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Lecture 8: Item-to-item; Page Rank

SI583: Recommender Systems



Item-Item Collaborative Filtering

High-level approach:

For each item X find similar items Y,Z...

For user Joe, recommend items most similar to items Joe has already liked



Users-by-Items Matrix

$$R = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$



Normalize the Rows for User-User Algorithm

$$X_{iJ} = R_{iJ} - R_i$$

$$X = \begin{vmatrix} 1/3 & -2/3 & 1/3 \\ -2/3 & 1/3 & 1/3 \\ 1/3 & -2/3 & 1/3 \\ -2/3 & 1/3 & 1/3 \end{vmatrix} \qquad R = \begin{vmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{vmatrix}$$

$$R = \begin{vmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{vmatrix}$$



Normalize the Columns for Item-Item Algorithm

$$W_{jk} = R_{jk} - W_k$$

$$X = \begin{bmatrix} 0.5 & -5 & 0 \\ -5 & 0.5 & 0 \\ 0.5 & -5 & 0 \\ -5 & 0.5 & 0 \end{bmatrix}$$

$$X = \begin{vmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{vmatrix}$$



Alternative similarity measure for 0-1 entries: co-occurence

- When X has just 0 or 1 for each entry
- Instead of computing actual covariances from W, compute a similarity score based on count of co-occurrence in X
 - Co-occur(It1, It2) = 0
 - Co-occur(It1, It3) = 2

	1	U	1
= _	0	1	1
<u> </u>	1	0	1
	0	1	1



Generalization of co-occurrence similarity: Association Rules

- From a database of purchases, can find significant co-occurence rules, e.g., person who buys bread and butter => 90% chance of also buying milk
- It's possible to precompute these association rules (Agarwal et al)



User-User vs. Item-Item

- Compute pairwise correlations between users
- Compute pairwise correlations between items

$$X X^{T}$$

 $W^{\mathcal{T}}W$



Computational Complexity

- With n items, m users,
 - user-user algorithm (unoptimized): about m²n operations
 - item-item algorithm (unoptimized): about mn² operations
- #items may be < #users</p>
- item-item similarities may be stable over long periods of time => batch computing leads to less inaccuracy



Predicted Scores for Target Item

- User-user
 - Weighted average of other user's ratings of this item
 - Weights taken from user-user similarities
- Item-item
 - Weighted average of this user's ratings of other items
 - Weights taken from item-item similarities



Finding Items from Items

- Item-item algorithm
 - Single starting item
 - Find other items with highest correlation
 - Starting from a group of items
 - Union of results for each item
 - (Why are association rules better than the itemitem similarity matrix?)
- User-user algorithm
 - **-??**



Finding Users from Users

- User-user algorithm
 - Find other users with highest correlation
- Item-item algorithm
 - **-??**



Web search as a recommender

- Use links between pages as implicit "ratings"
- No separate categories of users, items
 - can't easily use user-user algorithm, etc.
- How are the "best" pages for a query recommended?



Model

- Page is a node
- html link defines a directional link in the graph
- Terminology
 - If A has an html to B
 - A has an outgoing link to B
 - B has an incoming link from A



PageRank

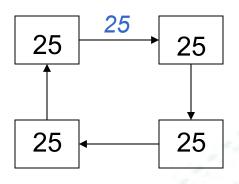
- Google's big original idea [Brin &Page, 1998]
- Idea: ranking is based on "random web surfer":
 - start from any page at random
 - pick a random link from the page, and follow it
 - repeat!
 - ultimately, this process will converge to a <u>stable distribution</u> over pages (with some tricks...)
 - most likely page in this stable distribution is ranked highest

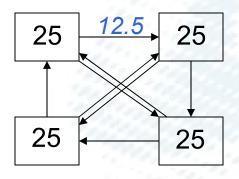
Strong points:

- Pages linked to by many pages tend to be ranked higher (not always)
- A link ("vote") from a highly-ranked page carries more weight
- Relatively hard to manipulate



PageRank, examples



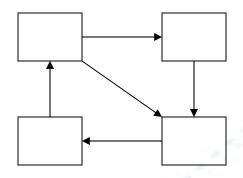


Final distribution properties:

- (a) Total weight = 100%
- (b) Weight of node is divided among outgoing links.
- (c) Weight of node is sum of incoming link weights.



PageRank, examples



Final distribution properties:

- (a) Total weight = 100%
- (b) Weight of node is divided among outgoing links.
- (c) Weight of node is some of incoming links



PageRank, mathematically

- Let the stable probabilities be x_i for page i, $x_i >= 0$
- For each i,j, define a_{ii} as
 - If j links to i, $a_{ij} = (1/number of links of j)$
 - If j does not link to i, $a_{ii} = 0$
- Form A =square matrix of a_{ii} for all i, j.
- Then, the PageRank probabilities satisfy

$$Ax = x$$

x is the eigenvector of the link matrix, with eigenvalue
1

* May need to modify A slightly to ensure unique solution



Finding the PageRank eigenvector

- One approach: solve linear equation $(A-I)x = (0\ 0\ 0\ ..0\ 0)^T$
- Alternative "power method" is more efficient in practice:
 - Start with an arbitrary X
 - Compute A^x , A^2x , ... A^tx (t large)
 - A^tx is approximately proportional to the correct solution!



Aside: why the power method works (optional)

- Known: the link matrix A has
 - eigenvalue 1 for the correct eigenvector v*
 - all other eigenvalues λ have $|\lambda|$ <1
- Known: any x can be expressed as a sum of eigenvectors of A

$$\mathbf{x} = \mathbf{a}_0 \mathbf{v}^* + \mathbf{a}_1 \mathbf{v}_1 + \mathbf{a}_2 \mathbf{v}_2 + \dots$$

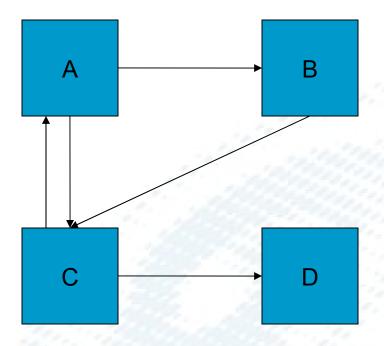
Multiplying by A t times,

$$A^{t}x = a_{0}v^{*} + a_{1}(\lambda_{1})^{t}v_{1} + a_{2}(\lambda_{2})^{t}v_{2} + ...$$

but $(\lambda_i)^t$ etc. are very close to 0 for large t



A Sample Graph



$$A = \begin{bmatrix} 0 & 0 & .5 & 0 \\ .5 & 0 & 0 & 0 \\ .5 & 1 & 0 & 0 \\ 0 & 0 & .5 & 0 \end{bmatrix}$$

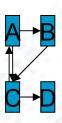


Handling Loops

- Let E be a set of "source" weight ranks
 - At each node, random surfer goes to nodes with probabilities in E
- Each node's final rank is a scaled multiple of
 - It's source rank PLUS
 - The sum of the rank on its backlinks
- Scale it such that the sum of final ranks is 1



A Sample Graph

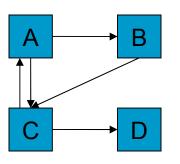


$$A = \begin{bmatrix} 0 & 0 & .5 & 0 \\ .5 & 0 & 0 & 0 \\ .5 & 1 & 0 & 0 \\ 0 & 0 & .5 & 0 \end{bmatrix}$$

$$E = \begin{array}{c} .1 \\ .1 \\ .1 \\ .1 \end{array}$$



Some Intuitions



- Will D's Rank be more or less than ¼?
- Will C's Rank be more or less than B's?
- How will A's Rank compare to D's?



Mathematical Expression

$$R' = c (AR + E)$$
 $||P'|| = |P_i| = 1$



Power Method Algorithm

Multiply by A, and then normalize so that the sum is 1

$$R_{i+1} = \frac{AR_i + E}{|AR_i + E|}$$



Before the First Iteration



- r1 .3
- r2 .1
- r3 .3
- r4 .1



First Iteration

AR+E

r1 .25

r2 .25

r3 .35

r4 .25

Normalize so sum is 1 (divide by 1.1)

r1 .22727273

r2 .22727273

r3 .31818182

r4 .22727273



Second Iteration

AR+E

- r1 .25909091
- r2 .21363636
- r3 .44090909
- r4 .25909091

Normalized (divide by 1.17)

- r1 .22093023
- r2 .18217054
- r3 .37596899
- r4 .22093023



Third Iteration

AR+E

- r1 .2879845
- r2 .21046512
- r3 .39263566
- r4 .2879845

Normalized (divide by 1.18)

- r1 .24424721
- r2 .17850099
- r3 .3330046
- r4 .24424721



What If More Weight in E?

Try (1 1 1 1) instead of (.1 .1 .1 .1)

```
r1 .23825503
```

- r2 .2360179
- r3 .28747204
- r4 .23825503

Try (10 10 10 10)

- r1 .24848512
- r2 .24845498
- r3 .25457478
- r4 .24848512



Personalized PageRank

- Pick E to be some sites that I like
 - My bookmarks
 - Links from my home page
- Rank flows more from these initial links than from other pages
 - But much of it may still flow to the popular sites, and from them to others that are not part of my initial set

