq \frac{\sigma_r(\hat{\Phi}) \sqrt{\Omega_{(r)} + \Omega_{(1)}}}{8\sqrt{s(2+\epsilon)(\Omega_{(1)})}} \), we have

\[
\text{MR}(\hat{\Omega}_{1:r}) \leq \frac{64(2+\epsilon)s^2}{\sigma_r^2(\hat{\Phi})}\epsilon_{\text{Lmp}}.
\]

(7)

Additionally, if \( \epsilon_{\text{Lmp}} \leq \frac{\sigma_r(\hat{\Phi})}{8\sqrt{s(2+\epsilon)(\Omega_{(1)})}} \), then \( \text{MR}(\hat{\Omega}_{1:r}) = 0 \).

The difference of using \( T \) and \( S_r \) to construct features for lumpable and aggregatable cases is also reflected in the comparison between Theorem 5 and 7. They are almost identical with an additional \( \gamma_3 \) term. This term describes how much the lumpability perturbation \( \epsilon_T \) on \( T \) affects the row equalities of its spectrum-related matrix \( S_r \) in Lemma 6. The assumption on the existence of \( T_0 \) with informative spectrum guarantees (i) the partition \( \Omega_{1:r} \) information is carried by the leading eigenvectors of \( T_0 \) as introduced in Lemma 6, and (ii) this information can still be preserved in \( S_r \) as long as \( T \) is close to \( T_0 \). Saying this, the theory for clustering in the lumpable case may not hold for an arbitrary lumpable \( T \), but only those close to Markov matrices with informative spectra.

5. Approximation Guarantees

With perturbation \( \epsilon_A, \epsilon_B, \epsilon_T \), the reduced \( \tilde{\Sigma} \) may not be equivalent to the original \( \Sigma \) as in Proposition 4. In this case, if certain approximation guarantees can be established, they can serve verification purposes such as safety (Julius and Pappas, 2009) and invariance (Soudjani and Abate, 2011) evaluations. In this section, we show that when \( \Sigma \) is autonomous and stable, the reduced system \( \tilde{\Sigma} \) can provably well approximate the original system \( \Sigma \) in terms of the trajectory difference \( \|x_t - \hat{x}_t\| \).

Note that autonomous MJSs also include MJSs under mode-dependent state-feedback controller \( K_{1:s} \) such that the closed-loop dynamics \( x_{t+1} = (A_{\omega_t} + B_{\omega_t} K_{\omega_t})x_t \) is autonomous. Autonomous system reduction is also considered in Abate et al. (2011); Tkachev and Abate (2014); Bian and Abate (2017). Since we have shown in Theorem 5 and 7 that \( \text{MR}(\hat{\Omega}_{1:r}) = 0 \) when perturbation \( \epsilon_A, \epsilon_B, \epsilon_T \) are small, in the following, we assume \( \Omega_{1:r} = \hat{\Omega}_{1:r} \) for simplicity.

The stability under consideration is the mean-square stability (MSS) that is typically studied in MJS Costa et al. (2006). We say \( \Sigma \) is mean-square stable (MSS), if there exists \( \Sigma_{\infty} \) such that \( \lim_{t \to \infty} \mathbb{E}[x_t^2 x_t^T] = \Sigma_{\infty} \). Define the augmented state matrix \( A \in \mathbb{R}^{sn^2 \times sn^2} \) with its \( ij \)-th \( n^2 \times n^2 \) block given by \( [A]_{ij} := \mathcal{T}(j, i) \cdot A_j \otimes A_j \) and let \( \rho \) denote the spectral radius of \( A \). Then, under the autonomous case, i.e., \( u_t = 0 \), \( \Sigma \) being MSS is equivalent to \( \rho < 1 \). Furthermore, we define \( \tau := \sup_{k \in \mathbb{N}} \|A^k\|/\rho^k \), which measures how fast \( \|A^k\|\hat{\tau} \) converges to \( \rho \). We let \( \bar{A} := \max_i \|A_i\| \) and \( \rho_0 := \frac{1 + \rho}{2} \). Then we have the following bounds on the trajectory difference.

**Theorem 8**  Consider \( \Sigma \) that is autonomous and MSS. Assume \( \tilde{\Omega}_{1:r} = \Omega_{1:r} \) in Algorithm 1. Suppose \( \Sigma \) and \( \hat{\Sigma} \) have the same initial states, i.e., \( x_0 = \hat{x}_0 \). Mode \( \omega_t \) of \( \tilde{\Sigma} \) is synchronous to \( \omega_t \) of \( \Sigma \), i.e., for all \( t \), if \( \omega_t \in \Omega_k \) then \( \hat{\omega}_t = k \). Then, when \( \epsilon_A \leq \min\{\bar{A}, \frac{1-\rho}{6\sqrt{\tau}A}\} \), we have,

\[
\mathbb{E}[\|x_t - \hat{x}_t\|^2] \leq 12rst\rho_0^{t-1}r^2\bar{A}\|x_0\|^2\epsilon_A.
\]

(8)

We require the synchrony between the mode sequences \( \omega_{0:t-1} \) and \( \hat{\omega}_{0:t-1} \) so that evaluating \( \|x_t - \hat{x}_t\| \) is meaningful as we discussed in Section 3.2. When \( \omega_t \) is observed while \( x_t \) is not, by forcing mode synchrony, we can estimate \( x_t \) with \( \hat{x}_t \). Since \( \rho_0 < 1 \) due to the MSS of \( \Sigma \), we know \( \rho_0 < 1 \), which implies that the trajectory difference \( \|x_t - \hat{x}_t\| \) converges to 0 exponentially with \( t \). Note that this also implies that \( \hat{\Sigma} \) is MSS, as otherwise the difference would not converge.

Proposition 4 provides a sanity check for Theorem 8 when \( \epsilon_A = 0 \), we have \( x_t = \hat{x}_t \).
6. Controller Design with Case Study on LQR

A typical class of controllers for MJS is known to be mode-dependent. A mode-dependent controller is essentially a collection of individual controllers such that each mode is associated with one, and the deployed controller switches with corresponding modes. With the reduced MJS Σ, we can design mode-dependent controller $\hat{K}_{1:r}$ for Σ and then associate every mode $i$ in Σ with $\hat{K}_k$ if $i \in \Omega_k$. Since Σ has a smaller scale than Σ, the computation cost may be reduced. In the following, with the classical infinite time linear quadratic regulator (LQR) problem as an example, we showcase the implementation of this idea and provide its suboptimality guarantees.

In the infinite LQR problems, given positive definite cost matrices $Q$ and $R$, we define quadratic cost $J_T = \mathbb{E} \left[ \sum_{t=0}^{T-1} \left( x_t^T Q x_t + u_t^T R u_t \right) + x_T^T Q x_T \right]$. The goal is to design inputs to minimize the infinite time average cost $\limsup_{T \to \infty} \frac{1}{T} J_T$ under Σ. The optimal solution is well studied in Costa et al. (2006). Mode-dependent controllers $K_{1:s}$ is obtained by solving $s$ coupled discrete algebraic Riccati equations each of which is parameterized by $A_i$, $B_i$, $T(i, :)$, $Q$, $R$ and resembles the Riccati equation used in LQR for LTI systems. Then, at time $t$, the optimal input is given by $u_t = K_{\omega_t} x_t$.

To design controllers with the reduced Σ, we can first compute controller $\hat{K}_{1:r}$ by solving LQR problem with Σ as the dynamics. This requires solving $r$ coupled Riccati equations each of which is parameterized by $\hat{A}_i$, $\hat{B}_i$, $T(i, :)$, $\hat{Q}$, $\hat{R}$. To apply $\hat{K}_{1:r}$ to the original Σ, we simply let $u_t = \hat{K}_k x_t$ if $\omega_t = k$. Since the number of coupled Riccati equations is the same as the number of modes, the computation cost for Σ is $O(s)$ while only $O(r)$ for $\hat{Σ}$, thus saving is prominent when $r \ll s$. Additionally, we have the following suboptimality guarantees.

**Theorem 9** Assume system Σ is mean square stabilizable and has additive Gaussian process noise $\mathcal{N}(0, \sigma_w^2 I_n)$ that is independent of the mode switching. Assume $\Omega_{1:r} = \Omega_{1:r}$ in Algorithm 1. Let $J^*$ and $\hat{J}$ respectively denote the infinite time average cost incurred by the optimal controller $K_{1:s}$ and controller $\hat{K}_{1:r}$ (at time $t$, $u_t = \hat{K}_k x_t$ if $\omega_t = \hat{\Omega}_k$). Then, there exists constant $\bar{\epsilon}_{A,B}$, $\bar{\epsilon}_T$, $C_{A,B}$, and $C_T$, such that when $\max\{\epsilon_A, \epsilon_B\} \leq \bar{\epsilon}_{A,B}$ and $\epsilon_T \leq \bar{\epsilon}_T$, we have $\hat{J} - J^* \leq \sigma_w^2 \left(C_{A,B} \max\{\epsilon_A, \epsilon_B\} + C_T \epsilon_T\right)^2$.

Having additive noise means the MJS dynamics in (1) becomes $x_{t+1} = A_{\omega_t} x_t + B_{\omega_t} u_t + w_t$ where $w_t \sim \mathcal{N}(0, \sigma_w^2 I_n)$. Constants $\bar{\epsilon}_{A,B}$, $\bar{\epsilon}_T$, $C_{A,B}$, and $C_T$ only depend on the original Σ and cost matrices $Q$ and $R$ with the exact expressions provided in Du et al. (2022).

7. Numerical Experiments

In this section, we present experiment results to evaluate the main results in the paper. System Σ is randomly generated with desired levels of perturbation $\epsilon_A$, $\epsilon_B$, $\epsilon_T$ under the uniform partition $\Omega_{1:r}$ on $[s]$, i.e., $|\Omega_i| = |\Omega_j|$ for any $i, j \in [r]$. We first evaluate the clustering performance of Algorithm 1. Then, we evaluate the performance of LQR controller designed with the reduced MJS Σ.

To evaluate Algorithm 1, we fix $n = 5$, $p = 3$, $r = 4$ and record the average clustering error $CE = \min_{h \in H} \sum_{k \in [r]} \left| \left\{ i : i \in \Omega_k, i \notin \hat{\Omega}_{h(k)} \right\} \right|$, i.e., the number of misclustered modes under the best cluster labeling, over 100 runs. Figure 2 presents the clustering performances under different number of modes $s$ and perturbation $\epsilon_A$, $\epsilon_B$ and $\epsilon_T$. In Figure 2(a), we set perturbation $\epsilon_A = \epsilon_B$ and hyper-parameters $\alpha_A = \frac{1}{\max \{\|A\|\}}$, $\alpha_B = \frac{1}{\max \{\|A\|\}}$, and $\alpha_T = 0$ in Algorithm 1. The clustering error goes up with increasing $s$ and $\epsilon_A$, $\epsilon_B$, and we can indeed see when the perturbation is small ($< 0.5$), there are no misclustered modes. The impact of $\epsilon_T$ is shown in Figure 2(b) (lumpable case P1) and 2(c) (aggregatable case P2), where we set $\alpha_A = \alpha_B = 0, \alpha_T = 1$. 


Next we implement the idea of designing LQR controllers for the original $\Sigma$ through the reduced $\hat{\Sigma}$ as discussed in Section 6. We let $r = 3$, $n = 10$, $p = 5$, process noise variance $\sigma_w^2 = 0.1$, initial state $x_0 = 1$, and perturbation $\epsilon_A = \epsilon_B = \epsilon_T$. In Figure 3(a), the finite time LQR with horizon $T = 500$ is considered, and the plot shows time-averaged sub-optimality $(\hat{J}_t - J^*_t)/t$ averaged over 4000 runs. At $t = 500$, since the behavior has reached stationary state, $(\hat{J}_t - J^*_t)/t$ can be viewed as the infinite time suboptimality $\hat{J} - J^*$ discussed in Theorem 9. We can easily see the trend that larger perturbations result in larger suboptimality. Figure 3(b) shows the time to compute controllers via Riccati iterations using $\Sigma$ and $\hat{\Sigma}$. We see when $s$ is large, $\Sigma$ needs significantly more time than $\hat{\Sigma}$. Since the optimal infinite time LQR controllers also can be obtained through Riccati iterations until convergence, on each curve we use circles to mark the time needed when the controller difference between two adjacent iterations falls below $10^{-12}$.

8. Conclusion and Future Work

In this work, we propose a clustering-based method that reduces the mode complexity of an MJS. The reduced MJS provably approximates the original MJS in terms of trajectory difference and controller optimality. Several potential extensions include: (i) stronger approximation metrics that allow for inputs; (ii) stability of the reduced MJS; (iii) partial observation case, i.e., the state $x_t$ is observed through $y_t = C_{\omega_t}x_t$ for some mode-dependent output matrices $C_{1:s}$.

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References


Appendix A. Proof for Clustering Results

We first provide several supporting lemmas. The first result provides a perturbation result for left singular vector space.

**Lemma 10 (Singular Vectors Perturbation Bound)** Consider two arbitrary matrices $\Phi, \hat{\Phi} \in \mathbb{R}^{s \times q}$. Let $U, \hat{U} \in \mathbb{R}^{s \times r}$ respectively denote the top-$r$ left singular vectors of $\Phi$ and $\hat{\Phi}$ with $U^T U = \hat{U}^T \hat{U} = I_r$. Then

$$\min_{O \in \mathcal{O}(r)} \| \bar{U} O - U \|_F \leq \frac{2\sqrt{2} \| \Phi - \Phi \|_F}{\sigma_r(\Phi) - \sigma_{r+1}(\Phi)} ,$$

(9)

where $\mathcal{O}(r)$ denotes the orthogonal group with dimension $r$, i.e., the set of all $r$-dimensional orthonormal matrices, and $\sigma_r(\Phi)$ denotes the $r$th largest singular value of $\Phi$.

**Proof** One can simply see this by combining Lemma 10 (which can be easily shown to hold for Frobenius norm as) and Lemma 11 in Du et al. (2019). Note that we need to replace the spectrum in Lemma 10 of Du et al. (2019) with Frobenius norm to complete this proof, and it is easy to show Lemma 10 of Du et al. (2019) also holds for Frobenius norm.

Note that the basis for a subspace is only unique up to rotation transformation, hence when we study the difference of subspace, we consider the rotation that gives the best match between their corresponding basis. The next result says if a matrix has certain columns being identical, then its left singular vectors share the same identity properties.

**Lemma 11 (Lemma 12 in Du et al. (2019))** Consider a matrix $\Phi \in \mathbb{R}^{s \times q}$ and a partition $\Omega_{1:r}$ on $[s]$ such that for any $i, i' \in \Omega_k$, $\hat{\Phi}(i,:) = \hat{\Phi}(i', :).$ Assume $\text{rank}(\Phi) = r$. Let $\bar{U} \in \mathbb{R}^{s \times r}$ denote the top-$r$ left singular vectors of with $\bar{U}^T \bar{U} = I_r$. Then for any $i \in \Omega_k$ and $j \in \Omega_l$,

$$\| \bar{U}(i,:) - \bar{U}(j,:) \| = \begin{cases} 0 & \text{if } k = l \\ \sqrt{\frac{1}{|\Omega_k|} + \frac{1}{|\Omega_l|}} & \text{if } k \neq l \end{cases} .$$

(10)

This result says, in the clustering analysis, if one has no idea of the inter-cluster distance in the original feature space as given by data matrix $\Phi$, the left singular vectors can provide such distances, and the distances only depend on the cluster sizes. This itself is double-edged: no matter how close two clusters are, their corresponding rows in the left singular vectors are fixed as in (10); and when certain rows of different clusters are linearly dependant (rank assumption is violated), no matter how faraway two clusters are, this result will not hold. The differences in left singular vectors decrease with the cluster size, on the bright side, in square root sense. Another choice is matrix $\tilde{U} := \bar{U} \text{diag}(\sigma_1(\Phi), \ldots, \sigma_r(\Phi))$, i.e., the projection of $\Phi$ onto its top-$r$ right singular vector space, and in this case, (10) gives $\sigma_r(\tilde{\Phi}) \sqrt{|\Omega_k|^{-1} + |\Omega_l|^{-1}} \leq \| \bar{U}(i,:) - \bar{U}(j,:) \| \leq \sigma_1(\tilde{\Phi}) \sqrt{|\Omega_k|^{-1} + |\Omega_l|^{-1}}$ if $k \neq l$ and 0 otherwise.

The next lemma provides a preliminary result on the performance of k-means when it is applied to a data matrix with feature dimension same as the number of clusters. Note that the matrices constructed with the top-$r$ left singular vectors in Lemma 10 and Lemma 11 fall into this category.

**Lemma 12 (Lemma 5.3 in Lei et al. (2015))** Consider two arbitrary matrices $\bar{U}, \tilde{U} \in \mathbb{R}^{s \times r}$ with $\Delta \bar{U} := \| \bar{U} - \tilde{U} \|_F$. Suppose there exists a partition $\Omega_{1:r}$ on $[s]$ such that for any $i, i' \in \Omega_k$,
Define the inter-cluster distance for cluster $k$ as $\delta_k := \min_{i \in [r]} \min_{\Omega_k, j \in \Omega_k} \|\hat{U}(i, :) - \hat{U}(j, :)\|$. Let $\{\Omega_{1, r}, \hat{c}_{1, r}\}$ be an $(1 + \epsilon)$ solution to the k-means problem on $\hat{U}$. Then, when $\Delta_{\hat{U}} \leq \frac{\min_k \sqrt{\Omega_k / \delta_k}}{\sqrt{8(2 + \epsilon)}}$, we have

$$
\min_{\hat{h} \in \mathcal{H}} \sum_{k \in [r]} \left| \left\{ i : i \in \Omega_k, i \notin \hat{\Omega}_{h(k)} \right\} \right| \cdot \delta_k^2 \leq 8(2 + \epsilon) \Delta_{\hat{U}}^2,
$$

(11)

where $\mathcal{H}$ is the set of all bijections from $[r]$ to $[r]$.

Finally, by combining Lemma 10, 11, and 12, we can obtain theoretical guarantee on the performance of k-means when it is applied to the left singular vectors of the data matrix, which is the key lemma we will use to show Theorem 5 and Theorem 7.

**Lemma 13 (Approximate k-means error bound)** Consider two arbitrary matrices $\Phi, \tilde{\Phi} \in \mathbb{R}^{s \times q}$ with $\Delta_{\Phi} := \|\Phi - \tilde{\Phi}\|_F$. Suppose there exists a partition $\Omega_{1, r}$ on $[s]$ such that for any $i, i' \in \Omega_k$, $\Phi(i, :) = \Phi(i', :)$. Assume $\text{rank}(\tilde{\Phi}) = r$. Let $\hat{U} \in \mathbb{R}^{s \times r}$ denote the top-$r$ left singular vectors of $\tilde{\Phi}$ with $\hat{U}^\top \hat{U} = \mathbf{I}_r$. Let $\{\hat{\Omega}_{1, r}, \hat{c}_{1, r}\}$ be an $(1 + \epsilon)$ solution to the k-means problem on $\hat{U}$. Then, when $\Delta_{\tilde{\Phi}} \leq \frac{\alpha_r(\Phi)\sqrt{|\Omega_{1, r}| + |\Omega_{1, s}|}}{8\sqrt{(2 + \epsilon)|\Omega_{1, r}|}}$, we have

$$
\text{MR}(\hat{\Omega}_{1, r}) \leq \frac{64(2 + \epsilon)}{\sigma_{\alpha_r}(\Phi)} \Delta_{\tilde{\Phi}}^2.
$$

(12)

**Proof** Let $\bar{U} \in \mathbb{R}^{s \times r}$ denote the top-$r$ left singular vectors of $\Phi$ with $\bar{U}^\top \bar{U} = \mathbf{I}_r$. Applying Lemma 11 and noticing that $\text{rank}(\bar{U}) = r$, we know there exists $\mathbf{O}^* \in \mathcal{O}(r)$ such that

$$
\|\bar{U} \mathbf{O}^* - \bar{U}\|_F \leq \frac{2\sqrt{2}\Delta_{\Phi}}{\sigma_{\alpha_r}(\Phi)}
$$

(13)

Note that $\|\bar{U} \mathbf{O}^*(i, :) - [\bar{U} \mathbf{O}^*](j, :)\| = \|(\hat{U}(i, :) - \hat{U}(j, :)) \mathbf{O}^*\| = \|\hat{U}(i, :) - \bar{U}(j, :)|. Then applying Lemma 11, we know for any $i \in \Omega_k$ and $j \in \Omega_l$

$$
\|\bar{U} \mathbf{O}^*(i, :) - [\bar{U} \mathbf{O}^*](j, :)\| = \begin{cases} 0 & \text{if } k = l, \\ \sqrt{\frac{1}{|\Omega_k|}} + \frac{1}{|\Omega_l|} & \text{if } k \neq l. \end{cases}
$$

(14)

Then, for any $k \in [r]$, let $\delta_k := \min_{i \in [r]} \min_{\Omega_k, j \in \Omega_k} \|\bar{U} \mathbf{O}^*(i, :) - [\bar{U} \mathbf{O}^*](j, :)\|$, we see $\delta_k \geq \sqrt{\frac{1}{|\Omega_k|}} + \frac{1}{|\Omega_k|}$. With (14), we can apply Lemma 12 to matrix $\{\bar{U} \mathbf{O}^*, \bar{U}\}$ and obtain the following: when $\|\bar{U} \mathbf{O}^* - \bar{U}\|_F \leq \frac{\min_k \sqrt{|\Omega_k|/\delta_k}}{\sqrt{8(2 + \epsilon)}}$, we have $\min_{\hat{h} \in \mathcal{H}} \sum_{k=1}^r \|\{ i : i \in \Omega_k, i \notin \hat{\Omega}_{h(k)} \} \| \cdot \delta_k^2 \leq 8(2 + \epsilon)\|\bar{U} \mathbf{O}^* - \bar{U}\|_F^2 \leq \frac{64(2 + \epsilon)}{\sigma_{\alpha_r}(\Phi)} \Delta_{\tilde{\Phi}}^2$. Finally, plugging in the upper bound for $\|\bar{U} \mathbf{O}^* - \bar{U}\|_F$ in (13) and the bounds for $\delta_k$, this gives that when $\Delta_{\tilde{\Phi}} \leq \frac{\alpha_r(\Phi)\sqrt{|\Omega_{1, r}| + |\Omega_{1, s}|}}{8\sqrt{(2 + \epsilon)|\Omega_{1, r}|}}$, we can guarantee

$$
\text{MR}(\hat{\Omega}_{1, r}) = \min_{\hat{h} \in \mathcal{H}} \sum_{k=1}^r \|\{ i : i \in \Omega_k, i \notin \hat{\Omega}_{h(k)} \} \| \cdot \frac{1}{|\Omega_k|} \leq \min_{\hat{h} \in \mathcal{H}} \sum_{k=1}^r \|\{ i : i \in \Omega_k, i \notin \hat{\Omega}_{h(k)} \} \| \cdot \frac{1}{|\Omega_k|} \leq \frac{64(2 + \epsilon)}{\sigma_{\alpha_r}(\Phi)} \Delta_{\tilde{\Phi}}^2.
$$

\[\blacksquare\]
A.1. Proof for Theorem 5

Proof [Proof for Theorem 5] Recall \( \Phi \) defined in Line 3 of Algorithm 1 and its averaged version \( \bar{\Phi} \) defined in Section 4.1. Let \( i' := \arg \max_{i \in [s]} \| \Phi(i, :) - \Phi(i, :) \| \) and suppose \( i' \in \Omega_k \). Then, we have

\[
\| \bar{\Phi} - \Phi \|_F \leq \sqrt{s} \left( \alpha_A \| A_{i'} - \frac{1}{|\Omega_k|} \sum_{i \in \Omega_k} A_i \|_F + \alpha_B \| B_{i'} - \frac{1}{|\Omega_k|} \sum_{i \in \Omega_k} B_i \|_F + \alpha_T \| T(i', :) - \frac{1}{|\Omega_k|} \sum_{i \in \Omega_k} T(i, :) \| \right) \]  (15)

By definition of \( \epsilon_A, \epsilon_B, \) and \( \epsilon_T \) in Problem P2 and triangle inequality, we have

\[
\| \bar{\Phi} - \Phi \|_F \leq \sqrt{s} \left( \alpha_A \epsilon_A + \alpha_B \epsilon_B + \alpha_T \epsilon_T \right) = \sqrt{s} \epsilon_{Agg}, \]  (16)

where \( \epsilon_{Agg} := \alpha_A \epsilon_A + \alpha_B \epsilon_B + \alpha_T \epsilon_T \). By construction, in \( \bar{\Phi} \), rows that belong to the same cluster have the same rows, thus we could apply Lemma 13 to \{ \( \Phi, \hat{\Phi} \) \} and obtain that when \( \sqrt{s} \epsilon_{Agg} \leq \frac{\sigma_{r}(\Phi)\sqrt{|\Omega(1)| + |\Omega(2)|}}{8\sqrt{2+\epsilon}|\Omega(1)|} \), we have \( \text{MR}(\Omega_{1:r}) \leq \frac{64(2+\epsilon)^{s}2}{\sigma_{r}(\Phi)^{2}} \epsilon_{Agg} \). \( \blacksquare \)

A.2. Proof for Theorem 7

We first provide a perturbation result regarding Markov matrices.

Lemma 14 (Section 3.6 in Cho and Meyer (2001)) Consider two Markov matrices \( \mathbf{T}, \mathbf{T}_0 \in \mathbb{R}^{s \times s} \) and their stationary distributions \( \pi, \pi_0 \in \mathbb{R}^s \), then \( \| \pi - \pi_0 \|_1 \leq \gamma_1 \| \mathbf{T} - \mathbf{T}_0 \|_{\infty} \), where \( \gamma_1 := \sum_{i=2}^{s} \frac{1}{1-\lambda_i(\mathbf{T})} \).

When the difference \( \| \mathbf{T} - \mathbf{T}_0 \| \) is small, we can further have the following corollary

Corollary 15 In Lemma 14, let \( \pi_{\min} := \min_i \pi(i), \pi_{\max} := \max_i \pi(i) \). Suppose \( \| \mathbf{T} - \mathbf{T}_0 \|_{\infty} \leq \frac{\pi_{\min}}{\gamma_1} \), then we have

\[
\max_i |\pi(i) - \pi_0(i)| \leq \frac{\pi_{\min}}{2}, \quad \min_i \pi_0(i) \geq \frac{\pi_{\min}}{2}, \quad \max_i \pi_0(i) \leq \pi_{\max} + \frac{\pi_{\min}}{2} \]  (17)

\[
\max_i |\pi(i)^{-\frac{1}{2}} - \pi_0(i)^{-\frac{1}{2}}| \leq (\sqrt{2} - 1)\gamma_1 \frac{\pi_{\min}}{2} \| \mathbf{T} - \mathbf{T}_0 \|_{\infty} \]  (18)

\[
\max_i |\pi(i)^{\frac{1}{2}} - \pi_0(i)^{\frac{1}{2}}| \leq \left( 1 - \frac{\sqrt{2}}{2} \right)\gamma_1 \frac{\pi_{\min}}{2} \| \mathbf{T} - \mathbf{T}_0 \|_{\infty} \]  (19)

Proof Since \( 1^T \pi = 1^T \pi_0 = 1 \), we have \( \max_i |\pi(i) - \pi_0(i)| \leq \frac{1}{2} \| \pi - \pi_0 \|_1 \leq \frac{\pi_{\min}}{2} \| \mathbf{T} - \mathbf{T}_0 \|_{\infty} \leq \frac{\pi_{\min}}{2} \). Then using triangle inequality, we could show (17). Note that the LHS of (18) is equivalent to \( \max_i \frac{|\pi_0(i) - \pi(i)|}{\sqrt{\pi(i)\pi_0(i)}} \), then plugging in the results in (17) can show (18). Note that the LHS of (19) is equivalent \( \max_i \frac{|\pi(i) - \pi_0(i)|}{\sqrt{\pi(i) + \sqrt{\pi_0(i)}}} \), then plugging in the results in (17) can show (19). \( \blacksquare \)
Lemma 16  Consider an ergodic Markov matrix $T \in \mathbb{R}^{s \times s}$ with stationary distribution $\pi$ and a partition $\Omega_{1:r}$ such that it is approximately lumpable as in (4) with perturbation $\epsilon_T$. Assume there exists an ergodic and reversible $T_0 \in \mathcal{L}(T, \Omega_{1:r}, \epsilon_T)$ that has informative spectrum. Using $T$ and $\pi$, construct $S_r \in \mathbb{R}^{s \times r}$ as in Line 7 of Algorithm 1. And define $S_r \in \mathbb{R}^{s \times r}$ such that for any $i \in [s]$ (suppose $i \in \Omega_k$), $S_r(i,:):=\frac{1}{|\Omega_k|} \sum_{i' \in \Omega_k} S_r(i',:).$ Let $\pi_{\min} := \max_i \pi(i)$, $\pi_{\max} := \min_i \pi(i)$, $\gamma_1 := \sum_{i=2}^{s} \frac{1}{\text{tr}(T_i^i)},$ and $\gamma_2 := \min \{\sigma_r(H) - \sigma_{r+1}(H), 1\}$ where $H$ is defined in Algorithm 1. Assume perturbation $\epsilon_T \leq \frac{\pi_{\min}}{\gamma_1}$. Then, for any $i \in [s]$ (suppose $i \in \Omega_k$), we have

$$\|S_r - \bar{S}_r\|_F \leq \frac{16\sqrt{\gamma_1 \gamma_2}}{\sqrt{|\Omega_k|}} \|T\|_F \epsilon_T. \tag{20}$$

Proof  We will start with analyzing $T_0$ and use it as a bridge to prove the claim. Let $\pi_0 \in \mathbb{R}^s$ denote the stationary distribution of $T_0$. Since $T_0$ is ergodic, we know $\pi_0$ is strictly positive. By definition of reversibility, we know $\text{diag}(\pi_0)T_0 = T_0^\dagger \text{diag}(\pi_0)$, and this further gives $\text{diag}(\pi_0)^{\frac{1}{2}} T_0 \text{diag}(\pi_0)^{-\frac{1}{2}} = \text{diag}(\pi_0)^{-\frac{1}{2}} T_0^\dagger \text{diag}(\pi_0)^{-\frac{1}{2}}$. Let $H_0 := \text{diag}(\pi_0)^{-\frac{1}{2}} T_0^\dagger \text{diag}(\pi_0)^{-\frac{1}{2}}$, then we see $H_0$ is symmetric. Let $W_{0,r} \in \mathbb{R}^{s \times r}$ denote the top $r$ left singular vectors of $H_0$, by spectrum theorem, we know the columns of $W_{0,r}$ also gives the top $r$ eigenvectors of $H_0$. Let $S_{0,r} := \text{diag}(\pi_0)^{-\frac{1}{2}} W_{0,r}$, by definition of $H_0$, we see the columns of $S_{0,r}$ are the top $r$ eigenvectors of $T_0$. Then, by Lemma 6 and the definition of informative spectrum, we know for any $i, i' \in \Omega_k$, we have $S_{0,r}(i,:) = S_{0,r}(i',:)$.

Recall in Algorithm 1, $W_r$ denotes the top $r$ left singular vectors of $H := \text{diag}(\pi)^{\frac{1}{2}} T \text{diag}(\pi)^{\frac{1}{2}}$ and let $S_r := \text{diag}(\pi)^{-\frac{1}{2}} W_r$. Let $O^* := \min_{O \in \mathcal{O}(r)} \| W_{0,r} O - W_r \|_F$, where $\mathcal{O}(r)$ is the set of all $r \times r$ orthonormal matrices. Similarly, we have for $i, i' \in \Omega_k$, we have $[S_{0,r}O^*](i,:) = [S_{0,r}O^*](i',:).$ Using this property, for any $i \in [s]$ (suppose $i \in \Omega_k$), we have

$$S_r(i,:) - \bar{S}_r(i,:) = \frac{|\Omega_k| - 1}{|\Omega_k|} S_r(i,:) - \frac{1}{|\Omega_k|} \sum_{i', i'' \in \Omega_k, i' \neq i, i'' \neq i} S_r(i',:).$$

(21)

Now, WLOG, we assume $\{1, \ldots, |\Omega_1|\} = \Omega_1$, $\{1, \ldots, |\Omega_1| + 1, \ldots, |\Omega_1| + |\Omega_2|\} = \Omega_2$, \ldots and define block diagonal matrices $D, P \in \mathbb{R}^{s \times s}$ both with $r$ diagonal blocks such that their $k$-th diagonal blocks $[D]_k, [P]_k \in \mathbb{R}^{[\Omega_k] \times [\Omega_k]}$ are given by

$$[D]_k = \frac{|\Omega_k| - 1}{|\Omega_k|} I_{[\Omega_k]}, \quad [P]_k = \frac{1}{|\Omega_k|} (I_{[\Omega_k]} - I_{[\Omega_k]})^T I_{[\Omega_k]} - I_{[\Omega_k]}).$$

(22)

Then, stacking (21), we see $S_r - \bar{S}_r = D(S_r - S_{0,r}O^*) + P(S_{0,r}O^* - S_r)$. Note that for a arbitrary matrix $E$, we have $\|PE\|_F^2 = \text{tr}(P^T EEG^T) \leq \text{tr}(D^T EEG^T) = \|DE\|_F^2$ where the inequality holds since for each diagonal block we have $[P]_k^T [P]_k \leq [D]_k^T [D]_k$. Therefore,

$$\|S_r - \bar{S}_r\|_F \leq 2\|D(S_r - S_{0,r}O^*)\|_F + 2\max_k \left| \frac{|\Omega_k| - 1}{|\Omega_k|} \right| \|S_r - S_{0,r}O^*\|_F \leq 2\|S_r - S_{0,r}O^*\|_F. \tag{23}$$
Note that when the aforementioned partition assumption "\{1, \ldots, |\Omega_1|\} = \Omega_1, \ldots" is not true, one can simply left multiply matrix \(S_r - \bar{S}_r\) by a row permutation matrix such that after permutation, row 1 to \(|\Omega_1|\) belong to \(\Omega_1\), etc. Then, (23) still holds since Frobenius norm is invariant under row permutations. From (23), we see to complete the proof, it suffices to study \(\|S_r - S_{0,r}O^*\|_F\).

\[
\|S_r - S_{0,r}O^*\|_F = \|\text{diag}(\pi)^{-\frac{1}{2}}(W_r - W_{0,r}O^*) + \text{diag}(\pi)^{-\frac{1}{2}} - \text{diag}(\pi_0)^{-\frac{1}{2})W_{0,r}O^*\|_F
\leq \frac{1}{\sqrt{\pi_{\min}}} \|W_r - W_{0,r}O^*\|_F + \sqrt{r} \max_i |\pi(i)^{\frac{1}{2}} - \pi_0(i)^{\frac{1}{2}}|
\]

In (24), for term \(\|W_r - W_{0,r}O^*\|_F\), according to Lemma 10, we know \(\|W_r - W_{0,r}O^*\|_F \leq \frac{2\sqrt{2}}{\sigma_1(H) - \sigma_{r+1}(H)} \|H - H_0\|_F\). Plugging this result and the upper bound for \(\max_i |\pi(i)^{\frac{1}{2}} - \pi_0(i)^{\frac{1}{2}}|\) in (18) into (24), we have

\[
\|S_r - S_{0,r}O^*\|_F \leq \frac{2\sqrt{2}}{(\sigma_1(H) - \sigma_{r+1}(H))\pi_{\min}^{0.5}} \|H - H_0\|_F + \frac{(\sqrt{2} - 1)\gamma_1\sqrt{r}}{\pi_{\min}^{1.5}} \epsilon_T.
\]

Next, we evaluate term \(\|H - H_0\|_F\). By definitions of \(H\) and \(H_0\), we have

\[
\|H - H_0\|_F \leq \|\text{diag}(\pi)^{\frac{1}{2}} - \text{diag}(\pi_0)^{\frac{1}{2})\cdot T\cdot \text{diag}(\pi)^{-\frac{1}{2}}\|_F
+ \|\text{diag}(\pi_0)^{\frac{1}{2}}\cdot T\cdot (\text{diag}(\pi)^{-\frac{1}{2}) - \text{diag}(\pi_0)^{-\frac{1}{2})}\|_F
+ \|\text{diag}(\pi_0)^{\frac{1}{2}}\cdot (T - T_0)\cdot \text{diag}(\pi_0)^{-\frac{1}{2})\|_F
\leq (\max_i |\pi(i)^{\frac{1}{2}} - \pi_0(i)^{\frac{1}{2})|) \cdot \|T\|_F \cdot (\min_i \pi(i)^{\frac{1}{2}})
+ (\max_i \pi_0(i)^{\frac{1}{2}) \cdot \sqrt{s} \|T - T_0\|_\infty \cdot (\min_i \pi_0(i)^{\frac{1}{2}})
\]

Plugging in the results in Corollary 15 gives

\[
\|H - H_0\|_F \leq 2.56\gamma_1 \sqrt{s} \pi_{\max}^{0.5} \pi_{\min}^{-1.5} \|T\|_F \epsilon_T.
\]

Applying (27) to (25), we see

\[
\|S_r - S_{0,r}O^*\|_F \leq \frac{8\gamma_1 \sqrt{s} \sqrt{s} \pi_{\max} \|T\|_F \epsilon_T}{\gamma_2 \pi_{\min}^{2}}
\]

where \(\gamma_2 := \min\{\sigma_1(H) - \sigma_{r+1}(H), 1\}\). Finally, by plugging (28) into (23), we have \(\|S_r - \bar{S}_r\|_F \leq \frac{16\gamma_1 \sqrt{s} \sqrt{s} \pi_{\max} \|T\|_F \epsilon_T}{\gamma_2 \pi_{\min}^{2}}\), which concludes the proof.

Now we are ready to present the main proof for Theorem 7.

\textbf{Proof [Proof for Theorem 7]} Recall \(\Phi\) defined in Line 8 of Algorithm 1 and its averaged version \(\bar{\Phi}\) defined in Section 4.1. Let \(i' := \arg \max_{i \in [s]} \|A - \Phi(i, n^2 + np + 1 : n^2 + np + r)\|\) and suppose \(i' \in \Omega_k\). Then, we have

\[
\|\Phi - \bar{\Phi}\|_F \leq \sqrt{s} \left(\alpha_A \left\|A_{i'} - \frac{1}{|\Omega_k|} \sum_{i \in \Omega_k} A_i\right\|_F + \alpha_B \left\|B_{i'} - \frac{1}{|\Omega_k|} \sum_{i \in \Omega_k} B_i\right\|_F\right) + \alpha_T \|S_r - \bar{S}_r\|_F,
\]
where \( \tilde{S} \) is defined the same as in Lemma 16. By definition of \( \epsilon_A \) and \( \epsilon_B \) in Problem P1 and Lemma 16, we have

\[
\| \tilde{\Phi} - \Phi \|_F \leq \sqrt{s} \left( \alpha_A \epsilon_A + \alpha_B \epsilon_B + \alpha_T \frac{16\gamma_1 \sqrt{\pi \max} \| T \|_F}{\gamma_2 \pi_\min^2} \epsilon_T \right) = \sqrt{s} \epsilon_{Lmp}
\]

(30)

where \( \epsilon_{Lmp} := \alpha_A \epsilon_A + \alpha_B \epsilon_B + \alpha_T \frac{16\gamma_1 \sqrt{\pi \max} \| T \|_F}{\gamma_2 \pi_\min^2} \epsilon_T \). By construction, in \( \tilde{\Phi} \), rows that belong to the same cluster have the same rows, thus we could apply Lemma 13 to \( \{ \tilde{\Phi}, \Phi \} \) and obtain that when \( \sqrt{s} \epsilon_{Lmp} \leq \frac{\gamma_0(\Phi)\sqrt{\max(\|\Omega_1\|, \|\Omega_2\|)}}{8\sqrt{2+\epsilon})\|\Omega_1\|} \), we have \( \text{MR} (\hat{\Omega}_{1,\tau}) \leq \frac{64(2+\epsilon)s_0^2}{\sigma_T^2(\Phi)} \epsilon_{Lmp} \). \( \blacksquare \)

Appendix B. Proof for Approximation Guarantees Theorem 8

We first provide several supporting lemmas.

Lemma 17 (Mania et al. (2019), Lemma 5) Consider two matrices \( \mathcal{A} \) and \( \tilde{\mathcal{A}} \) with \( \| \mathcal{A} - \tilde{\mathcal{A}} \| \leq \epsilon_A \). Let \( \rho \) be the spectral radius of \( \mathcal{A} \), and let \( \tau := \sup_{k \in \mathbb{N}} \| \tilde{\mathcal{A}}^k \| / \rho^k \). Then, we have

\[
\| \tilde{\mathcal{A}}^t \| \leq \tau (\epsilon_A + \rho)^t, \quad \| \mathcal{A}^t - \tilde{\mathcal{A}}^t \| \leq \tau^2 (\epsilon_A + \rho)^{t-1} \epsilon_A.
\]

(31)

Lemma 18 (Sattar et al. (2021)) Consider an autonomous MJS(\( A_{1,s}, 0, T \)). Define matrix \( \mathcal{A} \in \mathbb{R}^{sn^2 \times sn^2} \) with its \( ij \)-th \( n^2 \times n^2 \) block given by \( [\mathcal{A}]_{ij} := T(j, i) \cdot A_j \otimes A_j \). Let \( \Sigma_t^{(i)} := \mathbb{E}[x_t x_t^\top 1(\omega_t = i)] \) and \( s_t := [\text{vec}(\Sigma_t^{(1)})^\top, \ldots, \text{vec}(\Sigma_t^{(s)})^\top]^\top \). Then, we have

\[
s_t = \mathcal{A}^t s_0
\]

(32)

Recall in Sec 5, for \( \Sigma \), i.e., MJS(\( A_{1,s}, B_{1,s}, T \)), we define the augmented state matrix \( \tilde{\mathcal{A}} \in \mathbb{R}^{sn^2 \times sn^2} \) with its \( ij \)-th \( n^2 \times n^2 \) block given by \( [\tilde{\mathcal{A}}]_{ij} := T(j, i) \cdot A_j \otimes A_j \). We let \( \rho \) denote the spectral radius of \( \tilde{\mathcal{A}} \) and let \( \tau := \sup_{k \in \mathbb{N}} \| \tilde{\mathcal{A}}^k \| / \rho^k \). The next lemma is regarding the stability of the augmentation of two MJS with the same \( \mathcal{A} \) matrix.

Lemma 19 Construct matrix \( \tilde{\mathcal{A}} \in \mathbb{R}^{4sn^2 \times 4sn^2} \) with its \( ij \)-th \( 4n^2 \times 4n^2 \) block given by \( [\tilde{\mathcal{A}}]_{ij} := T(j, i) \cdot \begin{bmatrix} A_j & A_j \\ A_j & A_j \end{bmatrix} \). Then, (i) \( \tilde{\mathcal{A}} \) has spectral radius \( \rho \), (ii) \( \| \tilde{\mathcal{A}}^k \| = \| \mathcal{A}^k \| \), and (iii) \( \sup_{k \in \mathbb{N}} \| \tilde{\mathcal{A}}^k \| / \rho^k = \tau \).

Proof One can simply see these results by noticing there exists a permutation matrix \( P \) such that \( P \tilde{\mathcal{A}} P^\top = I_4 \otimes \mathcal{A} \), where \( I_4 \) denotes the \( 4 \times 4 \) identity matrix. \( \blacksquare \)

Now we are ready to present the proof for Theorem 8.

Proof [Proof for Theorem 8] First, we construct two autonomous switched systems:

\[
\tilde{\Pi} := \begin{cases} \dot{x}_{t+1} = \tilde{\mathcal{A}} \omega_t x_t, \quad \omega_{t+1} = \omega_T \end{cases}, \quad \Pi := \begin{cases} \dot{x}_{t+1} = \mathcal{A} \omega_t x_t, \quad \omega_{t+1} = \omega_T \end{cases}
\]

(33)

where \( \omega_t \) is the mode index for the original MJS \( \Sigma \), and for \( i \in [s] \) (suppose \( i \in \Omega_k \)),

\[
\tilde{\mathcal{A}}_i = \begin{bmatrix} A_i & A_i \\ A_i & A_i \end{bmatrix}, \quad \mathcal{A}_i = \begin{bmatrix} A_i & \tilde{A}_k \end{bmatrix}.
\]

(34)
Since $\omega_t$ of $\Sigma$ follows Markov chain $T$, systems $\bar{\Pi}$ and $\bar{\Pi}$ can be viewed as $\text{MJS}(\bar{\mathbf{A}}_{1,s}, 0, \bar{T})$ and $\text{MJS}(\bar{\mathbf{A}}_{1,s}, 0, \bar{T})$ respectively with $\bar{T} = \bar{T}$.

$\bar{\Pi} = \text{MJS}(\bar{\mathbf{A}}_{1,s}, 0, \bar{T})$ and $\bar{\Pi} = \text{MJS}(\bar{\mathbf{A}}_{1,s}, 0, \bar{T})$ where $\bar{T} = \bar{T} = \bar{T}$ and for all $i \in \Omega_t$. We then define observations for $\bar{\Pi}$ and $\bar{\Pi}$: $y_t = C\bar{x}_t$ and $\bar{y}_t = C\bar{x}_t$ where $C = C = [I_n, -I_n]$. We set their initial states as $\bar{x}_0 = [x_0^T, x_0^T]^T$, $\bar{x}_0 = [x_0^T, \bar{x}_0^T]^T$, where $x_0$ and $\bar{x}_0$ are the initial states of $\Sigma$ and $\bar{\Sigma}$ respectively. Due to the shared initial conditions for $\Sigma$ and $\bar{\Sigma}$, we have, for all $t$, $\bar{x}_t = [x_t^T, \bar{x}_t^T]^T$ and $\bar{x}_t = [x_t^T, \bar{x}_t^T]^T$, thus $y_t = 0$ and $\bar{y}_t = x_t - \bar{x}_t$. We define $\bar{\Sigma}_t = E[\bar{x}_t, \bar{x}_t]^T$ and $\bar{\Sigma}_t = E[\bar{x}_t, \bar{x}^T_t]$, then we have $E[\|x_t - \bar{x}_t\|^2] = E[\|\bar{y}_t, \bar{y}_t\|] = E[\|\bar{y}_t, \bar{y}_t\|] = \text{tr}(C^T \bar{\Sigma}_t) - \text{tr}(C^T \bar{\Sigma}_t) = \text{tr}(C^T(C(\bar{\Sigma}_t - \bar{\Sigma}_t)))$. Since $C^T C \succeq 0$, we further have

$$E[\|x_t - \bar{x}_t\|^2] \leq \text{tr}(C^T C)\|\bar{\Sigma}_t - \bar{\Sigma}_t\| = 2n\|\bar{\Sigma}_t - \bar{\Sigma}_t\|. \quad (35)$$

We let $\bar{\Sigma}_t^{(i)} := E[\bar{x}_t, \bar{x}_t^T 1_{\{\omega_t = i\}}]$, $\bar{\Sigma}_t^{(i)} := E[\bar{x}_t, \bar{x}_t^T 1_{\{\omega_t = i\}}]$, $\bar{s}_t := [\vec{\Sigma}_t^{(1)}, \ldots, \vec{\Sigma}_t^{(n)}]^T$ and $\bar{s}_t := [\vec{\Sigma}_t^{(1)}, \ldots, \vec{\Sigma}_t^{(n)}]^T$. Note that $\vec{\Sigma}_t = [I_{4n^2}, \ldots, I_{4n^2}]\bar{s}_t$ and $\vec{\Sigma}_t = [I_{4n^2}, \ldots, I_{4n^2}]\bar{s}_t$, thus we have $\|\bar{\Sigma}_t - \bar{\Sigma}_t\| = \|\bar{\Sigma}_t - \bar{\Sigma}_t\| = \sqrt{s}\|\bar{s}_t - \bar{s}_t\|$. Plugging this into (35), we have

$$E[\|x_t - \bar{x}_t\|^2] \leq 2n\sqrt{s}\|\bar{s}_t - \bar{s}_t\|. \quad (36)$$

By Lemma 18, we have $\bar{s}_t = \bar{\mathcal{A}}^t\bar{s}_0$ and $\bar{s}_t = \bar{\mathcal{A}}^t\bar{s}_0$, where $\bar{\mathcal{A}} \in \mathbb{R}^{4n^2 \times 4n^2}$ is constructed such that its $ij$-th $4n^2 \times 4n^2$ block given by $[\bar{\mathcal{A}}]_{ij} = \bar{T}(j, i)\bar{\mathcal{A}}_j \otimes \bar{\mathcal{A}}_j$, and $\bar{\mathcal{A}}$ is constructed similarly. Since $\bar{s}_0 = \bar{s}_0 = [\vec{\Sigma}_t^{(1)}, \ldots, \vec{\Sigma}_t^{(n)}]^T$, $\vec{\Sigma}_t = [I_{4n^2}, \ldots, I_{4n^2}]\bar{s}_t$ and $\vec{\Sigma}_t = [I_{4n^2}, \ldots, I_{4n^2}]\bar{s}_t$, thus we have $\|\bar{\Sigma}_t - \bar{\Sigma}_t\| = \|\bar{\Sigma}_t - \bar{\Sigma}_t\| = \sqrt{s}\|\bar{s}_t - \bar{s}_t\|$. Plugging this into (35), we have

$$E[\|x_t - \bar{x}_t\|^2] \leq 4n\sqrt{s}\|\bar{s}_t - \bar{s}_t\|. \quad (37)$$

Now, it only suffices to evaluate $\|\bar{\mathcal{A}}^t - \bar{\mathcal{A}}^t\|$. When $t = 1$, we have

$$\|\bar{\mathcal{A}} - \bar{\mathcal{A}}\| = s\max_i \|\bar{\mathcal{A}}_{i1}, \ldots, \bar{\mathcal{A}}_{i8}\| - s\max_i \|\bar{\mathcal{A}}_{i1}, \ldots, \bar{\mathcal{A}}_{i8}\|$$

$$= s\max_i \|\bar{T}(1, i) \cdot (\bar{\mathcal{A}}_1 \otimes \bar{\mathcal{A}}_1 - \bar{\mathcal{A}}_1 \otimes \bar{\mathcal{A}}_1), \ldots, \bar{T}(s, i) \cdot (\bar{\mathcal{A}}_s \otimes \bar{\mathcal{A}}_s - \bar{\mathcal{A}}_s \otimes \bar{\mathcal{A}}_s)\|$$

$$\leq s\max_i \|\bar{\mathcal{A}}_1 \otimes \bar{\mathcal{A}}_1 - \bar{\mathcal{A}}_1 \otimes \bar{\mathcal{A}}_1\| ||T(1, i)\|$$

$$\leq s\max_j \|\bar{\mathcal{A}}_j \otimes \bar{\mathcal{A}}_j - \bar{\mathcal{A}}_j \otimes \bar{\mathcal{A}}_j\|$$

$$\leq s\max_j \|\bar{\mathcal{A}}_j \otimes (\bar{\mathcal{A}} - \bar{\mathcal{A}}_j) + (\bar{\mathcal{A}}_j - \bar{\mathcal{A}}_j) \otimes \bar{\mathcal{A}}_j\|$$

$$\leq s\max_j \|\bar{\mathcal{A}}_j \|(\|\bar{\mathcal{A}}_j - \bar{\mathcal{A}}_j\|) + \|\bar{\mathcal{A}}_j - \bar{\mathcal{A}}_j\| \|\bar{\mathcal{A}}_j\|.$$  

(38)
where the second inequality holds since we assume $\epsilon_A \leq \frac{1-\rho}{6\sqrt{\pi} \rho}$ and let $\rho_0 := \frac{1+\rho}{2}$. Finally, plugging this upper bound for $\|A^i - A^f\|$ into (37), we have

$$\mathbb{E}[\|x_t - \hat{x}_t\|^2] \leq 12nst\rho_0^{-1} - T^2 \|x_0\|^2 \epsilon_A,$$

which concludes the proof.

### Appendix C. Proof for Suboptimality Guarantees Theorem 9

In this appendix, for completeness, we provide not only the proof for the infinite time LQR result Theorem 9, but also the discussion for the finite time case. In LQR problems, given positive definite cost matrices $Q$ and $R$, we define quadratic cost $J_T = \mathbb{E}\left[\sum_{t=0}^{T-1} (x_t^\top Q x_t + u_t^\top R u_t) + x_T^\top Q x_T\right]$, and the goal is to design inputs $u_{t:T-1}$ to minimize $J_T$ in the finite time horizon setting and $\limsup_{T \to \infty} \frac{1}{T} J_T$ in the infinite setting under MJS dynamics $\Sigma$. To ease the exposition, we let $S_i^+ := \{X_{1:i} \in [s], X_t \in \mathbb{R}^{nxn}, X_i \succeq 0\}$ and define the following mappings for $X_{1:i} \in S_i^+$: for $i \in [s]$, let

$$\varphi_i(X_{1:i}) := \sum_{j \in [s]} T(i,j) X_j \quad (40)$$

$$\mathcal{R}_i(X_{1:i}) := Q + A_i^\top \varphi_i(X_{1:i}) A_i - A_i^\top \varphi_i(X_{1:i})^\top B_i (R + B_i^\top \varphi_i(X_{1:i}) B_i)^{-1} B_i^\top \varphi_i(X_{1:i}) A_i \quad (41)$$

$$\mathcal{K}_i(X_{1:i}) := -(R + B_i^\top \varphi_i(X_{1:i}) B_i)^{-1} (B_i^\top \varphi_i(X_{1:i}) A_i) \quad (42)$$

According to Costa et al. (2006), to compute the optimal inputs for the finite time horizon LQR, we first compute a set of matrices $P^{(t)}_{1:i}$ for every $t$ using backward iteration: $\forall i \in [s]$, $P^{(T)}_{i} = Q$, and for $t = T-1, T-2, \ldots$, $P^{(t)}_{i} = \mathcal{R}_i(P^{(t+1)}_{1:i})$, then, if $\omega_t = i$ at time $t$, we let input $u_t = K^{(t)}_{i} x_t$, where $K^{(t)}_{i} = \mathcal{K}_i(P^{(t+1)}_{1:i})$. For the infinite horizon LQR problem, to guarantee its solvability, we assume the following.

**Assumption A1** $\Sigma$ is mean-square stabilizable. Cost matrices $Q \succ 0$ and $R \succ 0$.

Under A1, the solution to the infinite time horizon LQR is the following: we first solve for the coupled Riccati equations $P_t = \mathcal{R}_t(P_{1:s}), \forall t \in [s]$, and then if $\omega_t = i$ at time $t$, we let input $u_t = K_i x_t$ where $K_i = \mathcal{K}_i(P_{1:i})$. Particularly, the Riccati solution $P_{1:s}$ is unique among $S_i^+$ with $P_i \succ 0$ for all $i$, and $K_{1:s}$ stabilizes $\Sigma$. It is also known from Costa et al. (2006) that as the time horizon goes to infinity, the solution to the finite time setting converges to the infinite time setting, i.e., as $T \to \infty$, for any finite $t$, $P^{(T-t)}_{1:i} = P_{1:i}$ and $K^{(T-t)}_{1:i} = K_{1:i}$. Thus, one way to solve for coupled Riccati equations is to apply $\mathcal{R}_{1:i}$ recursively until convergence.

As we discussed at the beginning of this section, we may also use the reduced MJS $\hat{\Sigma}$ as a surrogate to compute controllers for $\Sigma$. To do this, we solve for the LQR controller $\hat{K}^{(1:T)}_{1:r}$ for the finite time setting (or $\hat{K}^{(1:T)}_{1:r}$ for the infinite time setting) in the same way as $K^{(1:T)}_{1:r}$ except replacing the MJS $\Sigma$ with $\hat{\Sigma}$. To apply these controllers to the original MJS $\Sigma$, if at time $t$, $\omega_t \in \Omega_k$, we let $u_t = K^{(t)}_k x_t$ for the finite time setting (or $u_t = K_k x_t$ for the infinite time setting). While
solving the LQR problems, each iteration computes a $K$ matrix for every mode, thus the complexity for $\Sigma$ is $O(s)$ while only $O(r)$ for $\hat{\Sigma}$.

Next, we analyze the sub-optimality when applying controllers computed with $\hat{\Sigma}$. To begin with, we first construct a fictional MJS by expanding the reduced $\Sigma$: we let $\hat{\Sigma} := \text{MJS}(\hat{A}_{1:s}, \hat{B}_{1:s}, \hat{T})$ such that $T \in \mathcal{L}(T, \hat{\Omega}_{1:r}, \hat{\tau}_T)$, and for all $i \in [s]$ (suppose $i \in \hat{\Omega}_k$), $\hat{A}_i = \hat{A}_k$, $\hat{B}_i = \hat{B}_k$. Note that by definition of $\mathcal{L}(T, \hat{\Omega}_{1:r}, \hat{\tau}_T)$, we can find $T$ via solving a feasibility linear programming problem. Particularly, if it is the aggregatable case $P2$, if suffices to let $T(i,:) := [\hat{\Omega}_k]^{-1} \sum_{i \in \hat{\Omega}_k} T(i,:)$ if $i \in \hat{\Omega}_k$. Note that by construction, $\hat{\Sigma}$ is mode-reducible with respect to $\hat{\Omega}_{1:r}$ and can be reduced to $\hat{\Sigma}$. According to Proposition 4, $\hat{\Sigma}$ has the same dynamics as $\Sigma$. Since $\Sigma$ has the same number of modes as $\hat{\Sigma}$, thus we can use $\hat{\Sigma}$ as a bridge to compare $\Sigma$ and $\hat{\Sigma}$.

Similar to the notations for $\Sigma$, for $\hat{\Sigma}$ we define $\hat{\varphi}_{1:r}, \hat{R}_{1:r}, \hat{K}_{1:r}, \hat{P}_{1:r}^{(1:T)}, \hat{K}_{1:r}^{(1:T)}, \hat{P}_{1:r}$, and $\hat{K}_{1:r}$. Particularly, $\hat{P}_{1:r}$ denotes the Riccati solution such that $\hat{P}_i = \hat{R}_i(\hat{P}_{1:r})$, and $\hat{K}_{1:r}$ is computed such that $\hat{K}_i = \hat{K}_i(\hat{P}_{1:r})$. For $\hat{\Sigma}$, we similarly define $\bar{\varphi}_{1:s}, \bar{R}_{1:s}, \bar{K}_{1:s}, \bar{P}_{1:s}^{(1:T)}, \bar{K}_{1:s}^{(1:T)}, \bar{P}_{1:s}$, and $\bar{K}_{1:s}$. In terms of LQR solutions, the relation between $\Sigma$ and $\hat{\Sigma}$ is given below.

**Lemma 20** For the finite time LQR controllers, we have, for any $t$ and any $i$ belonging to any $\hat{\Omega}_k$, $\hat{K}_{i}^{(t)} = \hat{K}_k^{(t)}$. For the infinite time LQR controllers, assuming the Riccati solution $\hat{P}_{1:s}$ uniquely exists among $S^+_s$, and $\hat{K}_{1:s}$ stabilizes $\hat{\Sigma}$, then (i) the Riccati solution $\hat{P}_{1:r}$ exists and uniquely exists among $S^+_r$; (ii) $\hat{K}_k = \hat{K}_i$ for any $i$ belonging to any $\hat{\Omega}_k$.

**Proof** For the finite time case, we first show via induction that, for any $t$ and any $i$ belonging to any $\hat{\Omega}_k$, $\hat{P}_i^{(t)} = \hat{P}_k^{(t)}$. For $T$, we have $\hat{P}_i^{(T)} = \hat{P}_k^{(T)} = Q$. Suppose for $t$, we have $\hat{P}_i^{(t)} = \hat{P}_k^{(t)}$. Then we consider time $t-1$. From the construction of $\hat{\Sigma}$, for $i \in \hat{\Omega}_k$, we have $\sum_{j \in \hat{\Omega}_k} \hat{T}(i,j) = \hat{T}(k,l)$, which further implies $\bar{\varphi}_i(\hat{P}_{1:s}^{(t)}) = \sum_{j \in [r], j \in \hat{\Omega}_k} \hat{T}(i,j) \hat{P}_j^{(t)} = \sum_{j \in [r]} \hat{T}(k,l) \hat{P}_l^{(t)} = \bar{\varphi}_k(\hat{P}_{1:s}^{(t)})$. Also note that $\hat{A}_i = \hat{A}_k$, $\hat{B}_i = \hat{B}_k$, then by definition of the $R$ operator, it is easy to see $\bar{R}_i(\hat{P}_{1:s}^{(t)}) = \bar{R}_k(\hat{P}_{1:s}^{(t)}).$ Since $\hat{P}_{1:s}^{(t-1)} = \bar{R}_i(\hat{P}_{1:s}^{(t)})$ and $\hat{P}_{1:s}^{(t-1)} = \bar{R}_k(\hat{P}_{1:s}^{(t)})$, we have $\hat{P}_i^{(t-1)} = \hat{P}_k^{(t-1)}$. This completes the induction and shows for any $t$ and any $i$ belonging to any $\hat{\Omega}_k$, $\hat{P}_i^{(t)} = \hat{P}_k^{(t)}$. Since during the induction, we have shown $\bar{\varphi}_i(\hat{P}_{1:s}^{(t+1)}) = \bar{\varphi}_k(\hat{P}_{1:s}^{(t+1)})$, and $\hat{A}_i = \hat{A}_k$, $\hat{B}_i = \hat{B}_k$, by definition of the $K$ operator, we have $\hat{K}_i(\hat{P}_{1:s}^{(t+1)}) = \hat{K}_k(\hat{P}_{1:s}^{(t+1)})$. Finally, since $\hat{K}_i^{(t)} = \hat{K}_i(\hat{P}_{1:s}^{(t+1)})$ and $\hat{K}_k^{(t)} = \hat{K}_k(\hat{P}_{1:s}^{(t+1)})$, we have $\hat{K}_i^{(t)} = \hat{K}_k^{(t)}$.

Now we consider the infinite time case. Under the given assumptions, by (Costa et al., 2006, Proposition A.23), we know for any $t$, $\lim_{T \to \infty} \hat{P}_i^{(T-t)} = \hat{P}_i$. Using the result for finite time case that for any $i, i' \in \hat{\Omega}_k$, $\hat{P}_{1:r}^{(T-t)} = \hat{P}_{1:r}^{(T-t)}$, we further have $\hat{P}_i = \hat{P}_{i'}$. Then, if we set $\hat{P}_{1:r}$ similar to the finite time case, we can verify $\hat{P}_i = \bar{R}_i(\hat{P}_{1:r})$ for all $i \in [r]$, which shows the existence of the Riccati solution $\hat{P}_{1:r}$. Suppose there exists another Riccati solution $\check{O}_{1:r} \in S_+^r$ such that $\check{O}_{1:r} \neq \hat{P}_{1:r}$ yet also satisfies $\check{O}_i = \bar{R}_i(\check{O}_{1:r})$. Define $\hat{O}_{1:s}$ such that $\hat{O}_i = \hat{O}_k$ if $i \in \hat{\Omega}_k$, then we see $\check{O}_{1:s} \in S_+^s$, $\check{O}_{1:s} \neq \hat{P}_{1:s}$ yet also satisfies $\hat{O}_i = \bar{R}_i(\hat{O}_{1:s})$, which violates the assumption regarding the uniqueness of Riccati solution $\hat{P}_{1:s}$. Therefore, $\hat{P}_{1:r}$ uniquely exists, and (i) is proved. Finally, (ii) can be shown simply by noticing that $\check{K}_k = \check{K}_k(\hat{P}_{1:r}) = \check{K}_k(\hat{P}_{1:s}) = \hat{K}_k$, where the second equality follows from similar arguments in the finite time case.

Other than illustrating the relations between $\hat{\Sigma}$ and $\hat{\Sigma}$ in terms LQR solutions, Lemma 20 also implies that if an MJS is mode-reducible, there is redundancy among its mode individual controllers,
and these controllers can be computed using the mode-reduced MJS, which justifies our discussion at the beginning of this section. The results for infinite time case in Lemma 20 is more involved as solutions $\hat{P}_{1,r}$ and $\hat{P}_{1,s}$ to Riccati equations are not defined as explicitly as their finite time counterparts $\hat{P}_{1,t}^{(i)}$ and $\hat{P}_{1,t}^{(s)}$, thus their existence and uniqueness are of importance and need separate discussions. For MJS that are not exactly mode-reducible, i.e., $\Sigma$ in $P_1$ and $P_2$, controllers computed using the reduced MJS $\Sigma$ will not be optimal if we apply them to $\Sigma$ due to the perturbation $\epsilon_A, \epsilon_B,$ and $\epsilon_T$. We are able to analyze this resulting suboptimality for the infinite time case with recent advances in MJS-LQR perturbation analysis Du et al. (2022).

**Theorem 21 (Complete Version of Theorem 9)** Assume $A_l$ holds for $\Sigma$, and $\Sigma$ has additive Gaussian noise $N(0, \sigma_w^2 I_n)$ that is independent of the mode switching. Let $J^*$ and $\hat{J}$ respectively denote the cost incurred by the optimal controller $K_{1,s}$ and controller $\hat{K}_{1,r}$ (at time $t$, $u_t = \hat{K}_k(t)x_t$ if $\omega_t = \hat{\Omega}_k$). Then, there exists constant $\epsilon_{A,B}, \epsilon_T, C_{A,B},$ and $C_T$ (depend on $\Sigma, Q,$ and $R$), such that when $\max\{\epsilon_A, \epsilon_B\} \leq \epsilon_{A,B}$ and $\epsilon_T \leq \epsilon_T$, the Riccati solution $\hat{P}_{1,r}$ exists and uniquely exists among $S^+_r$, thus $\hat{K}_{1,r}$ exists as well, and $\hat{J} - J^* \leq \sigma_w^2 (C_{A,B} \max\{\epsilon_A, \epsilon_B\} + C_T \epsilon_T)^2$.

**Proof** We will start with the discussion between $\Sigma$ and $\tilde{\Sigma}$, and use the relation between $\tilde{\Sigma}$ and $\hat{\Sigma}$ given in Lemma 20 to prove the claim.

Comparing $\Sigma$ and $\tilde{\Sigma}$, one can see $\|A_i - A_i\| \leq \epsilon_A, \|B_i - B_i\| \leq \epsilon_B,$ and $\|T - T\|_\infty \leq \epsilon_T$. Then, from Du et al. (2022) we know when $\max\{\epsilon_{A,B}, \epsilon_T\} \leq \tilde{\epsilon}_{A,B}$ and $\epsilon_T \leq \tilde{\epsilon}_T$, the Riccati solution $\hat{P}_{1,s}$ uniquely exists among $S^+_r$, thus $\hat{K}_{1,s}$ exists, and the cost $\hat{J}$ when applying $\hat{P}_{1,s}$ to $\Sigma$ has suboptimality $\hat{J} - J^* \leq \sigma_w^2 (C_{A,B} \max\{\epsilon_A, \epsilon_B\} + C_T \epsilon_T)$. Using Lemma 20, we know $\hat{P}_{1,r}$ exists and uniquely exists $S^+_r$, and $\hat{K}_k = \hat{K}_i$ for any $i$ belonging to any $\hat{\Omega}_k$, which implies applying $\hat{K}_{1,s}$ is equivalent to applying $\hat{K}_{1,r}$ as in the theorem statement. Thus $\hat{J} = \hat{J}$, and $\hat{J} - J^* = \hat{J} - J^* \leq \sigma_w^2 (C_{A,B} \max\{\epsilon_A, \epsilon_B\} + C_T \epsilon_T).$