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VECTOR AND PARALLEL IMPLEMENTATION OF
THE CONJUGATE GRADIENT FFT METHOD FOR
THE SOLUTION OF LARGE ELECTROMAGNETIC
SCATTERING PROBLEMS ON SUPERCOMPUTERS

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Abstract

There is a growing interest in the application of the Conjugate Gradient FFT (CGFFT) method to electromagnetic radiation and scattering problems primarily because of the low storage requirement associated with this method. The CGFFT lends itself well to efficient execution in vectorized fashion and in this study, a vector-concurrent form of the CGFFT method suitable for implementation on parallel, multiprocessor supercomputer systems is applied to the problem of scattering from electrically large planar structures. The vectorizable nature of the CGFFT algorithm is exploited by identifying the major processes involved in a given iteration including the Fast Fourier Transforms, vector dot products and norm calculations. Since for a given iteration there is no data recurrence in the operations involved, these processes may be vectorized to increase the computation speed.

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1 Introduction

Computational electromagnetics relies heavily on vector-oriented algorithms to simulate complex problems. With the computer technology approaching the limits of semiconductor speeds, the exploitation of parallel processing has emerged in order to meet the processing demands of computationally intensive applications in electromagnetics. In such applications machine instructions are vectorized and distributed across different vector processors for concurrent execution as opposed to the traditional approach where the computers are limited to sequential processing of data on a single scalar processing unit.

In this report, an example of vector and parallel processing as applied to an iterative technique, namely, the conjugate gradient method [1,2,3] is explored. In general, iterative techniques become more attractive for the solution of operator equations arising in electromagnetic problems as the size of the problem increases. This is mainly because the iterative schemes often involve only the multiplication of the matrices with vectors and thus do not require an explicit storage of the system matrix. Also, since these methods avoid the process of matrix inversion, they are numerically more stable for ill-conditioned systems.

In electromagnetic scattering, the pertaining integral operators are usually formulated as convolutions over the unknown current density. In these situations, the Fast Fourier Transform can be used to evaluate the convolution integrals efficiently. When combined with FFT, the conjugate gradient method requires roughly $4N(1 + \log_2 N)$ operations per iteration. The required number of iterations before the conjugate gradient method can yield a reasonable accuracy is often a fraction

of the total number of unknowns. This depends primarily on the distribution of the dominant eigenvalues of the operator projected onto the system matrix.

The CGFFT lends itself to efficient execution in vectorized fashion and in this report, a vector-concurrent form of the CGFFT method suitable for implementation on parallel, multiprocessor systems is applied to the problem of scattering from electrically large planar structures. Optimized algorithms are implemented on IBM 3090-600E supercomputer and on an Alliant FX/8 mini-supercomputer with both vector and vector-concurrent capabilities.

Recently [4], it was shown that the convergence of the CGFFT may be improved by the incorporation of subdomain basis functions in the expansion of the unknown current distributions. The improved convergence was attributed to a more accurate representation of the current distribution in the spectral domain. Implementation of this formulation in conjunction with the optimized algorithm will result in a drastic improvement in the efficiency of the method. Numerical results are presented to demonstrate the corresponding improvements in the processing speed.

2 The Conjugate Gradient FFT Method

The integral equations arising in the study of electromagnetic scattering problems can be written in a general operator form as (Fig. 1)

$$A[\mathbf{J}] = \mathbf{E}^i \quad (1)$$

where \mathbf{E}^i denotes the excitation field and \mathbf{J} is the unknown current density vector.

A typical form of such integral equations is

$$\mathbf{E}^i(\mathbf{r}) = \bar{\eta}(\mathbf{r}) \cdot \mathbf{J}(\mathbf{r}) + \int_{v'} \bar{\Gamma}(|\mathbf{r} - \mathbf{r}'|) \cdot \mathbf{J}(\mathbf{r}') dv', \quad (2)$$

where $\bar{\Gamma}$ is the associated dyadic Green's function and $\bar{\eta}$ is a tensor specific to the electrical properties of the scatterer.

The integral in (2) implies a convolution over the geometrical limits of the target. Thus, introducing the forward and inverse Fourier transform pair¹

$$\tilde{g} \iff g = \mathcal{F}^{-1}\{\tilde{g}\}, \quad (3)$$

equation (2) can be alternatively written as [4]

$$\mathbf{E}^i = \bar{\eta} \cdot \mathbf{J} + \mathcal{F}^{-1}\{\bar{\Gamma} \cdot \tilde{f}\tilde{\mathbf{J}}\} = A[\mathbf{J}]. \quad (4)$$

In the above, \tilde{f} is the Fourier transform of the basis function used in the expansion of the current density and $\hat{\mathbf{J}}$ is the discrete Fourier transform of its sampled train so that

$$\tilde{\mathbf{J}} = \tilde{f}\hat{\mathbf{J}}. \quad (5)$$

Clearly, this formulation avoids the generation of the square matrix corresponding to operator A and thus implies a storage requirement of $\mathcal{O}(N)$ as compared to

¹See Appendix I for the definitions of the Fourier transform relations used.

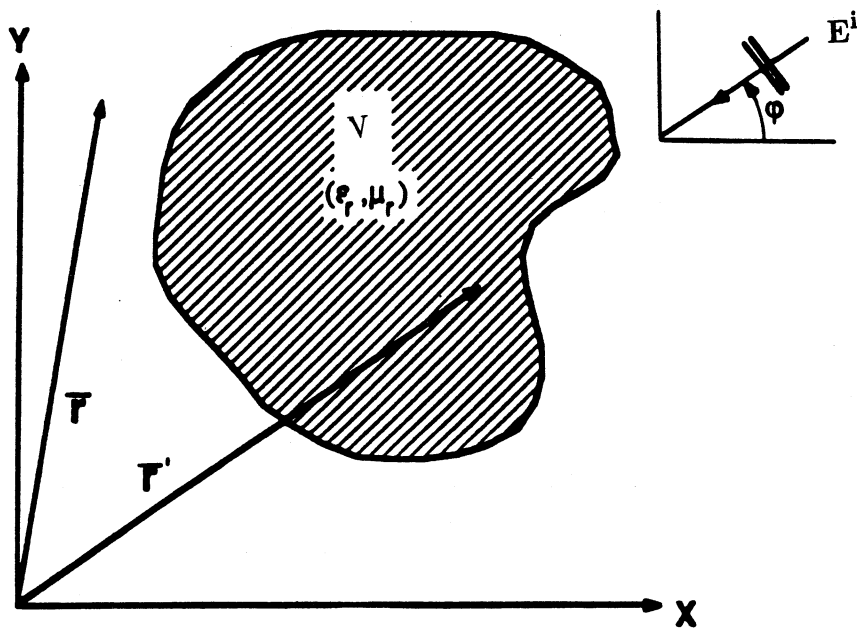


Fig. 1 Plane wave scattering from an arbitrary composite material target.

$\mathcal{O}(N^2)$ required for an implementation of the corresponding direct method. This storage economy has made the formulation suitable for large radiating systems. Several iterative methods for the solution of (4) are available. In the conjugate gradient method, the unknown current density is calculated by minimizing a functional through a finite number of successive searches in an iterative fashion. A conjugate gradient algorithm suitable for numerical computation is [2,3]

$$R_0 = A[J_0] - E^i$$

$$P_0 = -b_{-1}A^a[R_0]$$

Main Iteration Loop

For $n \geq 0$ Do

$$t_n = \frac{1}{\|A[P_n]\|^2}$$

$$J_{n+1} = J_n + t_n P_n$$

$$R_{n+1} = R_n + t_n A[P_n] \tag{6}$$

$$b_n = \frac{1}{\|A^a[R_n]\|^2}$$

$$P_{n+1} = P_n - b_n A^a[R_{n+1}]$$

$$\frac{|R_{n+1}|}{\|E^i\|} \stackrel{?}{\leq} tolerance$$

Repeat If Necessary.

In the above, R is the residual vector (the functional to be minimized) and P is the search direction. The norm and the adjoint operator are defined in terms of an inner product as

$$|g|^2 = \langle g, g \rangle \tag{7}$$

and

$$\langle A[g], f \rangle = \langle g, A^a[f] \rangle. \quad (8)$$

It may be shown that[5]

$$A^a[\mathbf{J}] = \bar{\eta}^*(\mathbf{r}) \cdot \mathbf{J}(\mathbf{r}) + \int_{v'} \bar{\Gamma}^*(|\mathbf{r} - \mathbf{r}'|) \cdot \mathbf{J}(\mathbf{r}') dv', \quad (9)$$

where * denotes the complex conjugate.

Each iteration in (6) requires two convolution operations, two norm calculations and three scalar products. The iterations are repeated until a preset tolerance on the residual is met.

3 Optimization of the CGFFT Algorithm

Before an assessment of the vectorizability of the CGFFT algorithm, a brief review of some general concepts in vector and parallel processing will be presented. This discussion is followed by an overview of the computer systems used in this study.

3.1 Vectorization

The crux of parallel computing is the process of vectorization and distribution of code among multiple processors. Typically, a vector instruction is capable of operating on 32 to 128 elements of data at once, depending on the machine used, resulting in two to four times gain in speed over the corresponding sequential scalar instruction. Vectorization requires some degree of independence in the access of data by the code. A dependence occurs when two statements—or iterations of the same statement—refer to the same storage location. Some data dependences *inhibit* vectorization; they are called *recurrence*. By changing the structure of the code, it may be possible to eliminate a recurrence and vectorize the modified code.

As an example of a vectorizable code, consider the following simple loop

```
DO I=1,N
    C(I)=A(I)+B(I)
END DO.
```

In scalar mode, registers for scalar arithmetic hold only one element at a time. To add two vectors A and B, each element in the vector B has to be added individually to the appropriate element in the vector A, and then assigned to the appropriate

element in vector C. The sequence of machine instructions generated for such a scalar computation would be

1. Load counter N $\langle 1 \rangle$
2. Load A(I) into scalar register through the loop $\langle 1 \rangle$
3. Add B(I) to the scalar register $\langle 3 \rangle$
4. Store C(I) from the scalar register $\langle 1 \rangle$
5. Decrement counter by 1 and repeat $\langle 1 \rangle$.

The time cycles required for each operation are shown in the brackets. Since in this particular case, the calculations in a given value of the Do loop index do not depend on a previous value of the index, there is no data recurrence in the loop and the operation may be carried out in vector mode as follows

1. Load counter N $\langle 1 \rangle$
2. Load A(1)-A(128) into vector register $\langle 1 \times 128 \rangle$
3. Add B(1)-B(128) to the vector register $\langle 3 + 127 \rangle$
4. Store C(1)-C(128) from the vector register $\langle 1 \times 128 \rangle$
5. Decrement counter by 128 and repeat $\langle 1 \rangle$.

When the instruction to add vectors is issued, the first elements of the vectors, A(1) and B(1), are transmitted to the add functional unit at time t_0 . Since the functional unit time for an addition is three clock periods, the first result (A(1) + B(1)) will

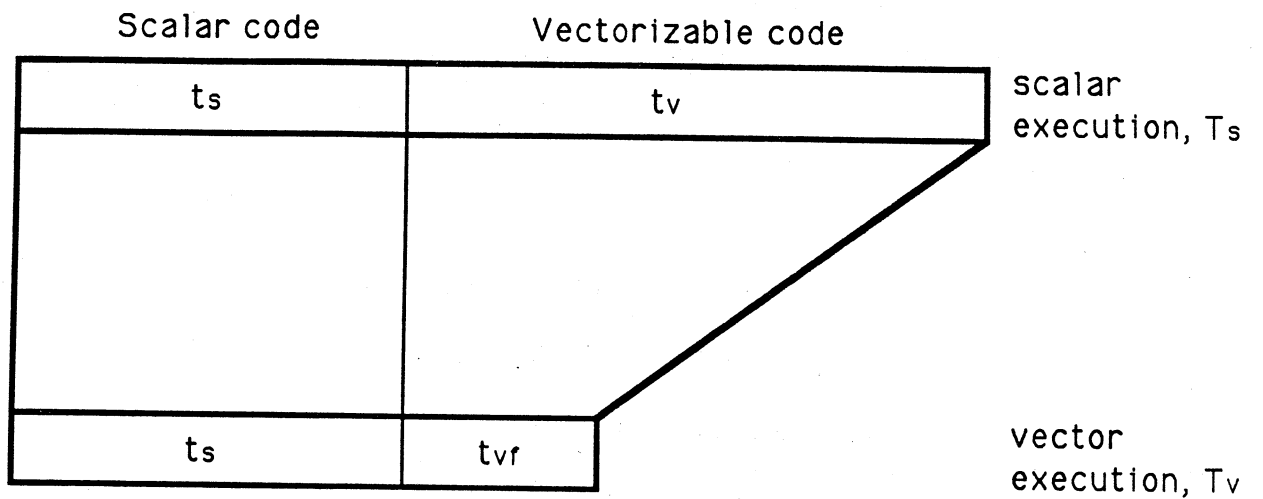
be produced at t_2 . It is important to note that while elements A(1) and B(1) were entering the functional unit, elements A(2) and B(2) were being transmitted to it. Due to full segmentation, the result of $A(2) + B(2)$ is produced at t_3 . From this point on, the process of one pair of vector register elements entering the functional unit and a result being produced every clock cycle continues until all elements of the vectors are added. So the addition of the first 128 pair of elements is accomplished in $127 + 3 = 130$ time cycles as opposed to the corresponding scalar processing time of $128 \times 3 = 384$ cycles. This process is called *pipelining* and is the major underlying process of optimization. If in the above example $N=256$, the scalar time would be 1537 clock cycles while a vector implementation would yield 775 cycles.

In order to measure the improvement in the speed of a vectorized program, several parameters are defined. The *program speedup* is defined as the increase in the speed of execution when a code is run in vector mode relative to that in the scalar mode. Therefore, referring to figure 2, if a code runs for T_s seconds in scalar mode and for T_v seconds in vector mode (i.e. after optimization), the corresponding program speedup is given by

$$\text{program speedup} = T_s / T_v.$$

Also, the *vector content* of a program is that percent of the scalar code which vectorizes. Thus, if for a given code, the scalar portion which may not be vectorized runs in t_s seconds and that which is vectorizable runs in t_v seconds in scalar mode and in t_{vf} seconds in vector mode², respectively, we have

² t_{vf} would be the time the code actually spends in the vector facilities.



Vector Content= t_v/t_s

Vector Speedup= t_v/t_{vr}

Program Speedup= T_s/T_v

Fig. 2 Scalar and vector execution times in a typical vectorized code.

$$\text{vector content} = t_v / T_s \times 100$$

and

$$\text{vector speedup} = t_v / t_{vf}$$

Typically, programs with more than about 70% vector content, run 1.5 to 2 times faster in vector mode. Higher vector contents, however, are achievable as will be shown to be the case for the CGFFT algorithm.

3.2 Concurrency

Although it is increasingly expensive to make a single processor faster, fairly fast processors are inexpensive. Therefore using several relatively inexpensive processors in parallel is often more efficient than using a single fast processor. A multiple processor system can devote several processors to the execution of different parts of a single program simultaneously. The compiler inserts protective synchronization code into the optimized loops so that the multiple processors work together without interfering with each other. Synchronization is needed to prevent conflicts in the use of memory shared by parallel tasks and is considered the major cost of parallel processing. Optimization is suppressed whenever the possibility of a data dependence exists. Also, from Amdhal's law the increase in the speed of execution in concurrent mode as the number of processors is increased, approaches an upper limit set by the presence of the sequential constructs in the algorithm.

In general, the concepts of data dependence in parallel processing are the same as those in vector processing; the consequences of certain dependences, however, are different.

A parameter of interest when processing in concurrent mode is the *efficiency* of execution. Efficiency is a measure of parallelism in the algorithm. An efficient parallel tasking system makes possible a nearly linear speedup in performance as processors are added, if the algorithm is parallel. Thus, if $T^{(n)}$ denotes the time required to run on n processors, we may define

$$\text{concurrency speedup} = T^{(1)}/T^{(n)}$$

and

$$\text{efficiency}\% = T^{(n)}/(nT^{(1)}) \times 100.$$

3.3 Compiler Directives

Once a well written code is compiled, most vector and parallel constructs are recognized by the compiler and automatically optimized for efficient execution. In addition to automatic optimization, however, most vector compilers provide directives for additional control. Compiler directives are user supplied control structures to override decisions made by the compiler and to give additional information to it. Upon compilation, the directives are interpreted by the processor and converted to library calls to be executed in a more efficient manner. In particular, the associative transformation directive recognizes operations like dot products and norm computations as reduction functions and optimizes these otherwise non-vectorizable loops.

Some general guidelines for code optimization in various architectural levels are given in Appendix II and have been followed in optimizing the CGFFT algorithm in the present study.

3.4 The IBM 3090 supercomputer

The supercomputer used in this study was an IBM 3090-600E with six processors, each with a vector facility and 64 kilobytes cache memory. An overview of the system is given in figure 3. It has 256 megabytes of real memory with an expanded storage of 512 megabytes and utilizes VM/XA operating system. The cycle time is 17.2 ns delivering a theoretical peak performance of 116 MFLOPS (mega floating point operations per second)³. The vector facility provides 16 vector registers of 32 bit words (or 8 vector registers of 64 bit words) operating on up to 128 data elements.

The operating system of the IBM 3090 provides detailed timing information on the execution of the code. The total CPU time and the time fraction spent in the vector facility when program is run in vector mode are reported upon execution. This information can be used to estimate the vector content of the code. The overall vector performance of the IBM 3090 is depicted in figure 4 where the program and vector speedups are given for various vector contents.

3.5 The Alliant FX/8 Mini-supercomputer

The Alliant FX/8 is a high performance system with both parallel and vector computing capabilities. The FX/8 runs the Concentrix operating system which conforms to the external specifications of UNIX and fully supports Transmission Control Protocol/Internet Protocol (TCP/IP) and Network File System (NFS). The Alliant used in this study has four vector processors (expandable to 8), 56 Mbytes of main memory, 256 kbytes of cache memory and a bus which moves 188

³The Cray X-MP has a cycle time of 8.5 ns.

IBM 3090-600 Processor Unit Design

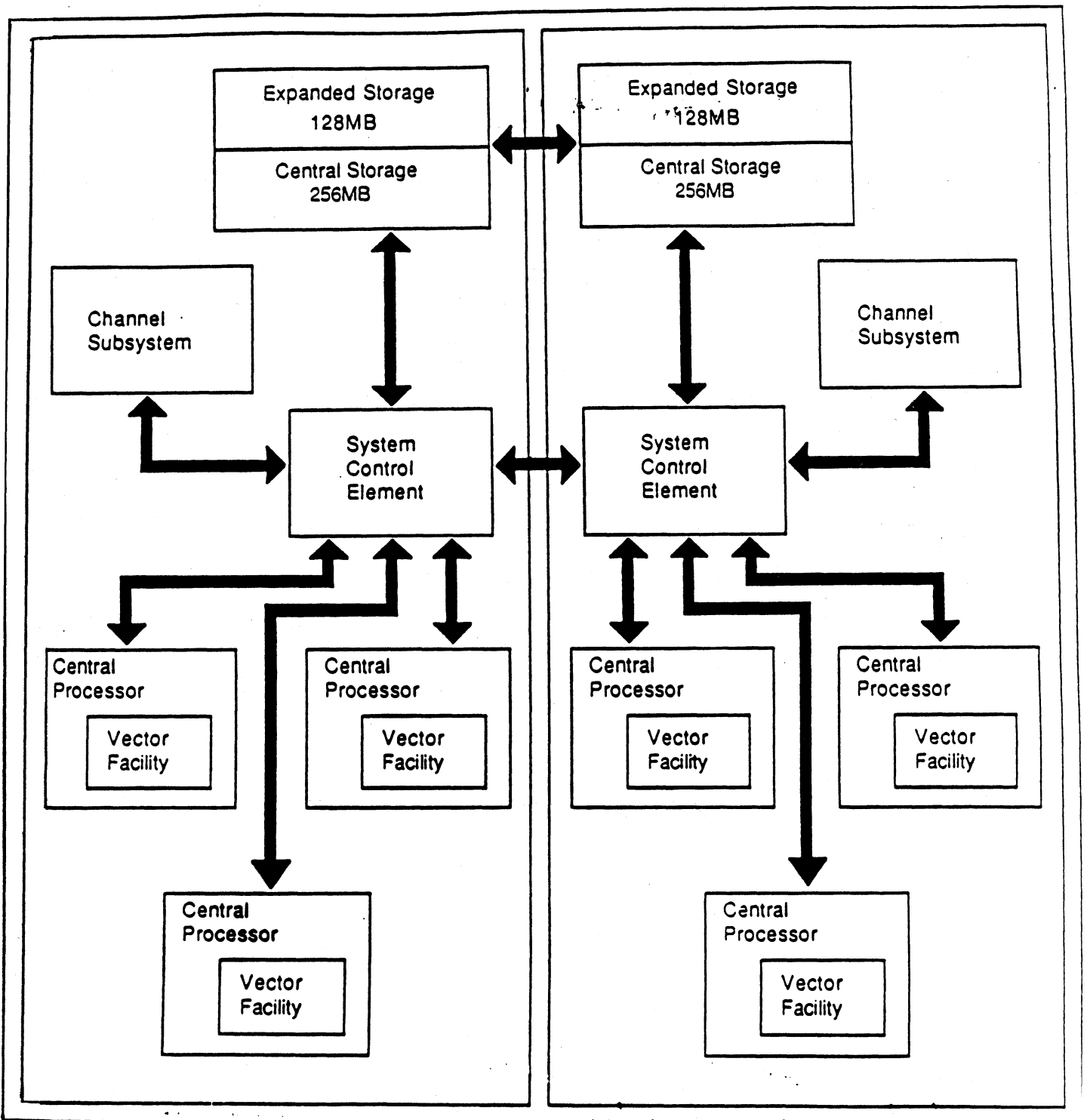


Fig. 3(a) An overview of the IBM 3090-600E supercomputer.

IBM 3090 Central Processor

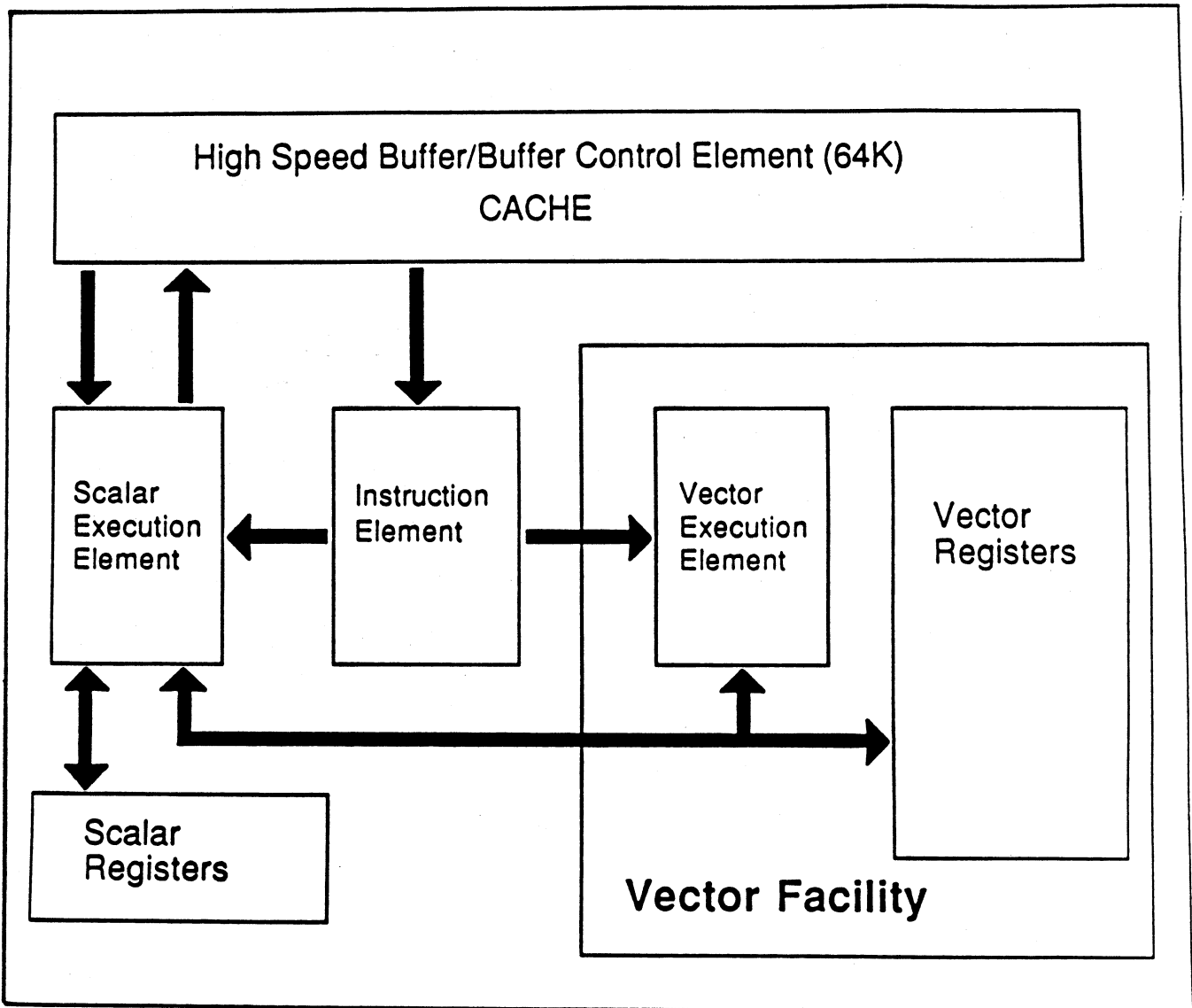


Fig. 3(b) IBM 3090- The central processor unit.

megabytes per second. The machine comprises of three main sections: vector unit, concurrency control unit and the FX/Fortran language processor which extends Fortran 77 to permit assignment and other operations on full arrays. The language extensions may be efficiently used to expose the inherent data independence which may have been masked by over-modularizing the program.

The Alliant's Computational Elements can run in three different modes as shown in figure 5, namely,

- Complex mode: all specified processors attack a single job,
- Detached mode: processors run independent jobs in parallel, and
- Dynamic complex mode: the complex attacks a single job and the detached processors run jobs concurrently.

The complex mode provides the optimum situation among the above modes since it provides maximum throughput for a single job; the user may specify the number of processors in the complex to achieve desired efficiency and speed. In the detached mode, the user specifies the number of CEs statically detached and dedicated to independent jobs. And in the dynamic mode, maximum flexibility is achieved for mixed mode production and development installations where user can favor detached or complex CEs to meet specific needs.

The Alliant's Fortran compiler automatically optimizes programs written in Fortran 77 for speed of execution through loop level concurrency and vectorization thus providing automatic parallelization of sequential code. The compiler's optimization strategy may be summarized as follows:

IBM 3090 Supercomputer Vector Performance

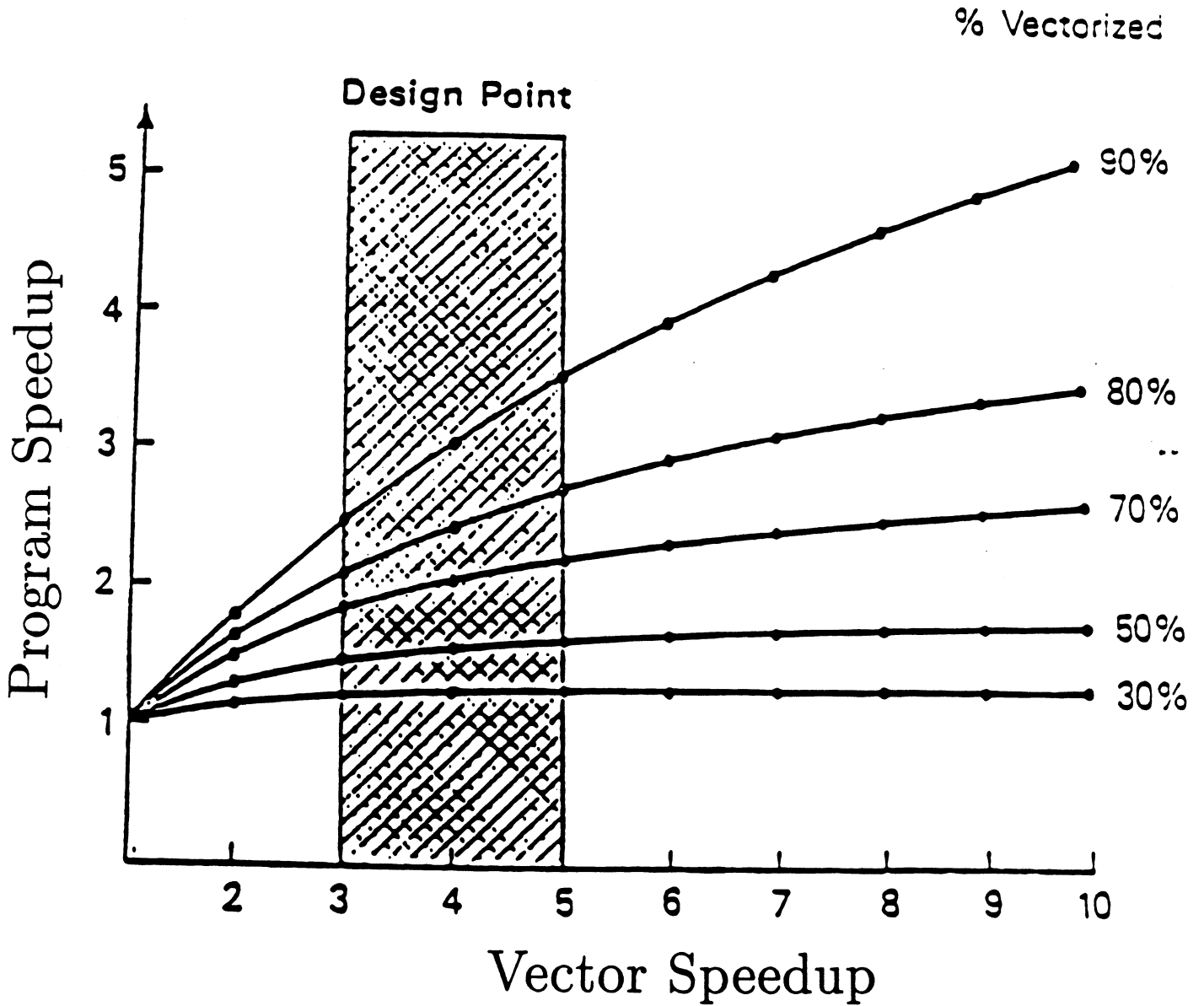
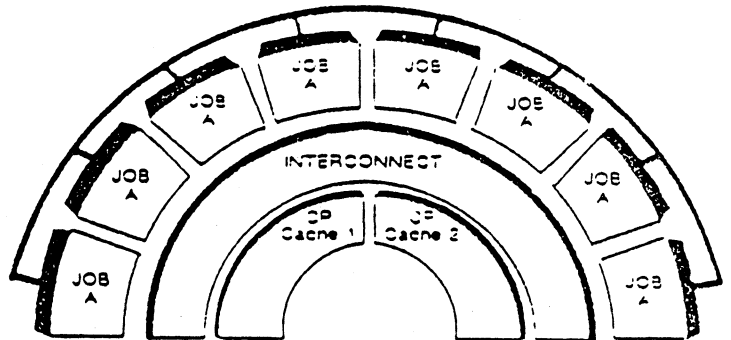


Fig. 4 IBM 3090- The vector performance.

Three Modes of Execution under Software Control

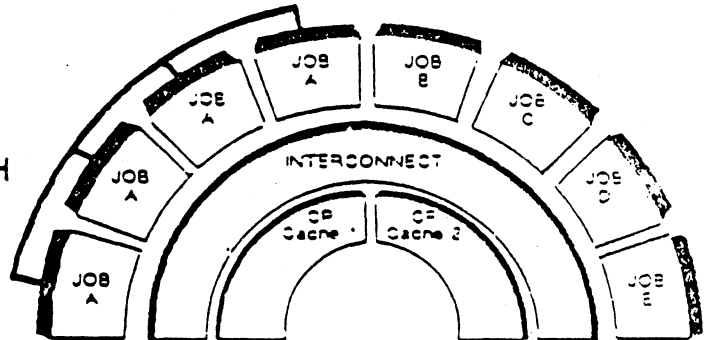
COMPLEX MODE:

OPTIMUM TIME-TO-SOLUTION



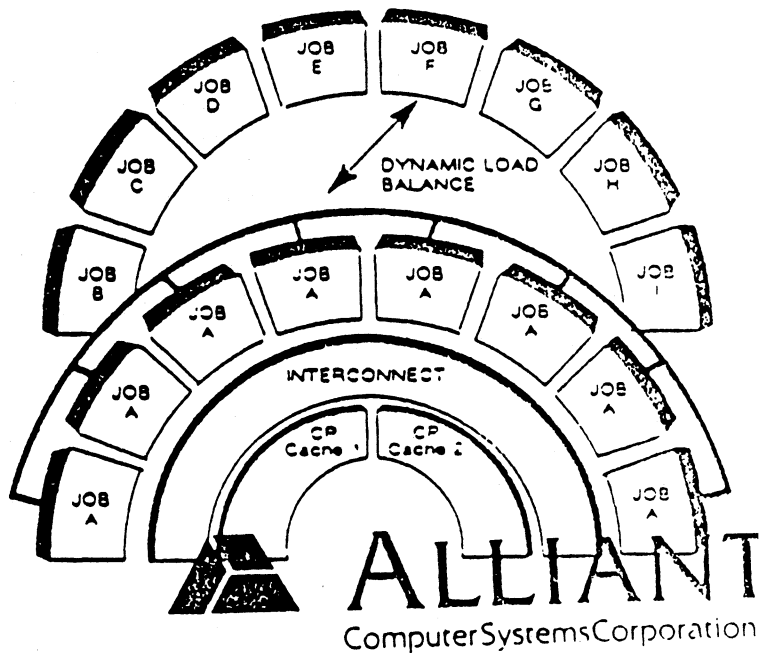
DETACHED MODE:

HIGH THROUGHPUT COMBINED WITH
PRODUCTION EXECUTION
(User specifies number of CEs
statically detached)



DYNAMIC COMPLEX MODE:

MAXIMUM FLEXIBILITY FOR MIXED
MODE PRODUCTION AND
DEVELOPMENT INSTALLATIONS
(User can favor detached or
complex to meet changing needs)



ALLIANT
Computer Systems Corporation

Fig. 5 Alliant FX/8 modes of operation.

For *each innermost loop*:

- If **vectorizable** then
 - if *next outer loop* is **parallelizable** then
 - * Concurrent-Outer Vector-Inner
 - else
 - * Vector-Concurrent
 - endif
- elseif **parallelizable** then
 - * Scalar-Concurrent
- else
 - * Not Optimized
- endif

The vector-concurrent execution in the complex mode provides the fastest performance since it utilizes all the available processors to attack a single task concurrently with vector operations performed on strides of data on each processor. The resulting high-performance, low-level parallelism can significantly boost the performance of compute-intensive operations such as Fourier transformations.

3.6 Optimized Algorithm

Here, the vectorizable nature of the CGFFT algorithm is exploited by identifying the major processes involved in a given iteration. From (6), it can be seen

that a considerable amount of computation time is spent in the calculation of the convolutions carried out in $A[P_n]$ and $A^a[R_n]$. Each convolution includes a pair of forward and inverse Fourier transform operations on the relevant components of the current density vector along with a Hadamard (element by element) multiplication of the current and the dyadic Green's function in the spectral domain. Since for a given current component there is no data recurrence at a particular point in an iteration, these operations may be vectorized to increase the speed of calculations. The same observation is true for the computation of the norms and the dot products used to update the current J , the residual vector R , and the search vector P . More importantly, since the FFT is a highly vectorizable algorithm, it plays a major role in the speed and efficiency of the optimized code.

In this study the scalar and optimized FFT routines available on the IBM 3090's ESSL library and the Alliant's Fortran math library were used in both scalar and vector modes. These routines are written in conjunction with assembler language and generate instructions appropriate for the architecture of the processors—they manage data to make efficient uses of the memory hierarchy.

4 Scattering from Large Planar Structures

The integral equation for a thin perfectly conducting strip of width w illuminated by a H-polarized plane wave of intensity H_0 incident at an angle ϕ_0 (Fig. 6) is given by

$$H_0 \sin \phi_0 e^{jk_0 x \cos \phi_0} = \frac{1}{4k_0} \left(k_0^2 + \frac{\partial^2}{\partial x^2} \right) \int_{-w/2}^{w/2} J_x(x') H_0^{(2)}(k_0 |x - x'|) dx' \quad (10)$$

where $H_0^{(2)}$ is the zeroth order Hankel function of the second kind and $J_x(x)$ is the unknown surface current density on the strip. By expanding the current in terms of sub-domain basis functions as

$$J_x(x) = \sum_{n=1}^N J_n f(x - x_n), \quad (11)$$

equation (10) takes the form

$$H_0 \sin \phi_0 e^{jk_0 x \cos(\phi_0)} = \frac{1}{4} \mathcal{F}^{-1} \{ \widetilde{\mathcal{W}}(f_x) \widehat{J} \} \quad (12)$$

where

$$\widetilde{\mathcal{W}}(f_x) = (1 - f_x^2) \widetilde{H}_0^{(2)}(f_x) \widetilde{f}(f_x), \quad (13)$$

$$\widetilde{H}_0^{(2)}(f_x) = \frac{2j}{k_0 \sqrt{f_x^2 - 1}}. \quad (14)$$

and f_x is the transform variable.

Consider now a perfectly conducting rectangular plate as shown in figure 7. The required integral equations for the current density are given by

$$E_x^i(x, y) = \left[\left(k_0^2 + \frac{\partial^2}{\partial x^2} \right) \Pi_x + \frac{\partial^2}{\partial x \partial y} \Pi_y \right] \quad (15)$$

$$E_y^i(x, y) = \left[\frac{\partial^2}{\partial x \partial y} \Pi_x + \left(k_0^2 + \frac{\partial^2}{\partial y^2} \right) \Pi_y \right] \quad (16)$$

where \mathbf{E}^i denotes the incident plane wave, $\mathbf{\Pi}$ is the electric Hertz vector defined as

$$\mathbf{\Pi}(\mathbf{r}) = \frac{jZ_0}{k_0} \int_{s'} \mathbf{J}(\mathbf{r}') G(\mathbf{r}; \mathbf{r}') ds', \quad (17)$$

and G is the free space Green's function. Introducing the two dimensional expansion

$$J(x, y) = \sum_{n=0}^{N-1} \sum_{m=0}^{M-1} J_{nm} f(x - x_n, y - y_m), \quad (18)$$

the system of equations (15) and (16) may be rewritten as

$$\mathbf{E}^i = jk_0 Z_0 \mathcal{F}^{-1} \{ \tilde{\mathcal{D}} \cdot \tilde{\mathbf{J}} \} \quad (19)$$

where $\tilde{\mathcal{D}}$ denotes the dyadic Green's function,

$$\tilde{\mathcal{D}} \equiv \begin{pmatrix} (1 - f_x^2) \tilde{G} & (-f_x f_y) \tilde{G} \\ (-f_x f_y) \tilde{G} & (1 - f_y^2) \tilde{G} \end{pmatrix}, \quad (20)$$

f_x and f_y are the transform variables, and

$$G(x, y) = \frac{e^{-jk_0 \sqrt{x^2 + y^2}}}{4\pi \sqrt{x^2 + y^2}} \iff \tilde{G}(f_x, f_y) = \frac{1}{2k_0 \sqrt{f_x^2 + f_y^2 - 1}}. \quad (21)$$

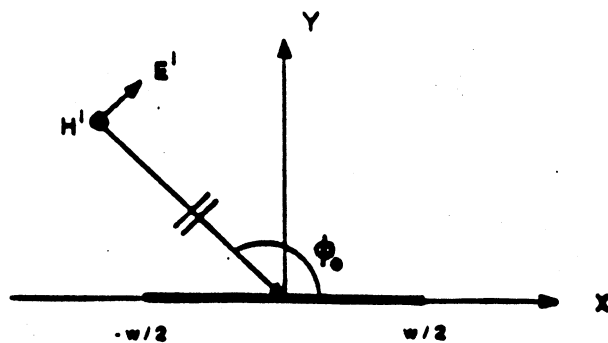


Fig. 6 Geometry of a perfectly conducting strip illuminated by a plane wave.

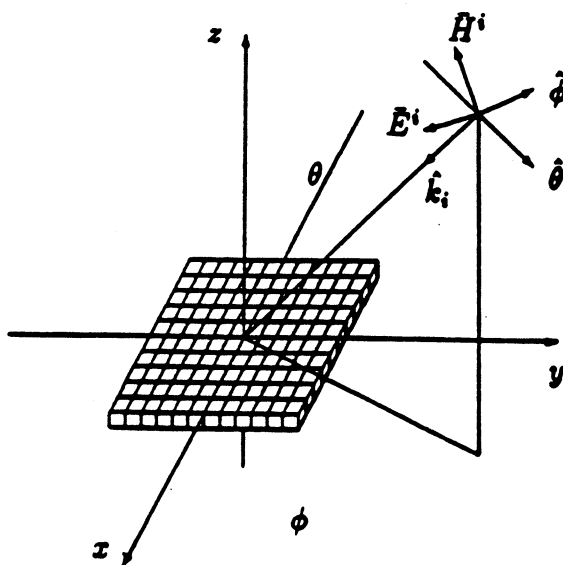


Fig. 7 Geometry of the perfectly conducting plate illuminated by a plane wave.

5 Results

Figure 8 shows the performance improvement of the CGFFT method when executed in vector-concurrent mode over the scalar mode for strips of various sizes at normal incidence. The actual FFT pad used in this experiment was of order one($\rho = 1$). Thus, the size of the pad varied from $N' = 32$ for the 1λ strip to $N' = 512$ for the 17λ one(15 unknowns/ λ). Clearly, the optimized algorithm provides a superior performance especially for larger problems. Figure 8 also includes the effect of incorporating the piecewise sinusoidal expansion functions into the CGFFT formulation which results in the improvement in the speed of the method by reducing the number of iterations required to meet the tolerance. This is also shown in figure 9 for both scalar and vector-concurrent processing modes. The magnitude of the surface current density on the 17λ strip is shown in figure 10 at normal incidence.

A detailed examination of the performance of the optimized code is given in Appendix III. The appendix lists the source code for each routine and enumerates the loops and array operations that were candidates for optimization. The information presented includes the location of the loop within the code, the structure of the loop, the nature of optimization applied and possible data dependences inhibiting the optimization of the loop. Clearly, all the important (CPU-intensive) loops are vectorized and processed concurrently on multiple processors.

The plate problem was examined in more detail. A $2\lambda \times 2\lambda$ plate was considered at normal incidence for E-polarization using 63×63 unknowns and 128×128 FFT pad. The vectorized algorithm was executed on the IBM 3090 under three

Performance of the Optimized CGFFT Algorithm

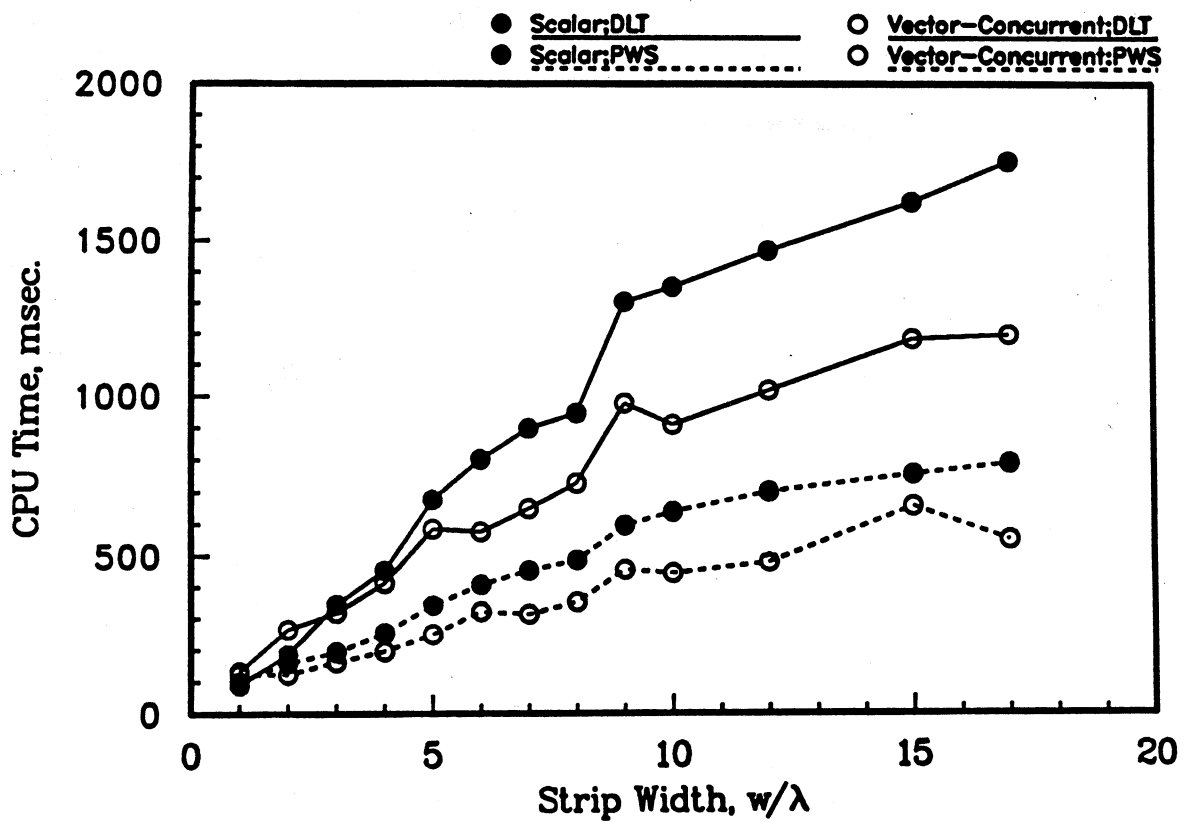


Fig. 8 Comparison of the CPU times associated with the scalar and vector-concurrent execution of the CGFFT method for various strips using different basis functions (15 unknowns/ λ and FFT pad of order 1).

Convergence of the Optimized CGFFT Algorithm

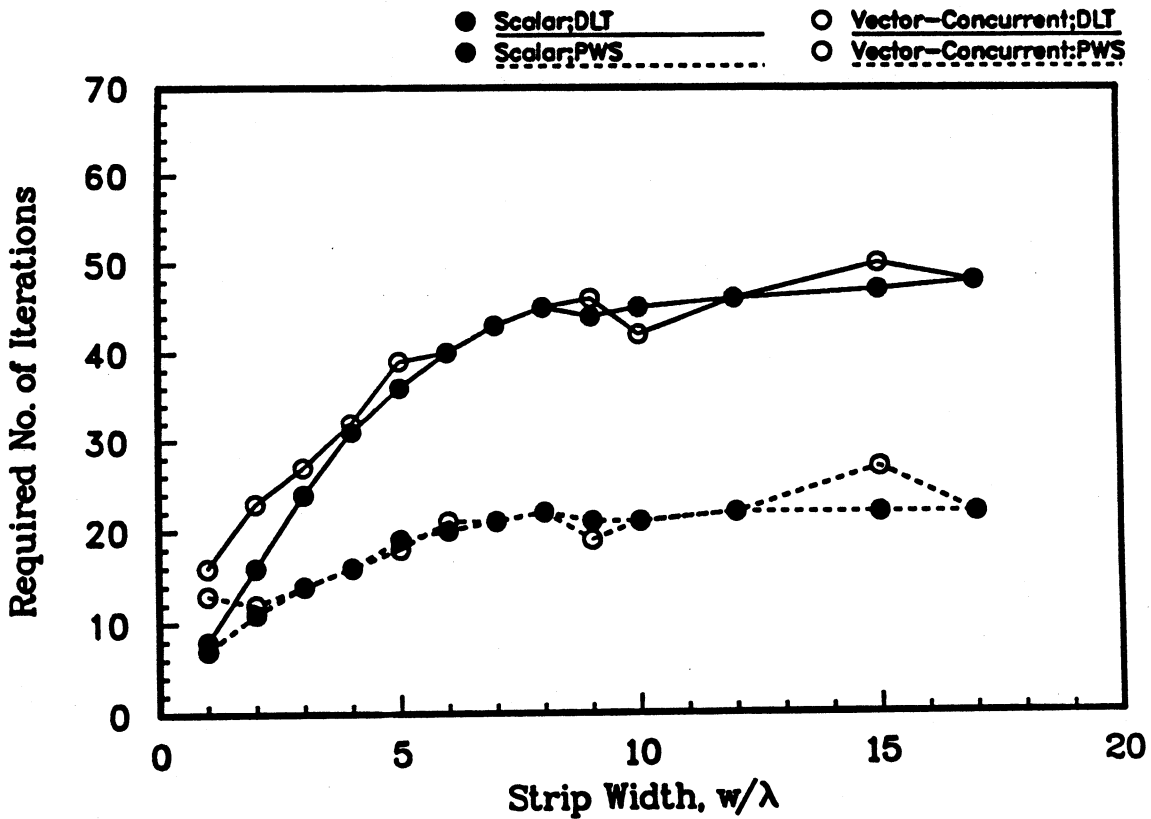


Fig. 9 A comparison of the required number of iterations associated with the convergence of the optimized CGFFT algorithm for various strips using different basis functions (15 unknowns/ λ ; FFT pad of order 1).

Strip Surface Current Density

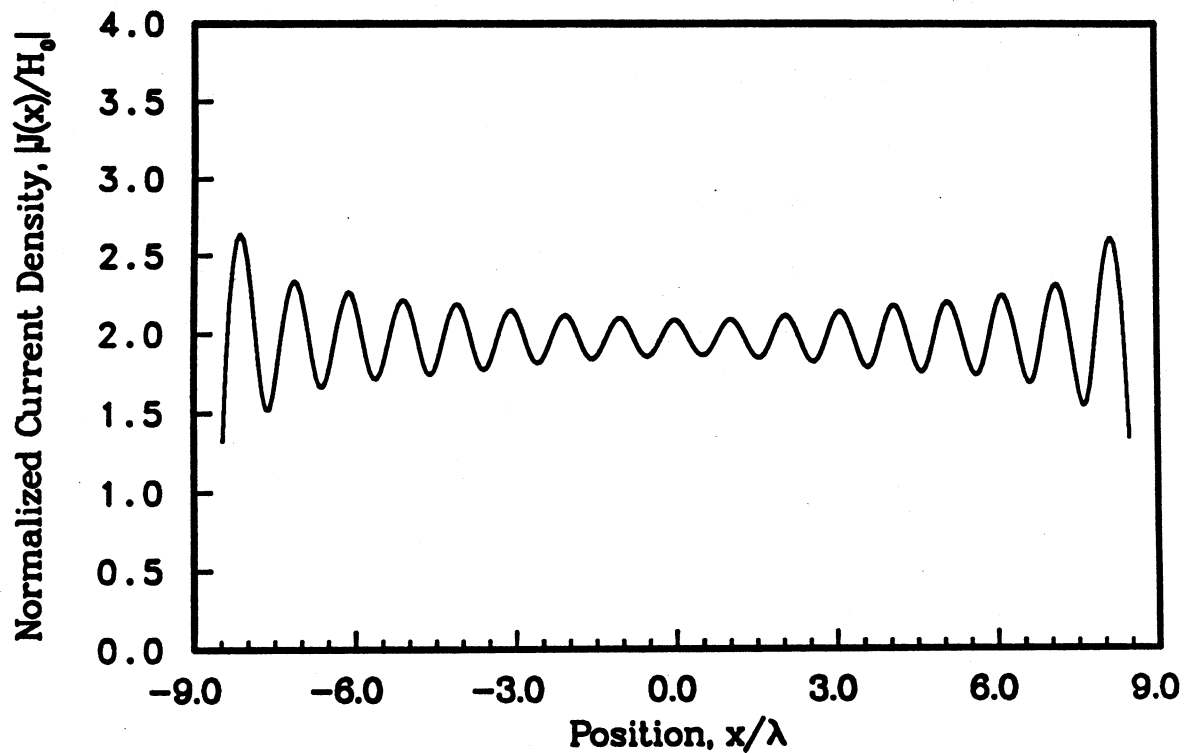


Fig. 10 Surface current density magnitude for the 17λ conducting strip illuminated by an H-pol. plane wave at normal incidence computed by the optimized CGFFT algorithm (15 unknowns/ λ) using PWS basis functions.

situations. In scalar mode using a scalar radix 2 FFT routine and in vector mode with scalar and vectorized FFT routines both provided by the IBM's ESSL library. The results are tabulated in Table 1 showing the total CPU time and the time spent in the vector facilities. The percent vectorizable code of 97% indicates a highly optimized algorithm resulting in a program speedup of more than four times. This is mostly attributed to the employed vectorized FFT routine. The performance point of the supercomputer for this particular case is given in figure 11.

Table 2 shows the corresponding vectorization data for the Alliant. The entries in the column denoted by MFLOPS are calculated based on the estimated number of operations from [5]. The distribution of the CPU time among the compute-intensive routines in the scalar and vector modes is shown in figure 12(a). Again, it is clear that calls to the FFT routine are the most significant factor in the solution process and about 90% of the CPU time is spent in these computations(Fig. 12(b)).

To examine the efficiency of the algorithm for vector-concurrent execution, the optimized code was run on one to four vector processors on the Alliant and the results are shown in Table 3. A speedup of 3.5 (over the vector execution on one processor) was achieved when employing four concurrent processors at an efficiency of 88%. This represents a speedup of more than 20 times (per iteration) over the sequential scalar execution.

To achieve near peak performance, vectors of length nR_v should be used, where R_v is the size of vector register (32 or 128) and n is the number of processors. In order to further evaluate the efficiency of the method when the size of the problem grows, the cases of $5\lambda \times 5\lambda$ and $10\lambda \times 10\lambda$ plates were also considered with the

corresponding results reported in Tables 4 and 5. The sampling density in these cases were $25 \text{ unknowns}/\lambda$ with a FFT pad of order one. A summary of the vector-concurrent executions for the plate problem is given in Table 6. The slight decrease in the efficiency, as the problem size becomes larger, is attributed to an increase in the complexity of the memory cache references. However, considering the fact that the cache memory size has been exceeded in all the cases reported here, the performance of the optimized algorithm is quite impressive. Figures 13 to 15 show the corresponding surface current distributions excited on the conducting plates.

CGFFT Code	Execution Mode		
	Scalar (scalar FFT)	Vector (scalar FFT)	Vector (vector FFT)
ELAPSED CPU, sec	148	129	34
VECTOR CPU, sec	-	9	30
VECTORIZABLE CODE, sec	-	28	144
VECTOR CONTENT	-	18.9%	97.3%
VECTOR SPEED UP	-	3.1	4.8
PROGRAM SPEED UP	-	1.15	4.35

Table 1: Performance of the scalar and vectorized 3-d code on the IBM 3090.

IBM 3090 Supercomputer Vector Performance

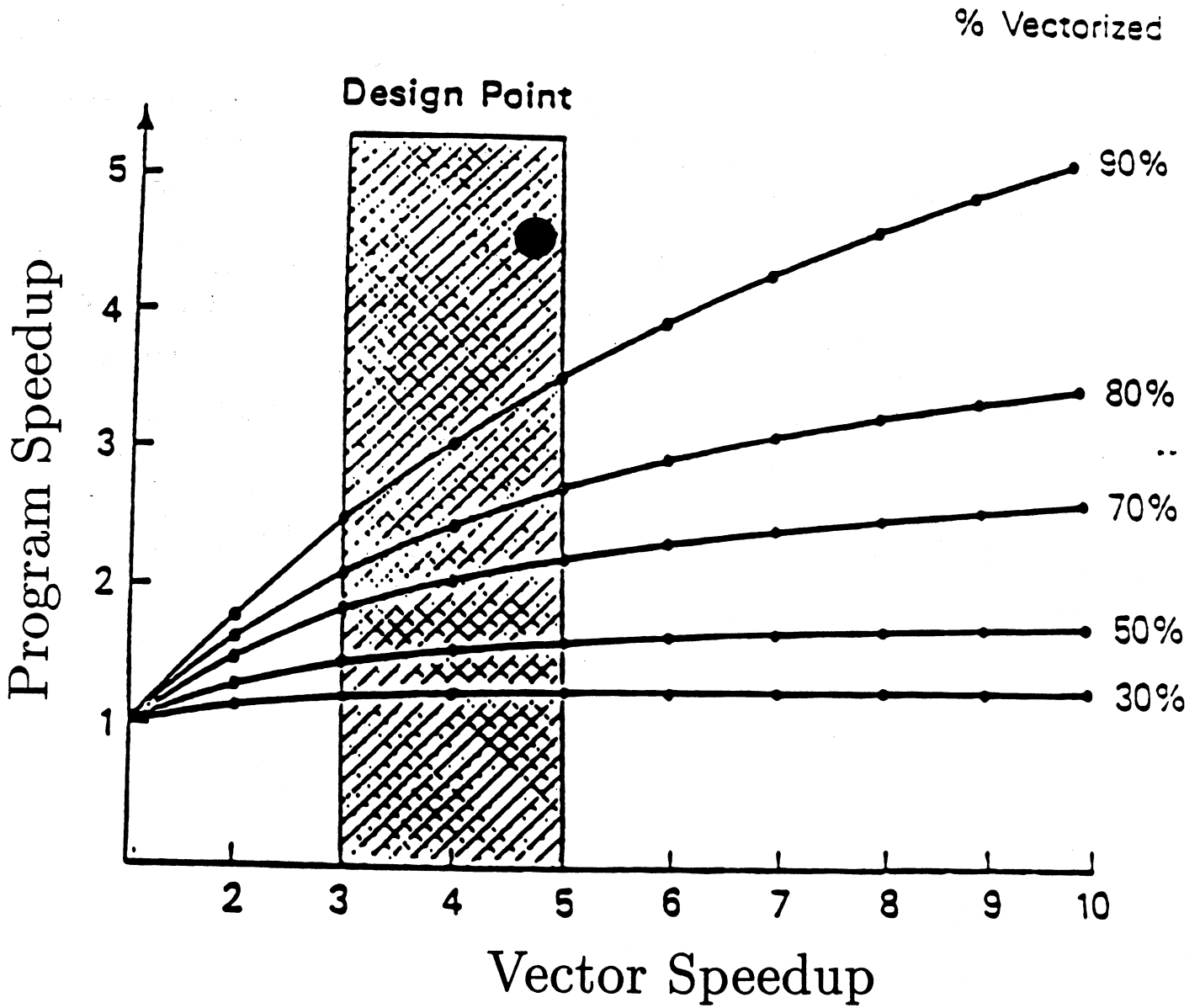


Fig. 11 Performance of the optimized CGFFT algorithm on the IBM 3090.

CGFFT Code	Execution Mode		
	Scalar (scalar FFT)	Vector (scalar FFT)	Vector (vector FFT)
INITIALIZATION, sec	2.39	0.72	0.80
CGFFT LOOP, sec	4307.0	1255.2	342.8
TOTAL CPU TIME, sec	4309.3	1255.9	343.6
ITERATIONS	111	111	59
PER ITERATION, sec	38.80	11.31	5.8
MEGAFLOPS	0.0527	0.1807	0.3512
PROGRAM SPEED UP	-	3.43	12.54
SPEED UP/ITERATION	-	3.43	6.69

Table 2: Performance of the scalar and vectorized 3-d code on the Alliant FX/8.

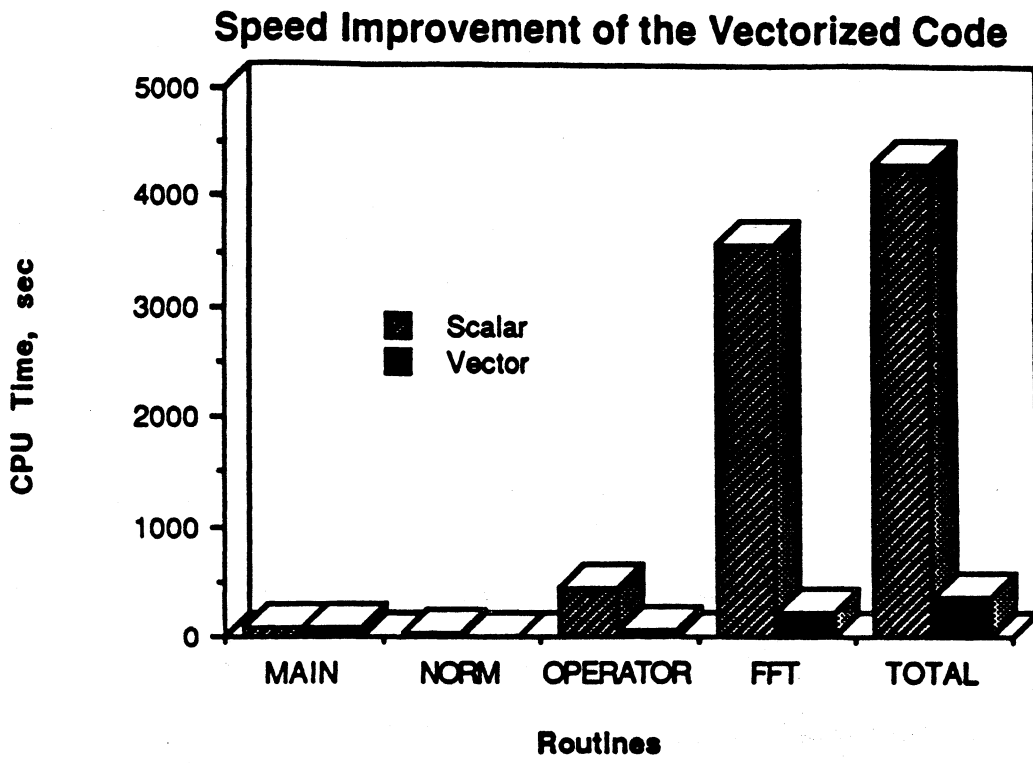


Fig. 12(a) Distribution of CPU time among the compute-intensive routines.

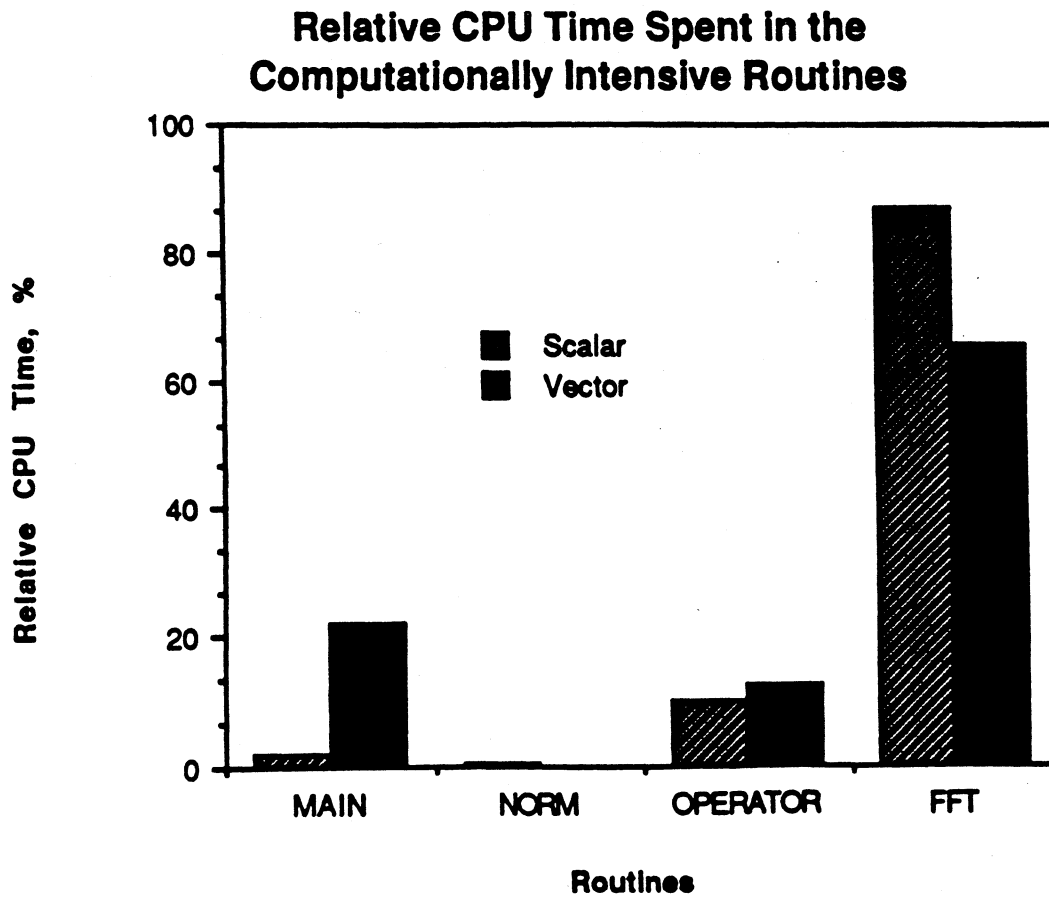


Fig. 12(b) Relative CPU time spent in the compute-intensive routines.

Optimized CGFFT	NO. of processors			
Performance	1	2	3	4
INITIALIZATION, sec	1.60	0.86	0.61	0.50
CGFFT LOOP, sec	407.4	211.0	146.7	115.6
TOTAL CPU TIME, sec	409.0	211.8	147.3	116.1
ITERATIONS	59	59	59	59
PER ITERATION, sec	6.91	3.58	2.49	1.96
MEGAFLOPS	0.2951	0.5697	0.8192	1.0340
SPEED UP	-	1.93	2.78	3.52
EFFICIENCY	-	96.5%	92.7%	88.0%

Table 3: Vector-Concurrent performance for a $2\lambda \times 2\lambda$ plate.

Optimized CGFFT	NO. of processors			
Performance	1	2	3	4
INITIALIZATION, sec	6.35	3.26	2.32	1.90
CGFFT LOOP, sec	1096.6	568.9	398.4	319.0
TOTAL CPU TIME, sec	1103.0	572.2	400.7	320.9
ITERATIONS	46	46	46	46
PER ITERATION, sec	23.98	12.44	8.71	6.98
MEGAFLOPS	0.3848	0.7417	1.0590	1.3226
SPEED UP	-	1.93	2.75	3.43
EFFICIENCY	-	96.5%	91.7%	85.8%

Table 4: Vector-Concurrent performance for a $5\lambda \times 5\lambda$ plate.

Optimized	NO. of processors			
Performance	1	2	3	4
INITIALIZATION, sec	24.55	13.60	9.10	7.23
CGFFT LOOP, sec	3673.3	1973.0	1349.8	1080.5
TOTAL CPU TIME, sec	3697.8	1986.6	1358.9	1087.7
ITERATIONS	38	38	38	38
PER ITERATION, sec	97.31	52.28	35.76	28.62
MEGAFLOPS	0.4223	0.7861	1.1492	1.4357
SPEED UP	-	1.86	2.72	3.40
EFFICIENCY	-	93.0%	90.7%	85.0%

Table 5: Vector-Concurrent performance for a $10\lambda \times 10\lambda$ plate.

Configuration			No. of processors*											
Plate	FFT SIZE		1			2			3			4		
	Pa	KBYTES	MBYTES	MFLOPS	Efficiency	MFLOPS	Speedup	Efficiency	MFLOPS	Speedup	Efficiency	MFLOPS	Speedup	Efficiency
2 λ × 2 λ	128 × 128	128	1.6	0.30	0.57	1.93	96.5%	0.82	2.78	92.7%	1.03	3.52	88.8%	
5 λ × 5 λ	256 × 256	512	6.0	0.38	0.74	1.93	96.5%	1.06	2.75	91.7%	1.32	3.43	85.8%	
10 λ × 10 λ	512 × 512	2048	24.0	0.42	0.79	1.86	93.0%	1.15	2.72	90.7%	1.44	3.40	85.0%	

*The speedup and efficiency are relative to execution on a single processor.

Table 6: A summary of the vector-concurrent performances of the optimized CGFFT algorithm for different configurations.

SCATTERING FROM A CONDUCTING PLATE
E-Pol., Plane Wave Incidence: 0 deg.

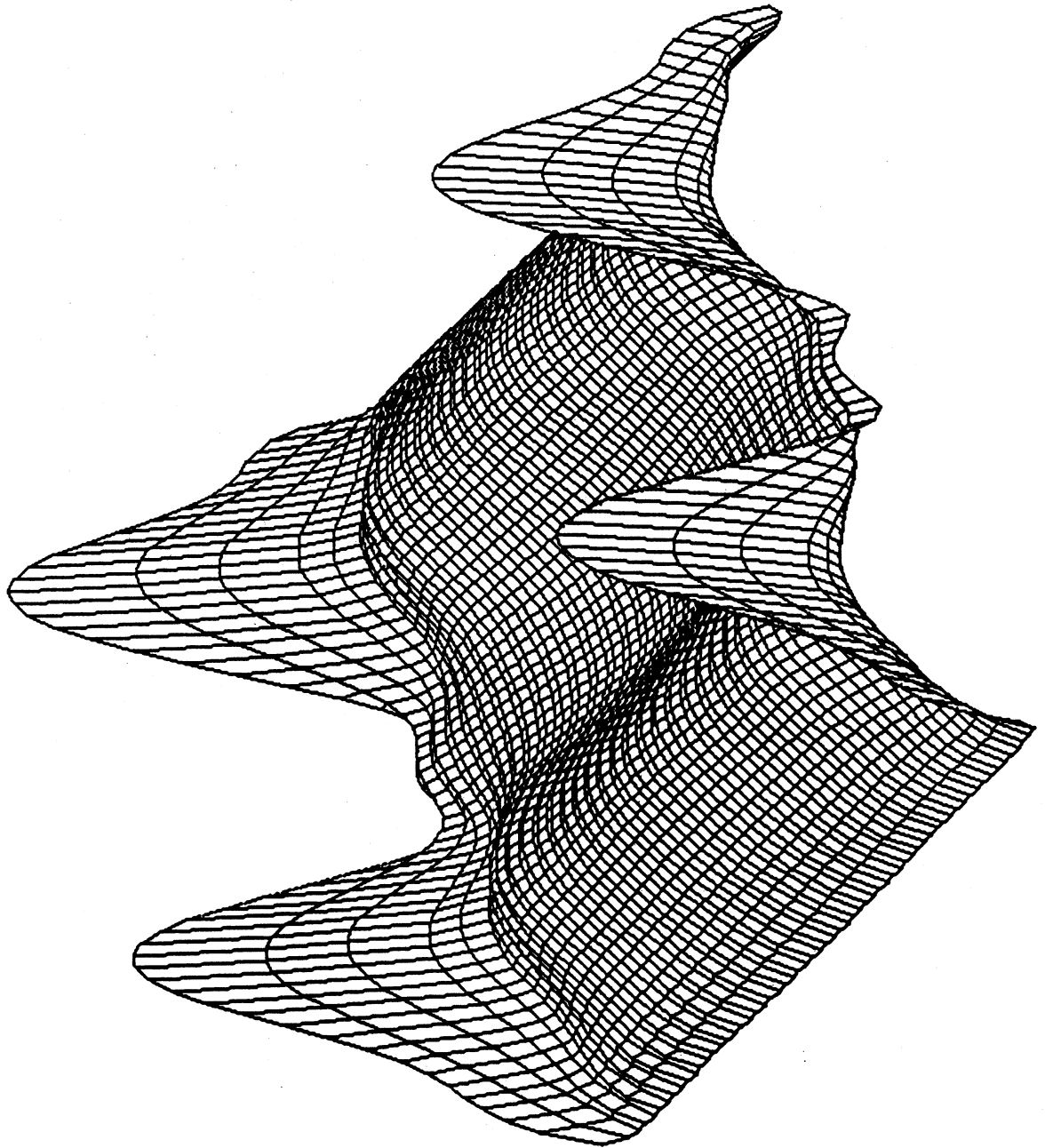
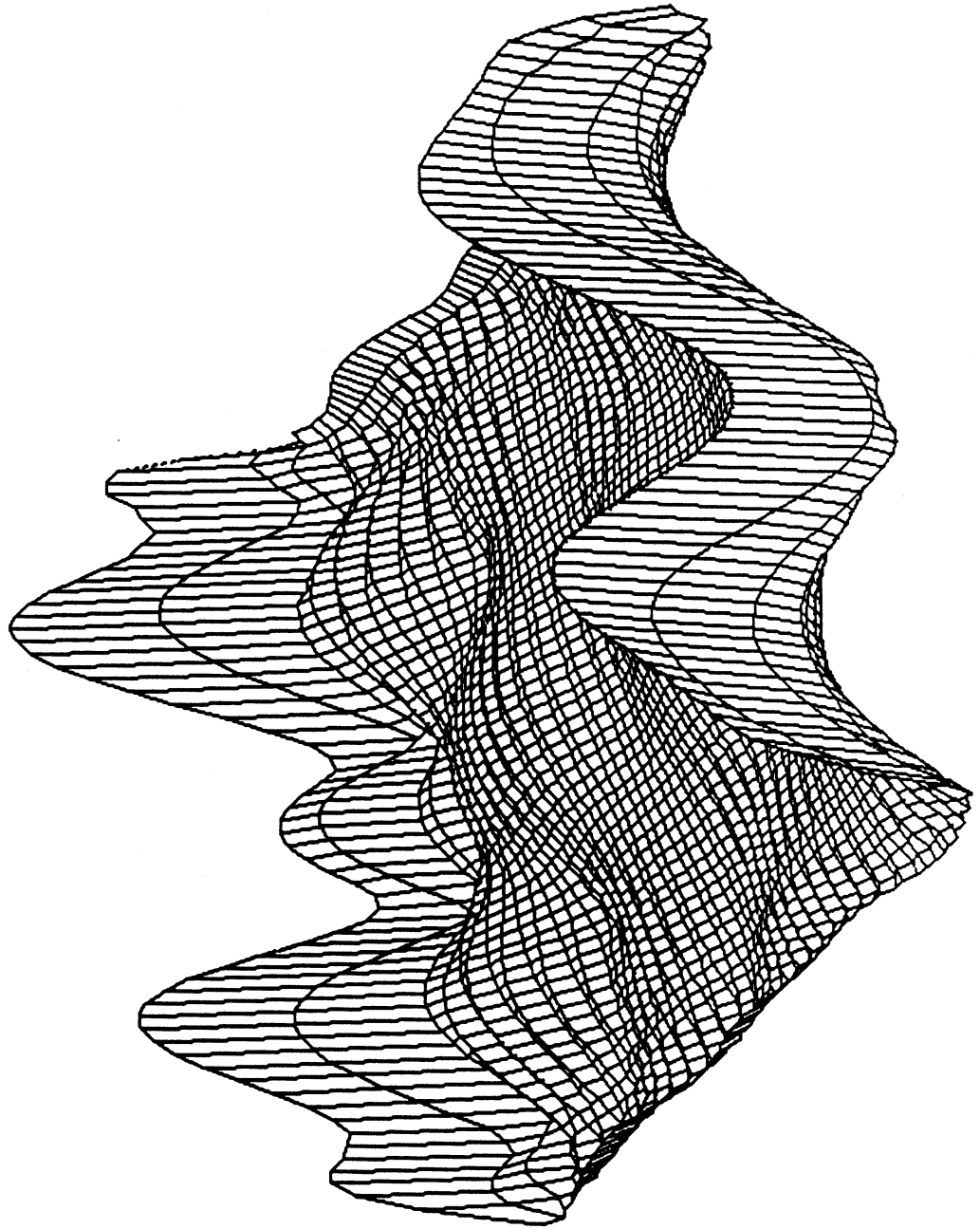


Fig. 13 The surface current density on a $2\lambda \times 2\lambda$ conducting plate illuminated by an E-polarized normally incident plane wave (63×63 unknowns and FFT pad of order 1).
(a) Like-polarized component, $|J_x(x, y)|$

SCATTERING FROM A CONDUCTING PLATE
E-Pol., Plane Wave Incidence: 90 deg.



(b) Cross-polarized component, $|J_y(x, y)|$

SCATTERING FROM A CONDUCTING PLATE
E-Pol., Plane Wave Incidence: 0 deg.

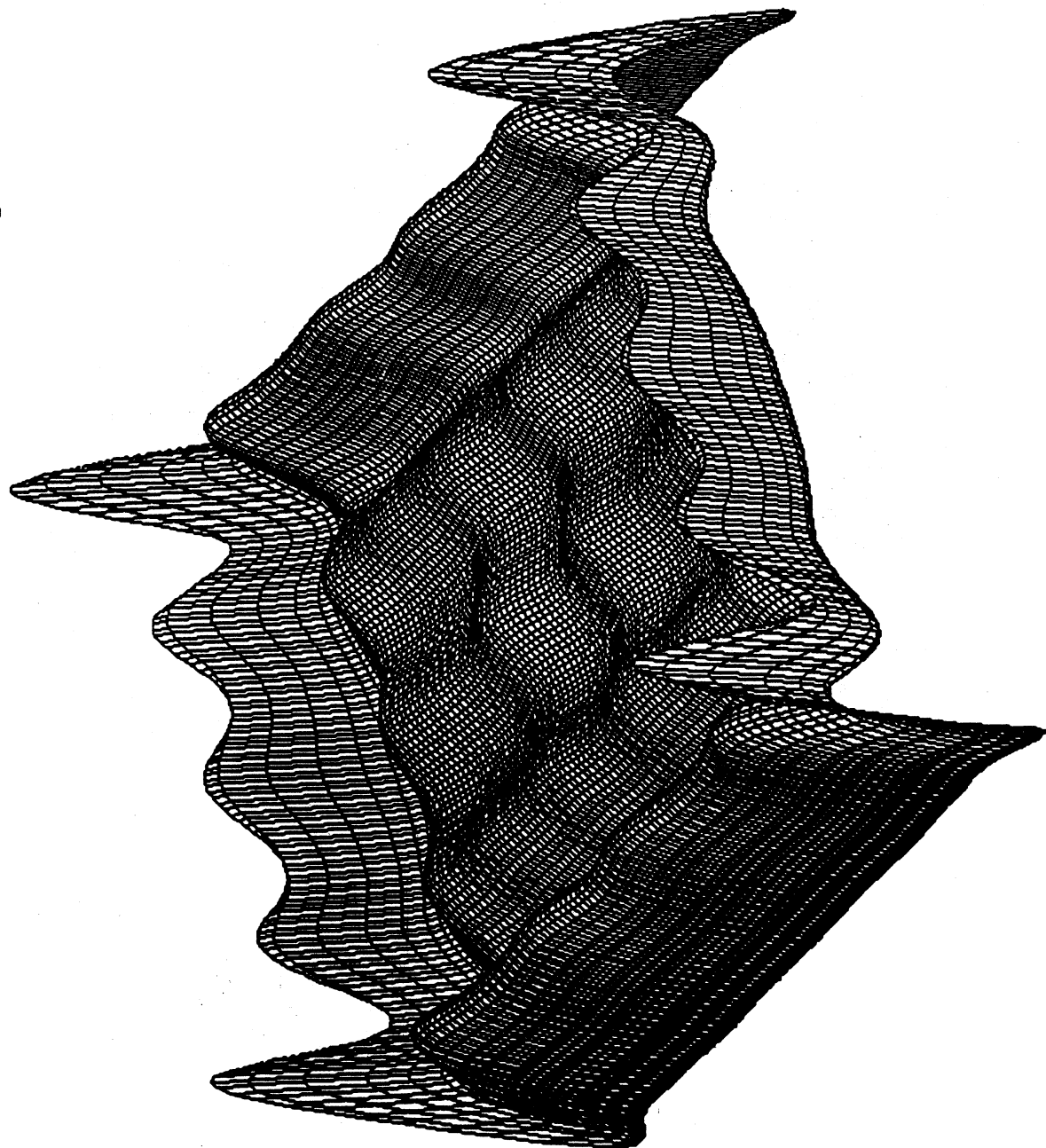
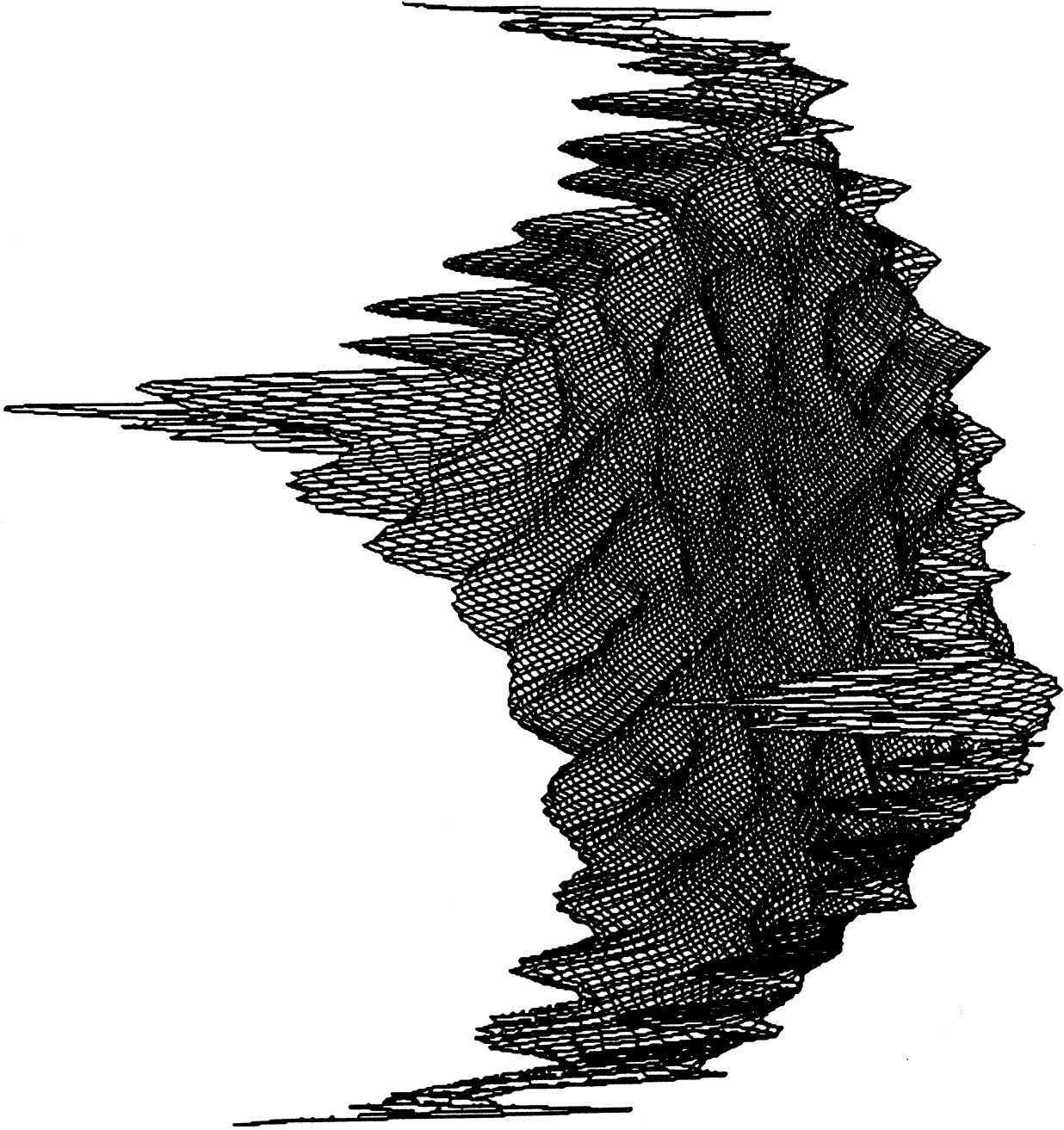


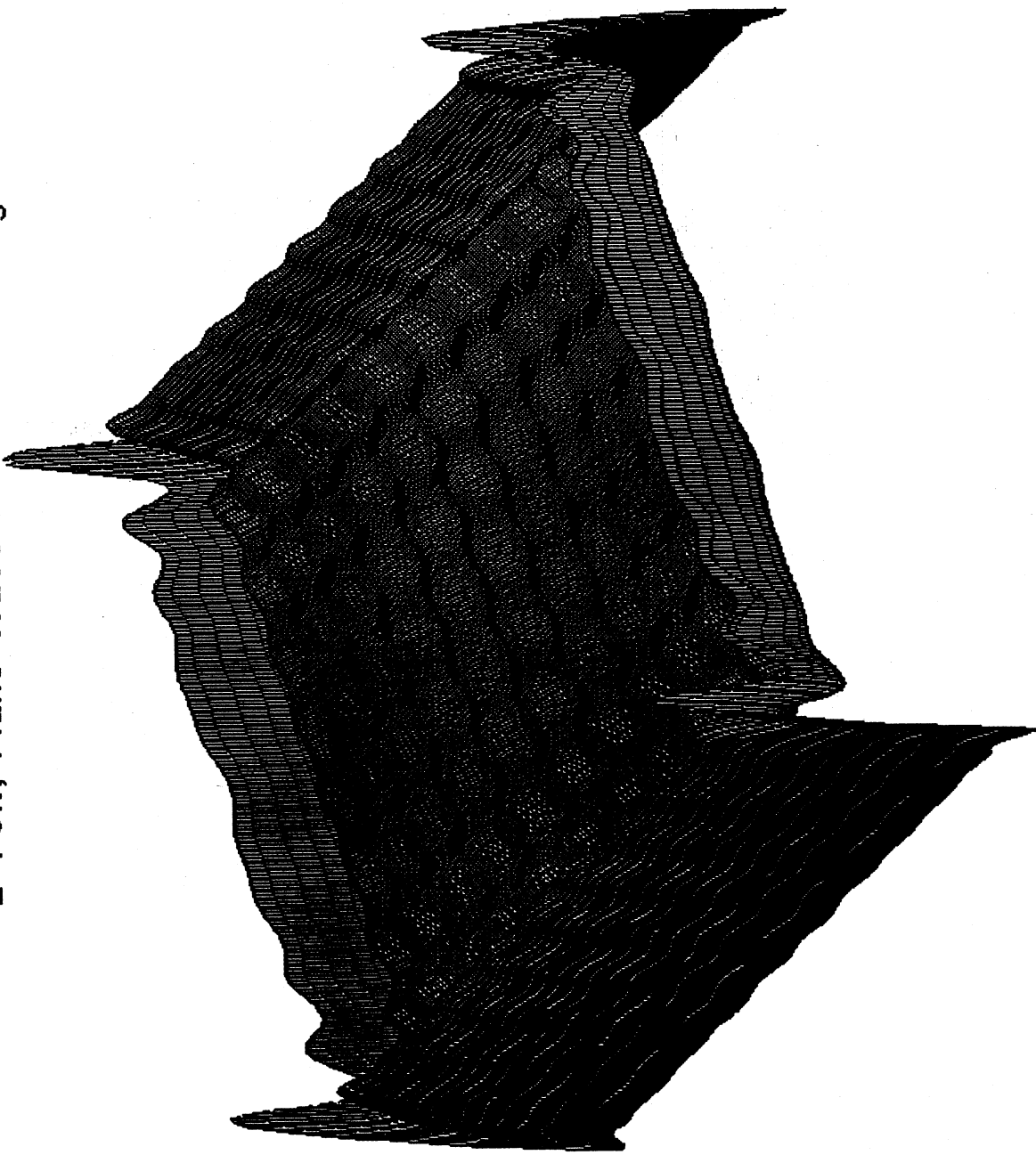
Fig. 14 The surface current density on a $5\lambda \times 5\lambda$ conducting plate illuminated by an E-polarized normally incident plane wave (125×125 unknowns and FFT pad of order 1).
(a) Like-polarized component, $|J_x(x, y)|$

SCATTERING FROM A CONDUCTING PLATE
E-Pol., Plane Wave Incidence: 0 deg.



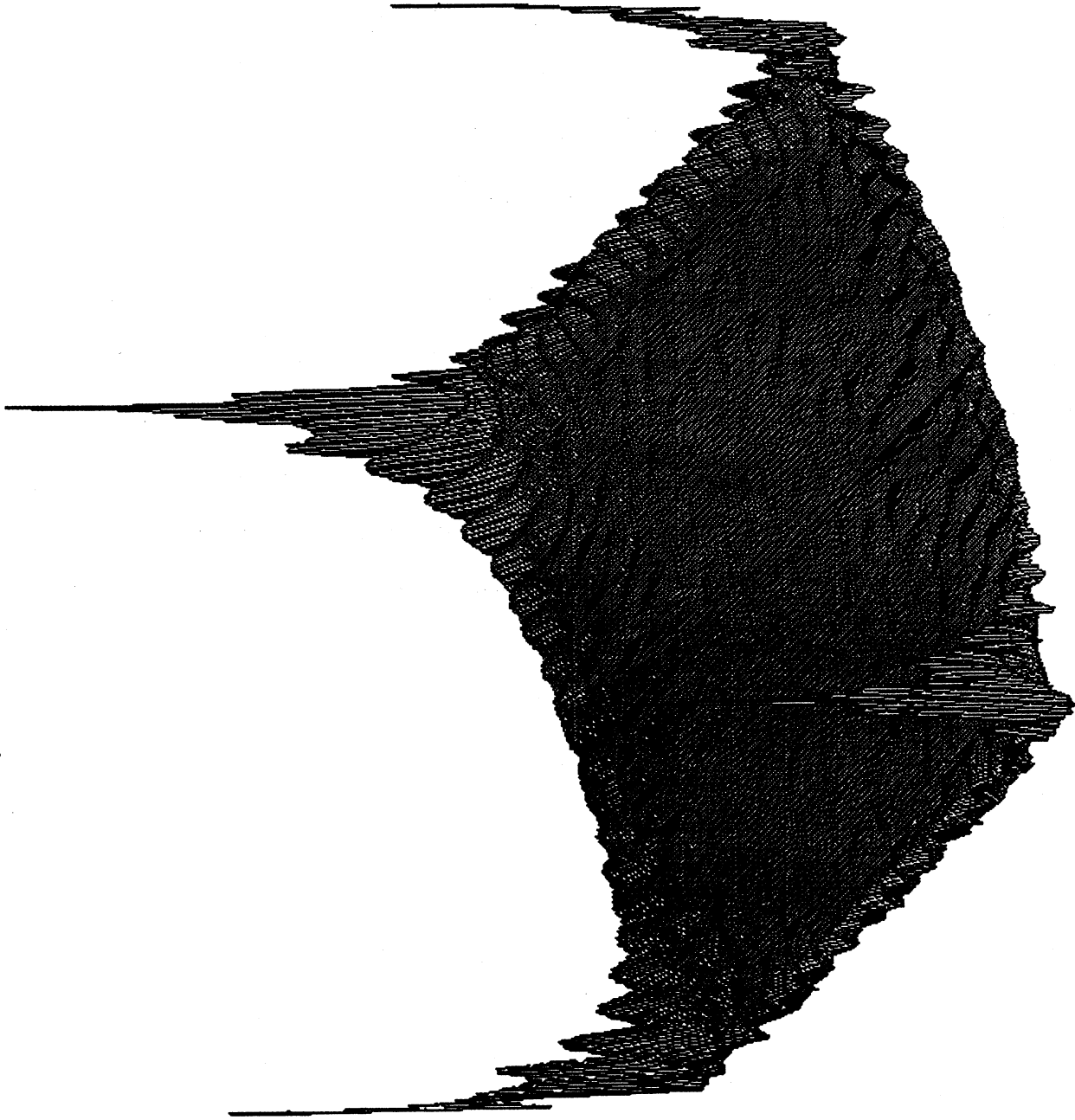
(b) Cross-polarized component, $|J_y(x, y)|$

SCATTERING FROM A CONDUCTING PLATE
E-Pol., Plane Wave Incidence: 0 deg.



(a) Like-polarized component, $|J_x(x, y)|$

E-Pol., Plane Wave Incidence: 0 deg.



(b) Cross-polarized component, $|J_y(x, y)|$

6 Conclusion

It was shown that the conjugate gradient FFT algorithm is suitable for vector-concurrent optimization and may be efficiently implemented on multi-processor computers. The FFT plays a crucial role in the speedup and the efficiency of such an application. As the size of the problem becomes larger, there is a corresponding degradation in the performance of the optimized code due to the complexity of the memory cache references. This complexity, however, has not proven restrictive in the examples considered because the CGFFT method does not suffer from the same memory requirements as the direct methods do. Thus, relatively large problems can be handled without considerable loss of efficiency and speed.

Although this study was concerned with automatic parallelization, which is limited to optimization of individual loops, parallelism at a larger granularity can be specified by the programmer to achieve a superior performance for more complex problems. An example of such an application is the problem of scattering by a dielectric plate of finite thickness [2] where the normal component of the current density is totally independent of the planar components and can be solved for by a dedicated processor in parallel with them.

Acknowledgements

The vector computations on the IBM 3090 were conducted on the Cornell National Supercomputer Facility (CNSF), a resource of the Center for Theory and Simulation in Science and Engineering. The vector-concurrent computations on the Alliant FX/8 were carried out on the University of Michigan's Computer Aided Engineering Network (CAEN).

The author is grateful to Professors John L. Volakis and R. J. Lomax for their support and encouragement in the course of this work.

7 Appendix I: The Fourier Transform Pair

The forward and inverse Fourier transform pair are defined for one and two dimensional cases, respectively, as follows:

$$\tilde{g}(f_x) = \int_{-\infty}^{\infty} g(x) e^{-j2\pi f_x x} dx, \quad (22)$$

$$g(x) = \int_{-\infty}^{\infty} \tilde{g}(f_x) e^{j2\pi f_x x} df_x, \quad (23)$$

and

$$\tilde{g}(f_x, f_y) = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} g(x, y) e^{-j2\pi(f_x x + f_y y)} dx dy, \quad (24)$$

$$g(x, y) = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} \tilde{g}(f_x, f_y) e^{j2\pi(f_x x + f_y y)} df_x df_y. \quad (25)$$

Since in practice the spatial variables are normalized with respect to λ , the free space wave number $k_0 = 2\pi$ and the following relations hold

$$\frac{\partial}{\partial x} \iff j2\pi f_x, \quad \frac{1}{k_0^2} \frac{\partial^2}{\partial x^2} \iff -f_x^2, \quad \frac{1}{k_0^2} \frac{\partial^2}{\partial x \partial y} \iff -f_x f_y.$$

For numerical computations the Discrete Fourier Transform (DFT) is introduced as

$$\hat{J}_k = \sum_{m=1}^N J_m e^{-2\pi j m k / N} \quad k = 1, \dots, N, \quad (26)$$

where each component of $\hat{\mathbf{J}}$ is in effect the Fourier transform of a periodic waveform, one period of which is assumed to be the N samples of the original function. The two dimensional DFT is defined similarly with summations over both dimensions.

The discrete Fourier transforms are carried out by a radix 2 FFT routines. The sampling intervals and the size of the FFT pads are chosen so that the Nyquist criterion in both the spatial and frequency domains as well as the requirements of

linearity in the convolutions are satisfied. Thus, the period N' of the array to be transformed is chosen so that[6]

$$N' = 2^\gamma : N' > N_{Nyquist}, N' > 2 \times N - 1 \quad (27)$$

where N is the number of unknowns and γ is an integer. In practice γ is chosen according to the rule

$$\gamma \geq \log_2(2N - 1) + \rho \quad (28)$$

where ρ is an integer setting the *order* of the FFT pad dimension to ensure adequate frequency sampling in the spectral domain when performing the inverse transform operation. The array elements beyond the scatterer's physical range are set to zero.

8 Appendix II: Optimization Techniques

ARCHITECTURAL TECHNIQUES	PROGRAMMER ACTIONS [†]
SCALAR PIPELINING	Use compiler switches (global optimization).
VECTOR PROCESSING	<ol style="list-style-type: none"> 1. Restructure loops and use compiler directives. 2. Maximize vector lengths by renesting, merging, unrolling loops; largest iterates to be inside. 3. Eliminate conditionals in loops & distribute them. 4. "Supervector" loops to fill vector registers. 5. Move I/O statements out of the loops. 6. Turn off vectorization for short loops, or some vectorized loops with conditionals, dependences.
CONCURRENCY	<ol style="list-style-type: none"> 1. Eliminate/relocate dependences, scalar carry-arounds. 2. Restructure for Concurrent-Outer Vector-Inner. 3. Create concurrently-callable subroutines. 4. Renest/merge loops for more concurrent iterations.
CACHE/MEMORY ACCESS	<ol style="list-style-type: none"> 1. Possible problems with strides. 2. Leftmost array dimension should be the largest. 4. Outer loop corresponds to the leftmost array index. 3. Process compact vectors, columns instead of rows. 4. Localize memory references.

[†]Extracted from IBM 3090 and Alliant FX/8 programming manuals

9 Appendix III: Listings

The source codes and compiler listings for the strip and plate programs are included in this appendix. The listings include informational messages pertaining to the optimization of the respective algorithms. For the plate program, the listings from both the Alliant FX/8 and the IBM 3090 are included.

Alliant FX/Fortran Compiler V4.0.28 (2.24D20) 15 Oct 1988 21:04:05 Page 1
 Source Listing of /u/barkeshli/strip.d/strip.f

```

1      PROGRAM STRIP
2      C      *****
3      C      MAIN
4      C      *****
5      C      TIMING PARAMETERS
6      REAL      ETIME1,ETIME2,ETIMEOH
7      REAL      USTIME1(2),USTIME2(2),UTIMEOH,STIMEOH
8      REAL*8    KX(512),C(512)
9      REAL      X(512),ETA(512),ECHO(360)
10     COMPLEX   J(512),E(512),A(512),R(512),P(512)
11     COMPLEX   DFGL(512),W(512),WORK(384),XJ,SUMIT,MYCSQRT
12     CHARACTER*1 POL,BASIS
13     CHARACTER*30 FILE(4)
14     COMMON    /DIM/MX,NX
15     COMMON    /BLK/DFGL,ETA
16     DATA XJ/(0.,1.),/PI/3.14159262,/TPI/6.28318524/,Z0/376.99111488/
17     C      TIMING OVERHEAD CALCULATION
18     C      TAKE LAST OF THREE TO ELIMINATE CACHE/PAGING EFFECTS
19     CVD$ NOVECTOR
20     CVD$ NOCONCUR
21     DO I=1,3
22         ETIME1=ETIME(USTIME1)
23         ETIME2=ETIME(USTIME2)
24         ETIMEOH=ETIME2-ETIME1
25         UTIMEOH=USTIME2(1)-USTIME1(1)
26         STIMEOH=USTIME2(2)-USTIME1(2)
27     ENDDO
28     C      -----
29     C      INITIALIZATION
30     C      -----
31     OPEN(UNIT=5,FILE='strip.input')
32     READ(5,'(A30)',END=3,ERR=4)(FILE(I),I=1,4)
33     READ(5,'(A1)',END=3,ERR=4)POL
34     READ(5,*,END=3,ERR=4)FMHZ,H0
35     READ(5,*,END=3,ERR=4)ISCAT,I1PHI,I2PHI,INC
36     READ(5,*,END=3,ERR=4)PHI0
37     READ(5,*,END=3,ERR=4)WIDTH,IDIEL,ET,ETC,ETE,TPR
38     READ(5,*,END=3,ERR=4)MX,IPAD
39     READ(5,*,END=3,ERR=4)ITMAX,TOL
40     READ(5,'(A1)',END=3,ERR=4)BASIS
41     CLOSE(5)
42     NX=2** (INT (ALOG (2.*MX-1)/ALOG (2.)))+IPAD)
43     WIDTH2=WIDTH/2.
44     DLTX=WIDTH/MX
45     DLTFX=1./(DLTX*NX)
46     NXH=NX/2+1
47     XL1=300./FMHZ
48     XK0=2.*PI/XL1

```

Alliant FX/Fortran Compiler V4.0.28 (2.24D20) 15 Oct 1988 21:04:05 Page 2
 Source Listing of /u/barkeshli/strip.d/strip.f

```

49      XK2=XK0*XK0
50      FACTOR=XK0/2.
51      C -----
52      C STATEMENT OF THE PROBLEM
53      C -----
54      WRITE(*,*)
55      WRITE(*,*)'SCATTERING FROM A STRIP'
56      WRITE(*,*)'FREQUENCY= ',FMHZ,' MHz, LAMBDA= ',XL1,' m'
57      WRITE(*,*)'STRIP RELATIVE LENGTH; W/LAMDA=',WIDTH/XL1
58      IF (IDIEL.EQ.1) THEN
59          WRITE(*,*)'UNIFORM RESISTIVITY: ETA=',ET
60      ELSE
61          WRITE(*,*)'NON-UNIFORM RESISTIVITY:'
62          WRITE(*,*)'ETA(C)=',ETC,' ETA(E)=',ETE,' TPR=',TPR
63      ENDIF
64      WRITE(*,*)POL,'-POLARIZATION',' PHI0=',PHI0
65      WRITE(*,*)'NO. OF UNKNOWNNS MX= ',MX,' PER LAMBDA= ',MX/WIDTH,
66      +      ' DLTX/LAMBDA= ',DLTX
67      WRITE(*,*)'FFT PAD ',NX,' OF ORDER ',IPAD
68      WRITE(*,*)'SCATTERING PATTERN COMPUTED:'
69      IF (ISCAT.EQ.1) THEN
70          WRITE(*,*)' BISTATIC'
71      ELSE
72          WRITE(*,*)' BACKSCATTERING'
73      ENDIF
74      WRITE(*,*)'TOLERANCE= ',TOL
75      WRITE(*,*)
76      C -----
77      C SPATIAL VARIABLE
78      C -----
79      DO I=1,MX
80          X(I)=-WIDTH2+(I-.5)*DLTX
81      ENDDO
82      C -----
83      C STRIP RESISTIVITY
84      C -----
85      C HOMOGENEOUS DIELECTRIC (UNIFORM)
86      IF (IDIEL.EQ.1) THEN
87          ETA(1:MX)=ET
88      C INHOMOGENEOUS DIELECTRIC (LINEAR OR PARABOLIC)
89      ELSE
90          ETA(1:MX)=ETC+ETE*ABS(X(1:MX)/WIDTH2)**TPR
91      ENDIF
92      C -----
93      C PHYSICAL OPTICS CURRENTS
94      C -----
95      SPH=SIN(PHI0*PI/180.)
96      CPH=COS(PHI0*PI/180.)

```

Alliant FX/Fortran Compiler V4.0.28 (2.24D20) 15 Oct 1988 21:04:05 Page 3
 Source Listing of /u/barkeshli/strip.d/strip.f

```

  97           IF (POL.EQ.'E') THEN
  98             J(1:MX)=SPH/(.5+ETA(1:MX)*SPH)*H0*CEXP(XJ*X(1:MX)*CPH)
  99           ELSE
 100             J(1:MX)=SPH/(ETA(1:MX)+.5*SPH)*H0*CEXP(XJ*X(1:MX)*CPH)
 101           ENDIF
 102 C -----
 103 C SPECTRAL VARIABLES AND TRANSFORM OF THE KERNEL
 104 C -----
 105           DO I = 1,NXH
 106             KX(I)=TPI*(I-1)*DLTFX
 107           ENDDO
 108           DO I = NXH+1,NX
 109             KX(I)=TPI*(I-NX-1)*DLTFX
 110           ENDDO
 111           C(1:NX)=DBLE(XK2)-KX(1:NX)*KX(1:NX)
 112 CVD$ CNCALL
 113           DO I=1,NX
 114             DFGL(I)=FACTOR/MYCSQRT(C(I))
 115           ENDDO
 116 C           DFGL(1:NX)=FACTOR/MYCSQRT(C(1:NX))
 117           IF (POL.EQ.'H') THEN
 118             DFGL(1:NX)=SNGL(C(1:NX))/XK2*DFGL(1:NX)
 119           ENDIF
 120           IF (BASIS.EQ.'C') THEN
 121 CVD$ CNCALL
 122             DFGL(1:NX)=DFGL(1:NX)*SINC(KX(1:NX)*DLTX/2.)
 123             ELSE IF (BASIS.EQ.'S') THEN
 124 CVD$ CNCALL
 125               DFGL(1:NX)=
 126 +             DFGL(1:NX)*2.*XK0*(COS(KX(1:NX)*DLTX)-COS(XK0*DLTX))/
 127 +             (DLTX*SNGL(C(1:NX))*SIN(XK0*DLTX))
 128             ENDIF
 129 C           INITIALIZE THE WORKING AREAS
 130             A(1:512)=(0.,0.)
 131             W(1:512)=(0.,0.)
 132             WORK(1:384)=(0.,0.)
 133             OPEN(UNIT=4,FILE=FILE(4))
 134             KK=PHI0
 135 C -----
 136 C START ON THE NEW ANGLE OF INCIDENCE
 137 C -----
 138 1           PH0=PHI0*PI/180.
 139             ITER=0
 140             RSS=0.
 141             CPH0=COS(PH0)
 142 C -----
 143 C INCIDENT FIELD
 144 C -----

```

Alliant FX/Fortran Compiler V4.0.28 (2.24D20) 15 Oct 1988 21:04:05 Page 4
 Source Listing of /u/barkeshli/strip.d/strip.f

```

145         IF (POL.EQ.'E') THEN
146             E(1:MX)=H0*CEXP(XJ*XK0*X(1:MX)*CPH0)
147         ELSE
148             SPH0=SIN(PH0)
149             E(1:MX)=H0*SPH0*CEXP(XJ*XK0*X(1:MX)*CPH0)
150         ENDIF
151         ENORM=SUM(CABS(E(1:MX))*CABS(E(1:MX)))
152     C     *-----*
153     C     BEGIN TIMING THE CGFFT LOOP
154     C     *-----*
155         ETIME1=ETIME(USTIME1)
156     C     -----
157     C     CGFFT INITIALIZATION
158     C     -----
159         CALL OPERATOR(A,J,W,WORK,-1)
160         R(1:MX)=A(1:MX)-E(1:MX)
161         P(1:MX)=(0.,0.)
162     C     -----
163     C     CGFFT MAIN ITERATION LOOP
164     C     -----
165     2     CONTINUE
166         CALL OPERATOR(A,R,W,WORK,1)
167         B=1./SUM(CABS(A(1:MX))*CABS(A(1:MX)))
168         P(1:MX)=P(1:MX)-B*A(1:MX)
169         CALL OPERATOR(A,P,W,WORK,-1)
170         T=1./SUM(CABS(A(1:MX))*CABS(A(1:MX)))
171         R(1:MX)=R(1:MX)+T*A(1:MX)
172         J(1:MX)=J(1:MX)+T*P(1:MX)
173         ITER=ITER+1
174         RSS=SQRT(SUM(CABS(R(1:MX))*CABS(R(1:MX)))/ENORM)
175     C     -----
176     C     END OF CGFFT MAIN ITERATION LOOP
177     C     -----
178     C     WRITE(*,*)'ITER= ',ITER,' RSS= ',RSS
179         WRITE(4,*)ITER,RSS
180         IF (RSS.LE.TOL) THEN
181             WRITE(*,*)'CONVERGENCE ACHIEVED'
182         ELSE IF (ITER.EQ.ITMAX) THEN
183             WRITE(*,*)'ITMAX EXCEEDED; CONVERGENCE CRITERION NOT MET.'
184         ELSE
185             GO TO 2
186         END IF
187     C     *-----*
188     C     END TIMING THE CGFFT LOOP
189     C     *-----*
190         ETIME2=ETIME(USTIME2)
191     C     -----
192     C     BISTATIC COMPUTATIONS

```

Alliant FX/Fortran Compiler V4.0.28 (2.24D20) 15 Oct 1988 21:04:05 Page 5
Source Listing of /u/barkeshli/strip.d/strip.f

```

193 C -----
194 CVD$ CNCALL
195     IF (ISCAT.EQ.1) THEN
196         DO K=I1PHI, I2PHI, INC
197             PH=K*PI/180.
198             CPH=COS (PH)
199             SUMIT=(0.,0.)
200             DO I=1, MX
201                 SUMIT=SUMIT+J (I) *CEXP (XJ*XK0*X (I) *CPH)
202             ENDDO
203             IF (POL.EQ.'E') THEN
204                 ECHO (K) = (XK0*DLTX*SINC (XK0*DLTX/2.*CPH) ) **2
205             +             *CABS (SUMIT) **2 / (8.*PI)
206             ELSE
207                 SPH=SIN (PH)
208                 ECHO (K) = (XK0*DLTX*SPH*SINC (XK0*DLTX/2.*CPH) ) **2
209             +             *CABS (SUMIT) **2 / (8.*PI)
210             ENDIF
211         ENDDO
212 C     SAVE CURRENT VALUES
213         OPEN (UNIT=1, FILE=FILE (1) )
214         WRITE (1, ' (A, /, I3) ') 'STRIP SURFACE CURRENT; INTEGRAL EQN', MX
215         WRITE (1, ' (2 (1X, F10.4) ) ') (X (I), CABS (J (I) ), I=1, MX)
216         CLOSE (1)
217 C     COMPUTE AND SAVE CURRENT SPECTRUM
218         CALL CFFT1D (NX, W, 0, WORK)
219         W (1:MX) = J (1:MX)
220         W (MX+1:NX) = (0., 0.)
221         CALL CFFT1D (NX, W, -1, WORK)
222         OPEN (UNIT=2, FILE=FILE (2) )
223         WRITE (2, ' (A, /, I3) ') 'CURRENT SPECTRUM', NX
224         WRITE (2, ' (2 (1X, F10.4) ) ') (KX (I), CABS (W (I) ), I=NXH+1, NX)
225         WRITE (2, ' (2 (1X, F10.4) ) ') (KX (I), CABS (W (I) ), I=1, NXH)
226         CLOSE (2)
227 C -----
228 C     BACKSCATTERING COMPUTATIONS
229 C -----
230         ELSE IF (ISCAT.EQ.2) THEN
231             CPH=COS (PH0)
232             SUMIT=(0.,0.)
233             DO I=1, MX
234                 SUMIT=SUMIT+J (I) *CEXP (XJ*XK0*X (I) *CPH)
235             ENDDO
236             IF (POL.EQ.'E') THEN
237                 ECHO (KK) = (XK0*DLTX*SINC (XK0*DLTX/2.*CPH) ) **2
238             +             *CABS (SUMIT) **2 / (8.*PI)
239             ELSE IF (POL.EQ.'H') THEN
240                 SPH=SIN (PH0)

```

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 Source Listing of /u/barkeshli/strip.d/strip.f

```

241          ECHO(KK)=(XK0*DLTX*SPH*SINC(XK0*DLTX/2.*CPH))**2
242      +          *CABS(SUMIT)**2/(8.*PI)
243          ENDIF
244          KK=KK+INC
245          PHI0=PHI0+INC
246          IF(PHI0.LE.FLOAT(I2PHI)) THEN
247              WRITE(*,*)'ANGLE OF INCIDENCE: ',PHI0
248              GOTO 1
249          ENDIF
250      ENDIF
251      WHERE(ECHO(I1PHI:I2PHI:INC).NE.0.)
252          ECHO(I1PHI:I2PHI:INC)=10.*ALOG10(ECHO(I1PHI:I2PHI:INC))
253      OTHERWISE
254          ECHO(I1PHI:I2PHI:INC)=1.E-15
255      END WHERE
256      OPEN(UNIT=3,FILE=FILE(3))
257      WRITE(3,*)'STRIP ECHOWIDTH; PHI0= ',PHI0
258      WRITE(3,*)(I2PHI-I1PHI)/INC+1
259      WRITE(3,'(2(1X,F10.4))')(FLOAT(K),ECHO(K),K=I1PHI,I2PHI,INC)
260      CLOSE(3)
261      CLOSE(4)
262      WRITE(*,*)'NORMAL TERMINATION'
263      WRITE(*,*)
264      PRINT *,'ITERATIONS: ',ITER
265      WRITE(*,*)'CPU TIME INFORMATION(msec.):'
266      PRINT *,'OVERHEADS (TOTAL,USER,SYSTEM)'
267      +      ,1000.*ETIMEOH,1000.*UTIMEOH,1000.*STIMEOH
268      PRINT *,'START TIME: ',1000.*ETIME1,1000.*USTIME1
269      PRINT *,'END TIME: ',1000.*ETIME2,1000.*USTIME2
270      PRINT *,'TOTAL TIME: ',1000.*(ETIME2-ETIME1-ETIMEOH)
271      PRINT *,'USER TIME: ',1000.*(USTIME2(1)-USTIME1(1)-UTIMEOH)
272      PRINT *,'SYSTEM TIME: ',1000.*(USTIME2(2)-USTIME1(2)-STIMEOH)
273      NOPER=(4.*NX*ALOG(FLOAT(NX))/ALOG(2.))+10.*MX+2)*ITER
274      PRINT *,'MEGAFLOPS: ',NOPER/((ETIME2-ETIME1-ETIMEOH)*1.0E6)
275      PRINT *,'USER MFLOPS: ',NOPER/((USTIME2(1)-USTIME1(1)-UTIMEOH)
276      +      *1.0E6)
277      GO TO 5
278      3 PRINT *,'END OF INPUT FILE; MORE INPUT EXPECTED'
279      4 PRINT *,'ERROR IN INPUT DATA; CHECK INPUT FILE'
280      5 STOP
281      END

```

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 Source Listing of /u/barkeshli/strip.d/strip.f

----- LOOP SUMMARY FOR ROUTINE STRIP -----								
LABEL	INDEX	START	END	NEST	COMMENT	CD?	DP?	ITERATIONS
	I	21	27	1	NO ACTION REQUESTED			3
	I	79	81	1	VECTOR-CONCURRENT			MX
		87	87	1	VECTOR-CONCURRENT			MX
		90	90	1	VECTOR-CONCURRENT			MX
		98	98	1	VECTOR-CONCURRENT			MX
		100	100	1	VECTOR-CONCURRENT			MX
	I	105	107	1	VECTOR-CONCURRENT			NXH
	I	108	110	1	VECTOR-CONCURRENT			NX-NXH
		111	111	1	VECTOR-CONCURRENT			NX
	I	113	115	1	CONCURRENT			NX
		118	118	1	VECTOR-CONCURRENT			NX
		125	125	1	VECTOR-CONCURRENT			NX
		130	131	1	VECTOR-CONCURRENT			512
		132	132	1	VECTOR-CONCURRENT			384
		146	146	1	VECTOR-CONCURRENT			MX
		149	149	1	VECTOR-CONCURRENT			MX
		151	151	1	VECTOR-CONCURRENT			MX
		160	161	1	VECTOR-CONCURRENT			MX
		167	167	1	VECTOR-CONCURRENT			MX
		168	168	1	VECTOR-CONCURRENT			MX
		170	170	1	VECTOR-CONCURRENT			MX
		171	172	1	VECTOR-CONCURRENT			MX
		174	174	1	VECTOR-CONCURRENT			MX
	K	196	211	1	CONCURRENT OUTER			
	I	200	202	2	VECTORIZED			MX
		219	219	1	VECTOR-CONCURRENT			MX
		220	220	1	VECTOR-CONCURRENT			NX-MX
	I	233	235	1	VECTOR-CONCURRENT			MX
		251	255	1	VECTOR-CONCURRENT		Y	
1	VECTOR							
1	NOT OPTMZD							
29	TOTAL							
		1			CONCURRENT			
		0			NOT INNER			
		29			EXAMINED			
		25			VECT-CONCR			
		0			TOO SHORT			
		28			OPTIMIZED			
		1			CON. OUTER			
		0			DEPENDENT			

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 Source Listing of /u/barkeshli/strip.d/strip.f

```

282 C *****
283 C             INTEGRO-DIFFERENTIAL OPERATOR
284 C *****
285     SUBROUTINE OPERATOR(A,Z,W,WORK,ISTAR)
286     COMPLEX    A(*),Z(*),DFGL(512),W(*),WORK(*)
287     REAL      ETA(512)
288     COMMON    /DIM/MX,NX
289     COMMON    /BLK/DFGL,ETA
290 C -----
291 C     FFT INITIALIZATION
292 C -----
293     CALL CFFT1D(NX,W,0,WORK)
294 C -----
295 C     FORWARD FFT
296 C -----
297     W(1:MX)=Z(1:MX)
298     W(MX+1:NX)=(0.,0.)
299     CALL CFFT1D(NX,W,-1,WORK)
300 C -----
301 C     DISCRETE CONVOLUTION
302 C -----
303     IF (ISTAR.EQ.-1) THEN
304     W(1:NX)=DFGL(1:NX)*W(1:NX)
305     ELSE
306     W(1:NX)=CONJG(DFGL(1:NX))*W(1:NX)
307     ENDIF
308 C -----
309 C     INVERSE FFT
310 C -----
311     CALL CFFT1D(NX,W,1,WORK)
312     A(1:MX)=ETA(1:MX)*Z(1:MX)+W(1:MX)
313     RETURN
314     END

```


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 Source Listing of /u/barkeshli/strip.d/strip.f

----- LOOP SUMMARY FOR ROUTINE OPERATOR -----						
LABEL	INDEX	START	END	NEST	COMMENT	CD? DP? ITERATIONS
		297	297	1	VECTOR-CONCURRENT	MX
		298	298	1	VECTOR-CONCURRENT	NX-MX
		304	304	1	VECTOR-CONCURRENT	NX
		306	306	1	VECTOR-CONCURRENT	NX
		312	312	1	VECTOR-CONCURRENT	MX
0	VECTOR	0	CONCURRENT	5	VECT-CONCR	0 CON. OUTER
0	NOT OPTMZD	0	NOT INNER	0	TOO SHORT	0 DEPENDENT
5	TOTAL	5	EXAMINED	5	OPTIMIZED	

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Source Listing of /u/barkeshli/strip.d/strip.f

```
315 C *****  
316 C           SQUARE ROOT OF A SIGNED REAL NUMBER  
317 C *****  
318     COMPLEX FUNCTION MYCSQRT(A)  
319     REAL*8  A  
320     COMPLEX XJ  
321     DATA  XJ/(0.,1.)/  
322     IF (A.GE.0D0) THEN  
323         MYCSQRT=SNGL(DSQRT(A))  
324     ELSE  
325         MYCSQRT=XJ*SNGL(DSQRT(-A))  
326     ENDIF  
327     RETURN  
328     END
```

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Source Listing of /u/barkeshli/strip.d/strip.f

```
329 C *****  
330 C SINC  
331 C *****  
332 REAL FUNCTION SINC(A)  
333 IF (A.NE.0) THEN  
334 SINC=SIN(A) /A  
335 ELSE  
336 SINC=1.  
337 ENDIF  
338 RETURN  
339 END
```


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 Source Listing of /u/barkeshli/slab.d/platevc.f

```

1      PROGRAM SLAB
2      C *****
3      C *SOLUTION OF A THIN DIELECTRIC SLAB BY COMPLEX CONJUGATE-FFT METHOD*
4      C * VECTOR-CONCURRENT IMPLEMENTATION ON ALLIANT SUPER-MINI COMPUTER *
5      C *****
6      C MAIN PROGRAM
7      C *****
8      C PARAMETER DESCRIPTION
9      C =====
10     C DLX,DLY.....RELATIVE LENGTH IN EACH DIMENSION
11     C IN TERMS OF LAMBDA
12     C UPL.....NO. OF UNKNOWNNS PER WAVELENGTH
13     C MX,MY.....NO. OF UNKNOWNNS PER DIMENSION
14     C NX,NY.....COMPOSITE NO. OF SAMPLES IN EACH
15     C DIMENSION SUCH THAT :
16     C N > 2*M | N = 2**L , L-INTEGERS
17     C E.....INCIDENT FIELD VECTOR
18     C JS.....SURFACE CURRENT VECTOR
19     C A.....INTEGRO-DIFFERENTIAL OPERATOR MX
20     C ER.....SLAB PERMITTIVITY
21     C TK.....SLAB RELATIVE THICKNESS
22     C R,P,T,B.....CG VARIABLES
23     C ITMAX.....MAX. NO. OF ITERATIONS
24     C TOL.....CONVERGENCE TOLERANCE
25     C RSS.....RELATIVE RESIDUAL ERROR
26     C ZT,W.....WORKING AREA
27     C ISORCE.....SOURCE TYPE
28     C 1 = E-POL
29     C 2 = H-POL
30     C TETA,PHI.....ANGLES OF INCIDENCE FOR PLANE WAVE
31     C 0 < TETA < PI
32     C 0 < PHI < 2*PI
33     C PSI.....POLARIZATION ANGLE
34     C PLANE WAVE INCIDENCE PSI = 90 E-POL, IDIM=2
35     C = 0 H-POL, IDIM=3
36     C =====
37     REAL ETIME1,ETIME2,ETIME3,ETIMEOH
38     REAL USTIME1(2),USTIME2(2),USTIME3(2),UTIMEOH,STIMEOH
39     INTEGER IRING(:)
40     REAL DFG(:)
41     COMPLEX JS(:),E(:),A(:),R(:),P(:)
42     COMPLEX ZT(:),AT(:),W(:)
43     COMPLEX XJ,RS,CRS,FCT,ER
44     COMMON /DIM/MX,MY,MT,MTI,NX,NY,NXH,NYH,NT,NTI,NG,NA,DLT,XK1,Z0
45     COMMON /PAR/RS,CRS,FCT
46     COMMON IBASIS
47     DATA PI/3.141592653589793/,TPI/6.28318530717959/,XJ/(0.,1.0)/
48     C -----
    
```

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 Source Listing of /u/barkeshli/slab.d/platevc.f

```

49 C ESTIMATE THE TIMING OVERHEAD
50 C -----
51 CVD$ NOVECTOR
52 CVD$ NOCONCUR
53 DO I=1,3
54 ETIME1=ETIME(USTIME1)
55 ETIME2=ETIME(USTIME2)
56 ETIMEOH=ETIME2-ETIME1
57 UTIMEOH=USTIME2(1)-USTIME1(1)
58 STIMEOH=USTIME2(2)-USTIME1(2)
59 ENDDO
60 C -----
61 C START TIMING
62 C -----
63 ETIME1=ETIME(USTIME1)
64 C =====
65 C INITIALIZATION
66 C =====
67 OPEN(UNIT=10)
68 READ(10,*)E0
69 READ(10,*)TETA,PHI,PSI,TMAX,INCT
70 READ(10,*)DLX,DLY,UPL,IPAD
71 READ(10,*)ER,TK
72 READ(10,*)IBASIS
73 READ(10,*)ITMAX,ITMIN,TOL
74 CLOSE(10)
75 Z0=120.*PI
76 XL1=1.
77 DLT=XL1/UPL
78 XK1=TPI/XL1
79 MX=DLX*UPL
80 MY=DLY*UPL
81 MT=MX*MY
82 MTI=2*MT
83 NX=2** (INT (ALOG (2.*MX-1) /ALOG (2.)) +IPAD)
84 NY=2** (INT (ALOG (2.*MY-1) /ALOG (2.)) +IPAD)
85 NXH=NX/2+1
86 NYH=NY/2+1
87 NT=NX*NY
88 NTI=2*NT
89 NG=4*NT
90 NA=NX+3*NY/4+1
91 NP=(TMAX-TETA)/INCT+1
92 IFILOFF=12
93 RS=-XJ*Z0/((ER-1)*XK1*TK)
94 CRS=CONJG(RS)
95 FCT=XJ*Z0/XK1
96 C -----

```

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 Source Listing of /u/barkeshli/slab.d/platevc.f

```

 97 C   ALLOCATE ARRAYS
 98 C   -----
 99       ALLOCATE (IRING (MT) , DFG (NG) )
100       ALLOCATE (JS (MTI) , E (MTI) , A (MTI) , R (MTI) , P (MTI) )
101       ALLOCATE (ZT (NTI) , AT (NTI) , W (NT) )
102 C   -----
103 C   FORCE CACHING
104 C   -----
105       IRING (1:MT)=0
106       DFG (1:NG)=0.
107       JS (1:MTI) = (0.,0.)
108       E (1:MTI) = (0.,0.)
109       A (1:MTI) = (0.,0.)
110       R (1:MTI) = (0.,0.)
111       P (1:MTI) = (0.,0.)
112       ZT (1:NTI) = (0.,0.)
113       AT (1:NTI) = (0.,0.)
114       W (1:NT) = (0.,0.)
115 C   =====
116 C   COMPUTE THE TRANSFORM OF THE GREEN'S FUNCTION
117 C   =====
118       CALL DIFREQ (DFG, IRING, IRMAX)
119       OPEN (UNIT=11)
120 10    CONTINUE
121       ITER = 0
122       RSS = 0.
123       WRITE (11, 200) TETA, PHI, PSI
124 C   =====
125 C   CALCULATE THE INCIDENT FIELD
126 C   =====
127       CALL INCIDENCE (E, E0, TETA, PHI, PSI)
128       ENORM=SUM (CABS (E (1:MTI) ) *CABS (E (1:MTI) ) )
129 C   -----
130 C   BEGIN TIMING THE CGFFT LOOP
131 C   -----
132       ETIME2=ETIME (USTIME2)
133 C   =====
134 C   ZEROTH ORDER RESIDUAL CALCULATION
135 C   =====
136       CALL OPERAT (A, -1, JS, ZT, DFG, IRING, IRMAX, AT, W)
137       R (1:MTI) = A (1:MTI) -E (1:MTI)
138       P (1:MTI) = (0.,0.)
139 C   =====
140 C   CGFFT MAIN LOOP
141 C   =====
142 30    CONTINUE
143       CALL OPERAT (A, 1, R, ZT, DFG, IRING, IRMAX, AT, W)
144       B=1./SUM (CABS (A (1:MTI) ) *CABS (A (1:MTI) ) )

```

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 Source Listing of /u/barkeshli/slab.d/platevc.f

```

145      P(1:MTI) = P(1:MTI)-B*A(1:MTI)
146      CALL OPERAT(A,-1,P,ZT,DFG,IRING,IRMAX,AT,W)
147      T=1./SUM(CABS(A(1:MTI))*CABS(A(1:MTI)))
148      R(1:MTI) = R(1:MTI)+T*A(1:MTI)
149      JS(1:MTI) = JS(1:MTI)+T*P(1:MTI)
150      ITER = ITER+1
151      RSS = SQRT(SUM(R(1:MTI)*R(1:MTI))/ENORM)
152      WRITE(11,300)ITER,RSS
153      IF((RSS.LE.TOL).AND.(ITER.GT.ITMIN)) THEN
154          WRITE(11,400)
155      ELSE IF(ITER.GE.ITMAX) THEN
156          WRITE(11,500)
157      ELSE
158          GO TO 30
159      END IF
160      90 CONTINUE
161      C -----
162      C END TIMING THE CGFFT LOOP
163      C -----
164      ETIME3=ETIME(USTIME3)
165      CALL OUTPUT(IFILOFF,1,JS,ZT,TETA,SIG,ITER,RSS)
166      IFILOFF=IFILOFF+2
167      TETA = TETA+INCT
168      IF(TETA.LE.TMAX) THEN
169          GO TO 10
170      ENDIF
171      C -----
172      C DEALLOCATE ARRAYS
173      C -----
174      DEALLOCATE(W,AT,ZT,P,R,A,E,JS,DFG,IRING)
175      PRINT *, 'USER CPU TIME INFORMATION (sec)'
176      PRINT *, '-----'
177      PRINT *, 'INITIALIZATION..', USTIME2(1)-USTIME1(1)-UTIMEOH
178      PRINT *, 'CGFFT LOOP.....', USTIME3(1)-USTIME2(1)-UTIMEOH
179      PRINT *, 'TOTAL.....', USTIME3(1)-USTIME1(1)-UTIMEOH
180      PRINT *, 'ITERATIONS.....', ITER
181      NOPER=(4.*NTI*ALOG(FLOAT(NTI))/ALOG(2.))+10.*MTI+2)*ITER
182      PRINT *, 'MEGAFLOPS.....', NOPER/((USTIME3(1)-USTIME1(1)
183      + -UTIMEOH)*1.0E6)
184      WRITE(11,*) 'PROGRAM TERMINATED.'
185      CLOSE(11)
186      WRITE(*,*) 'PROGRAM TERMINATED.'
187      200 FORMAT(29('='),/, 'TETA=', F4.1,
188      + 2X, 'PHI=', F4.1, 2X, 'PSI=', F4.1, /, ' ITER', 7X, ' RSS(R) ', /, 29('-'))
189      300 FORMAT(I4, 1X, G13.4)
190      400 FORMAT(' CONVERGENCE ACHIEVED.')
191      500 FORMAT(' ITMAX EXCEEDED; CONVERGENCE CRITERION NOT MET.')
192      600 FORMAT(1X, 'TETA', 3X, 'SIG0 dB',

```


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```
193          + 2X,'SIG0(PO) dB',2X,'PHI=',F5.1,2X,'PSI=',F5.1,/,I2)
194          STOP
195          END
```

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 Source Listing of /u/barkeshli/slab.d/platevc.f

-----		LOOP SUMMARY FOR ROUTINE SLAB				-----		
LABEL	INDEX	START	END	NEST	COMMENT	CD?	DP?	ITERATIONS
	I	53	59	1	NO ACTION REQUESTED			3
		105	105	1	VECTOR-CONCURRENT			MT
		106	106	1	VECTOR-CONCURRENT			NG
		107	111	1	VECTOR-CONCURRENT			MTI
		112	113	1	VECTOR-CONCURRENT			NTI
		114	114	1	VECTOR-CONCURRENT			NT
		128	128	1	VECTOR-CONCURRENT			MTI
		137	138	1	VECTOR-CONCURRENT			MTI
		144	144	1	VECTOR-CONCURRENT			MTI
		145	145	1	VECTOR-CONCURRENT			MTI
		147	147	1	VECTOR-CONCURRENT			MTI
		148	149	1	VECTOR-CONCURRENT			MTI
		151	151	1	VECTOR-CONCURRENT			MTI
0	VECTOR	0	CONCURRENT	12	VECT-CONCR	0	CON.	OUTER
1	NOT OPTMZD	0	NOT INNER	0	TOO SHORT	0	DEPENDENT	
13	TOTAL	13	EXAMINED	12	OPTIMIZED			

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 Source Listing of /u/barkeshli/slab.d/platevc.f

```

196 C *****
197 C                               OUTPUT
198 C *****
199     SUBROUTINE OUTPUT(IUNT,IFLG,Z,ZT,TETA,SIG,ITER,RSS)
200     COMPLEX Z(MTI),ZT(NTI)
201     COMMON /DIM/MX,MY,MT,MTI,NX,NY,NXH,NYH,NT,NTI,NG,NA,DLT,XK1,Z0
202 CVD$ NOVECTOR
203 CVD$ NOCONCUR
204     DO 10 I = 1,2
205         IUNTI = I+IUNT-1
206         OPEN(UNIT=IUNTI)
207         WRITE(IUNTI,*) TETA,S,ITER,RSS
208         WRITE(IUNTI,*) MX,MY
209         WRITE(IUNTI,*) (J,J=1,MX)
210         WRITE(IUNTI,*) (K,K=1,MY)
211         L1=(I-1)*MT+1
212         L2=I*MT
213         WRITE(IUNTI,20) (CABS(Z(L)),L=L1,L2)
214         CLOSE(IUNTI)
215     10 CONTINUE
216     20 FORMAT(10(1X,G10.4))
217     RETURN
218     END

```

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```
----- LOOP SUMMARY FOR ROUTINE OUTPUT -----
 LABEL  INDEX  START  END  NEST  COMMENT  CD?  DP?  ITERATIONS
   10    I      204   215   1  NO ACTION REQUESTED                2
  0 VECTOR      0 CONCURRENT  0 VECT-CONCR  0 CON. OUTER
  1 NOT OPTMZD  0 NOT INNER   0 TOO SHORT   0 DEPENDENT
  1 TOTAL      1 EXAMINED    0 OPTIMIZED
```

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 Source Listing of /u/barkeshli/slab.d/platevc.f

```

219 C *****
220 C             INCIDENT FIELD
221 C *****
222 C             PARAMETER DESCRIPTION
223 C =====
224 C ISORCE.....SOURCE TYPE:
225 C             1 : E-POL
226 C             2 : H-POL
227 C E0.....AMPLITUDE OF THE PLANE WAVE
228 C TETA,PHI.....ORIENTATION OF THE PLANE WAVE OR
229 C             ORIENTATION OF THE ANTENN     XIS
230 C PSI.....POLARIZATION ANGLE FOR
231 C             PLANE WAVE
232 C =====
233 C SUBROUTINE INCIDENCE (E, E0, TETA, PHI, PSI)
234 C COMPLEX E (MTI), XJ
235 C REAL     XP (MX), YP (MY)
236 C COMMON  /DIM/MX, MY, MT, MTI, NX, NY, NXH, NYH, NT, NTI, NG, NA, DLT, XK1, Z0
237 C DATA   PI/3.141592653589793/, XJ/(0.,1.)/
238 C MXO = MX/2+1
239 C MYO = MY/2+1
240 C TET = TETA*PI/180.
241 C PH  = PHI*PI/180.
242 C PS  = PSI*PI/180.
243 C ECS=E0*(COS(PS)*COS(TET)*COS(PH)-SIN(PS)*SIN(PH))
244 C ESC=E0*(COS(PS)*COS(TET)*SIN(PH)+SIN(PS)*COS(PH))
245 C XI  = XK1*SIN(TET)*COS(PH)
246 C YI  = XK1*SIN(TET)*SIN(PH)
247 CVD$R NODEPCHK
248 C DO 10 K=1,MY
249 C     YP(K)=(K-MYO)*DLT
250 C DO 20 J=1,MX
251 C     XP(J)=(J-MXO)*DLT
252 C     L1 = (K-1)*MX+J
253 C     E(L1)=ECS*CEXP(XJ*(XI*XP(J)+YI*YP(K)))
254 C     E(L2)=ESC*CEXP(XJ*(XI*XP(J)+YI*YP(K)))
255 C 20 CONTINUE
256 C 10 CONTINUE
257 C RETURN
258 C END
    
```

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 Source Listing of /u/barkeshli/slab.d/platevc.f

----- LOOP SUMMARY FOR ROUTINE INCIDENC -----									
LABEL	INDEX	START	END	NEST	COMMENT	CD?	DP?	ITERATIONS	
10	K	248	256	1	CONCURRENT OUTER			MY	
20	J	250	255	2	VECTORIZED			MX	
1	VECTOR	0	CONCURRENT	0	VECT-CONCR	1	CON.	OUTER	
0	NOT OPTMZD	0	NOT INNER	0	TOO SHORT	0	DEPENDENT		
2	TOTAL	2	EXAMINED	2	OPTIMIZED				

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Source Listing of /u/barkeshli/slab.d/platevc.f

```

259 C *****
260 C          TRANSFORM OF THE EFFECTIVE KERNEL
261 C *****
262       SUBROUTINE DIFREQ (DFG, IRING, IRMAX)
263       INTEGER IRING (MT)
264       REAL*8 X1 (NT), X2, G, KX (NX), KY (NY), FRX, FRY, C11, C121, C22, TPI, XK2
265       REAL DFG (NG)
266       COMMON /DIM/MX, MY, MT, MTI, NX, NY, NXH, NYH, NT, NTI, NG, NA, DLT, XK1, Z0
267       COMMON IBASIS
268       DATA PI/3.141592653589793/, TPI/6.28318530717959D0/
269       XK2=DBLE (XK1*XK1)
270       FRX=1.D0/DBLE (DLT*NX)
271       FRY=1.D0/DBLE (DLT*NY)
272 C =====
273 C COMPUTE THE SPECTRAL VARIABLES
274 C =====
275       DO 10 J=1, NXH
276           KX (J)=TPI*(J-1)*FRX
277 10      CONTINUE
278       DO 20 J=NXH+1, NX
279           KX (J)=TPI*(J-NX-1)*FRX
280 20      CONTINUE
281       DO 30 J=1, NYH
282           KY (J)=TPI*(J-1)*FRY
283 30      CONTINUE
284       DO 40 J=NYH+1, NY
285           KY (J)=TPI*(J-NY-1)*FRY
286 40      CONTINUE
287 C =====
288 C TRANSFORM OF THE GREEN'S FUNCTION
289 C =====
290 CVD$ NODEPCHK
291       DO 50 K=1, NY
292           C22=XK2-KY (K)**2
293           L=(K-1)*NX
294 CVD$ NODEPCHK
295       DO 60 J=1, NX
296           L1=L+J
297           L2=NT+L1
298           L3=NT+L2
299           L4=NT+L3
300           C11=XK2-KX (J)**2
301           C121=-KX (J)*KY (K)
302           X1 (L1) = KX (J)**2+KY (K)**2-XK2
303           X2=DABS (X1 (L1))
304           G = 1.D0/(2.*DSQRT (X2))
305           DFG (L1)=SNGL (C11*G)
306           DFG (L2)=SNGL (C121*G)

```

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Source Listing of /u/barkeshli/slab.d/platevc.f

```
307             DFG(L3)=SNGL(C121*G)
308             DFG(L4)=SNGL(C22*G)
309     60         CONTINUE
310     50         CONTINUE
311     C         =====
312     C         IDENTIFY THE GREEN'S FUNCTION RING
313     C         =====
314             INDX=0
315     CVD$ NOVECTOR
316     CVD$ NOCONCUR
317             DO 70 L1=1,NT
318                 IF (X1(L1).GT.0.D0) THEN
319                     ELSE
320                         INDX=INDX+1
321                         IRING(INDX)=L1
322                     END IF
323     70         CONTINUE
324             IRMAX=2*INDX
325     CVD$ NODEPCHK
326             IRING(INDX+1:IRMAX)=IRING(1:INDX)+NT
327             RETURN
328             END
```


Alliant FX/Fortran Compiler V4.0.28 (2.24D20) 5 Dec 1988 02:10:39 Page 13
 Source Listing of /u/barkeshli/slab.d/platevc.f

----- LOOP SUMMARY FOR ROUTINE DIFREQ -----								
LABEL	INDEX	START	END	NEST	COMMENT	CD?	DP?	ITERATIONS
10	J	275	277	1	VECTOR-CONCURRENT			NXH
20	J	278	280	1	VECTOR-CONCURRENT			NX-NXH
30	J	281	283	1	VECTOR-CONCURRENT			NYH
40	J	284	286	1	VECTOR-CONCURRENT			NY-NYH
50	K	291	310	1	CONCURRENT OUTER			NY
60	J	295	309	2	VECTORIZED			NX
70	L1	317	323	1	NO ACTION REQUESTED			NT
		326	326	1	VECTOR-CONCURRENT			INDX
1	VECTOR	0	CONCURRENT	5	VECT-CONCR	1		CON. OUTER
1	NOT OPTMZD	0	NOT INNER	0	TOO SHORT	0		DEPENDENT
8	TOTAL	8	EXAMINED	7	OPTIMIZED			

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 Source Listing of /u/barkeshli/slab.d/platevc.f

```

329 C *****
330 C             CONVOLUTION & DIFFERENTIATION
331 C *****
332 C     SUBROUTINE OPERAT (A, ISIGN, Z, ZT, DFG, IRING, IRMAX, AT, W)
333 C     REAL    DFG (NG)
334 C     INTEGER IRING (MT)
335 C     COMPLEX XJ, A (MTI), Z (MTI), ZT (NTI), AT (NTI), W (NT)
336 C     COMPLEX RS, CRS, FCT
337 C     COMMON  /DIM/MX, MY, MT, MTI, NX, NY, NXH, NYH, NT, NTI, NG, NA, DLT, XK1, Z0
338 C     COMMON  /PAR/RS, CRS, FCT
339 C     COMMON  IBASIS
340 C     DATA   XJ/(0.,1.)/
341 C     =====
342 C     FORWARD FOURIER TRANSFORM
343 C     =====
344 C     CALL SPECTRA (Z, ZT, -1, W)
345 C     =====
346 C     CARRY OUT THE CONVOLUTIONS IN THE TRANSFORM DOMAIN
347 C     =====
348 C     DO 10 I=1, NT
349 C         AT (I) = DFG (I) * ZT (I) + DFG (I+NT) * ZT (I+NT)
350 C     CONTINUE
351 C     DO 20 I=NT+1, NTI
352 C         AT (I) = DFG (I+NT) * ZT (I-NT) + DFG (I+NTI) * ZT (I)
353 C     CONTINUE
354 C     =====
355 C     CHOOSE THE RIGHT BRANCH OF THE GREEN'S FUNCTION
356 C     =====
357 C     CVD$ PERMUTATION (IRING)
358 C         AT (IRING (1:IRMAX)) = ISIGN * XJ * AT (IRING (1:IRMAX))
359 C     =====
360 C     INVERSE FOURIER TRANSFORM
361 C     =====
362 C     CALL SPECTRA (A, AT, 1, W)
363 C     IF (ISIGN.EQ.1) THEN
364 C         A (1:MTI) = CRS * Z (1:MTI) - FCT * A (1:MTI)
365 C     ELSE
366 C         A (1:MTI) = RS * Z (1:MTI) + FCT * A (1:MTI)
367 C     ENDIF
368 C     RETURN
369 C     END
    
```

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 Source Listing of /u/barkeshli/slab.d/platevc.f

----- LOOP SUMMARY FOR ROUTINE OPERAT -----								
LABEL	INDEX	START	END	NEST	COMMENT	CD?	DP?	ITERATIONS
10	I	348	350	1	VECTOR-CONCURRENT			NT
20	I	351	353	1	VECTOR-CONCURRENT			NTI-NT
		358	358	1	VECTOR-CONCURRENT			IRMAX
		364	364	1	VECTOR-CONCURRENT			MTI
		366	366	1	VECTOR-CONCURRENT			MTI
0	VECTOR	0	CONCURRENT	5	VECT-CONCR	0		CON. OUTER
0	NOT OPTMZD	0	NOT INNER	0	TOO SHORT	0		DEPENDENT
5	TOTAL	5	EXAMINED	5	OPTIMIZED			

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 Source Listing of /u/barkeshli/slab.d/platevc.f

```

370 C *****
371 C          FOURIER TRANSFORM VIA FFT
372 C *****
373          SUBROUTINE SPECTRA(Z,ZT,ISIGN,W)
374          COMPLEX Z(MTI),ZT(NTI),W(NT)
375          COMPLEX AUX(NA)
376          COMMON /DIM/MX,MY,MT,MTI,NX,NY,NXH,NYH,NT,NTI,NG,NA,DLT,XK1,Z0
377 C          =====
378 C          INITIALIZE FFT PARAMETERS
379 C          =====
380          CALL CFFT2D(NX,NY,W,NX,0,AUX)
381          IF (ISIGN.EQ.-1) THEN
382 CVD$ NOCONCUR
383          DO 10 I = 1,2
384 C          =====
385 C          PRECONDITIONING IN SPATIAL DOMAIN
386 C          =====
387          W(1:NT)=(0.,0.)
388          L1 = (I-1)*MT
389 CVD$ NODEPCHK
390          DO 20 K = 1,MY
391          L2 = L1+(K-1)*MX
392          L3 = (K-1)*NX
393 CVD$ NODEPCHK
394          DO 30 J = 1,MX
395          W(L3+J) = Z(L2+J)
396          30 CONTINUE
397          20 CONTINUE
398 C          =====
399 C          FOREARD FFT
400 C          =====
401          CALL CFFT2D(NX,NY,W,NX,-1,AUX)
Line 401 Informational message # 1342
          Concurrent and vector loop optimization inhibited by
          subroutine call without Cncall directive -- CFFT2D
402          L1 = (I-1)*NT
403          DO 40 J = 1,NT
404          ZT(L1+J) = W(J)
405          40 CONTINUE
406          10 CONTINUE
407          ELSE IF (ISIGN.EQ.1) THEN
408 CVD$ NOCONCUR
409          DO 50 I = 1,2
410          L1 = (I-1)*NT
411          DO 60 J = 1,NT
412          W(J) = ZT(L1+J)
413          60 CONTINUE
414 C          =====

```

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Source Listing of /u/barkeshli/slab.d/platevc.f

```
415 C    INVERSE FFT
416 C    =====
417          CALL CFFT2D(NX,NY,W,NX,1,AUX)
Line 417 Informational message # 1342
          Concurrent and vector loop optimization inhibited by
          subroutine call without Cncall directive -- CFFT2D
418          L1 = (I-1)*MT
419 CVD$ NODEPCHK
420          DO 70 K = 1,MY
421             L2 = (K-1)*NX
422             L3 = L1+(K-1)*MX
423 CVD$ NODEPCHK
424          DO 80 J = 1,MX
425             Z(L3+J) = W(L2+J)
426 80          CONTINUE
427 70          CONTINUE
428 50          CONTINUE
429          END IF
430          RETURN
431          END
```

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 Source Listing of /u/barkeshli/slab.d/platevc.f

----- LOOP SUMMARY FOR ROUTINE SPECTRA -----							
LABEL	INDEX	START	END	NEST	COMMENT	CD?	DP? ITERATIONS
10	I	383	406	1	NOT OPTIMIZED		2
		387	387	2	VECTOR-CONCURRENT		NT
20	K	390	397	2	CONCURRENT OUTER		MY
30	J	394	396	3	VECTORIZED		MX
40	J	403	405	2	VECTOR-CONCURRENT		NT
50	I	409	428	1	NOT OPTIMIZED		2
60	J	411	413	2	VECTOR-CONCURRENT		NT
70	K	420	427	2	CONCURRENT OUTER		MY
80	J	424	426	3	VECTORIZED		MX
2	VECTOR	0	CONCURRENT	3	VECT-CONCR	2	CON. OUTER
0	NOT OPTMZD	2	NOT INNER	0	TOO SHORT	0	DEPENDENT
9	TOTAL	7	EXAMINED	7	OPTIMIZED		

CORNELLF

USERID: M90Y ORIGIN: CORNELLE CREATED: 07/21/88 16:17:17
FILENAME: SLABVTM LISTING CLASS: A FORMAT: J
SPOOLID: 2110 RECS: 711 COPY: 1 DUPLICATE: 1

PRINTED AT: CORNELLD ID: OLN AT: 07/21/88 16:17:39

*
* THIS FILE WAS SENT BY THE COMMAND:
* PRT3812 SLABVTM LISTING A (OLIN
*

PROGRAM HAS TERMINATED; RC (0)
 ORTIAD
 PROGRAM SAMPLING INTERVAL WAS 4 MS; TOTAL NUMBER OF SAMPLES WAS 5279.
 IRECT SAMPLES:

PROGRAM UNIT	SAMPLES	%TOTAL
LAB	104	1.97
UTPUT	2	0.04
NCIDENCE	2	0.04
NORM	53	1.00
RING	2	0.04
IFREQ	9	0.17
PERAT	683	12.94 ***
PECTRA	452	8.56 **
#SQRT	5	0.09
FSCOC#	1	0.02
C#EXP	0	0.00
C#ABS	296	5.61 *
4I\$	1657	31.39 *****
CFTA\$	40	0.76
CFT\$	12	0.23
CFT2	3	0.06
4F\$	934	17.69 *****
2L\$	903	17.11 ***
IND\$	8	0.15
INDL\$	27	0.51
WF\$	23	0.44
W2L\$	8	0.15
WT\$	6	0.11
#ABSARY	43	0.06
UNKNOWN	1	0.02

81

PROGRAM SAMPLING INTERVAL WAS 4 MS; TOTAL NUMBER OF SAMPLES WAS 5279.

TATEMENT	SAMPLES	%UNIT	%TOTAL
UTPUT.ENTRY/EXIT	0	0.00	0.00
UTPUT.6	0	0.00	0.00
UTPUT.7	0	0.00	0.00
UTPUT.8	0	0.00	0.00
UTPUT.9	0	0.00	0.00
UTPUT.10	0	0.00	0.00
UTPUT.11	0	0.00	0.00
UTPUT.12	0	0.00	0.00
UTPUT.13	0	0.00	0.00
UTPUT.14	0	0.00	0.00
UTPUT.15	2	100.00	0.04
UTPUT.16	0	0.00	0.00
UTPUT.17/10	0	0.00	0.00
UTPUT.18	0	0.00	0.00
UTPUT.19	0	0.00	0.00

ORTIAD
 PROGRAM SAMPLING INTERVAL WAS 4 MS; TOTAL NUMBER OF SAMPLES WAS 5279.
 IRECT SAMPLES:

TATEMENT	SAMPLES	%UNIT	%TOTAL
NCIDENCE.ENTRY/EXIT	0	0.00	0.00
NCIDENCE.6	0	0.00	0.00
NCIDENCE.7	0	0.00	0.00
NCIDENCE.8	0	0.00	0.00
NCIDENCE.9	0	0.00	0.00
NCIDENCE.10	0	0.00	0.00
NCIDENCE.11	0	0.00	0.00

NCIDENCE.32 0 0.00 0.00
 NCIDENCE.33 0 0.00 0.00
 ORTIAD
 INVALID OPERAND SYNTAX SPECIFIED IN "LISTSAMP DIFREQ"
 ORTIAD
 ROGRAM SAMPLING INTERVAL WAS 4 MS; TOTAL NUMBER OF SAMPLES WAS 5279.
 IRECT SAMPLES:

TATEMENT	SAMPLES	%UNIT	%TOTAL
IFREQ.ENTRY/EXIT	0	0.00	0.00
IFREQ.7	0	0.00	0.00
IFREQ.8	0	0.00	0.00
IFREQ.9	0	0.00	0.00
IFREQ.10	0	0.00	0.00
IFREQ.11	0	0.00	0.00
IFREQ.12	0	0.00	0.00
IFREQ.13	0	0.00	0.00
IFREQ.14	0	0.00	0.00
IFREQ.15	0	0.00	0.00
IFREQ.16	0	0.00	0.00
IFREQ.17	0	0.00	0.00
IFREQ.18	0	0.00	0.00
IFREQ.19	0	0.00	0.00
IFREQ.20	1	11.11	0.02
IFREQ.21	0	0.00	0.00
IFREQ.22	0	0.00	0.00
IFREQ.23	0	0.00	0.00
IFREQ.24	0	0.00	0.00
IFREQ.25	0	0.00	0.00
IFREQ.26	0	0.00	0.00
IFREQ.27	0	0.00	0.00
IFREQ.28	0	0.00	0.00
IFREQ.29	1	11.11	0.02
IFREQ.30	0	0.00	0.00
IFREQ.31	0	0.00	0.00
IFREQ.32	1	11.11	0.02
IFREQ.33	0	0.00	0.00
IFREQ.34	1	11.11	0.02
IFREQ.35	0	0.00	0.00
IFREQ.36	0	0.00	0.00
IFREQ.37	0	0.00	0.00
IFREQ.38	0	0.00	0.00
IFREQ.39	4	44.44	0.08
IFREQ.40	1	11.11	0.02
IFREQ.41	0	0.00	0.00
IFREQ.42	0	0.00	0.00
IFREQ.43	0	0.00	0.00
IFREQ.44/20	0	0.00	0.00
IFREQ.45/10	0	0.00	0.00
IFREQ.46	0	0.00	0.00
IFREQ.47	0	0.00	0.00

ORTIAD
 ROGRAM SAMPLING INTERVAL WAS 4 MS; TOTAL NUMBER OF SAMPLES WAS 5279.
 IRECT SAMPLES:

TATEMENT	SAMPLES	%UNIT	%TOTAL
PERAT.ENTRY/EXIT	0	0.00	0.00
PERAT.9	0	0.00	0.00
PERAT.10	0	0.00	0.00
PERAT.11	0	0.00	0.00

PERAT.35 0 0.00 0.00
 PERAT.36 30 4.39 0.57
 PERAT.37 0 0.00 0.00
 PERAT.38 0 0.00 0.00
 PERAT.39 0 0.00 0.00
 PERAT.40 29 4.25 0.55
 PERAT.41 3 0.44 0.06
 PERAT.42 0 0.00 0.00
 PERAT.43 25 3.66 0.47
 PERAT.44 2 0.29 0.04
 PERAT.45 0 0.00 0.00
 PERAT.46 0 0.00 0.00
 PERAT.47 0 0.00 0.00

ORTIAD
 PROGRAM SAMPLING INTERVAL WAS 4 MS; TOTAL NUMBER OF SAMPLES WAS 5279.

IRECT SAMPLES:

TATEMENT	SAMPLES	%UNIT	%TOTAL
PECTRA.ENTRY/EXIT	0	0.00	0.00
PECTRA.5	0	0.00	0.00
PECTRA.6	0	0.00	0.00
PECTRA.7	0	0.00	0.00
PECTRA.8	53	11.73	1.00
PECTRA.9	30	6.64	0.57
PECTRA.10	0	0.00	0.00
PECTRA.11	0	0.00	0.00
PECTRA.12	0	0.00	0.00
PECTRA.13	0	0.00	0.00
PECTRA.14	2	0.44	0.04
PECTRA.15	32	7.08	0.61
PECTRA.16	8	1.77	0.15
PECTRA.17/20	2	0.44	0.04
PECTRA.18	0	0.00	0.00
PECTRA.19	0	0.00	0.00
PECTRA.20	0	0.00	0.00
PECTRA.21	0	0.00	0.00
PECTRA.22	117	25.88	2.22
PECTRA.23	27	5.97	0.51
PECTRA.24/10	0	0.00	0.00
PECTRA.25	0	0.00	0.00
PECTRA.26	0	0.00	0.00
PECTRA.27	0	0.00	0.00
PECTRA.28	0	0.00	0.00
PECTRA.29	108	23.89	2.05
PECTRA.30	18	3.98	0.34
PECTRA.31	0	0.00	0.00
PECTRA.32	0	0.00	0.00
PECTRA.33	0	0.00	0.00
PECTRA.34	0	0.00	0.00
PECTRA.35	0	0.00	0.00
PECTRA.36	0	0.00	0.00
PECTRA.37	2	0.44	0.04
PECTRA.38	37	8.19	0.70
PECTRA.39	13	2.88	0.25
PECTRA.40/70	3	0.66	0.06
PECTRA.41/50	0	0.00	0.00
PECTRA.42	0	0.00	0.00
PECTRA.43	0	0.00	0.00
PECTRA.44	0	0.00	0.00

ORTIAD
 PROGRAM SAMPLING INTERVAL WAS 4 MS; TOTAL NUMBER OF SAMPLES WAS 5279.

```

=====
C
PARAMETER ( IDIM=2, DLX=2, DLY=2., FMHZ=300., UPL=31.5, IPAD=1,
+          ITMAX=200, ITMIN=5, ICMAX=100, TOL=.04,
+          MX=63, MY=63, NX=128, NY=128,
+          MT=63*63, NT=128*128,
+          MTI=63*63*2, NTI=128*128*2, NTFG=128*128*(2+2) )
          NN(2), IRING(MT)
          INTEGER
          REAL
          COMPLEX
          +          JS(MTI), E(MTI), A(MTI), R(MTI), P(MTI),
          ZT(MTI), AT(MTI), W(MTI), RS(3), XJ, ER, JSG
          /DIM/ID, MXA, MYA, MTA, MTIA, NXA, NYA, NTA, NTIA,
          NTFG, DLTA, XK1A, ZOA
          COMMON
          +          COMMON
          IBASIS
=====

```

```

0002
0002
0002
0002
0002
0002
0003
0004
0005
0005
0006
0006
0007

```

REPORT(XLIST) VECTORIZATION ANALYSIS

LSN FLAG NESTING *.....*.....1.....*.....2.....*.....3.....*.....4.....*.....5.....*.....6.....*.....7.....*.....8 MESSAGES

DATA PI/3.141592653589793/,XJ/(0.,1.)/

```

0008 Z0=120.*PI
0009 XL1=300./FMHZ
0010 DLT=XL1/UPL
0011 XK1=2.*PI/XL1
0012 ID=IDIM
0013 MXA=MX
0014 MYA=MY
0015 MTA=MT
0016 MTIA=MTI
0017 NXA=NX
0018 NYA=NY
0019 NTA=NT
0020 NTIA=NTI
0021 NTF=NTFG
0022 DLTA=DLT
0023 XK1A=XK1
0024 Z0A=Z0
0025
0026 OPEN(UNIT=10)
0027 READ(10,*)ISORCE
0028 READ(10,*)TETA,PHI,PSI,TMAX,INCT
0029 READ(10,*)ANTL,ANTPX,ANTPY,ANTPZ
0030 READ(10,*)E0,C10
0031 READ(10,*)ER,TK
0032 READ(10,*)IBASIS
0033 CLOSE(10)

```

C READ IN THE INPUT INFORMATION

```

C
C | INITIALIZATION |
C |
C |

```

```

0034 RS(1) = -XJ*Z0/((ER-1)*XK1*TK)
0035 RS(2) = RS(1)
0036 RS(3) = -XJ*Z0*ER/((ER-1)*XK1*TK)
0037 CALL DIFREQ(DFG)
0038 CALL GRING(IRING,IRMAX)
0039 VECT +-----
0040 |
0041 JS(1) = (0.,0.)
0042 NP = (TMAX-TETA)/INCT+1
0043 IFILOFF=12
0044 OPEN(UNIT=11)
0045 CONTINUE
10 CONVERGENCE PARAMETERS
C ITER = 0
ICONT = 0
RSS = 0.
WRITE(11,200)TETA,PHI,PSI
CALL INCIDENCE(E,E0,TETA,PHI,PSI)
ENORM = VNORM(E)
C ZEROTH ORDER RESIDUAL CALCULATION
CALL OPERAT(A,-1,JS,RS,ZT,DFG,IRING,IRMAX,AT,W)
DO 20 I = 1,MTI
R(I) = A(I)-E(I)
P(I) = (0.,0.)

```

REPORT(XLIST) VECTORIZATION ANALYSIS

ISN FLAG NESTING *.....1.....2.....3.....4.....5.....6.....7.*.....8 MESSAGES

```

C
C
C | | MAIN ITERATION LOOP | |
C | |-----| |
30 CONTINUE
0057 CALL OPERAT(A,1,R,RS,ZT,DFG,IRING,IRMAX,AT,W)
0058 B = 1./VNORM(A)
0059 DO 35 I = 1,MTI
0060 VECT +-----|
0061 | P(I) = P(I)-B*A(I)
0063 | CALL OPERAT(A,-1,P,RS,ZT,DFG,IRING,IRMAX,AT,W)
0064 | T = 1./VNORM(A)
0065 VECT +-----|
0066 | R(I) = R(I)+T*A(I)
0067 | JS(I) = JS(I)+T*P(I)
0069 ITER = ITER+1
0070 ICONT = ICONT+1

C
C | | CONVERGENCE CHECK | |
C | |-----| |
C COMPUTE RELATIVE RESIDUAL ERROR
0071 RSS = SQRT(VNORM(R)/ENORM)
0072 WRITE(*,300)ITER,RSS
0073 WRITE(11,300)ITER,RSS
0074 IF((RSS.LE.TOL).AND.(ITER.GT.ITMIN)) THEN
0075 WRITE(11,400)
0076 ELSE IF(ITER.GE.ITMAX) THEN
0077 WRITE(11,500)
0078 ELSE IF(ICONT.EQ.ICMAX) THEN
0079 GO TO 90
0080 ELSE
0081 GO TO 30
0082 END IF
0083 CONTINUE
90 CCCCC CALL OUTPUT(12,1,JS,ZT,TETA,SIG,ITER,RSS)
C IF INTERMEDIATE RESULTS, RESET THE COUNTER.
IF(ICONT.EQ.ICMAX) THEN
ICONT = 0
GO TO 30
END IF
C IF (ISORCE.EQ.1) THEN
C IF DONE WITH THIS ANGLE OF INCIDENCE, STORE THE CONVERGED CURRENT AND
C COMPUTE EXCITED CURRENT AND SIGMAO FOR NEXT ANGLE OF
C INCIDENCE WITH THE LAST JS AS THE NEW SOLUTION GUESS.
CALL OUTPUT(IFILOFF,1,JS,ZT,TETA,SIG,ITER,RSS)
IFILOFF=IFILOFF+2
TETA = TETA+INCT
IF(TETA.LE.TMAX) THEN
GO TO 10
ENDIF
ENDIF
C WRITE(11,*)'PROGRAM TERMINATED.'

```

REPORT(XLIST) VECTORIZATION ANALYSIS

ISN FLAG NESTING *.....1.....2.....3.....4.....5.....6.....7.*.....8 MESSAGES

```

C DO I = 1,11
C CLOSE(I)
C
200 + FORMAT(29('='),/, 'TETA=',F4.1,
0097 2X, 'PHI=',F4.1,2X, 'PSI=',F4.1,/, 'ITER',7X, 'RSS(R)',/,29('-'))
0098 + FORMAT(14,1X,G13.4)
0099 + FORMAT(' CONVERGENCE ACHIEVED.')
0100 + FORMAT(' ITMAX EXCEEDED; CONVERGENCE CRITERION NOT MET.')
0101 + FORMAT(1X, 'TETA',3X, 'SIGO dB',
0102 2X, 'SIGO(PO) dB',2X, 'PHI=',F5.1,2X, 'PSI=',F5.1,/,12)
0103 STOP
END

```

```

0001 SUBROUTINE OUTPUT(IUNT,IFLG,Z,ZT,TETA,SIG,ITER,RSS)
0002 COMPLEX Z(*),ZT(*)
0003 CHARACTER*11 FMT1
0004 CHARACTER*16 FMTR
0005 COMMON /DIM/ID,MX,MY,MT,MTI,NX,NY,NT,NTI,NTF,DLT,XK1,ZO
0006 DO 10 I = 1,ID

```

I/O OPERATION UNANALYZABLE INTRINSI

```

0007 IUNT1 = I+IUNT-1
0008 OPEN(UNIT=IUNT1)
0009 WRITE(IUNT1,*)TETA,S,ITER,RSS
0010 WRITE(IUNT1,*)MX,MY
0011 WRITE(IUNT1,*)(J,J=1,MX)
0012 WRITE(IUNT1,*)(K,K=1,MY)
0013 L1=(I-1)*MT+1
0014 L2=I*MT
0015 WRITE(IUNT1,*)(CABS(Z(L)),L=L1,L2)
0016 CLOSE(IUNT1)
0017 CONTINUE
0018 RETURN

```

ISN FLAG NESTING *.....1.....2.....3.....4.....5.....6.....7.*.....8 MESSAGES

```

0001 SUBROUTINE INCIDENCE(E,E0,TETA,PHI,PSI)
0002 COMMON /DIM/ID,MX,MY,MT,MTI,NX,NY,NT,NTI,NTF,DLT,XK1,ZO
0003 COMPLEX E(*),XJ,G,RJK(7938)
0004 REAL XP(63),YP(63)
0005 DATA PI/3.141592653589793/,XJ/(0.,1.)/
0006 MXO = MX/2+1
0007 MYO = MY/2+1

```

```

C INCIDENT PLANE WAVE
C
0008 TET = TETA*PI/180.
0009 PH = PHI*PI/180.
0010 PS = PSI*PI/180.
0011 ECS=E0*(COS(PS)*COS(TET)*COS(PH)-SIN(PS)*SIN(PH))
0012 ESC=E0*(COS(PS)*COS(TET)*SIN(PH)+SIN(PS)*COS(PH))
0013 XI = XK1*SIN(TET)*COS(PH)
0014 YI = XK1*SIN(TET)*SIN(PH)
0015 DO 10 K=1,MY

```

VECT +-----

```

0003 COMMON /DIM/ID,MX,MY,MT,MTI,NX,NY,NT,NTI,NTF,DLT,XK1,ZO
0004 COMPLEX W(*)
0005 SUM = 0.DO
0006 VECT +-----
0007 |
0008 |
0009 |
0010 |
0011 |

```

VECTOR INTRINSIC FUNC
VECTOR SUM REDUCTION

```

0001 SUBROUTINE GRING(IRING,IRMAX)
0002 INTEGER IRING(*)
0003 REAL*8 X,KX,KY,FRX,FRY,PI,TPI,XK2
0004 COMMON /DIM/ID,MX,MY,MT,MTI,NX,NY,NT,NTI,NTF,DLT,XK1,ZO
0005 DATA PI/3.141592653589793/
0006 TPI=2.DO*PI
0007 XK2=DBLE(XK1*XK1)
0008 NXH=NX/2+1
0009 NYH=NY/2+1
0010 FRX=1.DO/DBLE(DLT*NX)
0011 FRY=1.DO/DBLE(DLT*NY)
0012 INDX=0

```

UNSP +-----

```

0013 DO 10 K = 1,NY
0014 IF(K.LE.NYH)KY = TPI*(K-1)*FRY
0016 IF(K.GT.NYH)KY = TPI*(K-NY-1)*FRY

```

DEPENDENT ON UNSUPPOR
CONDITIONAL SCALAR CO
DEPENDENT ON UNSUPPOR
CONDITIONAL SCALAR CO

RECR +-----

```

0019 DO 20 J = 1,NX
0020 IF(J.LE.NXH)KX = TPI*(J-1)*FRX
0022 IF(J.GT.NXH)KX = TPI*(J-NX-1)*FRX

```

```

X = KX**2+KY**2-XK2
IF(X.LT.0.DO)THEN
  INDX=INDX+1

```

SCALAR DEFINED BEFORE
CONDITIONAL SCALAR CO
SCALAR DEFINED BEFORE
CONDITIONAL SCALAR CO
DEPENDENT ON UNSUPPOR
DEPENDENT ON UNSUPPOR
SCALAR DEFINED BEFORE
CONDITIONAL SCALAR CO
SUBSCRIPT TOO COMPLEX
OPTIMIZER INDUCED DEP
CONDITIONAL NON-INDUC

```
IRING(INDX)=L1
```

```

0032 IRMAX=INDX
0033 RETURN
0034 END

```


REPORT(XLIST) VECIORIZATION ANALYSIS

ISN FLAG NESTING *.....1.....2.....3.....4.....5.....6.....7.*.....8 MESSAGES

```

0033 | _____ |
      | | |
0035 VECT +-----+
0036 | _____ |
0038 | _____ |
0039 VECT +-----+
0040 | _____ |
      | | |
0042 VECT +-----+
0043 | _____ |
0045 | _____ |
0046 | _____ |
0047 | _____ |

```

```

      A(I)=CRS(1)*Z(I)-FCT*A(I)
      DO 60 I=MT+1,MTI
      A(I)=CRS(2)*Z(I)-FCT*A(I)
      ELSE
      DO 70 I=1,MT
      A(I)=RS(1)*Z(I)+FCT*A(I)
      DO 80 I=MT+1,MTI
      A(I)=RS(2)*Z(I)+FCT*A(I)
      ENDIF
      RETURN
      END

```

NUMBER ISN FLAG VS FORTRAN VECTOR REPORT MESSAGES

ILX01291 0023 RECR THE LOOP VARIABLE OF THIS LOOP OR OF SOME NESTED LOOP AFFECTS THE LOOP VARIABLE OR AN AUXILIA INDUCTION VARIABLE USED BY SOME OTHER NESTED LOOP.

```

0001 SUBROUTINE SPECTRA(Z,ZT,ISIGN,W)
0002 /DIM/ID,MX,MY,MT,NTI,NX,NY,NT,NTF,DLT,XK1,ZO
0003 REAL*8 AUX1(20000),AUX2(20000)
0004 COMPLEX Z(*),ZT(*),W(*)
0005 IF(ISIGN.EQ.-1)THEN
0006 UNAN DO 10 I = 1, ID
          C PRECONDITIONING IN SPATIAL DOMAIN
          DO 15 II=1,NT
            W(II)=(0.,0.)
            L1 = (I-1)*MT
            DO 20 K = 1,MY
              DO 30 J = 1,MX
                W(L3+J) = Z(L2+J)
          10 CONTINUE
          CALL ZPADN(W)
          CALL SCFT2(1,W,1,NX,W,1,NX,NX,NY,1,1.0,AUX1,20000,AUX2,20000)
          CALL SCFT2(0,W,1,NX,W,1,NX,NX,NY,1,1.0,AUX1,20000,AUX2,20000)
          L1 = (I-1)*NT
          DO 40 J = 1,NT
            ZT(L1+J) = W(J)
          CONTINUE
          ELSE IF(ISIGN.EQ.1)THEN
            DO 50 I = 1, ID
              L1 = (I-1)*NT
              DO 60 J = 1,NT
                W(J) = ZT(L1+J)
            CALL SCFT2(1,W,1,NX,W,1,NX,NX,NY,-1,1.0,AUX1,20000,AUX2,20000)
            CALL SCFT2(0,W,1,NX,W,1,NX,NX,NY,-1,1.0,AUX1,20000,AUX2,20000)
            L1 = (I-1)*MT
            DO 70 K = 1,MY
              DO 80 J = 1,MX
                Z(L3+J) = W(L2+J)
          10 CONTINUE
          END IF
          RETURN
          END
0041
0042
0043
0044

```

USER FUNCTION OR SUBR

STRIDE UNKNOWN

USER FUNCTION OR SUBR

STRIDE UNKNOWN

NUMBER ISN FLAG VS FORTRAN VECTOR REPORT MESSAGES
 AFF5501 PROGRAM SAMPLING INTERVAL WAS 4 MS; TOTAL NUMBER OF SAMPLES WAS 5279.
 AFF5521 SUM OF DIRECT AND CALLED SAMPLES:
 AFF5641 PROGRAM UNIT
 AFF5651 SLAB
 AFF5651 OUTPUT
 AFF5651 INCIDENCE
 AFF5651 VNORM
 AFF5651 GRING
 AFF5651 DIFREQ
 AFF5651 OPERAT
 AFF5651 SPECTRA
 AFF5651 D#SQRT
 AFF5651 VFSCOC#
 AFF5651 VC#EXP
 AFF5651 VC#ABS
 AFF5651 F41\$ 87.22009.31.41
 AFF5651 SCFTAS
 AFF5651 SCFT\$ 87.21917.36.06
 AFF5651 SCFT2
 AFF5651 F4F\$
 AFF5651 F2L\$
 AFF5651 SINDS\$ 87.22009.33.35
 AFF5651 SINDL\$ 87.21918.20.27
 AFF5651 FWF\$ 87.21918.14.00
 AFF5651 FW2L\$
 AFF5651 FMT\$
 AFF5651 C#ABS
 AFF5651 *LIBRARY
 AFF5651 *UNKNOWN

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23	0.04	
27	0.27	*****+*****+*****+
32	90.09	*****+*****+*****+
38	77.15	*****+*****+*****+
	0.09	
	0.04	
	0.04	
	5.61	***
	31.39	*****+*****+*****+
	68.31	*****+*****+*****+
	68.54	*****+*****+*****+
	68.59	*****+*****+*****+
	17.69	*****+*****+*****+
	17.11	*****+*****+*****+
	0.15	
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