

USER'S GUIDE TO MATCOV

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Technical Report 80-5

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September 1980

Research effort is partially supported by the Department of Industrial and Operations Engineering of The University of Michigan and by the Air Force Office of Scientific Research, Air Force Systems Command, USAF, under grant no. AFOSR 78-3646. The United States Government is authorized to reproduce and distribute reprints for governmental purposes notwithstanding any copyright notation hereon. This paper was typed by Mrs. Geraldine Cox.



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## 1. INTRODUCTION

This report discusses a FORTRAN implementation of the algorithm presented in [1] for the minimum cost 1-matching/covering problem.

Given a network  $G = (N, A, c)$  where  $N$  is a nonempty set of  $n$  nodes,  $A$  is a set of  $m$  edges (an edge is an unordered pair of distinct nodes), and  $c = [c_{i;j}]$  is a vector of edge costs, the minimum cost 1-matching/covering problem is formulated as follows:

Find an  $m$ -vector  $x = [x_{i;j}]$

To minimize  $\sum_{i;j} c_{i;j} x_{i;j}$

Subject to 
$$\left\{ \begin{array}{ll} \leq 1 & \text{for } i \in N^{\leq} \\ = 1 & \text{for } i \in N^= \\ \geq 1 & \text{for } i \in N^{\geq} \\ \text{unconstrained} & \text{for } i \in N^0 \end{array} \right. \quad \Sigma_j x_{i;j}$$
$$x_{i;j} = 0, 1 \quad \text{for } (i;j) \in A$$

where  $N^0, N^{\leq}, N^=, N^{\geq}$  is a given partition of  $N$ .

It is assumed that the reader is already familiar with the solution of this problem, which is discussed in [1].

## 2. PROGRAM INPUT

Every node and edge of the network is uniquely identified by a user assigned node/edge number. Each pseudonode created during the execution of the program is automatically associated with a unique pseudonode number. These numbers need not be serial, but they should lie within the ranges specified later on.

The program starts by reading the network which is defined by the data below:

- a) node data (node number, subset the node belongs to);
- b) edge data (edge number, pair of nodes joined by the edge, and edge-cost).

Successive iterations of post-optimality can be carried out. Changes introduced at each iteration define a new network, and these changes are always made on the latest network. Changes in the network are classified into the following types:

- a) edge-cost change (edge number, and new edge-cost);
- b) partition change (node number, and new subset);
- c) edge elimination (edge number);
- d) node elimination (node number);
- e) new node introduction (node number, and subset);
- f) new edge introduction (edge number, pair of nodes joined by the edge, edge-cost).

In short, the program input has the following structure:

- i) network data (node and edge data);
  - ii) first post-optimality iteration (change data);
  - iii) second post-optimality iteration (change data);
- and so on.

Just as a remark, this implementation is able to work with multi-networks; i.e., it is possible to have two or more edges joining the same pair of nodes. Therefore, every edge must be uniquely identified by an user assigned edge number (the cost of two parallel edges may be the same).

The description of the input data is based on the variables given below.

title - string with at most 48 characters, chosen by the user to identify the problem.

n - number of nodes (  $n \leq 100$  ).

m - number of edges (  $m \leq 3000$  ).

option - program option (see section 5 for details)

option = 1 : program prints statistics only;

option = 2 : program prints final solution only (suggested for most of the applications);

option = 3 : program prints progress report, final solution, and statistics;

option = 4 : program prints input data, progress report, final solution, and statistics.

node - node number (integer between 1 and 100); must be unique for each node.

set - indicates the subset that node belongs to, under the code:

set = 0: node  $\in N^0$

set = 1: node  $\in N^1$

set = 2: node  $\in N^2$

set = 3: node  $\in N^3$

edge - edge number (integer between 1 and 3000); must be unique for each edge.

cost - edge cost (real number between -9999.9 and +9999.9).

node<sub>1</sub>, node<sub>2</sub> - numbers of the pair of nodes joined by edge.

k<sub>1</sub> - total number of edges which have the cost coefficient changed during a post-optimality iteration.

k<sub>2</sub> - total number of set changes.

k<sub>3</sub> - total number of edge eliminations.

k<sub>4</sub> - total number of node eliminations.

k<sub>5</sub> - total number of node introductions.

k<sub>6</sub> - total number of edge introductions.

The user must provide the input data in FORTRAN free format; i.e., blanks or commas can be used as separator of data, and each line below must start in a new line of the input file.



## Input Data

title  
m, n, option  
node, set (one line per node)  
edge, node<sub>1</sub>, node<sub>2</sub>, cost (one line per edge)  
k<sub>1</sub>, k<sub>2</sub>, k<sub>3</sub>, k<sub>4</sub>, k<sub>5</sub>, k<sub>6</sub>  
edge, newcost (one line per cost change)  
node, newset (one line per set change)  
edge (one line per edge elimination)  
node (one line per node elimination)  
node, set (one line per node introduction)  
edge, node<sub>1</sub>, node<sub>2</sub>, cost (one line per edge introduction)  
...  
0, 0, 0, 0, 0, 0 (last line)

The 2nd, 3rd., and 4th. lines define the original network. From the 5th. line to the 11th. line above, a post-optimality iteration is defined. Every post-optimal iteration should be given under such a general form. When several iterations are used, the network at the beginning of iteration  $k$  is the same as the network at the end of iteration  $k-1$ . Finally, the last line of the input data should be a line with six zeros separated by commas or blanks.

### 3. PROGRAM EXECUTION

This section is appropriate for users of the Michigan Terminal System (MTS) of the University of Michigan; users of different installations should prepare the necessary modifications.

#### Batch Mode

On batch mode, the deck of cards to be punched should have the following structure:

```
$SIG ccid  
password  
$RUN K45V:MATCOV T=time  
.  
.   input data  
.  
$ENDFILE  
$SIG
```

where ccid, and password are assigned for each user of MTS. As an order of magnitude for time, MATCOV is able to solve a network with 50 nodes and 1000 edges in less than 5 s.

### Example 1

Consider the network in Figure I. The input data to be used in the solution of these three problems is shown in Output I. Note that the original network has 4 nodes and 5 edges, the first modified network has 4 nodes and 5 edges, and the last network has 3 nodes and 2 edges. Output II shows the optimum solution for the original network. The first post-optimality iteration was performed with 1 edge cost change, 1 edge elimination, 1 node elimination, 1 node introduction, and 2 edge introduction; edge  $240 = (20;40)$  was automatically eliminated from the network because it became incident with a nonexisting node (node 40 was eliminated). The optimum solution for the first modified network is given in Output III. In the second post-optimality iteration, one edge and one node were explicitly removed from the network, and two edges were automatically removed because they were incident with eliminated node. The final solution for the second modified network, which is infeasible, is presented in Output IV; the last network corresponds to an infeasible minimum cost 1-matching/covering problem.

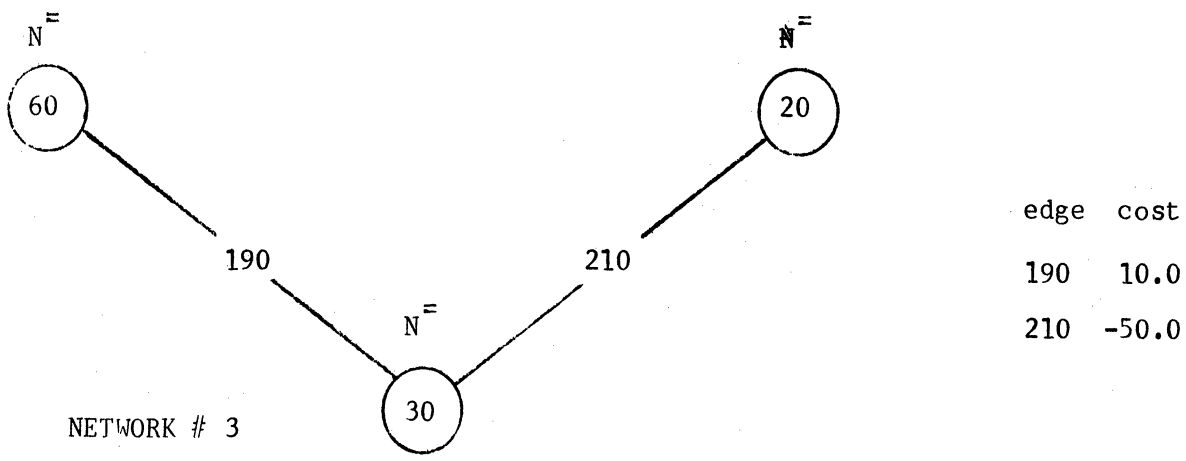
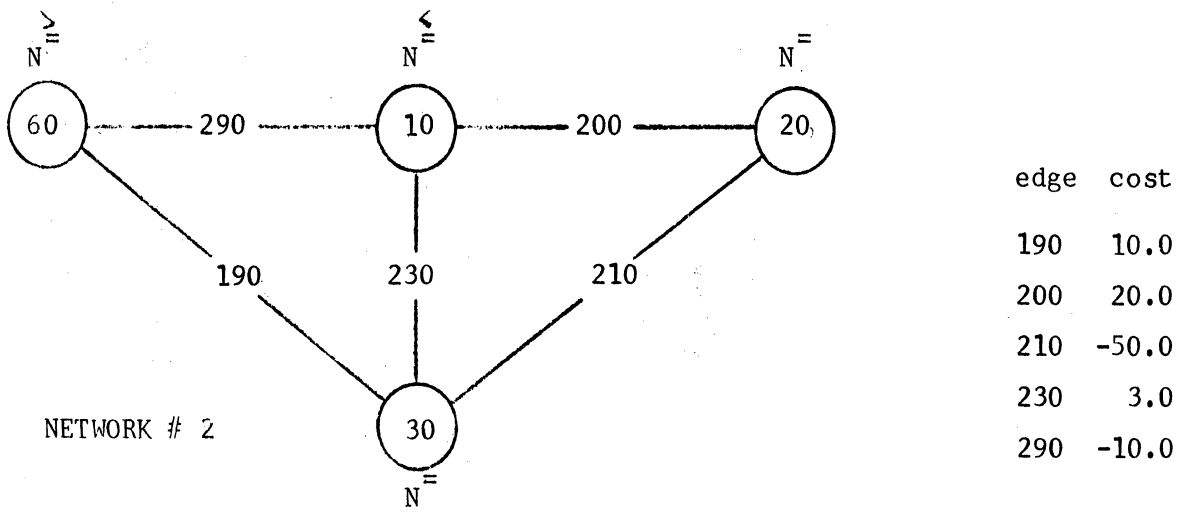
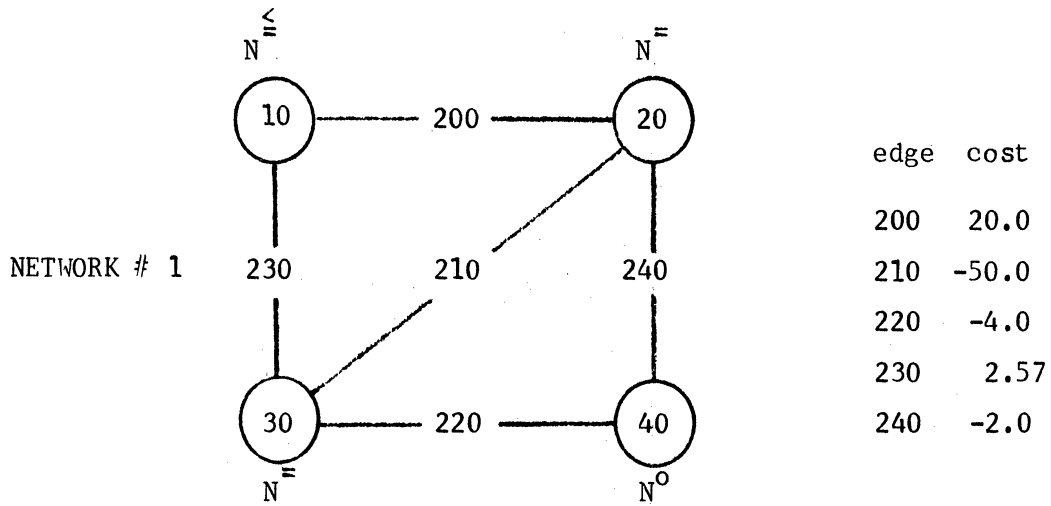


Figure I - Example 1

\$SI3 ccid  
 passwor1  
 sp1n K45V:HA^COV C=1

EXAMPLE 1

4	5	2		
10	2			
20	1			
40	0			
30	1			
240	20	40	-2.0	
220	30	40	-4.0	
200	20	10	20.0	
210	20	30	-50.0	
230	10	30	2.57	
1	0	1	1	2
230	3.0			
220				
40				
60	3			
290	10	60	-10.0	
190	60	30	+10.0	
0	1	0	1	2
60	1			
12				
0	0	0	0	0

\$ENDFILE

\$SI3

\*\*\*\*\* MINIMUM COST 1-MATCHING/COVERING PROBLEM \*\*\*\*\*

EXAMPLE 1

\*\*\*\*\* SOLUTION \*\*\*\*\*

NETWORK: 4 NODES, 5 EDGES

OBJECTIVE VALUE: -50.00

LIST OF NODES (NODE, SUP, PRICE) : 20 1 -25.0 40 3 0.0 30 1 -25.0

LIST OF EDGES (EDGE, NODE1, NODE2, COST) : 20 40 30 -4.0 200 10 20 20.0 210 30 20 -50.0 230 30 10 2.6

SOLUTION EDGES:  
210

\*\*\*\*\* POST-OPTIMALITY ANALYSIS \*\*\*\*\*

CHANGES: 1 EDGE COSTS, 2 NODE SETS  
 ELIMINATIONS: 1 EDGES, 1 NODES  
 INTRODUCTIONS: 1 NODES, 2 EDGES

\*\*\*\*\* SOLUTION \*\*\*\*\*

NETWORK: 4 NODES, 5 EDGES

OBJECTIVE VALUE: -60.00

LIST OF NODES (NODE, SET, PRICE) : 20 1 -25.0 30 1 -25.0 60 3 0.0  
 10 2 -10.0

LIST OF EDGES (EDGE, NODE1, NODE2, COST) : 230 30 10 3.0 290 10 60 -10.0 190 60 30 10.0  
 200 10 20 20.0 210 30 20 -50.0

SOLUTION EDGES: 230  
 210

\*\*\*\*\* POST-OPTIMILITY ANALYSIS \*\*\*\*\*

CHANGES: 0 EDGE COSTS, 1 NODE SETS  
ELIMINATIONS: 0 EDGES, 1 NODES  
INTRODUCTIONS: 0 NODES, 0 EDGES

\*\*\*\*\* SOLUTION \*\*\*\*\*

NETWORK: 3 NODES, 2 EDGES

= PROBLEM INFEASIBLE =

OBJECTIVE VALUE: -15.00

LIST OF NODES (NODE, SET, PRICE) : 30 1 -25.0 60 1 35.0  
20 1 -25.0

LIST OF EDGES (EDGE, NODE1, NODE2, COST) :  
210 30 20 -50.0 190 60 30 10.0

SOLUTION EDGES:  
210



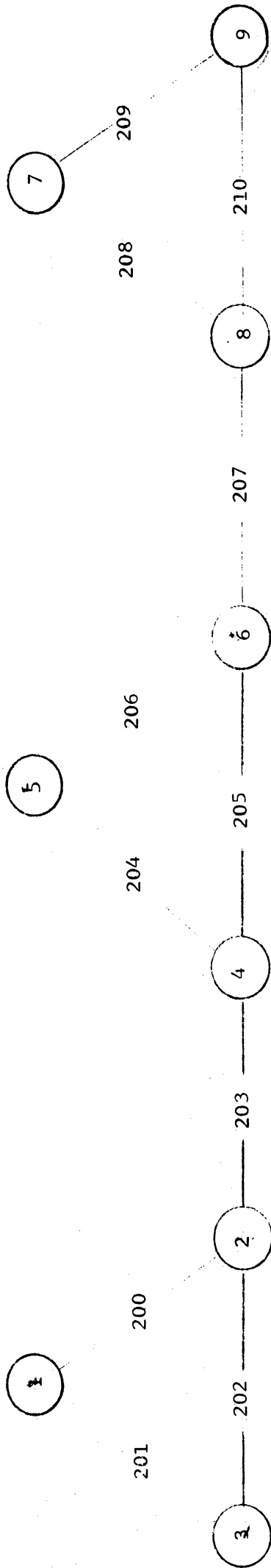
## Terminal Mode

On terminal mode, it is convenient to create a file containing the input data previously to the execution of the program. In the MTS statement below, infile indicates such a file, and outfile represents the file to which the output will be directed, such as \*PRINT\*. As an order of magnitude for time, MATCOV is able to solve a network with 50 nodes and 1000 edges in less than 5 s.

```
$RUN K45V:MATCOV 5=infile 6=outfile T=time
```

## Example 2

Consider the networks in Figure II. The input data for such an example is listed in Output V. All three networks have 9 nodes and 11 edges. In the first network, all nodes are in the subset  $N^=$ . During the first post-optimality iteration, node 9 is transferred to subset  $N^<=$ . During the second post-optimality iteration, node 9 is transferred to subset  $N^>=$ . The final solution for the original network, which is infeasible, is presented in Output VI; it contains 2 pseudonodes. Output VII shows the optimum solution for the first modified network, which contains just one pseudonode. Finally, the solution of the last network has two pseudonodes and it is shown in Output VIII.



edge 203 has cost 10.0

edge 207 has cost 80.0

remaining edges have cost 2.0

NETWORK # 1 - all nodes are in  $N^{\bar{}}$

NETWORK # 2 - node 9 is in  $N^{\bar{}}$ ; remaining nodes are in  $N^{\bar{}}$

NETWORK # 3 - node 9 is in  $N^{\bar{}}$ ; remaining nodes are in  $N^{\bar{}}$

Figure II - Networks of Example 2.

EXAMPLE 2

9 11 2

1 1  
 2 1  
 3 1  
 4 1  
 5 1  
 6 1  
 7 1  
 8 1  
 9 1

200 1 3 3.0  
 201 1 3 3.0  
 202 2 3 3.0  
 203 2 4 13.0  
 204 4 5 3.0  
 205 4 6 2.0  
 206 5 6 3.0  
 207 6 8 80.0  
 208 7 8 2.0  
 209 7 9 2.0  
 210 8 9 3.0

15

0 1 0 0 0  
 9 2  
 0 1 0 0 0  
 9 3  
 0 0 0 0 0

Output V - list of the input data corresponding to the Example 2

\*\*\*\*\* MINIMUM COST 1-MATCHING/COVERING PROBLEM \*\*\*\*\*

EXAMPLE 2

\*\*\*\*\* SOLUTION \*\*\*\*\*

NETWORK: 9 NODES, 11 EDGES

= PROBLEM INFASIBLE =

OBJECTIVE VALUE: 95.00

LIST OF NODES (NODE, SET, PRICE) :	2 1	1 0	3 1	1 0	4 1	1 0	5 1	1 0
1 1 1.0	7 1	1 0	8 1	1 0	9 1	1 0		
6 1 1.0								

LIST OF EDGES (EDGE, NODE1, NODE2, COST) :	2 0 2	3 2	2 0	2 0 3	4 2	10 0	2 0 4	5 4	2 0
200 2 1 2.0	201 3 1	2 0	202 3 2 2.0	203 4 2	10 0	204 5 4 2.0	205 6 4 2.0	206 6 5	2 0
207 8 6 90.0	208 8 7	2 0	209 9 7	2 0					
210 9 9 2.0									

PSEUDO PRICE SIMPLE BLOSSOM

101	9.0	2	3	1
103	78.0	0	7	9

SOLUTION EDGES: 203 206 201 208

Output VI - Final solution for Network # 1 of Example 2

\*\*\*\*\* POST-OPTIMALITY ANALYSIS \*\*\*\*\*

CHANGES: 2 EDGE COSTS, 1 NODE SPTS  
 ELIMINATIONS: 0 EDGES, 0 NODES  
 INTRODUCTIONS: 0 NODES, 0 EDGES

\*\*\*\*\* SOLUTION \*\*\*\*\*

NETWORK: 9 NODES, 11 EDGES

OBJECTIVE VALUE: 16.00

LIST OF NODES (NODE, SET, PRICE) :	2 1	1 0	3 1	1 0	4 1	1 0	5 1	1 0
1 1	1.0	7 1	1.0	3 1	1.0	4 1	1.0	1.0
6 1	1.0		1.0	4 1	1.0	9 2	0.0	

LIST OF EDGES (EDGE, NODE1, NODE2, COST) :	201	3 1	2 0	202	3 2	2 0	203	4 2	10 0	204	5 4	2 0
200 2 1	2.0	201 3 1	2.0	202 3 2	2.0	203 4 2	2.0	204 5 4	2.0	209 9 7	2.0	2.0
205 6 4	2.0	206 6 5	2.0	207 8 6	80.0	208 8 7	2.0					
210 9 9	2.0											

PSEUDO PRICE SIMPLE BLOSSOM  
 101 8.0 3 1

SOLUTION EDGES: 206 208 201

\*\*\*\*\* POST-OPTIMALITY ANALYSIS \*\*\*\*\*

CHANGES: 0 EDGE COSTS, 1 NODE SPTS  
 ELIMINATIONS: 0 EDGES, 0 NODES  
 INTRODUCTIONS: 0 NODES, 0 EDGES

\*\*\*\*\* SOLUTION \*\*\*\*\*

NETWORK: 9 NODES, 11 EDGES

OBJECTIVE VALUE: 18.00

LIST OF NODES (NODE, SPT, PRICE) :  
 1 1 1.0 2 1 1.0 3 1 1.0 4 1 1.0 5 1 1.0  
 6 1 1.0 7 1 1.0 8 1 1.0 9 3 -0.0

LIST OF EDGES (EDGE, NODE1, NODE2, COST) :  
 201 2 1 2.0 202 3 2 2.0 203 4 2 10.0 204 5 4 2.0  
 205 6 4 2.0 206 6 5 80.0 207 9 6 2.0 208 8 7 2.0  
 209 9 8 2.0

PSEUDO PRICE SIMPLE BLOSSOM  
 101 3.0 2 3 1  
 103 1.0 9 7 8

SOLUTION EDGES: 203 205 201 209 210

## 5. PROGRAM OUTPUT

There are several possible outputs in accordance with the value assigned by the user to the variable OPTION. They are classified into:

option = 1 : statistical output only;  
option = 2 : final solution output only;  
option = 3 : progress, final solution, and statistical outputs;  
option = 4 : input data, progress, final solution, and statistics;  
see section 9 for options providing debugging output.

The final solution output is printed at the end of each iteration; there is one for each network being solved (this output is produced by all options except by option 1). Such a report consists of the following information:

- a) message "PROBLEM IS INFEASIBLE", if appropriate;
- b) total number of nodes and total number of edges;
- c) dual objective value of the final solution (this is the optimum objective value whenever the problem is feasible);
- d) list of all nodes, comprising for each node,
  - i) node number,
  - ii) subset code (0, 1, 2, 3 for  $N^0$ ,  $N^=$ ,  $N^{\leq}$ ,  $N^{\geq}$ , respectively),
  - iii) node price (dual variable);

- e) list of pseudonodes (if appropriate), comprising for each pseudonode,
  - i) pseudonode number,
  - ii) pseudonode price (dual variable),
  - iii) list of nodes in the simple blossom;
  
- f) list of edges, comprising for each edge,
  - i) edge number,
  - ii) and iii) pair of nodes joined by the edge,
  - iv) edge cost;
  
- g) list of all edges in the (primal) solution.

Statistics are provided for every option, except by option 2.

It is expected to be used by researchers who are interested in comparing performance of codes. This report provides a table with the following information:

- a) Dual Solution Changes - number of times the dual solution change step was executed during the iteration;
  
- b) Paths rematched - number of alternating paths rematched during all augmentations of the iteration.
  
- c) Nodes Remated - number of nodes in those alternating paths;
  
- d) Nodes Scanned - number of nodes removed from the list of unscanned nodes;



- e) Edges Scanned - number of (equality) edges which were scanned;
- f) Nodes Labelled - number of outer and inner labels assigned;
- g) Shrinkings - number of simple blossom shrinkings;
- h) Unshrinkings - number of pseudonode unshrinkings;
- i) LINK/MATE's Used - number of storage positions used for LINK/MATE information (the value of P2 was estimated by this statistic);
- j) CPU MSEC (INPUT) - time in cpu msec used for initialization and input operations;
- k) CPU MSEC (SOLUTION) - time in cpu msec used for solving the problem;
- l) CPU MSEC (OUTPUT) - time in cpu msec used for performing the final solution transformation and printing the output.

The progress output is available by option 3 and option 4; it contains a number of messages indicating important marks that the algorithm achieves during the solution, such as the following:

- number of exposed nodes at initialization;
- node scanning;
- node labelling;
- augmentation;
- simple blossom shrinking;
- pseudonode unshrinking;

dual solution change;  
node/edge inspection;  
etc.

In this report, several symbols are used; among them we have the following:

$\langle R \rangle$  - indicates the root of an alternating tree

$\langle + \rangle$  - indicates an outer node

$\langle - \rangle$  - indicates an inner node

$\langle \rangle$  - indicates an unlabelled node

$\langle \Rightarrow \rangle$  - indicates an equality edge

edge = (node<sub>1</sub>, node<sub>2</sub>) - indicates an edge and the pair of nodes joined by the edge;

DUAL SOLUTION CHANGE: XXXX,X - indicates the increase in the dual objective value per exposed node (variable h in [1]).

Only option 4 provides a listing of the input data at the time data is entering; such output is useful for detection of errors in the input data.

## 5. ERROR MESSAGES

This section deals with the events that cause error messages and the subsequent halt in the execution of the program. It is expected that the message printed to be sufficient to lead to the error; however, since no indication is given to the input line the error was detected, it may be useful to rerun the program using option 4. For some errors, it may be necessary to augment the size of the program and recompile it (see section 10 for such cases).

The complete list of error messages is given below.

### UNEXPECTED END OF FILE

self explanatory.

### ILLEGAL PARAMETER

$n < 0,$        $k_1 < 0,$        $k_3 < 0,$        $k_5 < 0,$   
 $m < 0,$        $k_2 < 0,$        $k_4 < 0,$        $k_6 < 0.$

### ILLEGAL NODE

node  $< 1,$   
node  $> N2 (100),$   
duplication of an existing node,  
edge incident with a nonexisting node,  
elimination of a nonexisting node.

ILLEGAL EDGE

edge < 1

edge > M2 - L2 (3000)

duplication of an existing edge

elimination of a nonexistent edge.

ILLEGAL NODE SET

set < 0

set > 3

NODE STORAGE SPACE EXCEEDED [see section 10]

n > N2 (100)

introducing more than N2 (100) nodes.

EDGE STORAGE SPACE EXCEEDED [see section 10]

m > M2 - L2 (3000)

introducing more than M2-L2 (3000) edges.

LINK/MATE STORAGE SPACE EXCEEDED [see section 10]

HINK > P2 (1500)

ERROR NO. XXX

not used

## 6. CPU TIMES

Minimum cost 1-matching/covering problems were randomly generated by the program NET.OBJ (the source code called NET.FTN is listed in the appendix). The input for such a program should have the following structure (FORTRAN free format):

```
title
n, m, #sets, mincost, maxcost, seed
k1, k2, k3, k4, k5, k6
k1, k2, k3, k4, k5, k6
.
.
.
0, 0, 0, 0, 0, 0
```

where title, m, n, k<sub>1</sub> to k<sub>6</sub> have the meaning presented in section 2.

The parameter #sets indicates the number of subsets to be used as the partition for the set of nodes; #sets = 1 indicates that all the nodes should be in the set  $N^=$ ; #sets = 3 indicates that all nodes should be evenly distributed among the sets  $N^=$ ,  $N^{\leq}$ ,  $N^{\geq}$  (no node in  $N^0$ ); #sets = 4 indicates that all nodes should be evenly distributed among the sets  $N^=$ ,  $N^=$ ,  $N^=$ ,  $N^0$ . The parameters mincost, maxcost represent the range in which the edge costs are to be uniformly distributed. The parameter seed is used as an initialization number for the random number generator; different networks are generated just by changing seed. The number of post-optimality iterations to be generated are indicated by the number of nonnull lines with the parameters k<sub>1</sub> to k<sub>6</sub>. All problems are generated with program option = 1, which means that the execution of MATCOV will produce only statistics.

It should be noted here that the output produced by NET.OBJ can be used directly as input for MATCOV, except when there are node eliminations because this may imply that some edges become incident with nonexistent nodes; such edges are automatically eliminated by MATCOV from the network. Since NET.OBJ does not take into account such automatic edge eliminations, it may generate inconsistent changes afterwards.

Table I shows a set of average cpu times obtained from outputs of MATCOV under the statistic CPU MSEC (SOLUTION). These cpu times do not include the cpu time spent during the reading of the input data and during the printing of the final solution. All problems in Table I correspond to networks with 50 nodes; 1000 edges; edge costs (integers) generated in the intervals  $[-10,10]$ ,  $[-100,100]$ ,  $[100,200]$ ,  $[-200,-200]$ ; and partitions with 1, 3, 4 sets of nodes. Three problems were generated for each cell of Table I.

Table II presents the average cpu times obtained for problems with edge costs generated in the range  $[-100,100]$ ; partitions with 1, 3, 4 sets of nodes; and  $(\#nodes, \#edges)$  taking the values  $(20,200)$ ,  $(40,800)$ ,  $(60,1800)$ . Three problems were generated for each cell of Table II. Note that  $\#nodes = (\#edges)^2 / 2$  for all cases.

Partition	Edge Cost Range			
	[-10, 10]	[-100, 100]	[100, 200]	[-200, -100]
$N^=$	902	907	1064	887
$N^=, N^=, N^=$	485	505	1372	477
$N^=, N^=, N^=, N^0$	355	354	905	354

Table I - cpu msec (no readings, no writings) for solving minimum cost 1-matching/covering problems with 50 nodes and 1000 edges using AMDAHL 470/V7. Figures are the average of 3 different problems.

Partition	(# Nodes, # Edges)		
	(20, 200)	(40, 800)	(60, 1800)
$N^=$	63	472	1796
$N^=, N^=, N^=$	39	252	887
$N^=, N^=, N^=, N^0$	34	199	640

Table II - cpu msec (no readings, no writings) for solving minimum cost 1-matching/covering problems with edge cost ranging in the interval [-100,100] using AMDAHL 470/V7. Figures are the average of 3 different problems.

Table III indicates the average cpu times obtained for networks with 50 nodes, 1000 edges, edge cost range  $[-100,100]$ , and partitions with all 4 sets. Post-optimality analysis was applied to the problems in order to have 1000 edge cost changes. Due to the way problems are generated, some edges may have their edge costs changed more than one time in the same post-optimality iteration; this means that some edges may keep their edge costs during one iteration. Two post-optimality iterations were applied for each problem; each cell of Table III corresponds to the average of 3 problems. Cpu times spent for reading the input data, for initialization, and for preparing a post-optimality iteration are shown as CPU TIME (INPUT); they can be compared against the CPU TIME (SOLUTION) spent for solving the problems.



	CPU TIME (MSEC)	
	Input	Solution
Original Network	361	315
1000 edge cost changes	301	119
1000 edge cost changes	299	113

Table III - cpu times for input and solution of problems with 50 nodes and 1000 edges, edge cost range  $[-100,100]$  using AMDAHL 470/V8. Problems had twice 1000 edge cost changes. Figures are for the average of 3 different set of problems.

## 7. DATA STRUCTURE

Intenally, every node is identified by an unique number between N0 (1) and N2 (100); every pseudonode by a number between L0 (101) and L2 (150); and every edge by a number between M0 (151) and M2 (3150). So, the external and internal representation of nodes and pseudonodes is based on the same node/pseudonode number. However, the internal representation of an edge is obtained by the expression

$$\text{internal edge number} = \text{external edge number} + L2.$$

Moreover, each edge is decomposed into two arcs joining the same pair of nodes, but with different directions; one of the arcs is identified by the edge number corresponding to its edge; the other arc is identified by the complement of that number to  $M9 = 2.M2 + 1 = 6301$ .

$$\begin{aligned} \text{arc}_1 &= \text{edge} \\ \text{arc}_2 &= M9 - \text{edge} \end{aligned}$$

In this way, every arc is uniquely identified by an arc number in the range M0 (151) to  $M9 - L2 - 1$  (6150).

In the context of this section, all node/pseudonode/edge/arc numbers refer to the internal representation, unless it is explicitly sated as external.

## Main Variables

N - number of nodes

M - number of edges

N2 = 100 - maximum node number

L2 = 150 - maximum pseudonode number

M2 = 3150 - maximum edge number

P2 = 1500 - maximum position in LINK/MATE

INF = 9999.9 - maximum edge cost

TOLER =  $10^{-50}$  - tolerance for numerical comparisons

VALUE - dual objective value

ROOT\$ - number of exposed nodes

The variables NODE, PNODE (pseudonode), EDGE (edge or arc), IND (index), BLSM (node in a simple blossom), BASE (base of a simple blossom), APICE (apex of a blossom), BNODE or BCINOD (BCI of a blossom), PT (pointer) are used throughout the program with numerical appendices (e.g., NODE2); all these variables have always the meaning indicated, except NODE which is normally used as node or pseudonode, but it may indicate an edge or an arc. The variables ROOT\$, NODE\$, EDGE\$, SHRK\$, UNSH\$, AUGM\$, DUAL\$, RMAT\$, LABL\$, LINK\$, TIME1, TIME2, TIME3 have statistical usage. The values assigned to the variables PRINT, TRACE depend on the program option.

## List of existing node/edge numbers

LIST(ind) - three-stack vector containing all node/edge numbers which correspond to existing elements in the network.

node numbers:	head = N0	tail = N1
pseudonode numbers:	head = L0	tail = L1
edge numbers:	head = M0	tail = M1

The heads N0, L0, M0 are fixed during the execution; the tails N1, M1 are fixed for each post-optimality iteration; the tail L1 depends on the number of pseudonodes in the solution.

INDEX(node/edge) - vector of indices for LIST

N0 = INDEX(node) = N1	N0 = node = N2
L0 = INDEX(pnode) = L1	L0 = pnode = L2
M0 = INDEX(edge) = M1	M0 = edge = M2

INDEX(node/edge) = 0 implies a nonexisting node/edge

INDEX(node/edge) > 0 implies an index to LIST

for existing node/edge: LIST(INDEX(node/edge)) = node/edge

for a valid ind: INDEX(LIST(ind)) = ind

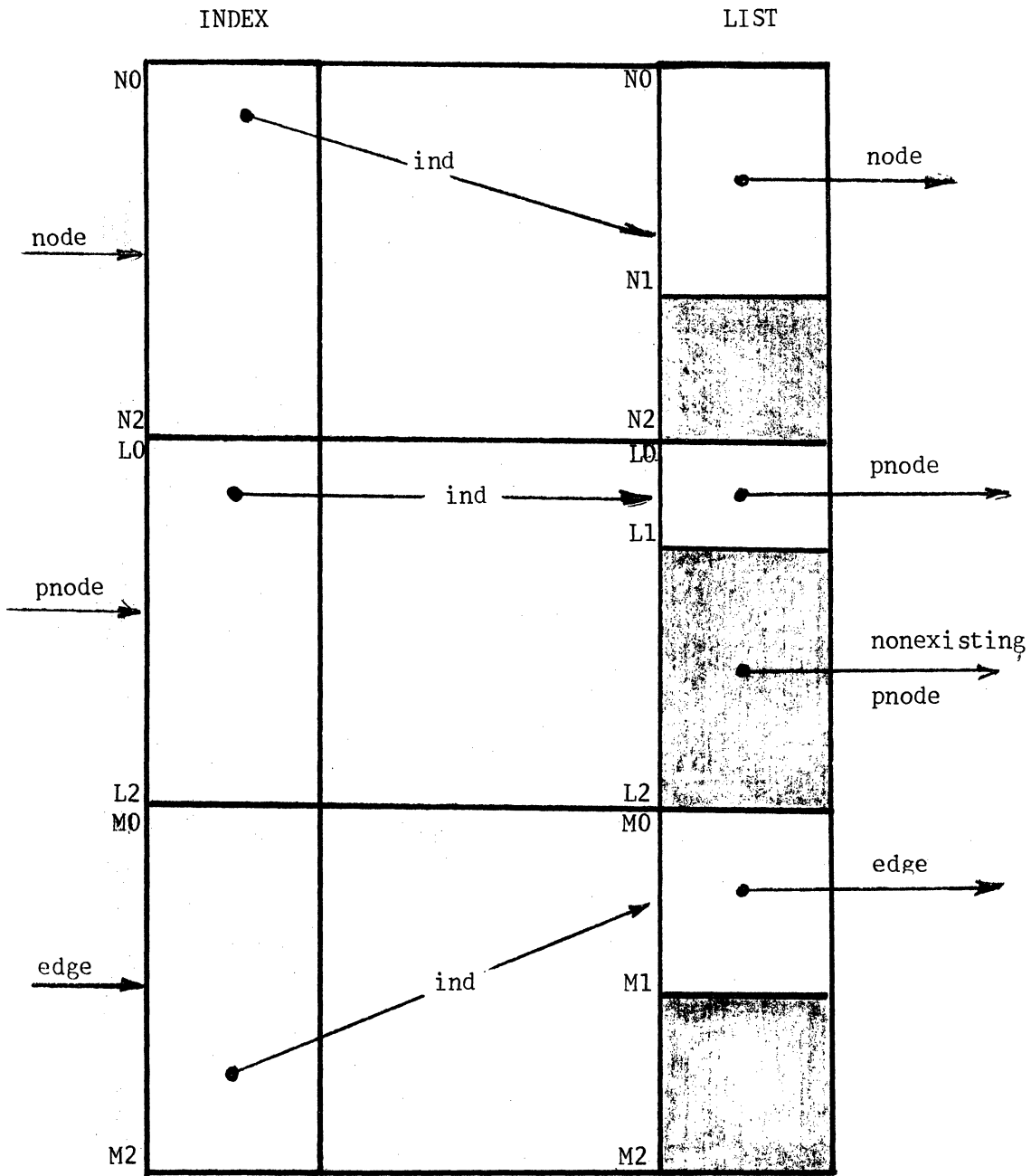


Figure III - list of existing node/edge numbers

## Network Specification:

SET(node) - indicates the partition subset that node belongs to,  
under the code:

- 0 - node of  $N^0$
- 1 - node of  $N^1$
- 2 - node of  $N^2$
- 3 - node of  $N^3$

PRICE(edge) - edge cost of edge

HEAD(arc) - indicates the node that arc is heading to.

in order to determine the two nodes joined by an edge, use the  
formulas:

$$\begin{aligned} \text{arc}_1 &= \text{edge} \\ \text{arc}_2 &= M9 - \text{edge} \\ \text{node}_1 &= \text{HEAD}(\text{arc}_1) \\ \text{node}_2 &= \text{HEAD}(\text{arc}_2) \end{aligned}$$

## Arc Lists

LINK/MATE(pt) - pair of vectors used for storing the following lists:

incidency-list(node) - contains all current equality arcs heading from node (incidency in the current equality network);

mating-list(node) - contains all current solution arcs heading from node;

internal-list(pnode) - contains equality arcs joining two nodes inside pnode

LINK(node) - head of

MATE(node) - head of mating-list(node)

LINK(pnode) - head of internal-list(pnode)

MATE(pnode) - base of pnode

MATE(pt) - contains the arc number of an element in one of those three lists

LINK(pt) - points to LINK/MATE indicating the next element in the list

Obs.: MATE(node) = 0 implies that there is no solution arc heading from node

MATE(node) > 0 implies that arc = MATE(node) is the only solution arc heading from node

MATE(node) < 0 implies that pt = -MATE(node) is the position in LINK/MATE containing the first element of mating-list(node).

if APEX(PSEUDO(node))  $\neq$  node then MATE(node) is meaningless; in words, the solution edges inside a pseudonode are determined only at the time the final solution is being printed.

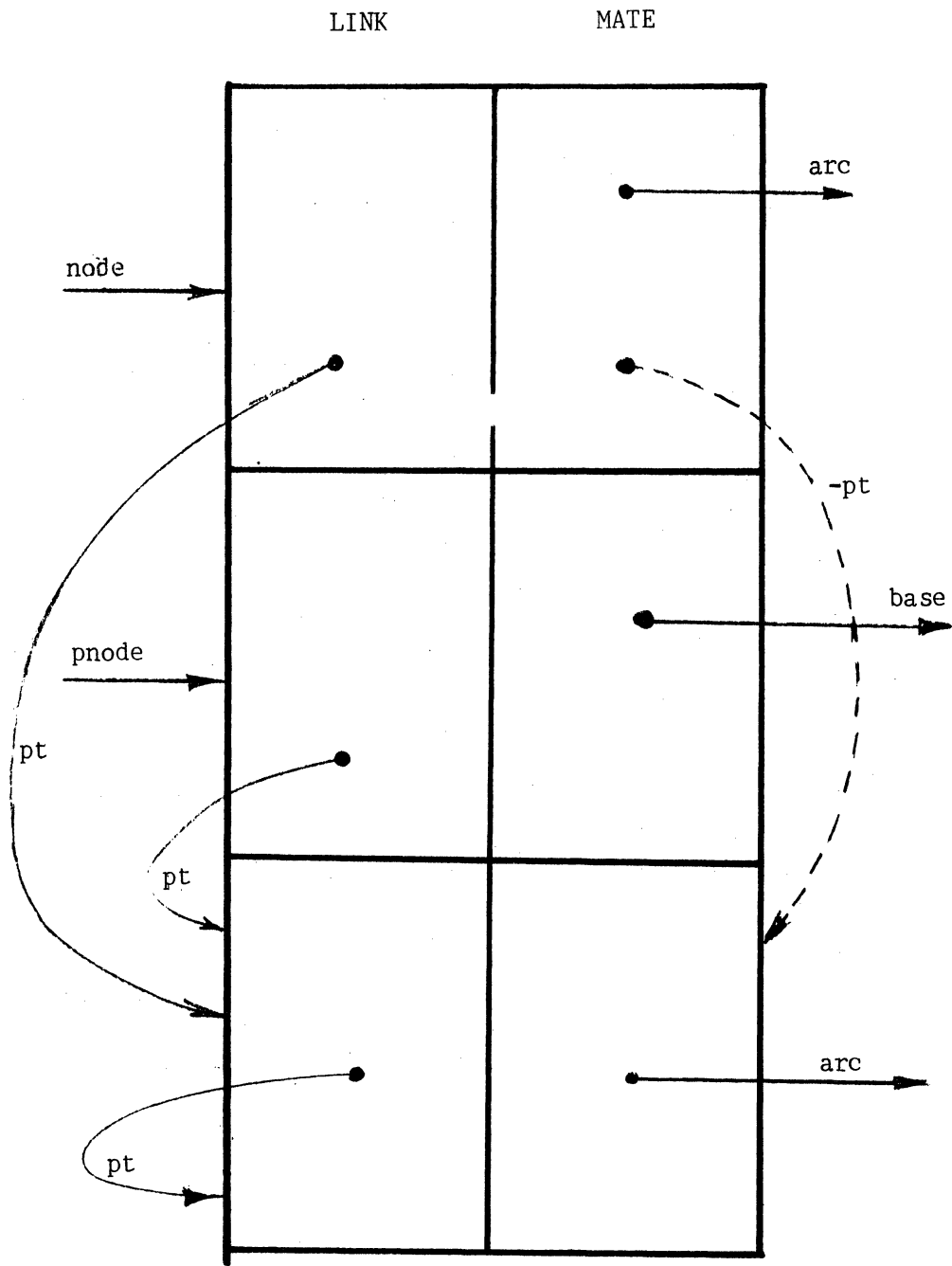


Figure IV - LINK/MATE pair of vectors



Labelling Structure (current nodes)

LABLO(node) - label of node (OUTER, INNER, UNLABD)

PREDA(node) - predecessor arc; this is the second arc on the path from node to the root; defined only for outer nodes

HEAD(node) - points to the apex of the root of the alternating tree

TREE(node) - single-linked list containing all nodes of an alternating tree. The head of such a list is the apex of the root of the alternating tree.

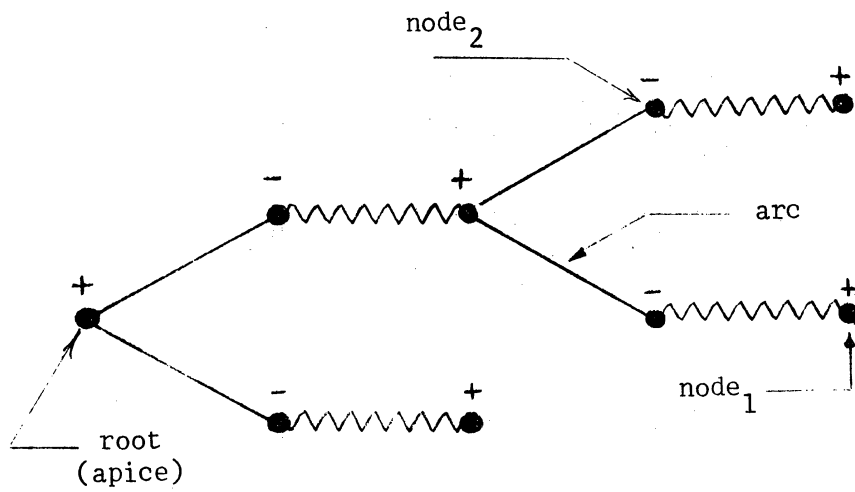
Simple Blossom Structure (noncurrent nodes)

LABLO(node) - pseudonode number of the simple blossom containing node.

HEAD(node) - node number of the next node in the cyclic list containing all nodes in a simple blossom

PREDA(node) - arc number of the arc heading from node to HEAD(node)

MATE(pnode) - base of the simple blossom associated with pnode



Alternating tree

LABLO(node<sub>1</sub>) = outer

LABLO(node<sub>2</sub>) = inner

LABLO(root) = outer

HEAD(node<sub>1</sub>) = apice

HEAD(node<sub>2</sub>) = apice

HEAD(root) = apice

apice = APEX(root)

PRED(node<sub>1</sub>) = arc

PRED(node<sub>2</sub>) = null

PRED(root) = null

Figure V - Labelling Structure

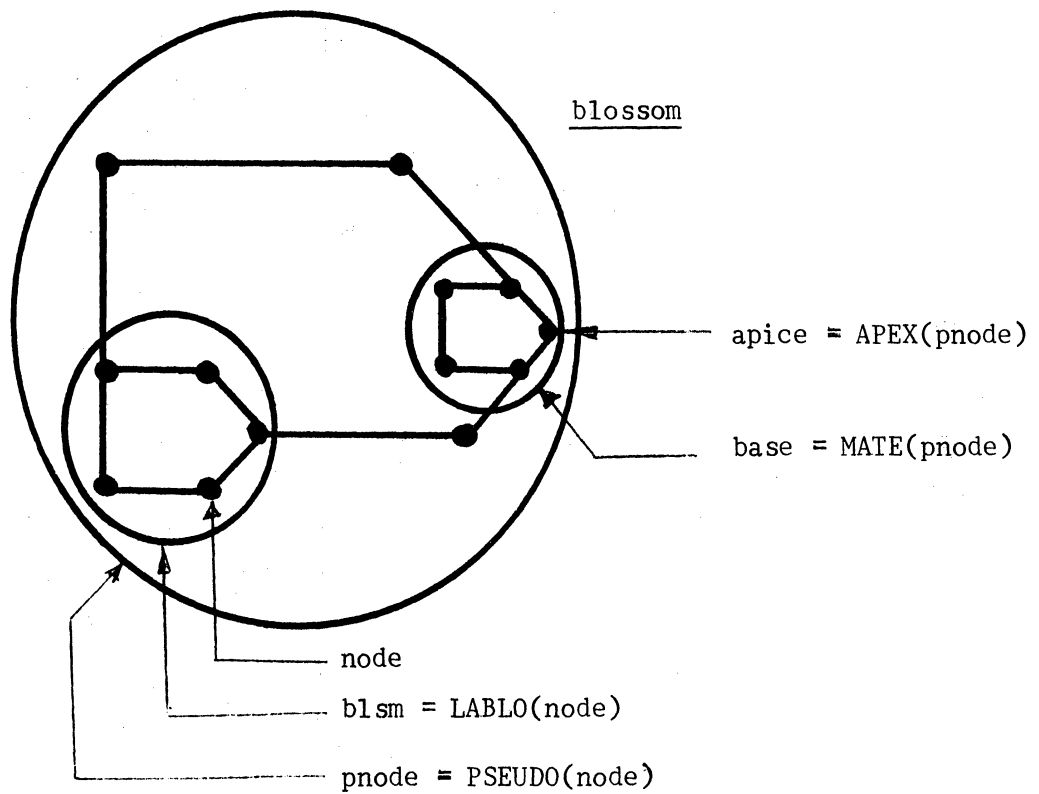
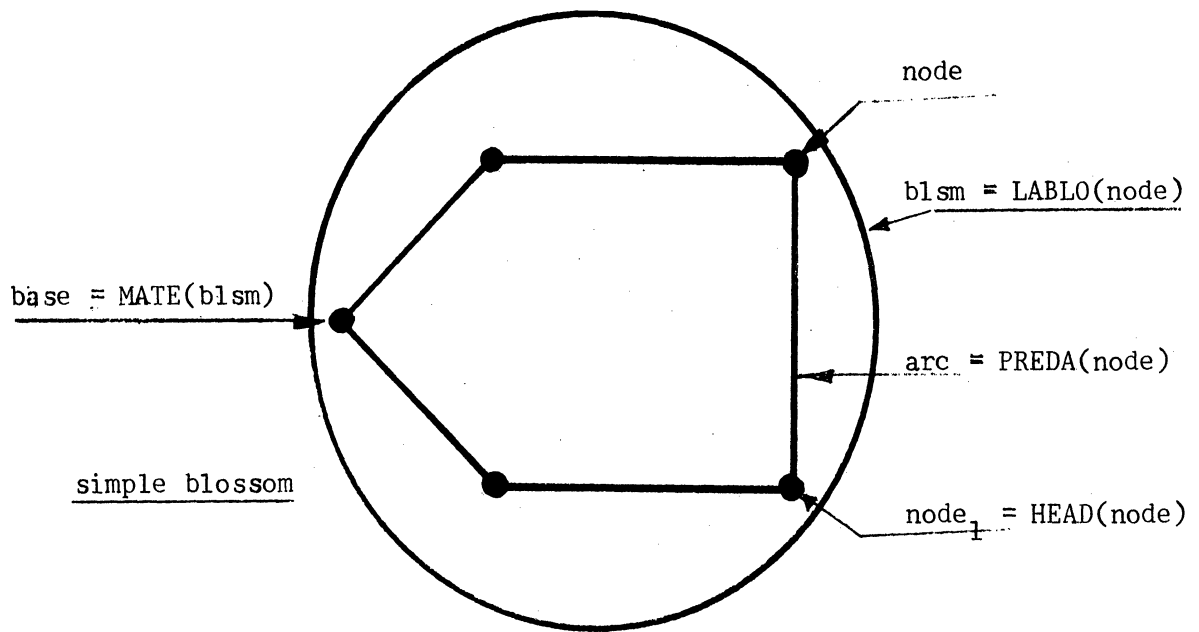


Figure VI - Blossom and Simple Blossom Structure

Blossom Structure (noncurrent nodes)

PSEUDO(node) - node in the current network which contains node;

PSEUDO(node) = node implies node is current

APEX(pnode) - points to an original node inside pnode;

for original nodes: APEX(node) = node

PRICE(node) - corresponds to the dual variables; let  $y(i)$ ,  $z(s)$  be

the dual variables defined in [1]; let  $\text{node} \equiv i$ ,  $\text{pnode} \equiv s$

$$\text{PRICE}(i) = \begin{cases} y(i) & \text{for } i \in N^0 \cup N^{\leq} \\ y(i) + \sum_s [ |z(s)| / 2 : s \text{ contains } i ] & \text{for } i \in N^{\leq} \\ y(i) + \sum_s [ |z(s)| : s \text{ contains } i ] & \text{for } i \in N^{\geq} \end{cases}$$

$$\text{PRICE}(s) = |z(s)| / 2 \quad \text{for } s \in S$$

At the time the final solution is to be printed, the values of  $y(i)$  and  $z(s)$  are computed from the above expressions and printed.

BCI(node) - points to the Blossom Constraint Identifier.

BCI(node) = node     $\text{node} \in N^{\leq} \cup N^{\geq} \cup N^0$

BCI(node) = null     $\text{node} \in N^{\leq}$

BCI(pnode) = node     $\text{node} \in N^{\leq} \cup N^{\geq}$

BCI(pnode) = null    all nodes inside pnode are in  $N^{\leq}$

MPRICE(node) - maximum value for PRICE(BCI(node))

$MPRICE(node) = 0$  for node  $\in N^{\leq} \cup N^0$

$MPRICE(node) = INF$  for node  $\in N^=$  or current node  $\in N^{>}$

$MPRICE(node) = \text{variable}$  for noncurrent node  $\in N^=$

$MPRICE(pnode) = MPRICE(BCI(node))$

For noncurrent node of  $N^{>}$  the value of  $MPRICE$  is assigned at the time the node becomes noncurrent by setting

$MPRICE(node) = 2.PRICE(node)$

### Special Lists

$NSCAN(pt)$  - cyclic queue of unscanned nodes

Head:  $HSCAN$

Tail:  $TSCAN$

Usable size:  $N9 = N + 1$

Nodes to be scanned are removed from the head of  $NSCAN$ ; nodes just outer labelled are introduced in the tail of  $NSCAN$ ; inner nodes of just erased alternating trees are introduced in head of  $NSCAN$ ; pnodes to be scanned just removed from  $NSCAN$  have the nodes in their simple blossoms introduced in the head of  $NSCAN$ .

$INSP(pt)$  stack of nodes/edges to be inspected after a dual solution change step have been performed

head:  $HINSP$

usable size :  $M + N$

`NORIG(pt)` - stack of nodes inside a blossom

`Head:` `HORIG`

`Usable size:`  $3 N / 2$

## 8. PROGRAM STRUCTURE

This section discusses the main characteristics of each subroutine in the FORTRAN program MATCOV.

### MATCOV (main program)

```
call GETNET
10 call GRWFOR
call PUTSOL
call POSTOP
go to 10
```

### Subroutine GETNET (Get network)

```
read the original network
initialize the variables
plant the alternating forest
special calls: MATCH
main divisions: i) initialization; ii) read node data and set
node variab.; iii) read edge data and set edge variables;
iv) plant alternating tree.
```

### Subroutine GRWFOR (Grow Forest)

```
solve a network with an alternating tree
controls the scanning, dual solution change, and inspection
processes;
starts unshrinking and augmentation processes;
checks both termination criteria: optimality and infeasibility;
special calls: SCAN, MINIM, RMATCH, NSHRINK
main divisions: i) scanning; ii) dual solution change; iii) inspection
```

### Subroutine PUTSOL (Put Solution)

```
transform the final solution: internal to external representation
prints the final solution
no special calls
main divisions: i) get solution edges and transform node prices;
ii) get simple blossoms and solution edges inside pseudonodes;
iii) print statistics
```

Subroutine POSTOP (Post-Optimality)

controls the post-optimality process;  
reads and performs the changes in the network  
plants; the alternating forest  
special calls: NMATCH, OPEN, DECRSE, MATCH, INCRSE, INCLUD  
main divisions: i) introduction; ii) edge cost changes; iii)  
node set changes; iv) edge elimination; v) node elimination;  
vi) node introduction; vii) edge introduction; viii) plants  
alternating forest

Subroutine SCAN (Scan arc)

Scan an equality arc whose tail is an outer node  
starts shrinking, unshrinking, and augmentation processes;  
assigns inner and outer labels;  
special calls: SHRINK, RMATCH, NSHRNK  
main divisions: i) initialization; ii) shrinking check; iii) double  
augmentation check; iv) type 2 augmentation check; v) type 0  
augmentation check; vi) type 1 augmentation check; vii) node  
labelling; viii) type 3 augmentation check; ix) unshrinking  
check.

Subroutine SHRINK (Shrink a simple blossom)

Given an edge joining two outer nodes of the same alternating  
tree, this subroutine finds the simple blossom containing that  
edge and shrinks it into a pseudonode  
special calls: ORIG, STPSDO, RMATCH  
main divisions: i) get simple blossom; ii) shrink simple blossom  
and determine BCI; iii) type 3 augmentation check.

Subroutine NSHRNK (Unshrink a pseudonode)

Given an inner labelled pseudonode with null pseudonode price ,  
this subroutine unshrinks it into its simple blossom.  
special calls: STPSDO, STAPEX  
main divisions: i) get alternating path in simple blossom; ii)  
unshrink pseudonode; iii) forest growing check;

Subroutine RMATCH (Rematch alternating path)

special calls: STAPEX  
main divisions: i) remate the nodes in the alternating path  
ii) erase the alternating tree



Subroutine MATCH

Match two nodes joined by a given edge

Subroutine NMATCH (Unmatch)

Eliminate the mate connections of two nodes joined by a given edge.

Subroutine MINIM (Minimization)

Checks and updates the inspection list

Subroutine STPSDO (Set Pseudo)

Update the variable PSEUDO of all nodes inside a pseudonode  
special call: ORIG

Subroutine STAPEX (Set Apex)

Update the basis and apices of all nodes inside a pseudonode with  
a new apex.

Subroutine ORIG (Original nodes)

Get the nodes in the simple blossom of a given pseudonode. It is  
useful for obtaining the original nodes inside a blossom.

Subroutine OPEN

Unshrink all pseudonodes containing a given noncurrent node  
Remove all solution edges incident with a given current node  
special calls; NMATCH, ORIG, STPSDO, STAPEX  
main divisions: i) open pseudonodes containing a noncurrent node;  
ii) remove the solution edges incident with a current node.

Subroutine INCLUD (Include edge)

Include a new edge in the network  
special calls: OPEN, DECRSE, MATCH

Subroutine DECRSE (Decrease node price)

Attempts to decrease the node price of a given current node by a given amount

Subroutine INCRSE (Increase node price)

Increases the price of a given current node to zero.  
Special calls: OPEN, DECRSE

Subroutine ERROR

Prints error messages.

Subroutine DUMP

Dumps node and edge tables

## 9. ADDITIONAL OUTPUTS

This section deals with program options which produces outputs with debugging purpose.

Options 5 to 9 activate the subroutine DUMP which prints tables with node, pseudonode, and edge internal information at specific points during the execution of the program, as follows:

option	subroutine	timming
5	PUTSOL	before and after each final solution transformation
6	POSTOP	before and after each post-optimality iteration
7	GETNET	after the initialization of the original network
8	GRWFOR	after each dual solution change step
9	GRWFOR	before and after each scanning step

Options corresponding to nonpositive integers causes messages indicating all subroutine calls and respective parameters to be printed during the periods of time given below. Also, dumps are provided at the beginning and at the end of these periods. The word trace is used as a reference to such actions.

option = 0 : traces while there are no exposed nodes during the solution of the original network;

option = -r : traces while there are r exposed nodes during the solution of the original network;

option = -N2 : traces since initialization until the first augmentation;

option = -(k.N2+r) : traces while there are r exposed nodes during the k-th post-optimality iteration;

option = -(k.N2) : traces since the beginning of the (k-1)-th post-optimality iteration until an augmentation occurs.

Finally, option =10 is useful in determining the period of time a trace is desirable. Such an option prints a message stating the number of exposed nodes every time an augmentation is performed.

A remark should be made here with respect to the dumps and traces which are produced by the options described in this section. The information is always given in terms of the internal representation and of the internal values. This means (cf. section 8):

- a) edges are internally identified by (edge number + L2); in the current version of the program, L2 = 150; so, an edge number 200 is internally represented by 350 = 200 + 150;
- b) node prices of original nodes inside a pseudonode are internally represented by:

$$\text{PRICE}(i) \begin{cases} = y(i) & \text{for } i \in N^{\leq} \\ = y(i) - \sum_s [ z(s) : s \text{ contains } i ] & \text{for } i \in N^{\equiv} \\ = y(i) - \sum_s [ 2z(s) : s \text{ contains } i ] & \text{for } i \in N^{\geq} \end{cases}$$

- c) solution edges inside a pseudonode are not provided.

## 10. MODIFYING THE PROGRAM SIZE

The design of MATCOV allows for modifications in the total number of nodes ( $N2 = 100$ ), in the total number of edges ( $M2-L2 = 3000$ ), in the total number of LINK/MATE's ( $P2 = 1500$ ), and in the range of the edge costs (  $[-INF,INF]$ ;  $INF = 9999.9$  ) to be easily performed, as explained below. Such modification should be made on the file PD.FTN which contains the source code in FORTRAN for MATCOV and it is listed in the appendix.

### a) Modifying the total number of nodes

Let  $n'$  denote the new maximum number of nodes that will replace the current value 100. In BLOCKDATA assign  $N2 = n'$ ,  $L2 = (3n'+1)/2$ . Every vector with dimension 100 should be redimensioned with  $n'$ . Every vector with dimension 101 should be redimensioned with  $(n'+1)$ . Every vector with dimension 150 should be redimensioned with  $(3n'+1)/2$ . It may be necessary to rewrite some FORMAT statements. It should be noted here that the maximum number of edges will be modified also; if necessary, follow the instructions in (b).

### b) Modifying the total number of edges

Let  $m'$  denote the new maximum number of edges that will replace the current value  $M2 - L2$  (3000). In BLOCKDATA assign  $M2 = L2 + m'$ . Every vector with dimension 3150 should be redimensioned to  $L2 + m'$ . Every vector with dimension 6150 should be redimensioned to  $L2 + 2m'$ . It may be necessary to rewrite some FORMAT statements.

c) Modifying the total number of LINK/MATE positions

Let  $p'$  denote the new maximum number of LINK/MATE positions that will replace the current value P2 (1500). In BLOCKDATA assign  $P2 = p'$ . Every vector with dimension 1500 should be redimensioned to  $p'$ .

d) Modifying the value of INF or TOLER

In BLOCKDATA assign INF (TOLER) to its new value. Note that this variable is in the current version of the program a REAL\*4 variable.

11. REFERENCE

C. Perin, "Matching and Edge Covering Algorithms", Ph.D. Dissertation, Department of Industrial and Operations Engineering, The University of Michigan, 1980.

## APPENDIX





```

WRITE (7,7)      K1, K2, K3, K4, K5, K6
IF              ( K1+ K2+ K3+ K4+ K5+ K6 .EQ. 0 ) STOP
C
C  GENERATE K1 CHANGES OF EDGE COSTS
IF (K1.EQ.0) GO TO 200
DO 120 K = 1, K1
110   EDGE = 151 + RAND(IX) * M
      IF (.NOT.SAVE(EDGE))          GO TO 110
      IF (.NOT.SAVE(HEAD1(EDGE))) GO TO 110
      IF (.NOT.SAVE(HEAD2(EDGE))) GO TO 110
      CST  = RAND(IX) * DELTA + MINCST
120   WRITE (7,8) EDGE, CST
C
C  GENERATE K2 CHANGES OF NODE SUBSETS
200 IF (K2.EQ.0) GO TO 300
DO 220 K = 1, K2
210   NODE      = 1 + RAND(IX) * N
      IF (.NOT.SAVE(NODE)) GO TO 210
      SET(NCDE) = RAND(IX) * NSETS + 1
      IF (SET(NODE).GT.3) SET(NODE) = 0
220   WRITE (7,5) NODE, SET(NODE)
C
C  GENRATE K3 EDGE DELETIONS
300 IF (K3.EQ.0) GO TO 400
IF (K3.GT.M1) GO TO 999
DO 320 K = 1, K3
310   EDGE      = 151 + RAND(IX)* M
      IF (.NOT.SAVE(EDGE))          GO TO 310
      IF (.NOT.SAVE(HEAD1(EDGE))) GO TO 310
      IF (.NOT.SAVE(HEAD2(EDGE))) GO TO 310
      SAVE(EDGE) = .FALSE.
320   WRITE (7,5) EDGE
      M1 = M1 - K3
C
C  GENERATE K4 NODE DELETIONS
400 IF (K4.EQ.0) GO TO 500
IF (K4.GT.N1) GO TO 999
DO 420 K = 1, K4
410   NODE      = 1 + RAND(IX) * N
      IF (.NOT.SAVE(NODE)) GO TO 410
      SAVE(NODE) = .FALSE.
420   WRITE (7,5) NODE
      N1 = N1 - K4
C
C  GENERATE K5 NEW NODE INTRODUCTIONS
500 IF (K5.EQ.0) GO TO 600
IF (N+K5.GT.100) GO TO 999
DO 510 K = 1, K5
      N          = N + 1
      NODE      = N
      SAVE(NODE) = .TRUE.
      SET(NCDE) = RAND(IX) * NSETS + 1
      IF (SET(NODE).GT.3) SET(NODE) = 0
510   WRITE (7,5) NODE, SET(NODE)
C
C  GENERATE K6 NEW EDGE INTRODUCTIONS
600 IF (K6.EQ.0) GO TO 100
IF (M+K6.GT.3000) GO TO 999
DO 630 K = 1, K6
      M          = M + 1

```

```

        EDGE          = M + 151
        SAVE(EDGE)    = .TRUE.
610     NODE1         = 1 + RAND(IX) * N
        IF (.NOT.SAVE(NODE1)) GO TO 610
620     NODE2         = 1 + RAND(IX) * N
        IF (.NOT.SAVE(NODE2)) GO TO 620
        IF (NODE1.EQ.NODE2) GO TO 620
        HEAD1(EDGE)  = NODE1
        HEAD2(EDGE)  = NODE2
        CST          = RAND(IX) * DELTA + MINCST
630     WRITE(7,5) EDGE, NODE1, NODE2, CST
C
C     RETURN FOR A NEW POST-OPTIMALITY ITERATION
C     GO TO 100
C
C     ERROR
999    WRITE(6,99)
        STOP
C
C     FORMATS
1     FORMAT(' ENTER TITLE')
2     FORMAT(48H12345678901234567890123456789012345678901234567890123456789012345678 )
3     FORMAT(' ENTER N, M, #SETS, MINCST, MAXCST, IX')
4     FORMAT(2I5,' 1')
5     FORMAT(3I5,F10.4)
6     FORMAT(' ENTER K1, K2, K3, K4, K5, K6')
7     FORMAT(6I5)
8     FORMAT(I5,F10.4)
99    FORMAT(' ***** ERROR *****')
        END
C
C     GENERATE RANDOM NUMBERS
        FUNCTION RAND(IX)
        IX = IX * 65539
        IF (IX.LT.C) IX = IX + 2147483647 + 1
        RAND = IX * .4656613E-9
        RETURN
        END

```







```

CALL GETNET
10 CALL GRWFOR
CALL PUTSOL
CALL POSTOP
GO TO 10

```

C

```

END
BLOCKDATA

```

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MINIMUM COST 1-MATCHING/COVERING PROBLEM
THE UNIVERSITY OF MICHIGAN
WRITTEN BY CLOVIS PERIN
AUGUST 1980

```

```

/BLOCK DATA/

```

```

DATA STRUCTURE DESCRIPTION

```

```

N = #NODES      M = #EDGES      L = MAX #PSEUDOS

```

```

LIST(IND) = NODE/EDGE,  NODE/EDGE EXISTS
LIST(IND) = NULL        NODE/EDGE DOES NOT EXIST

```

```

LIST( INDEX(NODE/EDGE) ) = NODE/EDGE  IFF NODE/EDGE EXISTS

```

```

EXISTING NODES:   FROM LIST(N0)  TO LIST(N1)
EXISTING PSEUDOS: FROM LIST(L0)  TO LIST(L1)
EXISTING EDGES:   FROM LIST(M0)  TO LIST(M1)

```

```

MAX NODE   = N2 >= N1
MAX PSEUDO = L2 >= L1
MAX EDGE   = M2 >= M1

```

```

VALUE = DUAL OBJECTIVE VALUE
ROOT$ = # EXPOSED NODES

```

```

SET(NODE) = 0, 1, 2, 3      (N0, N=, N<, N>, RESPECTIVELY)
LABLO(NODE) = OUTER, INNER, UNLABD   FOR CURRENT NODES
LABLO(NODE) = SIMPLE BLOSSOM        FOR NONCURRENT NODES

```

```

PSEUDO(NODE) - CURRENT NODE CONTAINING NODE (OR NODE)

```

```

APEX(NODE) - ORIGINAL NODE INSIDE NODE (OR NODE)

```

```

PRED(NODE) - PREDECESSOR ARC IN ALTERNATING PATH TO ROOT

```

```

TREE - LISTS OF NODES IN A TREE (HEADS = EXPOSED NODES)

```

```

BCI(NODE) - BLOSSOM CONSTRAINT IDENTIFIER (OR NODE)

```

```

MPRICE(NODE) - MAXIMUM PRICE OF BCI(NODE)

```

```

HEAD(NODE) = ROOT OF TREE
HEAD(ARC) = NODE POINTED BY THE ARC

```

```

PRICE(NODE) = NODE DUAL VARIABLE
PRICE(EDGE) = EDGE COST

```

```

NSCAN - CIRCULAR LIST OF UNSCANNED NODES

```

```

C
C   INSP - LIST OF NODE/EDGES TO BE INSPECTED
C
C   NORIG - LIST OF NODES IN A SIMPLE BLOSSOM (SEE ORIG)
C
C   LINK/MATE - LISTS OF EQUALITY/MATE EDGES
C           HEADS FOR CURRENT EQUALITY EDGES = LINK(NODE)
C           HEADS FOR NONCURRENT EQUALITY EDGES = LINK(PNODE)
C           HEADS FOR MATE EDGES = -MATE(NODE)
C           LINK(PT) = NEXT ELEMENT IN THE LIST
C           MATE(PT) = EDGE
C           PT <= HINK <= P2 = MAX ELEMENT
C
C   IMPLICIT INTEGER (A-Z)
C   REAL VALUE, DELTA, MCOST, TOLER, INF, PRICE, MPRICE
C   LOGICAL TRACE, PRINT
C   COMMON /NODES/ SET(100), LABLO(150), PSEUDO(150), APEX(150), PRED(
&           150), TREE(150), MPRICE(150), BCI(150)
C   COMMON /EDGES/ HEAD(6150), PRICE(3150), LIST(3150), INDEX(3150)
C   COMMON /LISTS/ NSCAN(101), INSP(3100), NORIG(100), LINK(1500), MAT
&           E(1500)
C   COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK
C   COMMON /PARMS/ N, NO, N1, N2, N9, L, L0, L1, L2, M, M0, M1, M2, M9, P0, P2
C   COMMON /VARBS/ VALUE, DELTA, MCOST, OPTION, TRACE, PRINT
C   COMMON /STATS/ ROOT$, NODE$, EDGES$, SHRK$, UNSH$, AUGM$, DUAL$, RMAT$
&           , LABL$, LINK$, TIME1, TIME2, TIME3
C   COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET
&           3, NULL
C   COMMON /PRNTS/ STACK(3000), PACK(100)
C
C   . . . . .
C
C   NODE VECTORS
C   DATA SET/100*0/, LABLO/150*0/, PSEUDO/150*0/, APEX/150*0/, PRED/15
&           0*0/, TREE/150*0/, MPRICE/150*0./, BCI/150*0/
C
C   NODE/EDGE VECTORS
C   DATA HEAD/6150*0/, PRICE/3150*0.0/
C
C   LIST OF EXISTING NODE/EDGES
C   DATA LIST/3150*0/, INDEX/3150*0/
C
C   LINK/MATE VECTORS
C   DATA LINK/1000*0/, MATE/1000*0/
C
C   HEAD/TAILS POINTERS
C   DATA HSCAN, TSCAN, HINSP, HORIG, HINK/5*0/
C
C   CONTROL VARIABLES
C   DATA VALUE, DELTA, MCOST, TRACE, PRINT/3*0.0, 2*.FALSE./
C
C   STATISTICAL VARIABLES
C   DATA ROOT$, NODE$, EDGES$, SHRK$, UNSH$, AUGM$, DUAL$, RMAT$, LABL$, LIN
&           K$/10*0/
C
C   CONSTANTS
C   DATA INF, TOLER, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET3, NULL/999
&           9.9, 1.E-50, 1., -1, 0, 0, 1, 2, 3, 0/
C
C   PROBLEM PARAMETERS

```



DATA N2, L2, M2, P2 / 100, 150, 3150, 1500 /  
END

SUBROUTINE GETNET

C  
C  
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C

MINIMUM COST 1-MATCHING/COVERING PROBLEM  
THE UNIVERSITY OF MICHIGAN  
WRITTEN BY CLOVIS PERIN  
AUGUST 1980

/GET NETWORK/  
READ NETWORK DATA  
INITIALIZE VARIABLES

IMPLICIT INTEGER (A-Z)  
REAL VALUE, DELTA, MCOST, TOLER, INF, PRICE, MPRICE  
LOGICAL TRACE, PRINT  
COMMON /NODES/ SET (100), LABLO (150), PSEUDO (150), APEX (150), PRED (150), TREE (150), MPRICE (150), BCI (150)  
COMMON /EDGES/ HEAD (6150), PRICE (3150), LIST (3150), INDEX (3150)  
COMMON /LISTS/ NSCAN (101), INSP (3100), NORIG (100), LINK (1500), MAT E (1500)  
COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK  
COMMON /PARMS/ N, NO, N1, N2, N9, L, LO, L1, L2, M, MO, M1, M2, M9, PO, P2  
COMMON /VARBS/ VALUE, DELTA, MCOST, OPTION, TRACE, PRINT  
COMMON /STATS/ ROOT\$, NODE\$, EDGES\$, SHRKS\$, UNSH\$, AUGM\$, DUAL\$, RMATS\$, LABL\$, LINK\$, TIME1, TIME2, TIME3  
COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET3, NULL  
COMMON /PRNTS/ STACK (3000), PACK (100)  
REAL MCOST, COST, TITLE (12)  
LOGICAL FREE (1) /'\*/

C  
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C  
C

READ PROBLEM PARAMETERS  
CALL TIME (0)  
READ (5, 1, END=999) TITLE  
1 FORMAT (12A4)  
READ (5, FREE, END=999) N, M, OPTION  
TRACE = .FALSE.  
PRINT = .FALSE.  
IF (OPTION.EQ.-N2) TRACE = .TRUE.  
IF (0.GT.OPTION.AND.OPTION.GE.-N2) PRINT = .TRUE.  
IF (OPTION.EQ.3.OR.OPTION.EQ.4) PRINT = .TRUE.  
WRITE (6, 2)  
2 FORMAT (1H1, 30X, '\*\*\*\*\* MINIMUM COST 1-MATCHING/COVERING PROBLEM  
EM \*\*\*\*\*'//)  
WRITE (6, 3) TITLE  
3 FORMAT (40X, 12A4)  
IF ( N.LE.0 .OR. M.LE.0 ) CALL ERROR (2)  
IF ( N.GT.N2 ) CALL ERROR (3)  
IF ( M.GT.M2-L2 ) CALL ERROR (4)

C  
C

INITIALIZE PARAMETRIC VARIABLES  
NO = 1  
N1 = N  
N9 = N + 1  
L = N / 2  
LO = N2 + 1  
L1 = N2  
MO = L2 + 1

```

M1 = L2 + M
M9 = 2 * M2 + 1
P0 = L2 + 1

C
C
10 INITIALIZE LINK/MATE STORAGE SPACE
DO 10 PT = P0, P2
LINK (PT) = PT + 1
LINK (P2) = NULL
HINK = M0

C
C
C READ NODE DATA
C INITIALIZE NODE VECTORS
DO 20 IND = N0, N1
READ (5, FREE, END=999) NODE, SET (NODE)
IF (TRACE. OR. OPTION. EQ. 4) WRITE (6, 4) NODE, SET (NODE)
4 FORMAT (' NODE DATA : ', 2I5)
IF (NODE. LT. N0. OR. NODE. GT. N2. OR. INDEX (NODE) . NE. NULL) CALL
& ERROR (6)
IF (SET (NODE) . LT. 0. OR. SET (NODE) . GT. 3) CALL ERROR (7)
INDEX (NODE) = IND
LIST (IND) = NODE
PSEUDO (NODE) = NODE
APEX (NODE) = NODE
LABLO (NODE) = UNLABD
ST = SET (NODE)
IF (ST. EQ. SET1. OR. ST. EQ. SET2) PRICE (NODE) = -INF/2
VALUE = VALUE + PRICE (NODE)
BCI (NODE) = NODE
IF (ST. EQ. SET1) BCI (NODE) = L2
IF (ST. EQ. SET1) MPRICE (NODE) = INF
20 CONTINUE
MPRICE (L2) = INF
H = 0

C
C
C READ EDGE DATA
C INITIALIZE EDGE VECTORS
C LINK EQUALITY EDGES
C MATCH EDGES IN A-
DO 40 IND = M0, M1
READ (5, FREE, END=999) EDGE, NODE1, NODE2, COST
IF (TRACE. OR. OPTION. EQ. 4) WRITE (6, 5) EDGE, NODE1, NODE2,
& COST
5 FORMAT (' EDGE DATA : ', 3I5, F7.1)
EDGE1 = EDGE + L2
IF (EDGE1. LT. M0. OR. EDGE1. GT. M2. OR. INDEX (EDGE1) . NE. NULL) C
& ALL ERROR (7)
IF (NODE1. LT. N0. OR. NODE1. GT. N2. OR. INDEX (NODE1) . EQ. NULL) C
& ALL ERROR (6)
IF (NODE2. LT. N0. OR. NODE2. GT. N2. OR. INDEX (NODE2) . EQ. NULL) C
& ALL ERROR (6)
IF (COST. GT. INF. OR. COST. LT. -INF) C
& ALL ERROR (9)
EDGE2 = M9 - EDGE1
HEAD (EDGE1) = NODE2
HEAD (EDGE2) = NODE1
PRICE (EDGE1) = COST
INDEX (EDGE1) = IND
LIST (IND) = EDGE1
COST = PRICE (EDGE1) - PRICE (NODE1) - PRICE (NODE2)
IF (COST. GT. +TOLER) GO TO 40

```

```

C      IF (COST.LT.-TOLER) GO TO 30
      "EDGE1" IS AN EQUALITY EDGE
      IF (HINK.EQ.NULL) CALL ERROR (5)
      PT          = HINK
      HINK        = LINK (HINK)
      MATE (PT)   = EDGE1
      LINK (PT)   = LINK (NODE1)
      LINK (NODE1) = PT
      IF (HINK.EQ.NULL) CALL ERROR (5)
      IF (HINK.GT.LINK$) LINK$ = HINK
      PT          = HINK
      HINK        = LINK (HINK)
      MATE (PT)   = EDGE2
      LINK (PT)   = LINK (NODE2)
      LINK (NODE2) = PT
      GO TO 40
C      "EDGE1" IS AN EGE OF A-
30     CALL MATCH (EDGE1)
      VALUE      = VALUE + PRICE (EDGE1)
      H          = H + 1
      STACK (H)  = EDGE1
40     CONTINUE
      IF (PRINT.AND.H.GT.1) WRITE (6,6) (STACK (I),HEAD (STACK (I)),HEAD (
&          M9-STACK (I)),PRICE (STACK (I)),I=1,H)
6     FORMAT (// ' EDGES OF A- (EDGE,NODE1,NODE2,COST) : ' /5 (I12,2I3,F6
&          .1))
C
C      INITIALIZE PSEUDONODE LIST
      DO 50 IND = L0, L2
50     LIST (IND) = IND
C
C      PLANT ALTERNATING FOREST
      DO 60 IND = N0, N1
      NODE          = LIST (IND)
      IF (SET (NODE).EQ.SET0.OR.MATE (NODE).NE.NULL) GO TO 60
      LABLO (NODE)  = OUTER
      HEAD (NODE)   = NODE
      HSCAN         = HSCAN + 1
      NSCAN (HSCAN) = NODE
      ROOT$         = ROOT$ + 1
60     CONTINUE
      LABL$ = ROOT$
      IF (PRINT) WRITE (6,7) ROOT$
7     FORMAT (// ' ALTERNATING FOREST PLANTED WITH',I4,' ROOTS')
C
      IF (OPTION.EQ.7) CALL DUMP
      CALL TIME (1,0,TIME1)
      IF (OPTION.NE.-ROOT$) RETURN
      TRACE = .TRUE.
      CALL DUMP
      RETURN
C
C      UNEXPECTED END-OF-FILE
999  CALL ERROR (1)
      STOP
      END

```

SUBROUTINE MINIM(NODE,NDELTA)

C  
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C

-----  
 MINIMUM COST 1-MATCHING/COVERING PROBLEM  
 THE UNIVERSITY OF MICHIGAN  
 WRITTEN BY CLOVIS PERIN  
 AUGUST 1980

/COMPUTE MINIMUM DELTA/  
 COMPARE DELTA AND NEWDELTA AND CONTROL INSPECTION LIST

IMPLICIT INTEGER (A-Z)  
 REAL VALUE,DELTA,MCOST,TOLER,INF,PRICE,MPRICE  
 LOGICAL TRACE,PRINT  
 COMMON /NODES/ SET (100),LABLO(150),PSEUDO(150),APEX(150),PRED(150),TREB(150),MPRICE(150),BCI(150)  
 COMMON /EDGES/ HEAD(6150),PRICE(3150),LIST(3150),INDEX(3150)  
 COMMON /LISTS/ NSCAN(101),INSP(3100),NORIG(100),LINK(1500),MATE(1500)  
 COMMON /PNTRS/ HSCAN,TSCAN,HINSP,HORIG,HINK  
 COMMON /PARMS/ N,NO,N1,N2,N9,L,LO,L1,L2,M,MO,M1,M2,M9,P0,P2  
 COMMON /VARBS/ VALUE,DELTA,MCOST,OPTION,TRACE,PRINT  
 COMMON /STATS/ ROOT\$,NODE\$,EDGES\$,SHRK\$,UNSH\$,AUGM\$,DUAL\$,RMTAT\$,LABL\$,LINK\$,TIME1,TIME2,TIME3  
 COMMON /CONTS/ TOLER,INF,OUTER,INNER,UNLABD,SET0,SET1,SET2,SET3,NULL  
 COMMON /PRNTS/ STACK(3000),PACK(100)  
 REAL NDELTA

C  
C  
C  
C  
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C  
C

.....  
 IF (TRACE) WRITE(6,1) NODE, NDELTA  
 1 FORMAT(' #TRACE# MINIM (' ,I4,' ,',F6.1,')')  
 NEWDELTA > DELTA -- RETURN  
 NEWDELTA = DELTA -- SAVE NODE/EDGE FOR INSPECTION  
 NEWDELTA < DELTA -- ERASE LIST, SAVE NODE/EDGE FOR INSPECTION  
 IF (NDELTA.GT.DELTA+TOLER) RETURN  
 IF (NDELTA.GE.DELTA-TOLER) GO TO 10  
 HINSP = 0  
 DELTA = NDELTA  
 10 HINSP = HINSP + 1  
 INSP(HINSP) = NODE  
 RETURN  
 END

SUBROUTINE MATCH(EDGE1)

```

C  -----
C
C  MINIMUM COST 1-MATCHING/COVERING PROBLEM
C  THE UNIVERSITY OF MICHIGAN
C  WRITTEN BY CLOVIS PERIN
C  AUGUST 1980
C
C  /MATCH NODES JOINED BY "EDGE1"/
C  FOR BOTH NODES DO:
C      IF "MATE(NODE) = NULL" -- "MATE(NODE) = EDGE1"
C      IF "NODE(NODE) > NULL" -- "-MATE(NODE)" BECOMES A HEAD
C      IF "NODE(NODE) < NULL" -- INTRODUCE "EDGE1" IN LIST
C
C  IMPLICIT INTEGER (A-Z)
C  REAL VALUE,DELTA,MCOST,TOLE,INF,PRICE,MPRICE
C  LOGICAL TRACE,PRINT
C  COMMON /NODES/ SET(100),LABLO(150),PSEUDO(150),APEX(150),PRED(
&      150),TREE(150),MPRICE(150),BCI(150)
C  COMMON /EDGES/ HEAD(6150),PRICE(3150),LIST(3150),INDEX(3150)
C  COMMON /LISTS/ NSCAN(101),INSP(3100),NORIG(100),LINK(1500),MAT
&      E(1500)
C  COMMON /PNTRS/ HSCAN,TSCAN,HINSP,HORIG,HINK
C  COMMON /PARMS/ N,N0,N1,N2,N9,L,L0,L1,L2,M,M0,M1,M2,M9,P0,P2
C  COMMON /VARBS/ VALUE,DELTA,MCOST,OPTION,TRACE,PRINT
C  COMMON /STATS/ ROOT$,NODE$,EDGE$,SHRKS$,UNSH$,AUGM$,DUAL$,RMAT$
&      ,LABL$,LINK$,TIME1,TIME2,TIME3
C  COMMON /CONTS/ TOLER,INF,OUTER,INNER,UNLABD,SET0,SET1,SET2,SET
&      3,NULL
C  COMMON /PRNTS/ STACK(3000),PACK(100)
C
C  . . . . .
C
C  IF (TRACE) WRITE(6,1) EDGE1
C  1 FORMAT(' #TRACE# MATCH (' ,I4,')')
C
C  DO 10 TO 40 FOR BOTH NODES JOINED BY "EDGE1"
C  EDGE = EDGE1
C  10 NODE = HEAD(M9-EDGE)
C  IF (MATE(NODE).NE.NULL) GO TO 20
C      NO MATES
C      MATE(NODE) = EDGE
C      GO TO 40
C  20 IF (MATE(NODE).LT.NULL) GO TO 30
C      ONE MATE -- "-MATE(NODE)" BECOMES A LIST HEAD
C      IF (HINK.EQ.NULL) CALL ERROR(5)
C      PT = HINK
C      HINK = LINK(HINK)
C      MATE(P) = MATE(NODE)
C      LINK(P) = NULL
C      MATE(NODE) = -PT
C  INTRODUCE "EDGE" IN THE MATE LIST WITH HEAD "-MATE(NODE)"
C  30 IF (HINK.EQ.NULL) CALL ERROR(5)
C  IF (HINK.GT.LINK$) LINK$ = HINK
C  PT = HINK
C  HINK = LINK(HINK)
C  MATE(P) = EDGE
C  LINK(P) = -MATE(NODE)
C  MATE(NODE) = -PT
C

```

```
C   GET SECOND NODE OR RETURN
40  IF (EDGE.NE.EDGE1) RETURN
    EDGE = M9 -EDGE1
    GO TO 10
C
    END
```

```

SUBROUTINE NMATCH(EDGE1)
C - - - - -
C
C MINIMUM COST 1-MATCHING/COVERING PROBLEM
C THE UNIVERSITY OF MICHIGAN
C WRITTEN BY CLOVIS PERIN
C AUGUST 1980
C
C /UNMATCH NODES JOINED BY "EDGE1"/
C FOR BOTH NODES DO:
C     IF EDGE IS NOT A MATE, RETURN
C     IF EDGE IS THE ONLY MATE, MAKE "MATE(NODE)" NULL
C     IF EDGE IS ONE OF THE MATES, REMOVE IT FROM THE LIST WITH
C         HEAD "-MATE(NODE)", AND IF SUCH A LIST CONTAINS NOW
C         JUST ONE EDGE, TRANSFORM "MATE(NODE)" INTO A POINTER
C
C IMPLICIT INTEGER (A-Z)
C REAL VALUE, DELTA, MCONST, TOLER, INF, PRICE, MPRICE
C LOGICAL TRACE, PRINT
C COMMON /NODES/ SET(100), LABLO(150), PSEUDO(150), APEX(150), PRED(
& 150), TREE(150), MPRICE(150), BCI(150)
C COMMON /EDGES/ HEAD(6150), PRICE(3150), LIST(3150), INDEX(3150)
C COMMON /LISTS/ NSCAN(101), INSP(3100), NORIG(100), LINK(1500), MAT
& E(1500)
C COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK
C COMMON /PARMS/ N, NO, N1, N2, N9, L, L0, L1, L2, M, M0, M1, M2, M9, P0, P2
C COMMON /VARBS/ VALUE, DELTA, MCONST, OPTION, TRACE, PRINT
C COMMON /STATS/ ROOT$, NODE$, EDGES$, SHRK$, UNSH$, AUGM$, DUAL$, RMAT$
& , LABL$, LINK$, TIME1, TIME2, TIME3
C COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET
& 3, NULL
C COMMON /PRNTS/ STACK(3000), PACK(100)
C
C . . . . .
C
C     IF (TRACE) WRITE(6,1) EDGE1
C     1 FORMAT(' #TRACE# NMATCH (' , I4, ') ')
C     IF (EDGE1.EQ.NULL) RETURN
C
C     DO 10 TO 50 FOR BOTH NODES JOINED BY "EDGE1"
C     EDGE = EDGE1
C     10 NODE = HEAD(M9-EDGE)
C     PT = MATE(NODE)
C     IF (PT.LT.NULL) GO TO 20
C     "MATE(NODE)" IS A POINTER
C     "MATE(NODE)" POINTS TO EDGE -- POINT IT TO NULL
C     IF (PT.EQ.EDGE) MATE(NODE) = NULL
C     GO TO 50
C
C     "-MATE(NODE)" IS A LIST HEAD
C     20 PT = -PT
C     IF (MATE(PT).NE.EDGE) GO TO 30
C     "EDGE" IS THE FIRST ELEMENT LIST
C     MATE(NODE) = -LINK(PT)
C     GO TO 40
C
C     "EDGE" IS NOT THE FIRST ELEMENT LIST
C     30 PT1 = PT
C     PT = LINK(PT1)
C     "EDGE" IS NOT IN THE LIST -- DONE
C     IF (PT.EQ.NULL) GO TO 50

```



```
C      IF (MATE (PT) .NE. EDGE) GO TO 30
      "EDGE" IS IN THE LIST
      LINK (PT1) = LINK (PT)
40     MATE (PT) = NULL
      LINK (PT) = HINK
      HINK      = PT

C
C      GET SECOND NODE OR RETURN
50     IF (EDGE .NE. EDGE1) RETURN
      EDGE = M9 - EDGE
      GO TO 10

C
      END
```

SUBROUTINE RMATCH(NODE)

```

C -----
C
C MINIMUM COST 1-MATCHING/COVERING PROBLEM
C THE UNIVERSITY OF MICHIGAN
C WRITTEN BY CLOVIS PERIN
C AUGUST 1980
C
C /REMATCH ALTERNATING PATH/
C "PSEUDO(NODE)" IS AN OUTER NODE
C REMATCH ITS ALTERNATING PATH
C ERASE ITS ALTERNATING TREE
C
C IMPLICIT INTEGER (A-Z)
C REAL VALUE, DELTA, MCONST, TOLER, INF, PRICE, MPRICE
C LOGICAL TRACE, PRINT
C COMMON /NODES/ SET(100), LABLO(150), PSEUDO(150), APEX(150), PRED(
& 150), TREE(150), MPRICE(150), BCI(150)
C COMMON /EDGES/ HEAD(6150), PRICE(3150), LIST(3150), INDEX(3150)
C COMMON /LISTS/ NSCAN(101), INSP(3100), NORIG(100), LINK(1500), MAT
& E(1500)
C COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK
C COMMON /PARMS/ N, NO, N1, N2, N9, L, L0, L1, L2, M, M0, M1, M2, M9, P0, P2
C COMMON /VARBS/ VALUE, DELTA, MCONST, OPTION, TRACE, PRINT
C COMMON /STATS/ ROOT$, NODE$, EDGES$, SHRKS$, UNSH$, AUGM$, DUAL$, RMAT$
& , LABL$, LINK$, TIME1, TIME2, TIME3
C COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET
& 3, NULL
C COMMON /PRNTS/ STACK(3000), PACK(100)
C
C . . . . .
C
C IF (TRACE) WRITE(6,1) NODE
1 FORMAT(' #TRACE# RMATCH (' ,I4, ' ) ')
AUGM$ = AUGM$ + 1
RMAT$ = RMAT$ + 1
C
C UPDATE "NODE" VECTORS
PNODE = PSEUDO(NODE)
IF (NODE.NE.PNODE) CALL STAPEX(NODE)
K = 1
STACK(K) = PNODE
C
C DO 10 FOR ALL OUTER NODES IN THE PATH
PNODE1 = PNODE
10 EDGE1 = PRED(PNODE1)
C "PRED(NODE1) = NULL" -- "PNODE1" IS A ROOT
IF (EDGE1.EQ.NULL) GO TO 100
C "PNODE2" IS INNER LABELLED
C "PNODE2" IS OUTER LABELLED
C UPDATE "NODE1" AND "NODE2" VECTORS
RMAT$ = RMAT$ + 2
EDGE2 = M9 - EDGE1
NODE1 = HEAD(EDGE1)
NODE2 = HEAD(EDGE2)
PNODE1 = PSEUDO(NODE1)
PNODE2 = PSEUDO(NODE2)
K = K + 4
STACK(K-3) = PNODE1
STACK(K-2) = MATE(NODE1)

```

```

        STACK(K-1) = PNODE2
        STACK(K)   = EDGE1
        MATE(NODE1) = EDGE2
        MATE(NODE2) = EDGE1
        IF (NODE1.NE.PNODE1) CALL STAPEX(NODE1)
        IF (NODE2.NE.PNODE2) CALL STAPEX(NODE2)
        GO TO 10

C
C   ERASE THE ALTERNATING TREE
C   DO 110 FOR EVERY NODE IN THE TREE
100 ROOT$ = ROOT$ - 1
    J     = 0
    NODE2 = HEAD(PNODE)
110 NODE1 = NODE2
    IF ( NODE1 .EQ. NULL ) GO TO 120
C     SAVE "NODE2"
C     UPDATE "NODE1" VECTORS
    NODE2      = TREE(NODE1)
    PNODE1     = PSEUDO(NODE1)
    TREE(NODE1) = NULL
    J          = J + 1
    PACK(J)    = NODE1
    IF ( PNODE1.NE.NODE1 ) GO TO 110
C   "NODE1" IS A CURRENT NODE
    LABEL      = LABELC(NODE1)
    LABLO(NODE1) = UNLABD
    HEAD(NODE1) = NULL
    PRED(NODE1) = NULL
    IF ( LABEL.NE.INNER ) GO TO 110
C   "NODE1" IS INNER LABELLED
C   STACK IT FOR SCANNING LATER
    NSCAN(TSCAN) = - NODE1
    TSCAN        = TSCAN - 1
    IF (TSCAN.LT.1) TSCAN = N9
    GO TO 110

C
C   PRINT PATH AND TREE
120 IF (PRINT) WRITE(6,2) (STACK(I),I=1,K)
    2 FORMAT(' REATCH ALTERN. PATH :',18I5/(24X,18I5))
    IF (PRINT) WRITE(6,3) (PACK(I),I=1,J)
    3 FORMAT(' ERASE ALTERNAT. TREE :',18I5/(24X,18I5))
    IF (OPTION.EQ.10) WRITE(6,4) ROOT$
    4 FORMAT(' #PRE-TRACE# ALTERNATING FOREST WITH',I4, '! ROOTS')
    IF (TRACE) GO TO 130
    IF (OPTION.NE.-ROOT$) RETURN
    TRACE = .TRUE.
    PRINT = .TRUE.
    RETURN
130 TRACE = .FALSE.
    PRINT = .TRUE.
    OPTION = 3
    RETURN
    END

```

SUBROUTINE SHRINK (EDGE)

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C

MINIMUM COST 1-MATCHING/COVERING PROBLEM  
THE UNIVERSITY OF MICHIGAN  
WRITTEN BY CLOVIS PERIN  
AUGUST 1980

/SHRINK SIMPLE BLOSSOM/  
"EDGE" JOINS TWO OUTER NODES IN THE SIMPLE BLOSSOM

IMPLICIT INTEGER (A-Z)  
REAL VALUE, DELTA, MFCOST, TOLER, INF, PRICE, MPRICE  
LOGICAL TRACE, PRINT  
COMMON /NODES/ SET (100), LABLO (150), PSEUDO (150), APEX (150), PRED (150), TREE (150), MPRICE (150), BCI (150)  
COMMON /EDGES/ HEAD (6150), PRICE (3150), LIST (3150), INDEX (3150)  
COMMON /LISTS/ NSCAN (101), INSP (3100), NORIG (100), LINK (1500), MAT E (1500)  
COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK  
COMMON /PARMS/ N, NO, N1, N2, N9, L, L0, L1, L2, M, M0, M1, M2, M9, P0, P2  
COMMON /VARBS/ VALUE, DELTA, MFCOST, OPTION, TRACE, PRINT  
COMMON /STATS/ ROOT\$, NODE\$, EDGE\$, SHRK\$, UNSH\$, AUGM\$, DUAL\$, RMAT\$, LABL\$, LINK\$, TIME1, TIME2, TIME3  
COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET3, NULL  
COMMON /PRNTS/ STACK (3000), PACK (100)  
REAL DD, DD1, DD2

C  
C  
C

IF (TRACE) WRITE (6, 1) EDGE  
1 FORMAT (' #TRACE# SHRINK (' , I4, ') ')  
SHRK\$ = SHRK\$ + 1  
LABL\$ = LABL\$ + 1

C  
C

GET FIRST ALTERNATING PATH  
H = 1  
PACK (H) = 0  
EDGE1 = EDGE  
10 NODE1 = PSEUDO (HEAD (EDGE1))  
H = H + 1  
PACK (H) = NODE1  
EDGE1 = PRED (NODE1)  
IF (EDGE1.NE.NULL) GO TO 10

C  
C

GET SECOND ALTERNATING PATH  
T = 50  
PACK (T) = -1  
EDGE2 = M9 - EDGE  
20 NODE2 = PSEUDO (HEAD (EDGE2))  
T = T - 1  
PACK (T) = NODE2  
EDGE2 = PRED (NODE2)  
IF (EDGE2.NE.NULL) GO TO 20

C  
C

ELIMINATE COMMON NODES IN THE PATHS  
30 H = H - 1  
T = T + 1  
IF (PACK (H) .EQ. PACK (T)) GO TO 30

```

C
C   CREATE "PNODE"
L1          = L1 + 1
PNODE       = LIST (L1)
INDEX (PNODE) = L1
BASE        = PACK (H+1)
LABLO (PNODE) = OUTER
PSEUDO (PNODE) = PNODE
PRICE (PNODE) = 0
LINK (PNODE) = NULL
MATE (PNODE) = BASE
APEX (PNODE) = APEX (BASE)
PRED (PNODE) = PRED (BASE)
ROOT        = HEAD (BASE)
HEAD (PNODE) = ROOT
TREE (PNODE) = TREE (ROOT)
TREE (ROOT) = PNODE

C
C   UPDATE NODE VECTORS OF THE FIRST PATH
PNODE1 = PACK (2)
40 IF (PNODE1.EQ.BASE) GO TO 50
    EDGE1      = MATE (APEX (PNODE1))
    PNODE2     = PSEUDO (HEAD (EDGE1))
    EDGE2     = PRED (PNODE1)
    PRED (PNODE1) = EDGE1
    HEAD (PNODE1) = PNODE2
    PNODE1     = PSEUDO (HEAD (EDGE2))
    PRED (PNODE2) = EDGE2
    HEAD (PNODE2) = PNODE1
    GO TO 40

C
C   UPDATE NODE VECTORS OF THE SECOND PATH
50 PNODE1      = PACK (49)
    EDGE2     = PRED (PNODE1)
    PRED (PNODE1) = EDGE
    HEAD (PNODE1) = PACK (2)
60 IF (PNODE1.EQ.BASE) GO TO 70
    EDGE1      = MATE (APEX (PNODE1))
    PNODE2     = PSEUDO (HEAD (EDGE1))
    HEAD (PNODE2) = PNODE1
    PRED (PNODE2) = M9 - EDGE1
    PNODE1     = PSEUDO (HEAD (EDGE2))
    HEAD (PNODE1) = PNODE2
    EDGE1      = PRED (PNODE1)
    PRED (PNODE1) = M9 - EDGE2
    EDGE2     = EDGE1
    GO TO 60

C
C   DETERMINE "BCI (PNODE) "
70 BCINOD = L2
    DD     = INF
    J      = 1
    STACK (J) = PNODE
    CALL ORIG (PNODE)
    DO 80 H = 1, HORIG
        NODE      = NORIG (H)
        LABLO (NODE) = PNODE
        J         = J + 1
        STACK (J) = NODE
        BNODE     = BCI (NODE)

```

```

      IF (SET (BNODE) .EQ. SET3 .AND. BNODE .EQ. NODE) MPRICE (BNODE) =
8         2 * PRICE (BNODE)
      DD1          = MPRICE (NODE) - PRICE (BNODE)
      IF (DD1 .GE. DD) GO TO 80
      DD          = DD1
      BCINOD     = BNODE
80      CONTINUE

C
C      PRINT SIMPLE BLOSSOM
      BCI (PNODE) = BCINOD
      MPRICE (PNODE) = MPRICE (BCINOD)
      IF (PRINT) WRITE (6, 2) (STACK (I), I=1, J).
2      FORMAT (' SHRINK SIMPLE BLOSSOM ', I6, ' WITH NODES ', I3, ' (BASE) ',
8         10I4 / (47X, 10I4))
C      SET PSEUDO OF ALL NODES INSIDE "PNODE"
      CALL STPSDO (PNODE)
      IF (DD .GT. 0) RETURN
C      TYPE 3 AUGMENTATION
      MATE (BCINOD) = NULL
      IF (PRINT) WRITE (6, 3) ROOT, PNODE
3      FORMAT (' TYPE 3 AUGMENTATION : ', I6, ' <R> ... ', I3, ' <+> ')
      CALL RMATCH (BCINOD)
      RETURN

END

```

SUBROUTINE NSHRNK (PNODE)

```

C -----
C
C MINIMUM COST 1-MATCHING/COVERING PROBLEM
C THE UNIVERSITY OF MICHIGAN
C WRITTEN BY CLOVIS PERIN
C AUGUST 1980
C
C /UNSHRINK "PNODE"/
C "PNODE" IS AN INNER LABELLED PSEUDONODE WITH NULL NODE PRICE
C "PNODE" "PNODE0" ARE MATED BY "EDGE0"
C "BLSM0" AND "BASE" ARE THE EXTREME OF THE ALTERNATING PATH
C TO BE CONSTRUCTED ON THE SIMPLE BLOSSOM OF "PNODE"
C
C IMPLICIT INTEGER (A-Z)
C REAL VALUE, DELTA, MCOST, TOLER, INF, PRICE, MPRICE
C LOGICAL TRACE, PRINT
C COMMON /NODES/ SET (100), LABLO (150), PSEUDO (150), APEX (150), PRED (
& 150), TREE (150), MPRICE (150), BCI (150)
C COMMON /EDGES/ HEAD (6150), PRICE (3150), LIST (3150), JINDEX (3150)
C COMMON /LISTS/ NSCAN (101), INSP (3100), NORIG (100), LINK (1500), MAT
& E (1500)
C COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK
C COMMON /PARMS/ N, NO, N1, N2, N9, L, L0, L1, L2, M, M0, M1, M2, M9, P0, P2
C COMMON /VARBS/ VALUE, DELTA, MCOST, OPTION, TRACE, PRINT
C COMMON /STATS/ ROOT$, NODE$, EDGES$, SHRK$, UNSH$, AUGM$, DUAL$, RMAT$
& E , LABL$, LINK$, TIME1, TIME2, TIME3
C COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET
& E 3, NULL
C COMMON /PRNTS/ STACK (3000), PACK (100)
C INTEGER LAB (3) /'<+>', '<->', '< >'/
C LOGICAL DONT
C
C . . . . .
C
C IF (TRACE) WRITE (6, 1) PNODE
C 1 FORMAT (' #TRACE# NSHRNK (' , I4, ') ')
C UNSH$ = UNSH$ + 1
C
C ELIMINATE "PNODE" FROM TREE
C NODE = HEAD (PNODE)
C 10 IF (TREE (NODE) .EQ. PNODE) GO TO 20
C NODE = TREE (NODE)
C GO TO 10
C 20 TREE (NODE) = TREE (PNODE)
C TREE (PNODE) = NULL
C DETERMINE "BASE", "PNODE0", AND "BLSM0"
C BASE = MATE (PNODE)
C PNODE0 = PSEUDO (HEAD (MATE (APEX (PNODE))))
C EDGE0 = PRED (PNODE0)
C BLSM0 = HEAD (M9-EDGE0)
C 30 IF (LABLO (BLSM0) .EQ. PNODE) GO TO 40
C BLSM0 = LABLO (BLSM0)
C GO TO 30
C
C DETERMINE NUMBER OF NODES BETWEEN "BASE" AND "BLSM0"
C 40 BLSM = BASE
C PT = 1
C 50 IF (BLSM .EQ. BLSM0) GO TO 60
C PT = PT + 1

```

```

        EDGE1 = PRED (BLSM)
        BLSM1 = BLSM
        BLSM = HEAD (BLSM)
        GO TO 50
C      EVEN NUMBER OF NODES -- RIGHT DIRECTION
60    IF (PT.EQ.PT/2*2.OR.PT.EQ.1) GO TO 80
C      ODD NUMBER -- REDIRECT THE SIMPLF BLOSSOM
70    EDGE2      = PRED (BLSM)
        BLSM2      = HEAD (BLSM)
        PRED (BLSM) = M9 - EDGE1
        HEAD (BLSM) = BLSM1
        EDGE1      = EDGE2
        BLSM1      = BLSM
        BLSM       = BLSM2
        IF (BLSM.NE.BLSM0) GO TO 70
C
C      UPDATE NODE VECTORS FOR NODLS IN THE ALTERNATING PATH
80    EDGE      = EDGE0
        BLSM1     = BLSM0
        ROOT     = HEAD (PNODE)
        BLSM2     = TREE (ROOT)
        TREE (ROOT) = BASE
        J        = 1
        STACK (J) = BLSM1
        PACK (J)  = 2
90    LABLO (BLSM1) = INNER
        LABL$      = LABL$ + 1
        CALL STPSDO (BLSM1)
        TREE (BLSM1) = BLSM2
        BLSM2       = HEAD (BLSM1)
        HEAD (BLSM1) = ROOT
        EDGE1       = PRED (BLSM1)
        PRED (BLSM1) = NULL
        IF (BLSM1.EQ.BASE) GO TO 100
        LABLO (BLSM2) = OUTER
        LABL$      = LABL$ + 1
        CALL STPSDO (BLSM2)
        EDGE2       = M9 - EDGE1
        NODE1       = HEAD (EDGE2)
        NODE2       = HEAD (EDGE1)
        MATE (NODE1) = EDGE1
        MATE (NODE2) = EDGE2
        CALL STAPEX (NODE1)
        CALL STAPEX (NODE2)
        TREE (BLSM2) = BLSM1
        BLSM1       = HEAD (BLSM2)
        HEAD (BLSM2) = RCOT
        EDGE2       = PRED (BLSM2)
        PRED (BLSM2) = EDGE
        EDGE        = M9 - EDGE2
        HSCAN       = HSCAN + 1
        IF (HSCAN.GT.N9) HSCAN = 1
        NSCAN (HSCAN) = BLSM2
        J           = J + 2
        STACK (J-1) = BLSM2
        STACK (J)   = BLSM1
        PACK (J-1)  = 1
        PACK (J)    = 2
        GO TO 90
C

```



C UPDATE NODE VECTORS FOR NODE NOT IN THE ALTERNATING PATH

```
100 PRED(PNODE0) = EDGE
    DONT          = .FALSE.
110 IF (BLSM2.EQ.BLSM0) GO TO 120
    BLSM1          = HEAD(BLSM2)
    LABLO(BLSM2) = UNLABD
    LABLO(BLSM1) = UNLABD
    CALL STPSDO(BLSM2)
    CALL STPSDO(BLSM1)
    EDGE1          = PRED(BLSM2)
    EDGE2          = N9 - EDGE1
    NODE2          = HEAD(EDGE1)
    NODE1          = HEAD(EDGE2)
    MATE(NODE2)   = EDGE2
    MATE(NODE1)   = EDGE1
    CALL STAPEX(NODE1)
    CALL STAPEX(NODE2)
    PRED(BLSM2)   = NULL
    PRED(BLSM1)   = NULL
    NSCAN(TSCAN) = -BISM2
    TSCAN         = TSCAN - 1
    IF (TSCAN.LT.1) TSCAN = N9
    NSCAN(TSCAN) = -BLSM1
    TSCAN         = TSCAN - 1
    IF (TSCAN.LT.1) TSCAN = N9
    J             = J + 2
    STACK(J-1)   = BLSM2
    PACK(J-1)    = 3
    STACK(J)     = BLSM1
    PACK(J)      = 3
    HEAD(BLSM2)  = NULL
    BLSM2        = HEAD(BLSM1)
    HEAD(BLSM1)  = NULL
    GO TO 110
```

C

C UPDATE THE EDGE VECTORS FOR EDGES LINKED TO PNODE

```
120 PT1          = LINK(PNODE)
    LINK(PNODE) = NULL
130 PT = PT1
    IF (PT.EQ.NULL) GO TO 150
        PT1          = LINK(PT)
        EDGE = MATE(PT)
        IF (EDGE.EQ.NULL) GO TO 140
        NODE          = HEAD(EDGE)
        LINK(PT)      = LINK(NODE)
        LINK(NODE)    = PT
        GO TO 130
140 LINK(PT) = HINK
    HINK      = PT
    GO TO 130
```

C

C UPDATE PSEUDONODE LIST

```
150 NODE          = LIST(L1)
    IND           = INDEX(PNODE)
    LIST(L1)      = PNODE
    LIST(IND)     = NODE
    INDEX(NODE)   = IND
    INDEX(PNODE) = NULL
    L1            = L1 - 1
    IF (PRINT) WRITE(6,2) PNODE, (STACK(I),LAB(PACK(I)))
```

```
2 FORMAT (' UNSHRINK PSEUDONODE : ',I4,' WITH NODES',10(I5,A3)/  
&      39X,10(I5,A3))  
RETURN  
END
```

SUBROUTINE STPSDO(PNODE)

```

C -----
C
C MINIMUM COST 1-MATCHING/COVERING PROBLEM
C THE UNIVERSITY OF MICHIGAN
C WRITTEN BY CLOVIS PERIN
C AUGUST 1980
C
C /SET PSEUDO FOR NODES INSIDE "PNODE"/
C
C IMPLICIT INTEGER (A-Z)
C REAL VALUE, DELTA, MCOST, TOLER, INF, PRICE, MPRICE
C LOGICAL TRACE, PRINT
C COMMON /NODES/ SET (100), LABLO (150), PSEUDO (150), APEX (150), PRED(
& 150), TREE (150), MPRICE (150), BCI (150)
C COMMON /EDGES/ HEAD (6150), PRICE (3150), LIST (3150), INDEX (3150)
C COMMON /LISTS/ NSCAN (101), INSP (3100), NORIG (100), LINK (1500), MAT
& E (1500)
C COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK
C COMMON /PARMS/ N, NO, N1, N2, N9, L, L0, L1, L2, M, M0, M1, M2, M9, P0, P2
C COMMON /VARBS/ VALUE, DELTA, MCOST, OPTION, TRACE, PRINT
C COMMON /STATS/ ROOT$, NODE$, EDGES$, SHRS$, UNSH$, AUGM$, DUAL$, RMT$
& , LABL$, LINK$, TIME1, TIME2, TIME3
C COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET
& 3, NULL
C COMMON /PRNTS/ STACK (3000), PACK (100)
C
C . . . . .
C
C IF (TRACE) WRITE(6,1) PNODE
1 FORMAT(' #TRACE# STPSDO (' ,I4,') ')
C PSEUDO(PNODE) = PNODE
C "PNODE" IS AN ORIGINAL NODE -- RETURN
C IF (PNODE.LE.N2) RETURN
C
C SET PSEUDO FOR ALL ORIGINAL NODES INSIDE "PNODE"
C CALL ORIG (PNODE)
10 IF (HORIG.EQ.0) RETURN
C NODE = NORIG (HORIG)
C HORIG = HORIG - 1
C PSEUDO (NODE) = PNODE
C IF (NODE.GT.N2) CALL ORIG (NODE)
C GO TO 10
C
C END

```

SUBROUTINE STAPEX (NODE)

C  
C  
C  
C  
C  
C  
C  
C  
C  
C  
C  
C  
C  
C

MINIMUM COST 1-MATCHING/COVERING PROBLEM  
THE UNIVERSITY OF MICHIGAN  
WRITTEN BY CLOVIS PERIN  
AUGUST 1980

/SET "NODE" AS APEX/  
"NODE" IS AN ORIGINAL NODE  
SET "NODE" AS APEX OF ALL PSEUDONODES  
SET "BASE" AS BASE OF ALL PSEUDONODES

IMPLICIT INTEGER (A-Z)  
REAL VALUE, DELTA, MCOST, TOLER INF, PRICE, MPRICE  
LOGICAL TRACE, PRINT  
COMMON /NODES/ SET (100), LABLO (150), PSEUDO (150), APEX (150), PRED (150), TREE (150), MPRICE (150), BCI (150)  
COMMON /EDGES/ HEAD (6150), PRICE (3150), LIST (3150), INDEX (3150)  
COMMON /LISTS/ NSCAN (101), INSP (3100), NORIG (100), LINK (1500), MAT E (1500)  
COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK  
COMMON /PARMS/ N, NO, N1, N2, N9, L, LO, L1, L2, M, MO, M1, M2, M9, PO, P2  
COMMON /VARBS/ VALUE, DELTA, MCOST, OPTION, TRACE, PRINT  
COMMON /STATS/ ROOT\$, NODE\$, EDGE\$, SHRKS\$, UNSH\$, AUGMS\$, DUAL\$, RMAT\$ E  
, LABL\$, LINK\$, TIME1, TIME2, TIME3  
COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET E  
3, NULL  
COMMON /PRNTS/ STACK (3000), PACK (100)

C  
C  
C

IF (TRACE) WRITE (6, 1) NODE  
1 FORMAT (' #TRACE# STAPEX (' , I4, ') ' )  
  
"NODE" IS CONSTANT FOR ALL PSEUDOS  
"BASE" IS VARIABLE  
"BLSM" IS THE SIMPLE BLOSSOM WITH BASE "BASE"  
BASE = NODE  
10 BLSM = LABLO (BASE)  
IF (BLSM.LE.N2) RETURN  
MATE (BLSM) = BASE  
APEX (BLSM) = NODE  
BASE = BLSM  
GO TO 10  
  
END



```

IF ( LABEL2 .EQ. UNLABD ) GO TO 20
C "NODE1", "NODE2" ARE BOTH OUTER NODES
IF ( ROOT1 .NE. ROOT2 ) GO TO 10
C "NODE1", "NODE2" ARE IN THE SAME TREE -- SHRINK
CALL SHRINK(EDGE1)
RETURN
C "NODE1", "NODE2" ARE IN DIFFERENT TREES -- AUGMENT
10 MATE(NODE1) = EDGE1
MATE(NODE2) = EDGE2
IF (PRINT) WRITE(6,2) ROOT1,PNODE1,PNODE2,ROOT2
2 FORMAT(' DOUBLE AUGMENTATION :',I6,'<R> ...',I3,'<+>',I4,
& '<+> ...',I3,'<R>')
CALL RMATCH(NODE1)
CALL RMATCH(NODE2)
RETURN

C
C "NODE2" IS UNLABELLED
20 IF ( PRICE(PNODE2).GT.TOLER .OR. PNODE2 .GT.N1 ) GO TO 30
IF ( SET(PNODE2).NE.SET0 .AND. SET(PNODE2).NE.SET3 ) GO TO 30
C "NODE2" CAN BE A TYPE 2 NODE-- AUGMENT
MATE(NODE1) = NULL
CALL MATCH(EDGE1)
IF (PRINT) WRITE(6,3) ROOT1,PNODE1,PNODE2
3 FORMAT(' TYPE 2 AUGMENTATION :',I6,'<R> ...',I3,'<+>',I4,
& '< >')
CALL RMATCH(NODE1)
RETURN

C
30 IF (MATE(APEX2).NE.NULL) GO TO 40
C "NODE2" IS A TYPE 0 NODE -- AUGMENT
MATE(NODE1) = EDGE1
MATE(NODE2) = EDGE2
IF (PRINT) WRITE(6,4) ROOT1,PNODE1,PNODE2
4 FORMAT(' TYPE 0 AUGMENTATION :',I6,'<R> ...',I3,'<+>',I4,
& '< >')
IF (NODE2.NE.PNODE2) CALL STAPEX(NODE2)
CALL RMATCH(NODE1)
RETURN

C
40 EDGE3 = MATE(APEX2)
APEX3 = HEAD(EDGE3)
IF ( MATE(APEX3).GE.NULL ) GO TO 50
C "NODE2" IS A TYPE 1 NODE -- AUGMENT
CALL NMATCH(EDGE3)
MATE(NODE1) = EDGE1
MATE(NODE2) = EDGE2
IF (PRINT) WRITE(6,5) ROOT1,PNODE1,PNODE2
5 FORMAT(' TYPE 1 AUGMENTATION :',I6,'<R> ...',I3,'<+>',I4,
& '< >')
IF (NODE2.NE.PNODE2) CALL STAPEX(NODE2)
CALL RMATCH(NODE1)
RETURN

C
C LABEL "NODE2" INNER
C LABEL "NODE3" OUTER
50 PNODE3 = PSEUDO(APEX3)
LABLO(PNODE2) = INNER
HEAD(PNODE2) = ROOT1
LABLO(PNODE3) = OUTER
TREE(PNODE2) = PNODE3

```

```

TREE (PNODE3) = TREE (ROOT1)
TREE (ROOT1) = PNODE2
PRED (PNODE3) = EDGE2
HEAD (PNODE3) = ROOT1
LABL$ = LABL$ + 2
IF (PRINT) WRITE (6,6) ROOT1,PNODE1,PNODE2, PNODE3
6 FORMAT(' NODE LABELING :      ',I6,'<R> ...',I3,'<+>',I4,'<->'
8      ,I4,'<+>')
C
NODE3 = BCI(PNODE3)
IF ( PRICE (NODE3).NE.MPRICE(PNODE3) ) GO TO 60
C      "NODE3" CAN BECOME A TYPE 0 NODE -- AUGMENT
MATE (NODE3) = NULL
IF (PRINT) WRITE (6,7) ROOT1,PNODE3
7      FORMAT(' TYPE 3 AUGMENTATION :',I6,'<R> ...',I3,'<+>')
CALL RMATCH (NODE3)
RETURN
C
60 IF ( PRICE (PNODE2).NE.0.0 .OR. PNODE2.LE.N ) GO TO 70
C      "PNODE2" IS A INNER PSEUDO WITH NULL NODE PRICE --
C      -- UNSHRINK IT
CALL NSHRNK (PNODE2)
RETURN
C
C      PUSH "PNODE3" INTO THE LIST OF UNSCANNED NODES
70 HSCAN = HSCAN + 1
IF (HSCAN.GT.N9) HSCAN = 1
NSCAN (HSCAN) = PNODE3
RETURN
END

```

SUBROUTINE GRWFOR

```

C -----
C
C MINIMUM COST 1-MATCHING/COVERING PROBLEM
C THE UNIVERSITY OF MICHIGAN
C WRITTEN BY CLOVIS PERIN
C AUGUST 1980
C
C
C /GROW ALTERNATING FOREST/
C SCAN UNSCANNED NODES
C PERFORM DUAL SOLUTION CHANGE
C CHECK TERMINATION OR INFEASIBILITY
C
C IMPLICIT INTEGER (A-Z)
C REAL VALUE, DELTA, MCOST, TOLER, INF, PRICE, MPRICE
C LOGICAL TRACE, PRINT
C COMMON /NODES/ SET (100), LABLO (150), PSEUDO (150), APEX (150), PRED(
& 150), TREE (150), MPRICE (150), BCI (150)
C COMMON /EDGES/ HEAD (6150), PRICE (3150), LIST (3150), INDEX (3150)
C COMMON /LISTS/ NSCAN (101), INSP (3100), NORIG (100), LINK (1500), MAT
& E (1500)
C COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK
C COMMON /PARMS/ N, NO, N1, N2, N9, L, LO, L1, L2, M, MO, M1, M2, M9, PO, P2
C COMMON /VARBS/ VALUE, DELTA, MCOST, OPTION, TRACE, PRINT
C COMMON /STATS/ ROOT$, NODE$, EDGES$, SHRK$, UNSH$, AUGM$, DUAL$, RMAT$
& , LABL$, LINK$, TIME1, TIME2, TIME3
C COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET
& 3, NULL
C COMMON /PRNTS/ STACK (3000), PACK (100)
C REAL COST
C
C . . . . .
C
C 10 IF (TRACE) WRITE (6, 1)
C 1 FORMAT (' #TRACE# GRWFOR')
C IF (OPTION.EQ.9) CALL DUMP
C
C DO 20 TO 80 FOR EVERY NODE IN "NSCAN"
C 20 IF (HSCAN.EQ.TSCAN) GO TO 100
C TSCAN = TSCAN + 1
C IF (TSCAN.GT.N9) TSCAN = 1
C NODE = NSCAN (TSCAN)
C LABEL = OUTER
C IF (NODE.GT.0) GO TO 30
C "NODE > 0" -- OUTER NODE
C "NODE < 0" -- EX-INNER NODE
C LABEL = UNLABD
C NODE = -NODE
C CHECK WHETHER "NODE" IS A DESTROYED PSEUDO
C 30 IF (INDEX(NODE).EQ.NULL) GO TO 20
C PNODE = PSEUDO (NODE)
C CHECK WHETHER LABEL IS STILL OUTER OR EX-INNER (UNLABELLED)
C IF (LABLO (PNODE).NE.LABEL) GO TO 20
C IF (PRINT) WRITE (6, 7) NODE
C 7 FORMAT (' SCAN NODE : ', I4)
C NODE$ = NODE$ + 1
C P" = NODE
C BIFURCATE NODE : PSEUDO
C IF (NODE.LE.N2) GO TO 50

```



```

C
C "NODE" IS A PSEUDONODE
C PUSH ITS SIMPLE BLOSSOM INTO "NSCAN"
C LABEL = 2 * LABEL - 1
C BASE = MATE(NODE)
C BLSM = BASE
40 NSCAN(TSCAN) = LABEL * BLSM
TSCAN = TSCAN - 1
IF (TSCAN.LT.1) TSCAN = N9
BLSM = HEAD(BLSM)
IF (BLSM.NE.BASE) GO TO 40
GO TO 20

C
C "NODE" IS AN ORIGINAL NODE
C SCAN ALL EQUALITY EDGES INCIDENT WITH "NODE"
50 PT1 = PT
60 PT = LINK(PT1)
IF (PT.EQ.NULL) GO TO 20
EDGE1 = MATE(PT)
C CHECK FOR DESTROYED EDGE
IF (EDGE1.EQ.NULL) GO TO 70
C DETERMINE DIRECTION OF ARC
IF (LABEL.EQ.UNLABD) EDGE1 = N9 - EDGE1
EDGE2 = N9 - EDGE1
EDGE = EDGE1
IF (EDGE.GT.EDGE2) EDGE = EDGE2
NODE1 = HEAD(EDGE2)
NODE2 = HEAD(EDGE1)
PNODE1 = PSEUDO(NODE1)
PNODE2 = PSEUDO(NODE2)
COST = PRICE(EDGE) - PRICE(NODE1) - PRICE(NODE2)
C CHECK OUTER >-- EDGE --> UNLABELLED
IF (LABLO(PNODE1).NE.CUTER) GO TO 50
IF (LABLO(PNODE2).EQ.INNER) GO TO 50
C CHECK FOR EQUALITY EDGE
IF (COST.GT.TOLER) GO TO 70
C CHECK FOR CURRENT EDGE
IF (PNODE1.EQ.PNODE2) GO TO 80
CALL SCAN(EDGE1)
C CHECK AUGMENTATION
IF (LABLO(PSEUDO(NODE1)).NE.OUTER.AND.LABEL.EQ.OUTER) GO
& TO 20
GO TO 50

C
C RETURN "LINK/MATE(PT)" TO HINK
70 LINK(PT1) = LINK(PT)
MATE(PT) = NULL
LINK(PT) = HINK
HINK = PT
GO TO 60

C
C TRANSFER "LINK/MATE(PT)" TO "PNODE1"
80 LINK(PT1) = LINK(PT)
LINK(PT) = LINK(PNODE1)
LINK(PNODE1) = PT
GO TO 60

C
C NO MORE AUGMENTATIONS, SHRINKINGS, UNSHRINKINGS, OR LABELLINGS
C START DUAL SOLUTION CHANGE
100 IF (TRACE) CALL DUMP

```

```

IF (OPTION.EQ.9) CALL DUMP
C CHECK TERMINATION
IF (ROOT$.EQ.0) RETURN
C INITIALIZE CONTROL VARIABLES
DELTA = INF
HINSP = 0

C
C CHECK ORIGINAL NODES
DO 110 IND = NO, N1
    NODE = LIST(IND)
    LABEL = LABLO(NODE)
    PNODE = PSEUDO(NODE)
    IF ( PNODE.NE.NODE ) GO TO 110
C ONLY CURRENT ORIGINAL NODES
C CHECK MAXIMUM INCREASE
    IF ( SET(NODE).EQ.SET2 .AND. LABEL.EQ.OUTER ) CALL MINIM
    & ( NODE, -PRICE(NODE) )
C CHECK MAXIMUM DECREASE
    IF ( SET(NODE).EQ.SET3 .AND. LABEL.EQ.INNER ) CALL MINIM
    & ( NODE, PRICE(NODE) )
110 CONTINUE

C
C CHECK PSEUDONODES
IF (L0.GT.L1) GO TO 130
DO 120 IND = L0, L1
    NODE = LIST(IND)
    LABEL = LABLO(NODE)
C CHECK MAXIMUM DECREASE
    IF ( LABEL.EQ.INNER ) CALL MINIM(NODE,PRICE(NODE))
    BNODE = BCI(NODE)
C CHECK MAXIMUM INCREASE
    IF ( LABEL.EQ.OUTER .AND. BNODE.LT.N2 ) CALL MINIM( NODE,
    & MPRICE(BNODE)-PRICE(BNODE) )
120 CONTINUE

C
C CHECK EDGES
130 DO 140 IND = M0, M1
    EDGE = LIST(IND)
    NODE1 = HEAD(EDGE)
    NODE2 = HEAD(M9-EDGE)
    PNODE1 = PSEUDO(NODE1)
    PNODE2 = PSEUDO(NODE2)
    IF (PNODE1.EQ.PNODE2) GO TO 140
C ONLY CURRENT EDGES
    LABEL1 = LABLO(PNODE1)
    LABEL2 = LABLO(PNODE2)
    IF ( LABEL1.NE.OUTER .AND. LABEL2.NE.OUTER ) GO TO 140
C "AT LEAST ONE OF "LABEL1" OR "LABEL2" IS OUTER
C CHECK MAXIMUM INCREASE
    COST = PRICE(EDGE) - PRICE(NODE1) - PRICE(NODE2)
    IF ( LABEL1.EQ.OUTER .AND. LABEL2.EQ.UNLABD )
    1 CALL MINIM(EDGE,COST)
    IF ( LABEL1.EQ.UNLABD .AND. LABEL2.EQ.OUTER )
    1 CALL MINIM(EDGE,COST)
    IF ( LABEL1.EQ.OUTER .AND. LABEL2.EQ.OUTER )
    1 CALL MINIM(EDGE,COST/2)
140 CONTINUE

C
C NO NODE/EDGE TO INSPECT -- RETURN
IF (HINSP.EQ.0) RETURN

```

```

C
C   UPDATE DUAL SOLUTION
   DUAL$ = DUAL$ + 1
   IF (PRINT) WRITE(6,2) DELTA
2   FORMAT(' DUAL SOLUTION CHANGE: ',F6.1)
C
C   UPDATE ORIGINAL NODE PRICES
   DO 210 IND = N0, N1
210   NODE          = LIST(IND)
      PRICE(NODE) = PRICE(NODE) + DELTA*LABLO(PSEUDO(NODE))
C
C   UPDATE PSEUDONODE PRICES
   IF (L0.GT.L1) GO TO 300
   DO 220 IND = L0, L1
      NODE = LIST(IND)
      IF (NODE.EQ.PSEUDO(NODE))
1         PRICE(NODE) = PRICE(NODE) + DELTA*LABLO(NODE)
220   CONTINUE
C
C   UPDATE DUAL OBJECTIVE VALUE
300  VALUE = VALUE + ROOT$*DELTA
C
C   START INSPECTING EVERY NODE/EDGE IN "INSP"
   DO 340 PT = 1, HINSP
      NODE = INSP(PT)
C      BIFURCATE NODE : EDGE
      IF (NODE.GE.M0) GO TO 330
      IF (PRINT) WRITE(6,3) NODE
3      FORMAT(' INSPECT NODE :           ',I6)
      LABEL = LABLO(NODE)
      IF (LABEL.EQ.INNER) GO TO 310
      IF (LABEL.NE.OUTER) GO TO 340
C      ONLY OUTER NODES -- AUGMENT
      BNODE          = BCI(NODE)
      MATE(BNODE)    = NULL
      IF (PRINT) WRITE(6,4) HEAD(NODE),NODE
4      FORMAT(' TYPE 3 AUGMENTATION : ',I6,'<R> ...',I3,'<+>
      ' )
      CALL RMATCH(BNODE)
      GO TO 340
C      ONLY INNER NODES
C      BIFURCATE NODE : PSEUDO
310   IF (NODE.GT.N2) GO TO 320
C      ONLY INNER ORIGINAL NODES -- AUGMENT
      EDGE          = PRED(HEAD(MATE(NODE)))
      CALL MATCH(EDGE)
      APEX1         = HEAD(EDGE)
      PNODE         = PSEUDO(APEX1)
      IF (PRINT) WRITE(6,5) HEAD(NODE),PNODE,NODE
5      FORMAT(' TYPE 2 AUGMENTATION : ',I6,'<R> ...',I3,'<+>
      ',I4,'< >')
      CALL RMATCH(APEX1)
      GO TO 340
C      ONLY INNER PSEUDONODES -- UNSHRINK
320   CALL NSHRNK(NODE)
      GO TO 340
C
C      ONLY EDGES
C      EQUALITY EDGE
330   EDGE1        = NODE

```

```

EDGE2      = M9 - EDGE1
NODE1      = HEAD (EDGE2)
NODE2      = HEAD (EDGE1)
IF (HINK.EQ.NULL) CALL ERROR(5)
PTR        = HINK
HINK       = LINK (HINK)
MATE (PTR) = EDGE1
LINK (PTR) = LINK (NODE1)
LINK (NODE1) = PTR
IF (HINK.EQ.NULL) CALL ERROR(5)
IF (HINK.GT.LINK$) LINK$ = HINK
PTR        = HINK
HINK       = LINK (HINK)
MATE (PTR) = EDGE2
LINK (PTR) = LINK (NODE2)
LINK (NODE2) = PTR
LABEL1     = LABLO (PSEUDO (NODE1) )
LABEL2     = LABLO (PSEUDO (NODE2) )
EDGE = EDGE1 - L2
IF (PRINT) WRITE (6,6) EDGE, NODE1, NODE2
6          FORMAT (' INSPECT EDGE :           ',I5,' = (' ,I3,' ; ' ,I3,' ) ')
C          CHECK   OUTER >-- EDGE --> OUTER
C          OUTER >-- EDGE --> UNLABD
C          UNLABD >-- EDGE --> OUTER
IF (LABEL1.EQ.INNER.OR. LABEL2.EQ.INNER) GO TO 340
IF (LABEL1.NE.OUTER.AND.LABEL2.NE.OUTER) GO TO 340
IF (LABEL1.EQ.OUTER) CALL SCAN (EDGE1)
IF (LABEL1.NE.OUTER) CALL SCAN (EDGE2)
340       CONTINUE
C
C CHECK TERMINATION
IF (TRACE) CALL DUMP
IF (OPTION.EQ.8) CALL DUMP
IF (ROOT$.GT.0) GO TO 10
RETURN
END

```

SUBROUTINE ORIG(PNODE)

```

C -----
C
C MINIMUM COST 1-MATCHING/COVERING PROBLEM
C THE UNIVERSITY OF MICHIGAN
C WRITTEN BY CLOVIS PERIN
C AUGUST 1980
C
C /GET SIMPLE BLOSSOM, EVENTUALLY ALL ORIGINAL NODES/
C PUSH INTO NORIG ALL NODES IN THE SIMPLE BLOSSOM OF "PNODE"
C IMPLICIT INTEGER (A-Z)
C RFAL VALUE, DELTA, MCONST, TOLER, INF, PRICE, MPRICE
C LOGICAL TRACE, PRINT
C COMMON /NODES/ SET (100), LABLO (150), PSEUDO (150), APEX (150), PRED (
& 150), TREE (150), MPRICE (150), BCI (150)
C COMMON /EDGES/ HEAD (6150), PRICE (3150), LIST (3150), INDEX (3150)
C COMMON /LISTS/ NSCAN (101), INSP (3100), NORIG (100), LINK (1500), MAT
& E (1500)
C COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK
C COMMON /PARMS/ N, NO, N1, N2, N9, L, LO, L1, L2, M, MO, M1, M2, M9, P0, P2
C COMMON /VARBS/ VALUE, DELTA, MCONST, OPTION, TRACE, PRINT
C COMMON /STATS/ ROOTS$, NODES$, EDGES$, SHRK$, UNSH$, AUGMS$, DUAL$, RMATS$
& , LABL$, LINK$, TIME1, TIME2, TIME3
C COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET
& 3, NULL
C COMMON /PRNTS/ STACK (3000), PACK (100)
C
C . . . . .
C
C IF (TRACE) WRITE (6,1) PNODE
C 1 FORMAT (' #TRACE# ORIG (' ,I4, ') ')
C
C "PNODE" IS ORIGINAL NODE? -- RETURN
C IF (PNODE.LE.N2) RETURN
C
C PUSH SIMPLE BLOSSOM INTO NORIG
C BASE = MATE (PNODE)
C NODE = BASE
C 10 HORIG = HORIG + 1
C NORIG (HORIG) = NODE
C NODE = HEAD (NODE)
C IF (NODE.NE.BASE) GO TO 10
C
C RETURN
C END

```

SUBROUTINE PUTSOL

```

C -----
C
C   MINIMUM COST 1-MATCHING/COVERING PROBLEM
C   THE UNIVERSITY OF MICHIGAN
C   WRITTEN BY CLOVIS PERIN
C   AUGUST 1980
C
C   /PUT SOLUTION/
C   TRANSFORM NODE PRICES
C   TRANSFORM TYPE E CONFIGURATION
C
C   IMPLICIT INTEGER (A-Z)
C   REAL VALUE, DELTA, MCONST, TOLER, INF, PRICE, MPRICE
C   LOGICAL TRACE, PRINT
C   COMMON /NODES/ SET (100), LABI (150), PSEUDO (150), APEX (150), PRED (
&      150), TREE (150), MPRICE (150), BCI (150)
C   COMMON /EDGES/ HEAD (6150), PRICE (3150), LIST (3150), INDEX (3150)
C   COMMON /LISTS/ NSCAN (101), INSP (3100), NORIG (100), LINK (1500), MAT
&      E (1500)
C   COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK
C   COMMON /PARMS/ N, NO, N1, N2, N9, L, L0, L1, L2, M, M0, M1, M2, M9, P0, P2
C   COMMON /VARBS/ VALUE, DELTA, MCONST, OPTION, TRACE, PRINT
C   COMMON /STATS/ ROOT$, NODE$, EDGES$, SHRKS$, UNSH$, AUGM$, DUAL$, RMAT$
&      , LABL$, LINK$, TIME1, TIME2, TIME3
C   COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET
&      3, NULL
C   COMMON /PRNTS/ STACK (3000), PACK (100)
C   INTEGER EDLIST (3000)
C   REAL SIGMA, MPRICE (100)
C
C   . . . . .
C
C   CALL TIME (1, 0, TIME2)
C   IF (OPTION.EQ.5) CALL DUMP
C   IF (OPTION.EQ.1) GO TO 310
C
C   PRINT HEADING
C   K = 0
C   WRITE (6, 1) N, M
1  FORMAT (//43X, '      ***** SOLUTION *****'//43X, 'NETWORK: ', I5
&      , ' NODES, ', I5, ' EDGES')
C
C   CHECK INFEASIBILITY
C   IF (DELTA.GE.INF) WRITE (6, 2)
2  FORMAT (//43X, '      = PROBLEM INFEASIBLE =')
C
C   PRINT DUAL OBJECTIVE VALUE
C   WRITE (6, 3) VALUE
3  FORMAT (//43X, '      OBJECTIVE VALUE: ', F9.2)
C
C   FOR ALL NODES
C   DO 70 IND = NO, N1
C       NODE = LIST (IND)
C       MPRICE (NODE) = PRICE (NODE)
C
C   CHECK CURRENCY
C   IF (PSEUDO (NODE).NE.NODE) GO TO 40
C   GET MATES (SOLUTION EDGES) OF CURRENT NODES
C   PT = -MATE (NODE)

```

```

IF (PT) 30, 70, 10
C      SEVERAL MATES
10    EDGE = MATE(PT)
      IF (EDGE.GT.M2) GO TO 20
          K      = K + 1
          STACK(K) = EDGE - L2
20    PT = LINK(PT)
      IF (PT.NE.NULL) GO TO 10
      GO TO 70
C      EXACTLY ONE MATE
30    IF (-PT.GT.M2) GO TO 70
          K      = K + 1
          STACK(K) = -PT - L2
          GO TO 70
C      NONCURRENT NODES
40    IF (NODE.NE.APEX(PSEUDO(NODE))) GO TO 50
C      IF APEX --> GET MATE (SOLUTION EDGE)
      EDGE = MATE(NODE)
      IF (EDGE.EQ.NULL.OR.EDGE.GT.M2) GO TO 50
          K      = K + 1
          STACK(K) = EDGE - L2
C      TRANSFORM NODE PRICES
50    IF (SET(NODE).EQ.SET2) GO TO 70
C      N< OK!
C      N= PRICE DECREASED ON EVERY PSEUDO
C      N> PRICE DOUBLE DECREASED ON EVERY PSEUDO
      BLSM = NODE
      SIGN = 1
      IF (SET(NODE).EQ.SET3) SIGN = 2
60    BLSM = LABLO(BLSM)
      IF (BLSM.LE.N2) GO TO 70
      NPRICE(NODE) = NPRICE(NODE) - SIGN * PRICE(BLSM)
      GO TO 60
70    CONTINUE
C
C      PRINT NODES AND EDGES
      WRITE(6,4) (LIST(I), SET(LIST(I)), NPRICE(LIST(I)), I=N0, M1)
4    FORMAT(//' LIST OF NODES (NODE,SET,PRICE) :'/5(I15,I2,F7.1))
      DO 80 IND = M0, M1
80   EDLIST(IND) = LIST(IND) - L2
      WRITE(6,5) (EDLIST(I), HEAD(LIST(I)), HEAD(M9-LIST(I)), PRICE(LIS
&      T(I)), I=M0, M1)
5    FORMAT(//' LIST OF EDGES (EDGE,NODE1,NODE2,COST) :'/5(I11,2I3,
&      F7.1))
C
C      CHECK EXISTENCE OF PSEUDONODES
      IF (L0.GT.L1) GO TO 200
      WRITE(6,6)
6    FORMAT(//' PSEUDO PRICE SIMPLE BLOSSOM')
C
C      CORRECT SOLUTION EDGES ON EVERY SIMPLE BLOSSOM
      DO 150 IND = L0, L1
          PNODE = LIST(IND)
          BASE = MATE(PNODE)
C      GET SIMPLE BLOSSOM
          I      = 0
          NODE  = BASE
110   I      = I + 1
          PACK(I) = NODE
          NODE  = HEAD(NODE)

```

```

IF (NODE.NE.BASE) GO TO 110
WRITE (6,7) PNODE, PRICE(PNODE), (PACK(J), J=1,I)
7  FORMAT (I7,F7.1,2X,20I5/(19X,20I5))
C  CHECK TYPE C SIMPLE BLOSSOM
   APICE = APEX (NODE)
   IF (BASE.GT.N2.OR.SET (APICE).NE.SET3.OR.MATE (APICE).NE.NU
8     LL.OR.NPRICE (APICE).GT.TOLER) GO TO 130
C  TYPE C SIMPLE BLOSSOM
C  GET SOLUTION EDGES
   BLSM1 = BASE
120  BLSM2 = HEAD (ELSM1)
      K = K + 1
      STACK (K) = PRED (BLSM1)
      IF (STACK (K).GT.M2) STACK (K) = M9 - STACK (K)
      STACK (K) = STACK (K) - L2
      IF (BLSM2.EQ.BASE) GO TO 150
      BLSM1 = HEAD (BLSM2)
      GO TO 120
C  TYPE A, TYPE B, TYPE D, OR ROOTED SIMPLE BLOSSOM
C  GET SOLUTION EDGES
130  BLSM1 = HEAD (BASE)
140  K = K + 1
      STACK (K) = PRED (BLSM1)
      IF (STACK (K).GT.M2) STACK (K) = M9 - STACK (K)
      STACK (K) = STACK (K) - L2
      BLSM1 = HEAD (HEAD (BLSM1))
      IF (BLSM1.NE.BASE) GO TO 140
150  CONTINUE
C
C  PRINT SOLUTION EDGES
200  WRITE (6,8) (STACK (I), I=1,K)
      8  FORMAT (// ' SOLUTION EDGES: ' / (10I12))
C
C  PRINT STATISTICS
300  IF (OPTION.EQ.2) RETURN
      WRITE (6,11)
11  FORMAT (///45X, ' = STATISTICS = ')
310  CALL TIME (1,0,TIME3)
      TIME3 = TIME3 - TIME2
      TIME2 = TIME2 - TIME1
      WRITE (6,12) DUAL$, AUGM$, RMAT$, LABL$, NODE$, EDGES$, SHRKS$, UNSH$, LI
8     NK$, TIME1, TIME2, TIME3
12  FORMAT (//
119X, I5, ' DUAL SOL. CHANGES ', I5, ' PATHS REMATCHED ',
2    I5, ' NODES REMATED ', /19X, I5, ' NODES LABELLED ',
3    I5, ' NODES SCANNED ', I5, ' EDGES SCANNED ', /
419X, I5, ' SHRINKINGS ', I5, ' UNSHRINKINGS ',
5    I5, ' LINK/MATES USED ', /19X, I5, ' CPU MSEC (INPUT) ',
5    I5, ' CPU MSEC (SOLUTION) ', I5, ' CPU MSEC (OUTPUT) ')
      IF (OPTION.EQ.5) CALL DUMP
      RETURN
      END

```





60 CONTINUE

C

C

DUMP PSEUDO INFORMATION

IF (L0.GT.L1) GO TO 200

WRITE(6,3)

3 FORMAT(//' PNODE BASE LABEL PSEUD HEAD PRED TREE APEX B  
8 CI PRICE MPRICE LINK(S)')

DO 120 IND = L0, L1

NODE = LIST(IND)

PT = LINK(NODE)

PACK(1) = 0

K = 0

110 IF (PT.EQ.NULL) GO TO 120

K = K + 1

PACK(K) = MATE(PT)

PT = LINK(PT)

GO TO 110

120 WRITE(6,2) NODE, MATE(NODE), LABLO(NODE), PSEUDO(NODE), HEAD(  
8 NODE), PRED(NODE), TREE(NODE), APEX(NODE), BCI(NODE)

C

C

DUMP EDGE INFORMATION

200 IF (J.EQ.0) RETURN

WRITE(6,4) (STACK(I), HEAD(STACK(I)), HEAD(M9-STACK(I)), I=1, J)

4 FORMAT(//' MATES : (EDGE, NODE1, NODE2) '/5(I9, 2I4))

RETURN

END

SUBROUTINE ERROR (I)

C  
C  
C  
C  
C  
C  
C  
C  
C  
C  
C

MINIMUM COST 1-MATCHING/COVERING PROBLEM  
THE UNIVERSITY OF MICHIGAN  
WRITTEN BY CLOVIS PERIN  
AUGUST 1980

/PRINT MESSAGE OF ERROR/

IMPLICIT INTEGER (A-Z)  
REAL VALUE, DELTA, MCOST, TOLER, INF, PRICE, MPRICE  
LOGICAL TRACE, PRINT  
COMMON /NODES/ SET (100), LABLO (150), PSEUDO (150), APEX (150), PRED (150), TREE (150), MPRICE (150), BCI (150)  
COMMON /EDGES/ HEAD (6150), PRICE (3150), LIST (3150), INDEX (3150)  
COMMON /LISTS/ NSCAN (101), INSP (3100), NORIG (100), LINK (1500), MAT E (1500)  
COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK  
COMMON /PARMS/ N, NO, N1, N2, N9, L, LO, L1, L2, M, MO, M1, M2, M9, P0, P2  
COMMON /VARBS/ VALUE, DELTA, MCOST, OPTION, TRACE, PRINT  
COMMON /STATS/ ROOT\$, NODE\$, EDGE\$, SHRK\$, UNSH\$, AUGM\$, DUAL\$, RMAT\$, LABL\$, LINK\$, TIME1, TIME2, TIME3  
COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET3, NULL  
COMMON /PRNTS/ STACK (3000), PACK (100)

C  
C  
C

GO TO (10, 20, 30, 40, 50, 60, 70, 80, 90, 100), I  
GO TO 110

C

10 WRITE (6, 1)  
1 FORMAT (' \*\*\*\*\* UNEXPECTED END OF FILE \*\*\*\*\*')  
STOP

C

20 WRITE (6, 2)  
2 FORMAT (' \*\*\*\*\* ILLEGAL PARAMETER \*\*\*\*\*')  
STOP

C

30 WRITE (6, 3)  
3 FORMAT (' \*\*\*\*\* NODE STORAGE SPACE EXCEEDED \*\*\*\*\*')  
STOP

C

40 WRITE (6, 4)  
4 FORMAT (' \*\*\*\*\* EDGE STORAGE SPACE EXCEEDED \*\*\*\*\*')  
STOP

C

50 WRITE (6, 5)  
5 FORMAT (' \*\*\*\*\* LINK/MATE STORAGE SPACE EXCEEDED \*\*\*\*\*')  
STOP

C

60 WRITE (6, 6)  
6 FORMAT (' \*\*\*\*\* ILLEGAL NODE \*\*\*\*\*')  
STOP

C

70 WRITE (6, 7)  
7 FORMAT (' \*\*\*\*\* ILLEGAL EDGE \*\*\*\*\*')  
STOP

```
      STOP
C
80 WRITE(6,8)
  8 FORMAT(' ***** ILLEGAL NODE SET *****')
  STOP
C
90 WRITE(6,9)
  9 FORMAT(' ***** ILLEGAL COST *****')
  STOP
C
100 CONTINUE
110 WRITE(6,12) I
  12 FORMAT(' ***** ERROR NO. ',I3,' *****')
  STOP
  END
```



```

OPTION = CPTION + N2
IF (OPTION.EQ.-N2) TRACE = .TRUE.
IF (0.GT.OPTION.AND.OPTION.GE.-N2) PRINT = .TRUE.

```

```

C
C   EDGE COST CHANGES
100 IF (K1.LT.1) GO TO 200
    DO 110 K = 1, K1
    IF (TRACE) CALL DUMP
C   GET EDGE AND NEW COST
    READ(5,FREE,END=999) EDGE, NPRICE
    IF (TRACE.OR.OPTION.EQ.4) WRITE(6,4) EDGE, NPRICE
4   FORMAT(' EDGE COST CHANGE :      ',I5,F7.1)
    STACK(K) = EDGE
    EDGE = EDGE + L2
    IF (EDGE.LT.M0.OR.EDGE.GT.M2.OR.INDEX(EDGE).EQ.NULL) CALL ERRO
&                                     R(7)
    IF (NPRICE.GT.INF.OR.NPRICE.LT.-INF) CALL ERRO
&                                     R(9)
    IF ( NPRICE.EQ.PRICE(EDGE) ) GO TO 110
    EDGE1 = EDGE
    EDGE2 = M9 - EDGE
    NODE2 = HEAD(EDGE1)
    NODE1 = HEAD(EDGE2)
    PRICE(EDGE) = NPRICE
    COST = PRICE(EDGE) - PRICE(NODE1) - PRICE(NODE2)
    CALL NMATCH(EDGE)
    IF (PSEUDO(NODE1).NE.PSEUDO(NODE2).AND.COST.GE.-TOLER) GO TO 1
&                                     10
C   EDGE IS OUT-OF-KILTER
C   TRY TO DECREASE PRICE(NODE1)
    CALL OPEN(NODE1)
    CALL DECRSE(EDGE1,COST)
    IF (COST.GT.-TOLER) GO TO 110
C   TRY TO DECREAS PRICE(NODE2)
    CALL OPEN(NODE2)
    CALL DECRSE(EDGE2,COST)
    IF (COST.GT.-TOLER) GO TO 110
C   EDGE OF A-
    CALL MATCH(EDGE)
    VALUE = VALUE + COST
110 CONTINUE
    IF (PRINT) WRITE(6,5) (STACK(K),HEAD(STACK(K)),HEAD(M9-STACK(K)
&                                     )),PRICE(STACK(K)),K=1,K1)
5   FORMAT('//' EDGE COST CHANGE (EDGE,NODE1,NODE2,NEWCOST) : '/'5 (I9,
&                                     2I3,F6.1))
C
C   SUBSET CHANGE
200 IF (K2.LT.1) GO TO 300
    DO 250 K = 1, K2
    IF (TRACE) CALL DUMP
C   GET NODE AND NEWSET
    READ(5,FREE,END=999) NODE, NSET
    IF (TRACE.OR.OPTION.EQ.4) WRITE(6,6) NODE, NSET
6   FORMAT(' CHANGE NODE SET :      ',2I5)
    STACK(K) = NODE
    IF (NSET.LT.SET0.OR.NSET.GT.SET3) CALL ERROR(8)
    IF (NODE.LT.N0.OR.NODE.GT.N2.OR.INDEX(NODE).EQ.NULL) CALL ERRO
&                                     R(6)
    IF ( NSET.EQ.SET(NODE) ) GO TO 250
C   UPDATE SET(NODE0, MPRICE(NODE), AND BCI(NODE)

```

```

J          = SET (NODE) + 1
SET(NODE) = NSET
MPRICE(NODE) = 0
IF (SET(NODE).EQ.SET1) MPRICE(NODE) = INF
IF (SET(NODE).EQ.SET3) MPRICE(NODE) = INF
BCI(NODE) = NODE
IF (SET(NODE).EQ.SET1) BCI(NODE) = L2
PNODE = PSEUDO(NODE)
C FORCE NODE TO BECOME CURRENT
IF (PNODE.NE.NCDE) CALL OPEN(NODE)
C BRANCH IN ACCORDANCE WITH OLD SUBSET
GO TO (210,220,230,240), J
C
C OLD SUBSET WAS SET0
C SET3 IS OK
210 IF ( NSET.EQ.SET3 ) GO TO 250
C SET1, SET2 : UNMATCH ALL MATES OF NODE
212 PT = MATE(NODE)
IF (PT) 214, 218, 216
214 EDGE = MATE(-PT)
CALL NMATCH(EDGE)
GO TO 212
216 CALL NMATCH(PT)
C UPDATE PRICE(NODE)
218 VALUE = VALUE- PRICE(NODE)
PRICE(NODE) = -INF/2
VALUE = VALUE + PRICE(NODE)
GO TO 250
C
C OLD SUBSET WAS SET1
220 IF (NSET.NE.SET0) GO TO 222
C NEW SET IS SET0
CALL OPEN(NODE)
IF (PRICE(NODE).LT.-TOLER) CALL INCRSE(NODE)
VALUE = VALUE - PRICE(NODE)
PRICE(NODE) = 0
GO TO 250
C
C NEW SET IS SET2
222 IF (NSET.NE.SET2) GO TO 224
IF (PRICE(NODE).LE.+TOLER) GO TO 250
CALL OPEN(NODE)
VALUE = VALUE - PRICE(NODE)
PRICE(NODE) = -INF/2
VALUE = VALUE + PRICE(NODE)
GO TO 250
C
C NEW SET IS SET3
224 IF (PRICE(NODE).LT.-TOLER) CALL INCRSE(NODE)
GO TO 250
C
C OLD SET WAS SET2
C NEW SET IS SET1 OK
230 IF ( NSET.EQ.SET1 ) GO TO 250
C NEW SET IS SET0 OR SET3
CALL OPEN(NODE)
CALL INCRSE(NODE)
GO TO 250
C
C OLD SET WAS SET3

```

```

240 IF (SET(NODE).EQ.SET0.AND.PRICE(NODE).LE.TOLER) GO TO 250
    IF (SET(NODE).EQ.SET1.AND.MATE(NODE).GE.NULL) GO TO 250
    IF (SET(NODE).EQ.SET2.AND.PRICE(NODE).LE.TOLER.AND.MATE(NODE).
&      GE.NULL) GC TO 250
C
C UNMATCH ALL MATES OF NODE
242 PT = MATE(NODE)
    IF (PT) 244, 248, 246
244     EDGE = MATE(-PT)
        CALL NMATCH(EDGE)
        GO TO 242
246 CALL NMATCH(PT)
C UPDATE PRICE(NODE)
248 VALUE = VALUE - PRICE(NODE)
    PRICE(NODE) = -INF/2
    IF ( SET(NODE).EQ.SET0 ) PRICE(NODE) = 0
    VALUE = VALUE + PRICE(NODE)
C
250 CONTINUE
    IF (PRINT) WRITE(6,7) (STACK(K),SET(STACK(K)),K=1,K2)
    7 FORMAT(//' NODE SET CHANGE (NODE,NEWSET) :'/10(I10,I2))
C
C EDGE ELIMINATION
300 IF (K3.LT.1) GO TO 400
    DO 310 K = 1, K3
    IF (TRACE) CALL DUMP
C GET EDGE
    READ(5,FREE,END=999) EDGE
    IF (TRACE.OR.OPTION.EQ.4) WRITE(6,8) EDGE
    8 FORMAT(' EXCLUDE EDGE :           ',I5)
    STACK(K) = EDGE
    EDGE = EDGE + L2
    IF (EDGE.LT.M0.OR.EDGE.GT.M2.OR.INDEX(EDGE).EQ.NULL) CALL ERRO
&      R(7)
C UPDATE POINTERS
    CALL NMATCH(EDGE)
    IND = INDEX(EDGE)
    LIST(IND) = NULL
    INDEX(EDGE) = NULL
    M = M - 1
C CHECK LINK/MATE POSITIONS
    DO 310 PT = P0, P2
        IF (MATE(PT).EQ.EDGE.CR.MATE(PT).EQ.M9-EDGE) MATE(PT) = N
&      ULL
310 CONTINUE
    IF (PRINT) WRITE(6,9) (STACK(K),HEAD(STACK(K)),HEAD(M9-STACK(K)
&      )),PRICE(STACK(K)),K=1,K3)
    9 FORMAT(//' EDGE EXCLUSION (EDGE,NODE1,NODE2,COST) :'/5(I9,2I3,F
&      6.1))
C
C NODE ELIMINATION
400 IF (K4.LT.1) GO TO 500
    DO 420 K = 1, K4
    IF (TRACE) CALL DUMP
C GET NODE
    READ(5,FREE,END=999) NODE
    IF (TRACE.OR.OPTION.EQ.4) WRITE(6,11) NODE
    11 FORMAT(' EXCLUDE NODE :           ',I5)
    STACK(K) = NODE
    IF (NODE.LT.N0.OR.NODE.GT.N2.OR.INDEX(NODE).EQ.NULL) CALL ERRO

```



```

E          R(6)
CALL OPEN(NODE)
DO 410 IND = M0, M1
    EDGE = LIST(IND)
    IF (HEAD(EDGE).NE.NODE.AND.HEAD(M9-EDGE).NE.NODE) GO TO 4
E          10
    IF (TRACE) WRITE(6,8) EDGE
    CALL NMATCH(EDGE)
    M          = M - 1
    INDEX(EDGE) = NULL
    LIST(IND)   = NULL
DO 410 PT = P0, P2
    IF (MATE(PT).EQ.EDGE.OR.MATE(PT).EQ.M9-EDGE) MATE(PT) = N
E          ULL
410 CONTINUE
    VALUE      = VALUE - PRICE(NODE)
    IND        = INDEX(NODE)
    LIST(IND)  = NULL
    INDEX(NODE) = NULL
420 N          = N - 1
    IF (PRINT) WRITE(6,12) (STACK(K),SET(STACK(K)),K=1,K4)
12 FORMAT(// ' NODE EXCLUSION (NODE,SET) : ' /10(I10,I2))
C
C NODE INTRODUCTION
500 IF (K5.LT.1) GO TO 600
    DO 510 K = 1, K5
    IF (TRACE) CALL DUMP
C GET NODE AND SUBSET
    READ(5,FREE,END=999) NODE, NSET
    IF (TRACE.OR.OPTION.EQ.4) WRITE(6,13) NODE, NSET
13 FORMAT(' INCLUDE NODE :          ',I4,I2)
    STACK(K) = NODE
    IF (NODE.LT.N0.OR.NODE.GT.N2.OR.INDEX(NODE).NE.NULL) CALL ERRO
E          R(6)
    IF (NSET.LT.0.OR.NSET.GT.3) CALL ERROR(3)
    IF (N1.GE.N2) CALL ERROR(2)
C UPDATE POINTERS
    N          = N + 1
    N1         = N1 + 1
    INDEX(NODE) = N1
    LIST(N1)   = NODE
    SET(NODE)  = NSET
    LABLO(NODE) = UNLABD
    PSEUDO(NODE) = NODE
    PRED(NODE) = NULL
    APEX(NODE) = NODE
    MPRICE(NODE) = 0
    IF (NSET.EQ.SET1) MPRICE(NODE) = INF
    MATE(NODE) = NULL
    BCI(NODE) = NODE
    IF (NSET.EQ.SET1) BCI(NODE) = L2
    PRICE(NODE) = 0
    IF (SET(NODE).EQ.SET1.OR.SET(NODE).EQ.SET2) PRICE(NODE) = -INF/2
    VALUE      = VALUE + PRICE(NODE)
    LINK(NODE) = NULL
510 HEAD(NODE) = NULL
    IF (PRINT) WRITE(6,14) (STACK(K),SET(STACK(K)),K=1,K5)
14 FORMAT(// ' NODE INCLUSION (NODE,SET) : ' /10(I10,I2))
C
C EDGE INTRODUCTION

```

```

600 IF (K6.LT.1) GO TO 700
DO 610 K = 1, K6
IF (TRACE) CALL DUMP
C GET EDGE, NODES, AND EDGE COST
READ(5,FREE, END=999) EDGE, NODE1, NODE2, NPRICE
IF (TRACE.OR.OPTION.EQ.4) WRITE(6,15) EDGE,NODE1,NODE2,NPRICE
15 FORMAT(' INCLUDE EDGE :           ',3I5,F7.1)
STACK(K) = EDGE
EDGE = EDGE + L2
IF (EDGE .LT.M0.OR.EDGE .GT.M2.OR.INDEX(EDGE) .NE.NULL) CALL E
& RROR(7)
IF (NODE1.LT.N0.OR.NODE1.GT.N2.OR.INDEX(NODE1) .EQ.NULL) CALL E
& RROR(6)
IF (NODE2.LT.N0.OR.NODE2.GT.N2.OR.INDEX(NODE2) .EQ.NULL) CALL E
& RROR(6)
IF (NPRICE.GT.INF.OR.NPRICE.LT.-INF) CALL E
& RROR(9)
IF (M1.GE.M2) CALL ERROR(4)
610 CALL INCLUD(EDGE,NODE1,NODE2,NPRICE)
IF (PRINT) WRITE(6,16) (STACK(K),HEAD(STACK(K)),HEAD(M9-STACK(
& K)),PRICE(STACK(K)),K=1,K6)
16 FORMAT(//' EDGE INCLUSION (EDGE,NODE1,NODE2,COST): '/5(I9,2I3,F
& 6.1))
C
C RE-INITIALIZE VARIABLES
700 ROOT$ = 0
NODE$ = 0
EDGE$ = 0
SHRK$ = 0
UNSH$ = 0
AUGM$ = 0
DUAL$ = 0
LABL$ = 0
RMAT$ = 0
HSCAN = 0
TSCAN = 0
C
C UPDATE LIST OF VALID NODE/EDGES
C PLANT ALTERNATING FOREST
I = NO
N1 = N
DO 720 IND = NO, N1
710 NODE = LIST(I)
I = I + 1
IF (NODE.EQ.NULL) GO TO 710
LIST(IND) = NODE
INDEX(NODE) = IND
PNODE = PSEUDO(NODE)
LABLO(PNODE) = UNLABL
HEAD(PNODE) = NULL
PRED(PNODE) = NULL
TREE(PNODE) = NULL
IF (MATE(NODE) .NE.NULL) GO TO 720
IF (SET(NODE) .EQ.SET0) GO TO 720
IF (SET(NODE) .EQ.SET2.AND.PRICE(NODE) .GE.-TOLER) GO TO 72
& 0
IF (NODE.NE.APEX(PNODE)) GO TO 720
IF (SET(NODE) .EQ.SET3.AND.PRICE(NODE) .GE.MPRICE(NODE)-TOL
& ER) GO TO 720
C EXPOSED NODE

```

```

LABLO (PNODE) = OUTER
HEAD (PNODE) = NODE
PRED (PNODE) = NULL
LINK (PNODE) = NULL
ROOT$      = ROOT$ + 1
HSCAN     = HSCAN + 1
IF (HSCAN.GT.N9) HSCAN = 1
NSCAN (HSCAN) = PNODE
720 CONTINUE
I = M0
M1 = M0 + M - 1
DO 740 IND = M0, M1
730     EDGE1 = LIST (I)
        I     = I + 1
        IF (EDGE1.EQ.NULL) GO TO 730
        LIST (IND) = EDGE1
740     INDEX (EDGE1) = IND
IF (TRACE) CALL DUMP
LABL$ = ROOT$
IF (PRINT) WRITE (6,17) ROOT$
17 FORMAT (/ ' ALTERNATING FOREST PLANTED WITH',I4,' ROOTS')
DELTA = 0
IF (OPTION.EQ.6) CALL DUMP
CALL TIME (1,0,TIME1)
IF (OPTION.NE.-ROOTS) RETURN
TRACE = .TRUE.
CALL DUMP
RETURN
C
C ERROR MESSAGE
999 CALL ERROR (1)
1000 STOP
END

```

SUBROUTINE INCIUD(EDGE1,NODE1,NODE2,NPRICE)

```

C - - - - -
C
C
C
C   MINIMUM COST 1-MATCHING/COVERING PROBLEM
C   THE UNIVERSITY OF MICHIGAN
C   WRITTEN BY CLOVIS PERIN
C   AUGUST 1980
C
C
C   /INCLUDE EDGE1 IN THE NETWORK/
C   NODE1 -- EDGE1 --> NODE2
C   PRICE(EDGE1) = NPRICE
C
C   IMPLICIT INTEGER (A-Z)
C   REAL VALUE,DELTA,MCOST,TOLFR,INF,PRICE,MPRICE
C   LOGICAL TRACE,PRINT
C   COMMON /NODES/ SET(100),LABLO(150),PSEUDO(150),APEX(150),PRED(
&           150),TREE(150),MPRICE(150),BCI(150)
C   COMMON /EDGES/ HEAD(6150),PRICE(3150),LIST(3150),INDEX(3150)
C   COMMON /LISTS/ NSCAN(101),INSP(3100),NORIG(100),LINK(1500),MAT
&           E(1500)
C   COMMON /PNTRS/ HSCAN,TSCAN,HINSP,HORIG,HINK
C   COMMON /PARMS/ N,N0,N1,N2,N9,L,L0,L1,L2,M,M0,M1,M2,M9,P0,P2
C   COMMON /VARBS/ VALUE,DELTA,MCOST,OPTION,TRACE,PRINT
C   COMMON /STATS/ ROOT$,NODE$,EDGE$,SHRK$,UNSH$,AUGM$,DUAL$,RMAT$
&           ,LABL$,LINK$,TIME1,TIME2,TIME3
C   COMMON /CONTS/ TOLER,INF,OUTER,INNER,UNLABD,SET0,SET1,SET2,SET
&           3,NULL
C   COMMON /PRNTS/ STACK(3000),PACK(100)
C   LOGICAL FREE(1) /'*/
C   REAL COST, NPRICE
C
C
C   . . . . .
C
C   IF (TRACE) WRITE(6,1) EDGE1, NODE1, NODE2, NPRICE
1  FORMAT(' #TRACE# INCLUD (' ,3(I4,','),F6.1,') ')
C
C   UPDATE VARIABLES
C   M           = M + 1
C   M1          = M1 + 1
C   LIST(M1)    = EDGE1
C   INDEX(EDGE1) = M1
C   EDGE2       = M9 - EDGE1
C   HEAD(EDGE1) = NODE1
C   HEAD(EDGE2) = NODE2
C   PRICE(EDGE1) = NPRICE
C   COST        = PRICE(EDGE1) - PRICE(NODE1) - PRICE(NODE2)
C   IF (COST.GE.0) RETURN
C
C   EDGE IS OUT-OF-KILTER
C   TRY TO DECREASE THE PRICE OF NODE1
C   CALL OPEN(NODE1)
C   CALL DECRSE(EDGE1,COST)
C   IF (COST.GE.0) RETURN
C   TRY TO DECREASE THE PRICE OF NODE2
C   CALL OPEN(NODE2)
C   CALL DECRSE(EDGE2,COST)
C   IF (COST.LT.-TOLER) CALL MATCH(EDGE1)
C   RETURN

```

C  
999 CALL ERROR(3)  
RETURN  
END



SUBROUTINE INCRSE(NODE)

C  
C  
C  
C  
C  
C  
C  
C  
C  
C  
C  
C

-----  
MINIMUM COST 1-MATCHING/COVERING PROBLEM  
THE UNIVERSITY OF MICHIGAN  
WRITTEN BY CLOVIS PERIN  
AUGUST 1980

/INCREASE PRICE(NODE) UP TO ZERO/  
CHECK EDGES INCIDENT WITH NODE

IMPLICIT INTEGER (A-Z)  
REAL VALUE, DELTA, MCOST, TOLER, INF, PRICE, MPRICE  
LOGICAL TRACE, PRINT  
COMMON /NODES/ SET (100), LABLO (150), PSEUDO (150), APEX (150), PRED (150), TREE (150), MPRICE (150), BCI (150)  
COMMON /EDGES/ HEAD (6150), PRICE (3150), LIST (3150), INDEX (3150)  
COMMON /LISTS/ NSCAN (101), INSP (3100), NORIG (100), LINK (1500), MAT E (1500)  
COMMON /PNTRS/ HSCAN, TSCAN, HINSP, HORIG, HINK  
COMMON /PARMS/ N, NO, N1, N2, N9, L, L0, L1, L2, M, M0, M1, M2, M9, P0, P2  
COMMON /VARBS/ VALUE, DELTA, MCOST, OPTION, TRACE, PRINT  
COMMON /STATS/ ROOT\$, NODE\$, EDGES\$, SHRKS\$, UNSH\$, AUGM\$, DUAL\$, RNATS\$, LABL\$, LINK\$, TIME1, TIME2, TIME3  
COMMON /CONTS/ TOLER, INF, OUTER, INNER, UNLABD, SET0, SET1, SET2, SET3, NULL  
COMMON /PRNTS/ STACK (3000), PACK (100)  
REAL COST

C  
C  
C

IF (TRACE) WRITE (6, 1) NODE  
1 FORMAT (' #TRACE# INCRSE (' , I4, ' ) ')

C  
C

INCREASE PRICE(NODE) TO ZERO  
VALUE = VALUE - PRICE(NODE)  
PRICE(NODE) = 0

C  
C

CHECK EDGES INCIDENT WITH NODE  
DO 10 IND = M0, M1  
EDGE = LIST (IND)  
NODE1 = HEAD (M9-EDGE)  
NODE2 = HEAD (EDGE)  
IF (NODE.NE.NODE1.AND.NODE.NE.NODE2) GO TO 10  
COST = PRICE (EDGE) - PRICE (NODE1) - PRICE (NODE2)  
IF (COST.GE.-TOLER) GO TO 10  
NEGATIVE COST  
NODE0 = NODE1 + NODE2 - NODE  
CALL OPEN (NODE0)  
CALL DECRSE (EDGE, COST)  
IF (COST.LT.-TOLER) CALL MATCH (EDGE)

C

10 CONTINUE  
RETURN  
END

```

SUBROUTINE OPEN(NODE)
C -----
C
C
C   MINIMUM COST 1-MATCHING/COVERING PROBLEM
C   THE UNIVERSITY OF MICHIGAN
C   WRITTEN BY CLOVIS PERIN
C   AUGUST 1980
C
C   /OPEN PSEUDO(NODE) /
C   UNSHRINK PSEUDO(NODE) OR UNMATCH MATE(NODE)
C
C   IMPLICIT INTEGER (A-Z)
C   REAL VALUE,DELTA,MCOST,TOLER,INF,PRICE,MPRICE
C   LOGICAL TRACE,PRINT
C   COMMON /NODES/ SET(100),LABLO(150),PSEUDO(150),APEX(150),PRED(
&      150),TREE(150),MPRICE(150),BCI(150)
C   COMMON /EDGES/ HEAD(6150),PRICE(3150),LIST(3150),INDEX(3150)
C   COMMON /LISTS/ NSCAN(101),INSP(3100),NORIG(100),LINK(1500),MAT
&      E(1500)
C   COMMON /PNTRS/ HSCAN,TSCAN,HINSP,HORIG,HINK
C   COMMON /PARMS/ N,NO,N1,N2,N9,L,LO,L1,L2,M,MO,M1,M2,M9,P0,P2
C   COMMON /VARBS/ VALUE,DELTA,MCOST,OPTION,TRACE,PRINT
C   COMMON /STATS/ ROOT$,NODE$,EDGE$,SHRK$,UNSH$,AUGM$,DUAL$,RMT$
&      ,LABL$,LINK$,TIME1,TIME2,TIME3
C   COMMON /CONTS/ TOLER,INF,OUTER,INNER,UNLABD,SET0,SET1,SET2,SET
&      3,NULL
C   COMMON /PRNTS/ STACK(3000),PACK(100)
C
C   . . . . .
C
C   IF (TRACE) WRITE(6,1) NODE
C   1 FORMAT(' #TRACE# OPEN (' ,I4,') ')
C
C   BRANCH IF NODE IS NONCURRENT
C   IF (PSEUDO(NODE).NE.NODE) GO TO 10
C   NODE IS CURRENT
C   IF (SET(NODE).EQ.SET0) RETURN
C   IF (SET(NODE).EQ.SET3.AND.PRICE(NODE).EQ.0) RETURN
C   UNMATCH MATE(NODE)
C   EDGE = MATE(NODE)
C   CALL NMATCH(EDGE)
C   RETURN
C
C   UNSHRINK PSEUDO(NODE)
C   10 PNODE = PSEUDO(NODE)
C   EDGE = MATE(APEX(PNODE))
C   UNMATCH MATE(NODE)
C   CALL NMATCH(EDGE)
C   IF (PNODE.EQ.NODE) RETURN
C   IF (TRACE) WRITE(6,2) PNODE
C   2 FORMAT(' #TRACE# OPEN PSEUDO : ' ,I4)
C
C   REDUCE NODE PRICES
C   HORIG = 0
C   CALL ORIG(PNODE)
C   DO 20 H = 1, HORIG
C   NODE2 = NORIG(H)
C   IF (NODE2.GT.N2) CALL ORIG(NODE2)
C   IF (NODE2.LE.N2) PRICE(NODE2) = PRICE(NODE2) - PRICE(PNOD

```



```

      &           E)
20      CONTINUE
      VALUE = VALUE - PRICE(PNODE)
C
C      UPDATE MATES, APICES, HEADS, AND PREDS
      BASE = MATE(PNODE)
      CALL STPSDO(BASE)
      LABLO(BASE) = UNLABD
      BLSM = HEAD(BASE)
30      EDGE1      = PRED(BLSM)
      EDGE2      = M9 - EDGE1
      NODE2      = HEAD(EDGE1)
      NODE1      = HEAD(EDGE2)
      MATE(NODE1) = EDGE1
      MATE(NODE2) = EDGE2
      BLSM1     = HEAD(BLSM)
      CALL STPSDO(BLSM)
      CALL STPSDO(BLSM1)
      LABLO(BLSM) = UNLABD
      LABLO(BLSM1) = UNLABD
      CALL STAPEX(NODE1)
      CALL STAPEX(NODE2)
      HEAD(BLSM) = 0
      PRED(BLSM) = 0
      BLSM      = HEAD(BLSM1)
      HEAD(BLSM1) = 0
      PRED(BLSM1) = 0
      IF ( BLSM .NE. BASE ) GO TO 30
      HEAD(BASE) = 0
      PRED(BASE) = 0
C
C      UPDATE PSEUDONODE LIST
      IND      = INDEX(PNODE)
      LIST(IND) = LIST(L1)
      LIST(L1) = PNODE
      INDEX(LIST(IND)) = IND
      INDEX(PNODE) = L1
      L1      = L1 - 1
      GO TO 10
C
      END

```

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